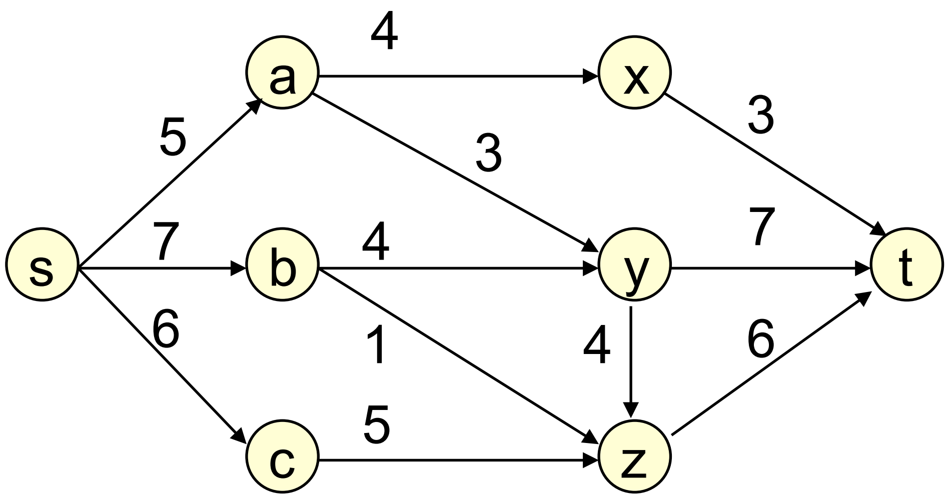
CSE 521 Algorithms

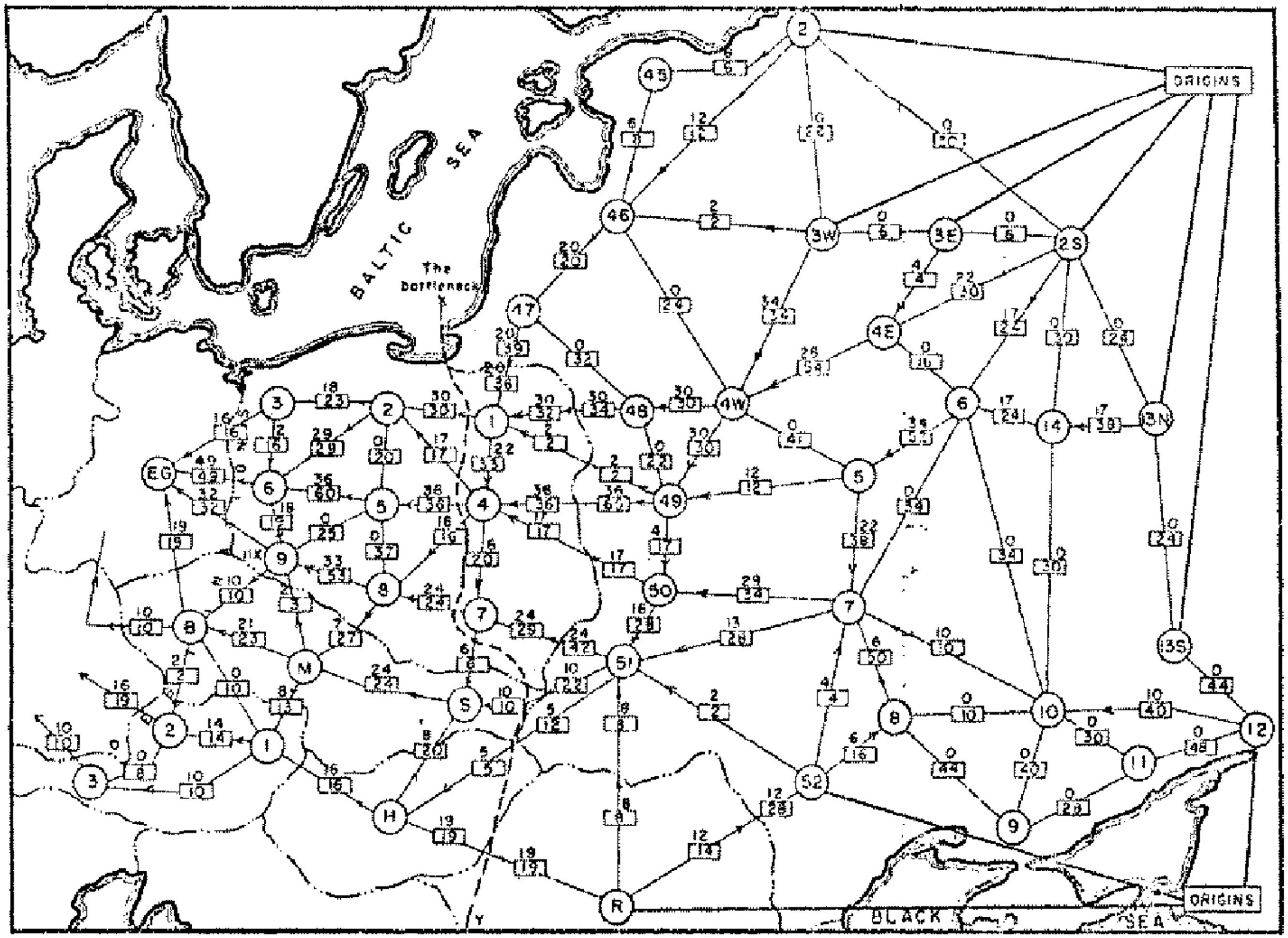
The Network Flow Problem

The Network Flow Problem



How much stuff can flow from s to t?

Soviet Rail Network, 1955



Reference: *On the history of the transportation and maximum flow problems*.
Alexander Schrijver in Math Programming, 91: 3, 2002.

Net Flow: Formal Definition

Given:

A digraph $G = (V,E)$

Two vertices s,t in V (source & sink)

A *capacity* $c(u,v) \geq 0$ for each $(u,v) \in E$ (and $c(u,v) = 0$ for all non-edges (u,v))

Find:

A *flow function* $f: V \times V \to R$ s.t., for all u,v:

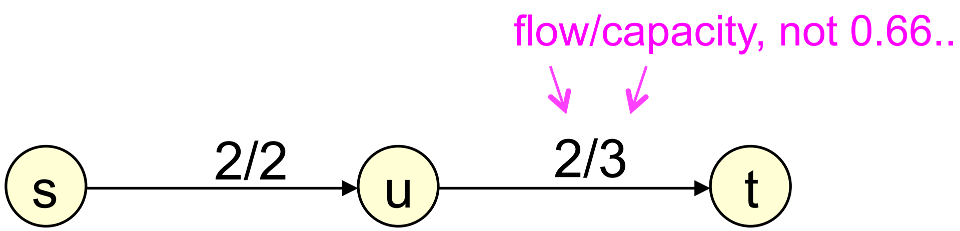
- $f(u,v) \leq c(u,v)$ [Capacity Constraint]
- $f(u,v) = -f(v,u)$ [Skew Symmetry]
- if $u \neq s,t$, $f(u,V) = 0$ [Flow Conservation]

Maximizing total flow $|f| = f(s,V)$

Notation:

$$f(X,Y) = \sum_{x \in X} \sum_{y \in Y} f(x,y)$$

Example: A Flow Function



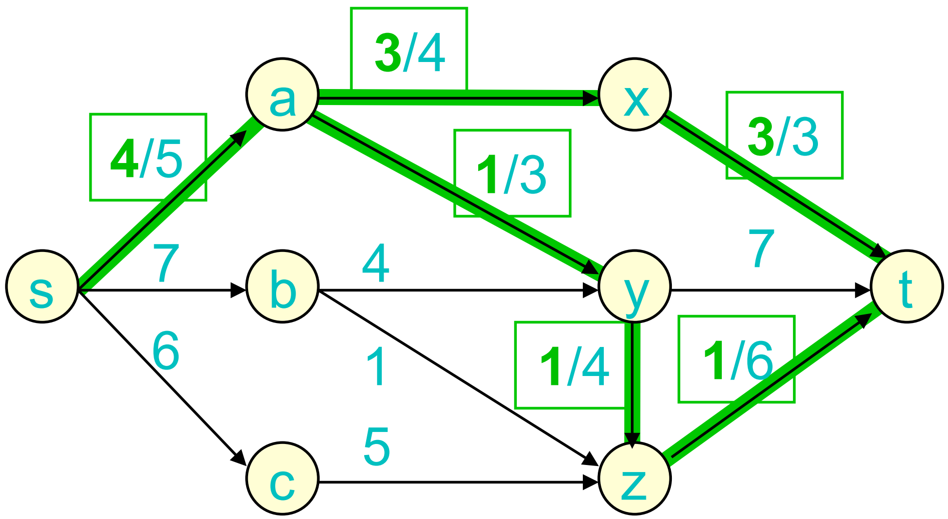
$f(s,u) = f(u,t) = 2$

$f(u,s) = f(t,u) = -2$ *(Why?)*

$f(s,t) = -f(t,s) = 0$ *(In every flow function for this G. Why?)*

$$f(u,V) = \sum_{v \in V} f(u,v) = f(u,s) + f(u,t) = -2 + 2 = 0$$

Example: A Flow Function



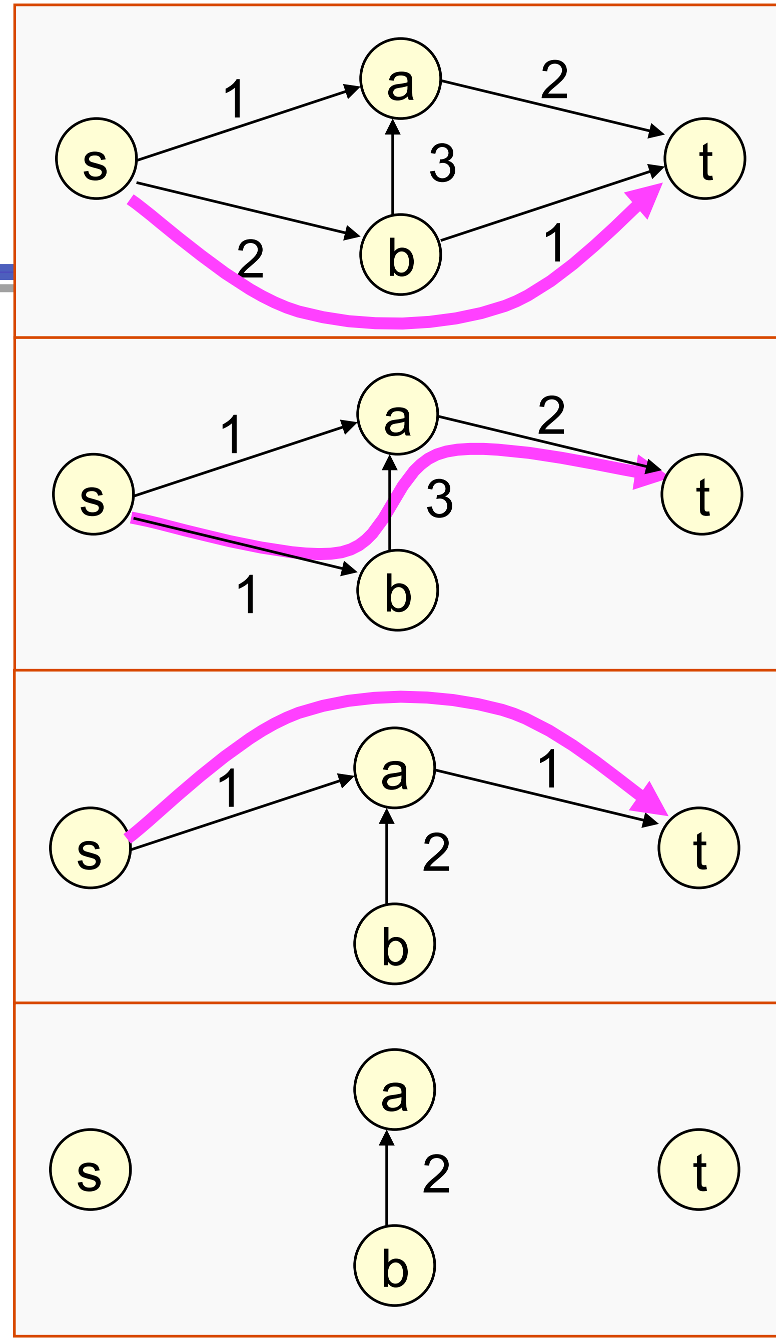
Not shown: f(u,v) if ≤ 0

Note: max flow ≥ 4 since f is a flow, |f| = 4

Max Flow via a Greedy Alg?

While there is an $s \to t$ path in G

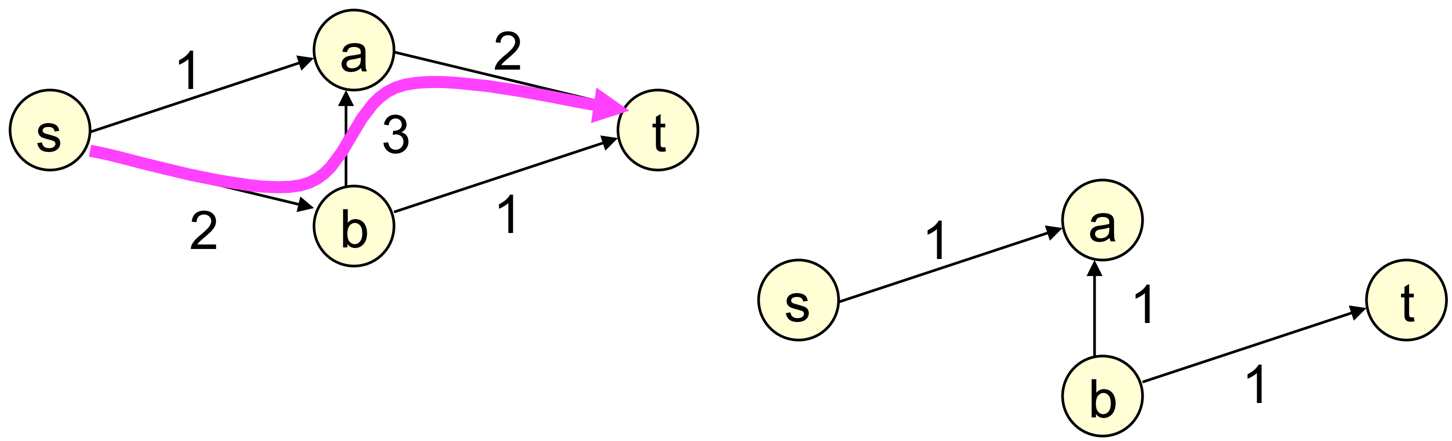
- Pick such a path, p
- Find c_p, the min capacity of any edge in p
- Subtract c_p from all capacities on p
- Delete edges of capacity 0



Max Flow via a Greedy Alg?

This does NOT always find a max flow:

If you pick s →b →a →t first,



Flow stuck at 2. But flow 3 possible.

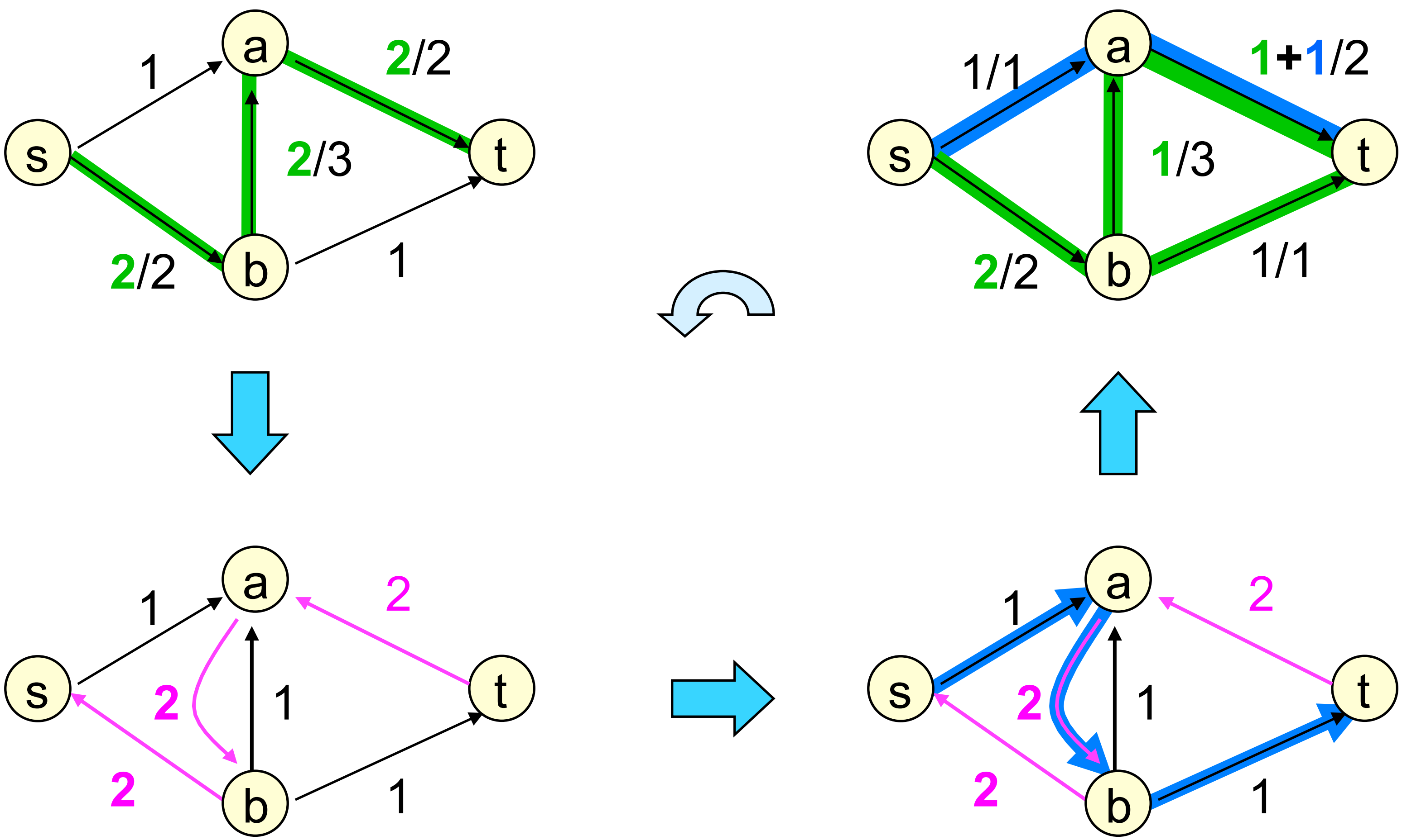
A Brief History of Flow

#	Year	Discoverer(s)	Bound
1	1951	Dantzig	$O(n^2mC)$
2	1955	Ford & Fulkerson	$O(nmC)$
3	1970	Dinitz; Edmonds & Karp	$O(nm^2)$
4	1970	Dinitz	$O(n^2m)$
5	1972	Edmonds & Karp; Dinitz	$O(m^2 \log C)$
6	1973	Dinitz;Gabow	$O(nm \log C)$
7	1974	Karzanov	$O(n^3)$
8	1977	Cherkassky	$O(n^2 \text{sqrt}(m))$
9	1980	Galil & Naamad	$O(nm \log^2 n)$
10	1983	Sleator & Tarjan	$O(nm \log n)$
11	1986	Goldberg &Tarjan	$O(nm \log (n^2/m))$
12	1987	Ahuja & Orlin	$O(nm + n^2 \log C)$
13	1987	Ahuja et al.	$O(nm \log(n \text{ sqrt}(\log C)/(m+2))$
14	1989	Cheriyan & Hagerup	$E(nm + n^2 \log^2 n)$
15	1990	Cheriyan et al.	$O(n^3/\log n)$
16	1990	Alon	$O(nm + n^{8/3} \log n)$
17	1992	King et al.	$O(nm + n^{2+\varepsilon})$
18	1993	Phillips & Westbrook	$O(nm(\log_{m/n} n + \log^{2+\varepsilon} n)$
19	1994	King et al.	$O(nm(\log_{m/(n \log n)} n)$
20	1997	Goldberg & Rao	$O(m^{3/2} \log(n^2/m) \log C)$; $O(n^{2/3} m \log(n^2/m) \log C)$
…	…	…	…

n = # of vertices
m= # of edges
C = Max capacity

Source: Goldberg & Rao, FOCS '97

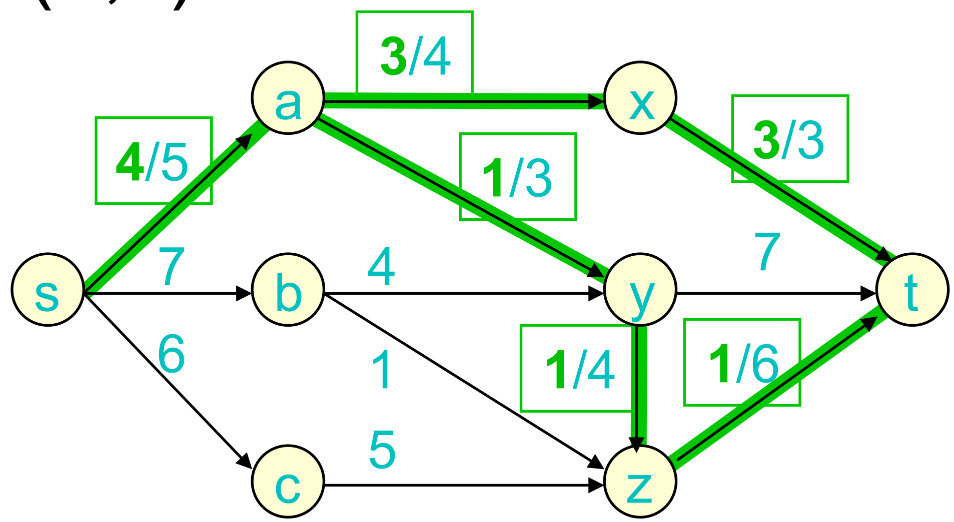
Greed Revisited



Residual Capacity

The *residual capacity* (w.r.t. f) of (u,v) is
$c_f(u,v) = c(u,v) - f(u,v)$

E.g.:



$c_f(s,b) = 7;$

$c_f(a,x) = 1;$

$c_f(x,a) = 3;$

$c_f(x,t) = 0$ (a saturated edge)

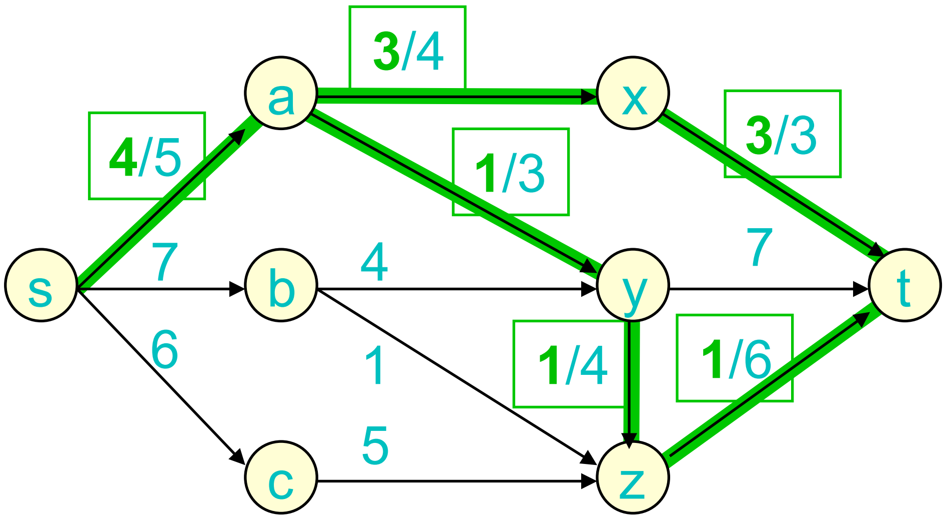
Residual Networks & Augmenting Paths

The *residual network* (w.r.t. f) is the graph $G_f = (V, E_f)$, where

$$E_f = \{ (u,v) \mid c_f(u,v) > 0 \}$$

An *augmenting path* (w.r.t. f) is a simple $s \rightarrow t$ path in G_f.

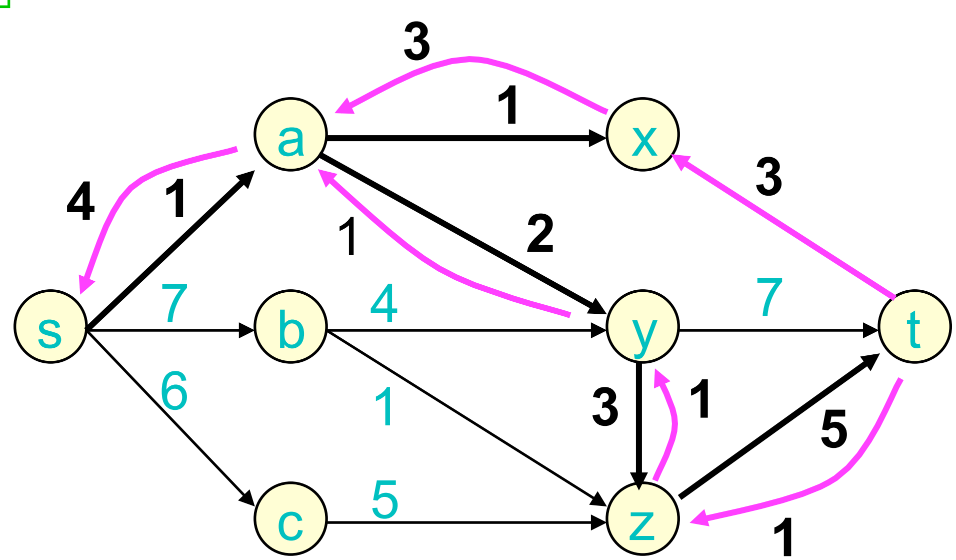
A Residual Network



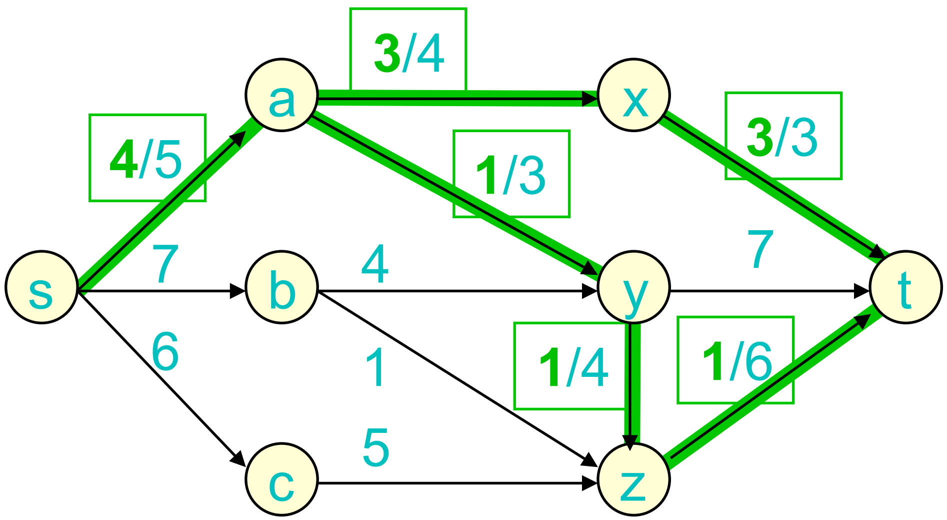
residual network: the graph

$G_f = (V, E_f)$, where

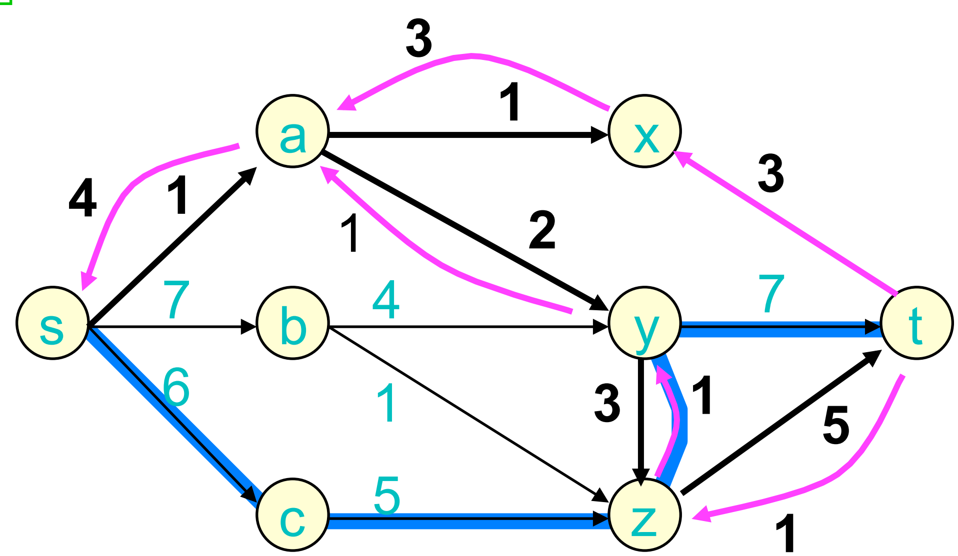
$E_f = \{ (u,v) \mid c_f(u,v) > 0 \}$



An Augmenting Path



augmenting path:
a simple s → t path in G_f.

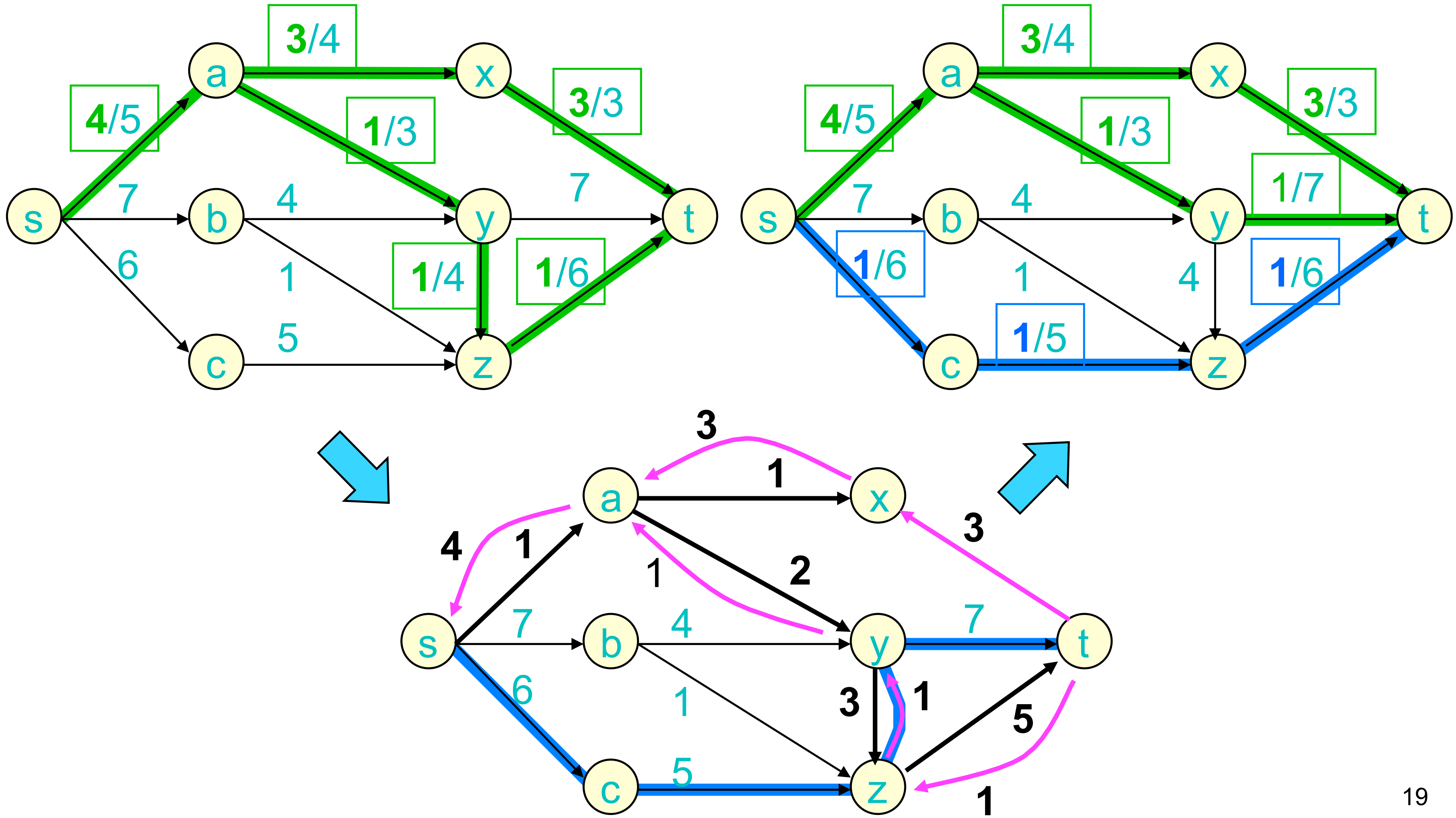


Lemma 1

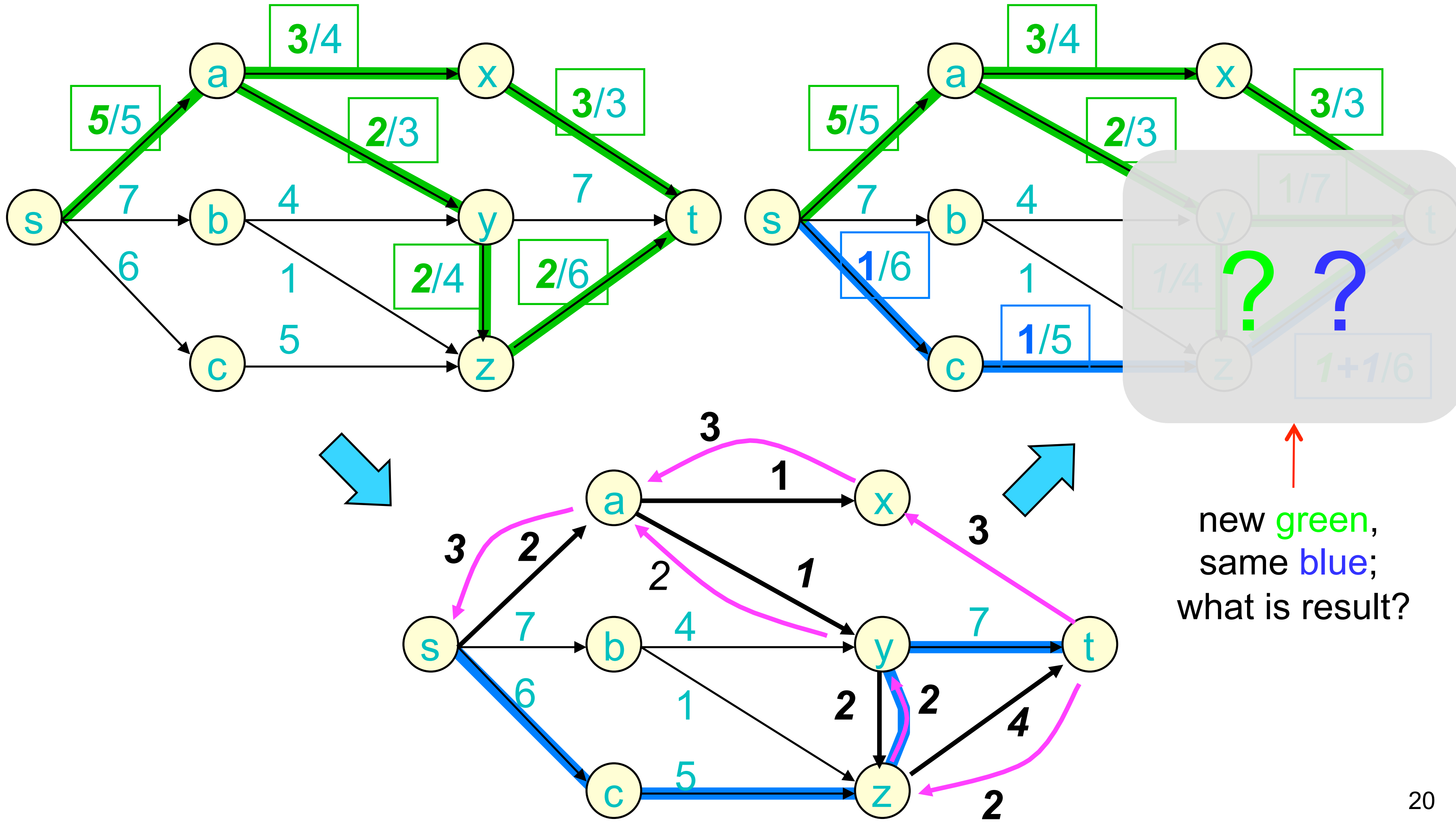
If f admits an augmenting path p, then f is not maximal.

Proof: “obvious” -- augment along p by c_p, the min residual capacity of p’s edges.

Augmenting A Flow



Augmenting A Flow



Lemma 1’: Augmented Flows are Flows

If f is a flow & p an augmenting path of capacity c_p, then f' is also a valid flow, where

$$f'(u,v) = \begin{cases} f(u,v) + c_p, & \text{if } (u,v) \text{ in path } p \\ f(u,v) - c_p, & \text{if } (v,u) \text{ in path } p \\ f(u,v), & \text{otherwise} \end{cases}$$

Proof:

a) Flow conservation – easy
b) Skew symmetry – easy
c) Capacity constraints – pretty easy

Lma 1': Augmented Flows are Flows

$$f'(u,v) = \begin{cases} f(u,v) + c_p, & \text{if } (u,v) \text{ in path } p \\ f(u,v) - c_p, & \text{if } (v,u) \text{ in path } p \\ f(u,v), & \text{otherwise} \end{cases}$$

f a flow & p an aug path of cap c_p, then f' also a valid flow.

Proof (Capacity constraints):

(u,v), (v,u) not on path: no change

(u,v) on path:

$$\begin{aligned} f'(u,v) &= f(u,v) + c_p \\ &\le f(u,v) + c_f(u,v) \\ &= f(u,v) + c(u,v) - f(u,v) \\ &= c(u,v) \end{aligned}$$

$$\begin{aligned} f'(v,u) &= f(v,u) - c_p \\ &< f(v,u) \\ &\le c(v,u) \end{aligned}$$

QED

Residual Capacity:

$0 < c_p \le c_f(u,v) = c(u,v) - f(u,v)$

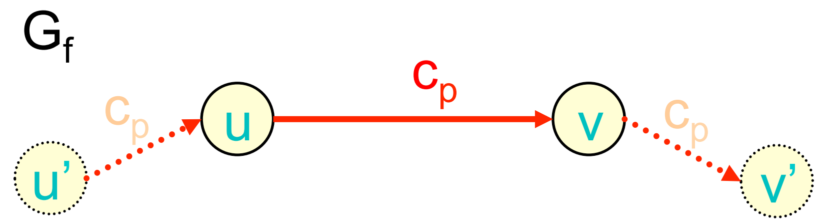
Cap Constraints:

$-c(v,u) \le f(u,v) \le c(u,v)$

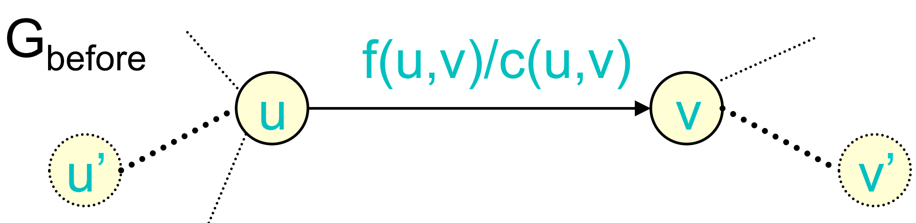
Lemma 1' Example – Case 1

Let (u,v) be any edge in augmenting path. Note

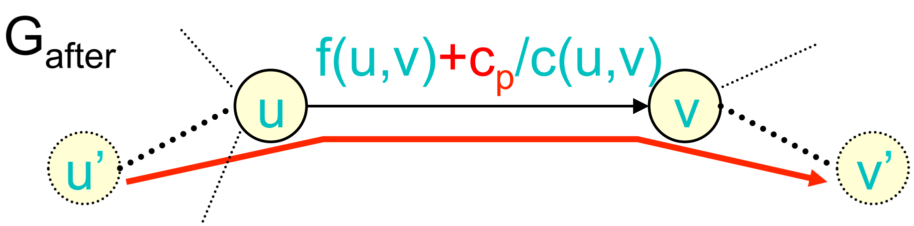
$c_f(u,v) = c(u,v) - f(u,v) \geq c_p > 0$



Case 1: $f(u,v) \geq 0$:



Add forward flow



Lemma 1' Example – Case 2

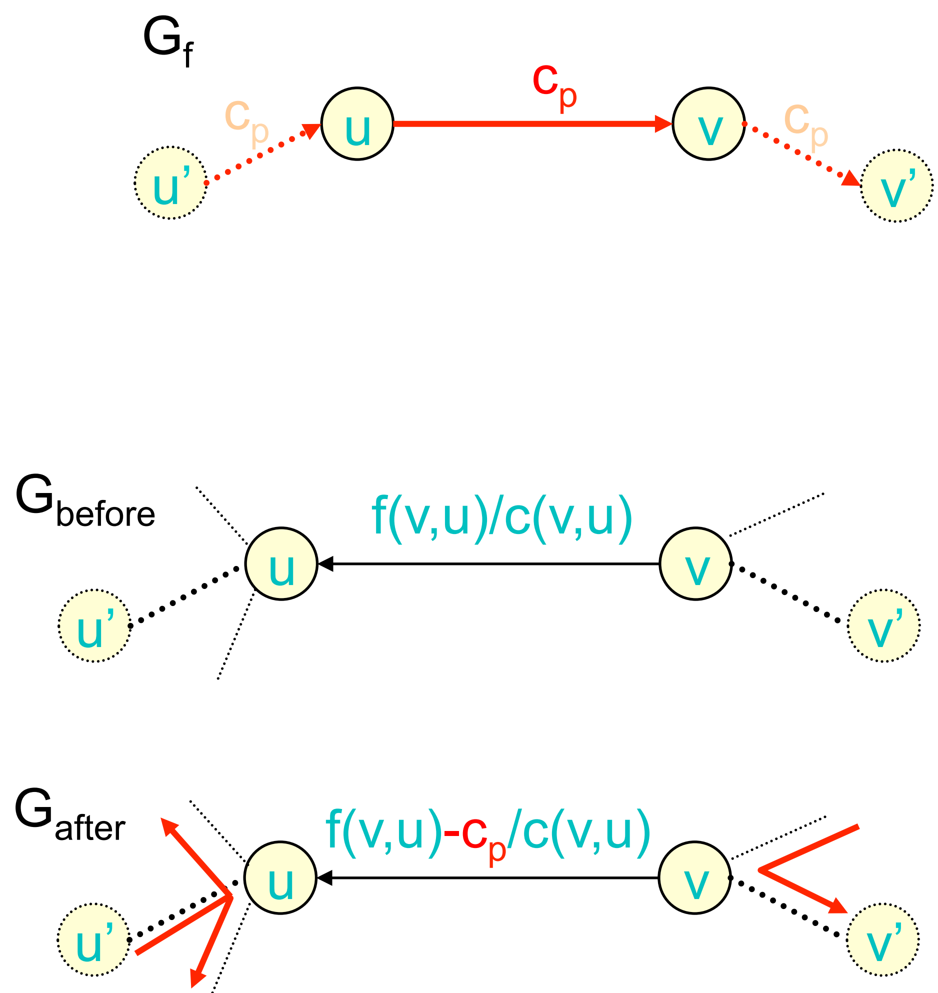
Let (u,v) be any edge in augmenting path. Note

$c_f(u,v) = c(u,v) - f(u,v) \geq c_p > 0$

Case 2: $f(u,v) \leq -c_p$:

$f(v,u) = -f(u,v) \geq c_p$

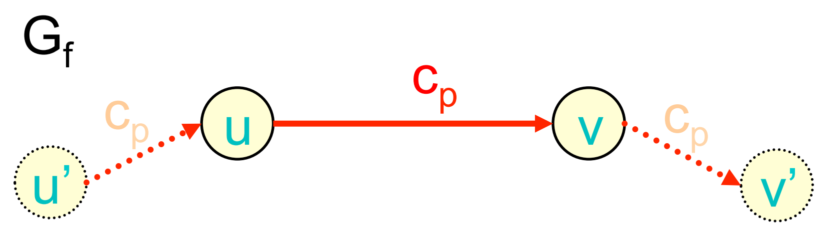
Cancel/redirect reverse flow



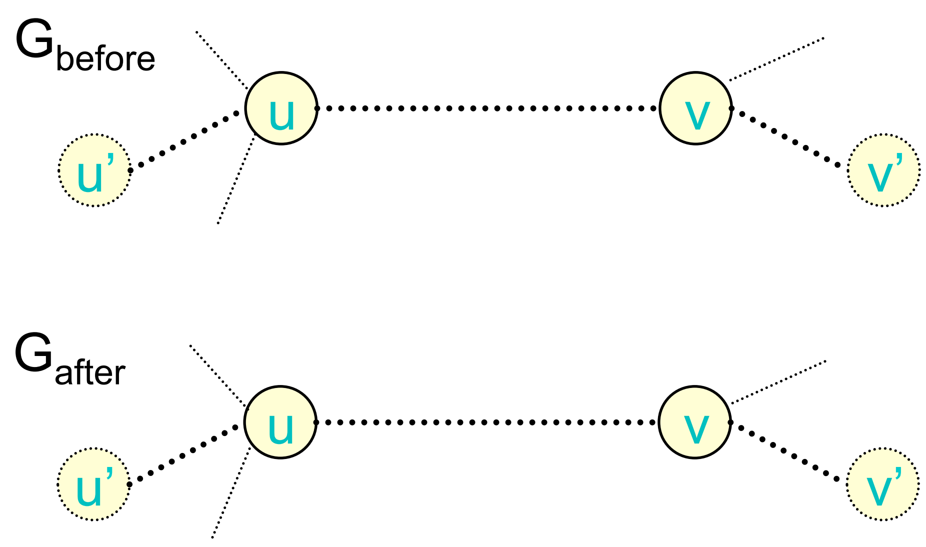
Lemma 1' Example – Case 3

Let (u,v) be any edge in augmenting path. Note

$c_f(u,v) = c(u,v) - f(u,v) \geq c_p > 0$



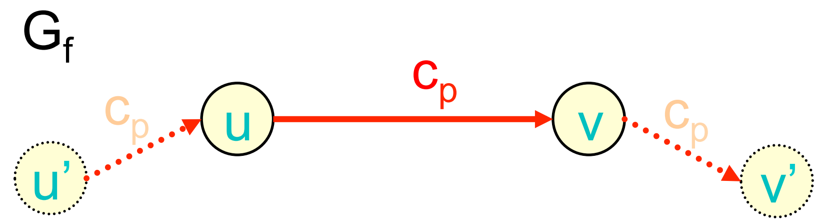
Case 3: $-c_p < f(u,v) < 0$:



???

[E.g., $c_p = 8$, $f(u,v) = -5$]

Lemma 1' Example – Case 3

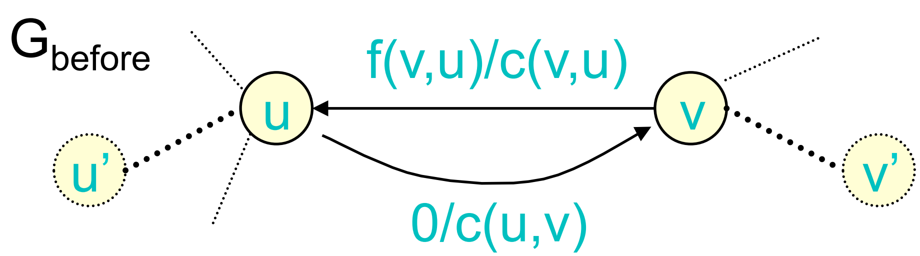


Let (u,v) be any edge in augmenting path. Note

$c_f(u,v) = c(u,v) - f(u,v) \geq c_p > 0$

Case 3: $-c_p < f(u,v) < 0$

$c_p > f(v,u) > 0$:

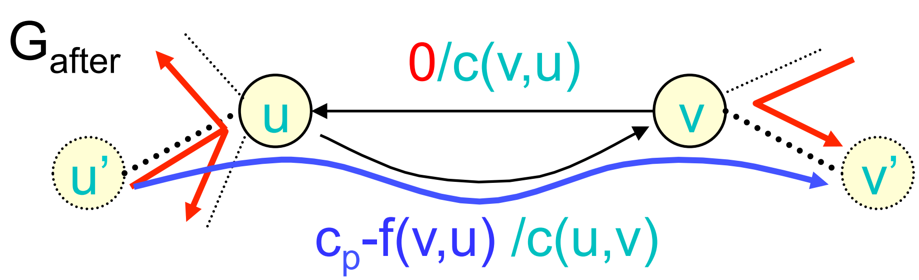


Both:

cancel/redirect reverse flow

and

add forward flow



Ford-Fulkerson Method

While G_f has an augmenting path, augment

Questions:

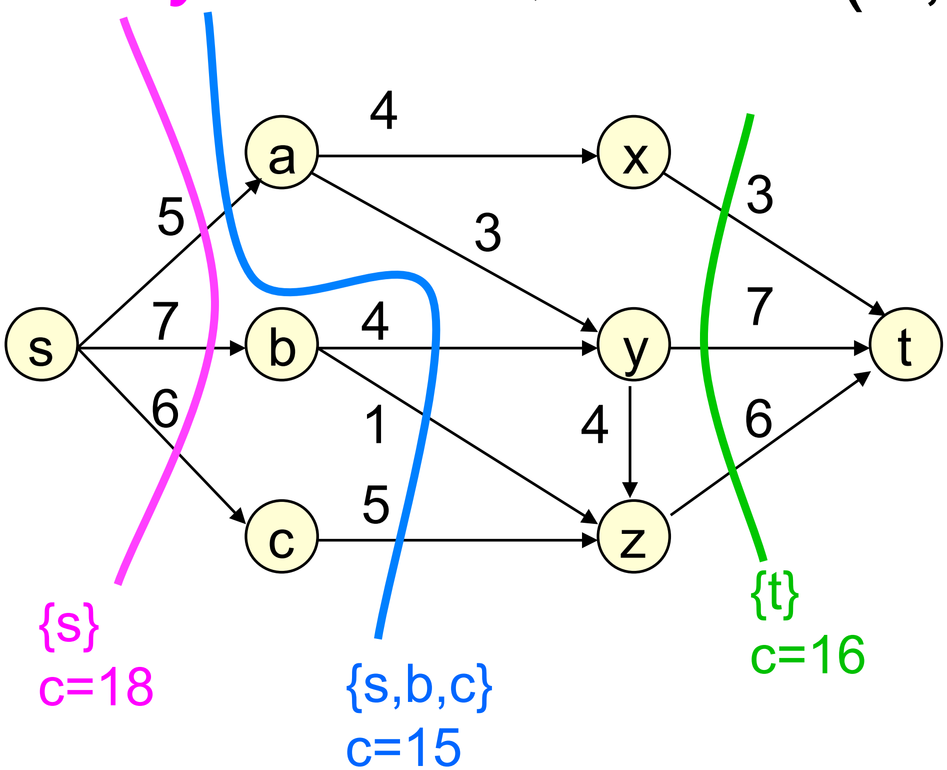
- Does it halt?
- Does it find a maximum flow?
- How fast?

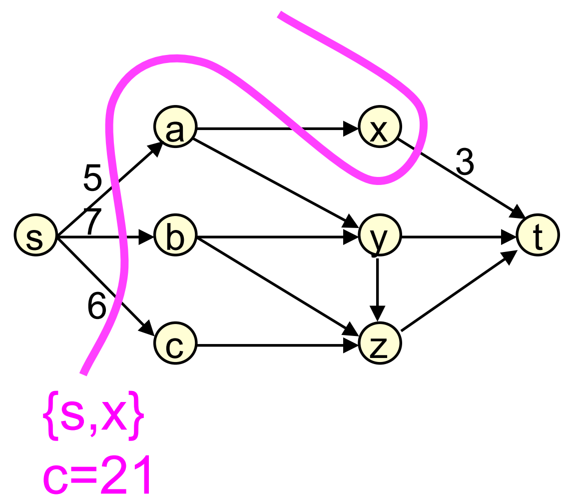
Cuts

A partition S,T of V is a *cut* if s ∈ S, t ∈ T.

Capacity of cut S,T is $c(S,T) = \sum_{u \in S, v \in T} c(u,v)$

sum of caps of edges *from S to T*



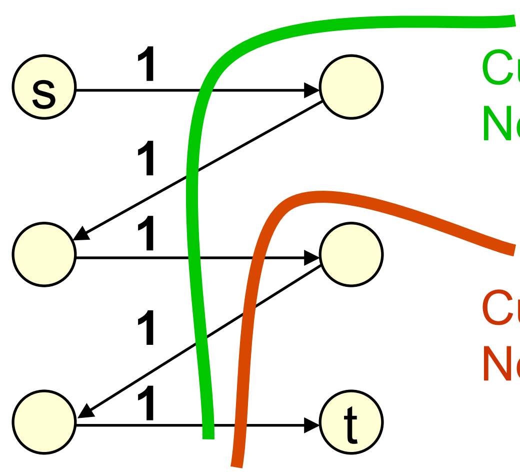


Lemma 2

For any flow f and any cut S,T,

- the net flow across the cut equals the total flow, i.e., $|f| = f(S,T)$, and
- the net flow across the cut cannot exceed the capacity of the cut, i.e. $f(S,T) \le c(S,T)$

Corollary:
Max flow ≤ Min cut



Cut Cap = 3
Net Flow = 1

Cut Cap = 2
Net Flow = 1

Lemma 2

For any flow f and any cut S,T,

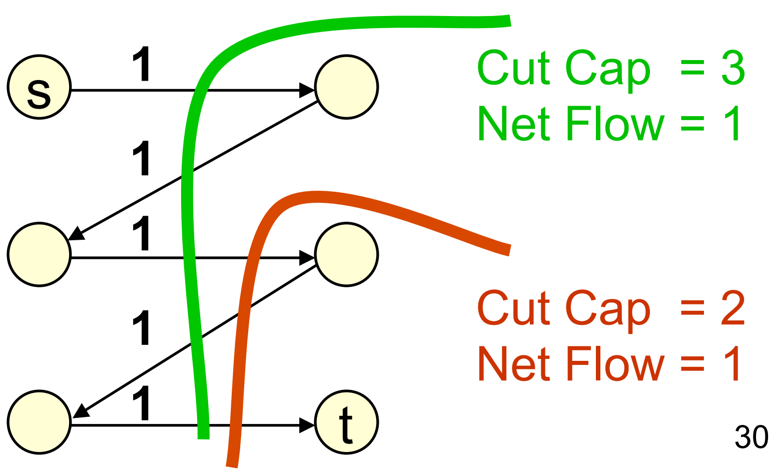
net flow across cut = total flow ≤ cut capacity

Proof:

Track a flow unit. Starts at s, ends at t.

crosses cut an odd # of times; net = 1.

Last crossing uses a forward edge totaled in C(S,T)



Max Flow / Min Cut Theorem

For any flow f, the following are equivalent

(1) $|f| = c(S,T)$ for some cut S,T (a min cut)

(2) f is a maximum flow

(3) f admits no augmenting path

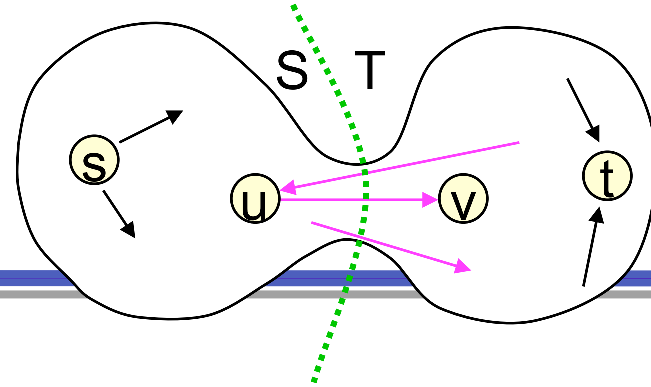
Proof:

(1) $\Rightarrow$ (2): corollary to lemma 2

(2) $\Rightarrow$ (3): contrapositive of lemma 1

(3) ⇒ (1)

(no aug) ⇒ (cut)



S = { u | ∃ an augmenting path wrt f from s to u }

T = V - S; s ∈ S, t ∈ T

For any (u,v) in S × T, ∃ an augmenting path from s to u, but not to v.

∴ (u,v) has 0 residual capacity:

(u,v) ∈ E ⇒ saturated   f(u,v) = c(u,v)

(v,u) ∈ E ⇒ no flow   f(u,v) = 0 = -f(v,u)

This is true for every edge crossing the cut, i.e.

$$|f| = f(S,T) = \sum_{u \in S} \sum_{v \in T} f(u,v) =$$

$$\sum_{u \in S, v \in T, (u,v) \in E} f(u,v) = \sum_{u \in S, v \in T, (u,v) \in E} c(u,v) = c(S,T)$$

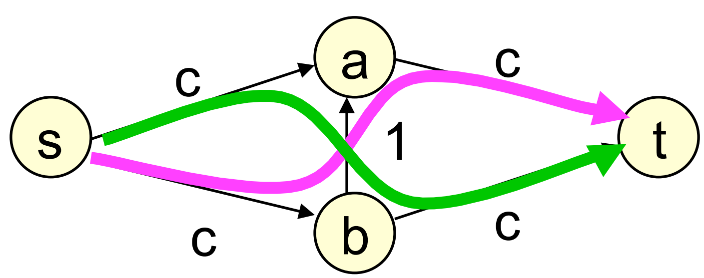
Idea: where's the bottleneck

Corollaries & Facts

If Ford-Fulkerson terminates, then it's found a max flow.

It will terminate if c(e) integer or rational (but may not if they're irrational).

However, may take exponential time, even with integer capacities:



c = 10^{99}, say

How to Make it Faster

Several ways. Three important ones:

- Edmonds-Karp '70; Dinitz '70
  - 1st "strongly" poly time alg. (next) $T = O(nm^2)$
- "Scaling" [Edmonds-Karp, '72; Dinitz '72]
  - do *largest* edges first; see text, and below.
  - if C = max capacity, $T = O(m^2 \log C)$
- Preflow-Push [Goldberg, Tarjan '86]
  - see text $T = O(n^3)$

Edmonds-Karp-Dinitz '70 Algorithm

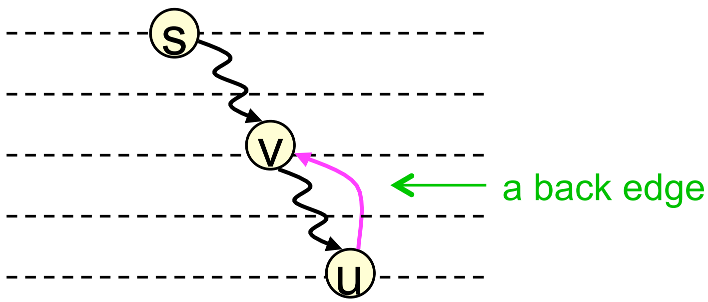
Use a shortest augmenting path
(via Breadth First Search in residual graph)

Time: $O(n\ m^2)$

BFS/Shortest Path Lemmas

Distance from s is never reduced by:

- Deleting an edge
  proof: no new (hence no shorter) path created
- Adding an edge (u,v), provided v is nearer than u
  proof: BFS is unchanged, since v visited before (u,v) examined

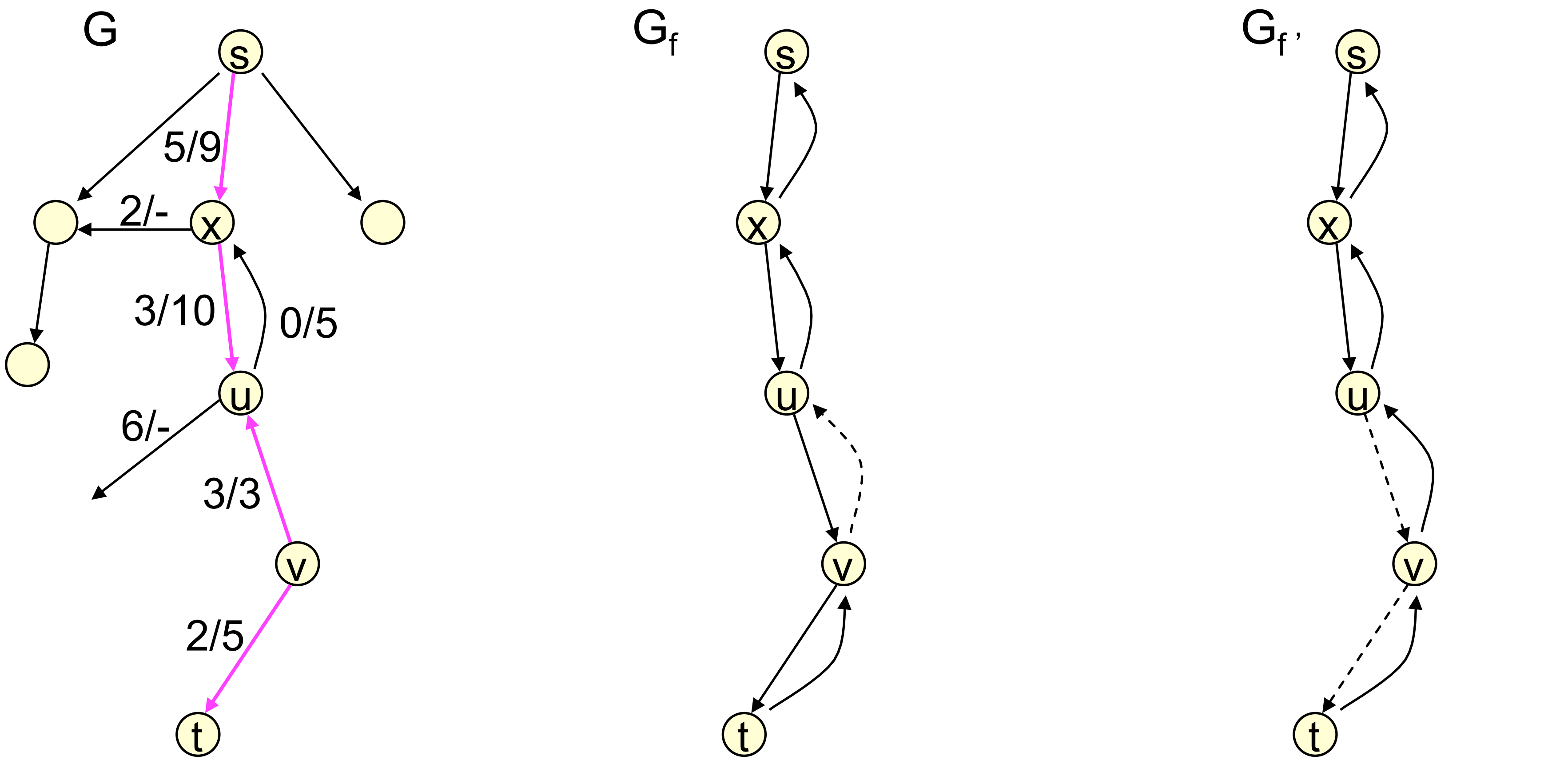


Lemma 3

Let f be a flow, G_f the residual graph, and p a shortest augmenting path. Then no vertex is closer to s in the new residual graph G_{f+p} after augmentation along p.

Proof: Augmentation only deletes edges, adds back edges

Augmentation vs BFS



Theorem 2

The Edmonds-Karp-Dinitiz Algorithm performs O(mn) flow augmentations

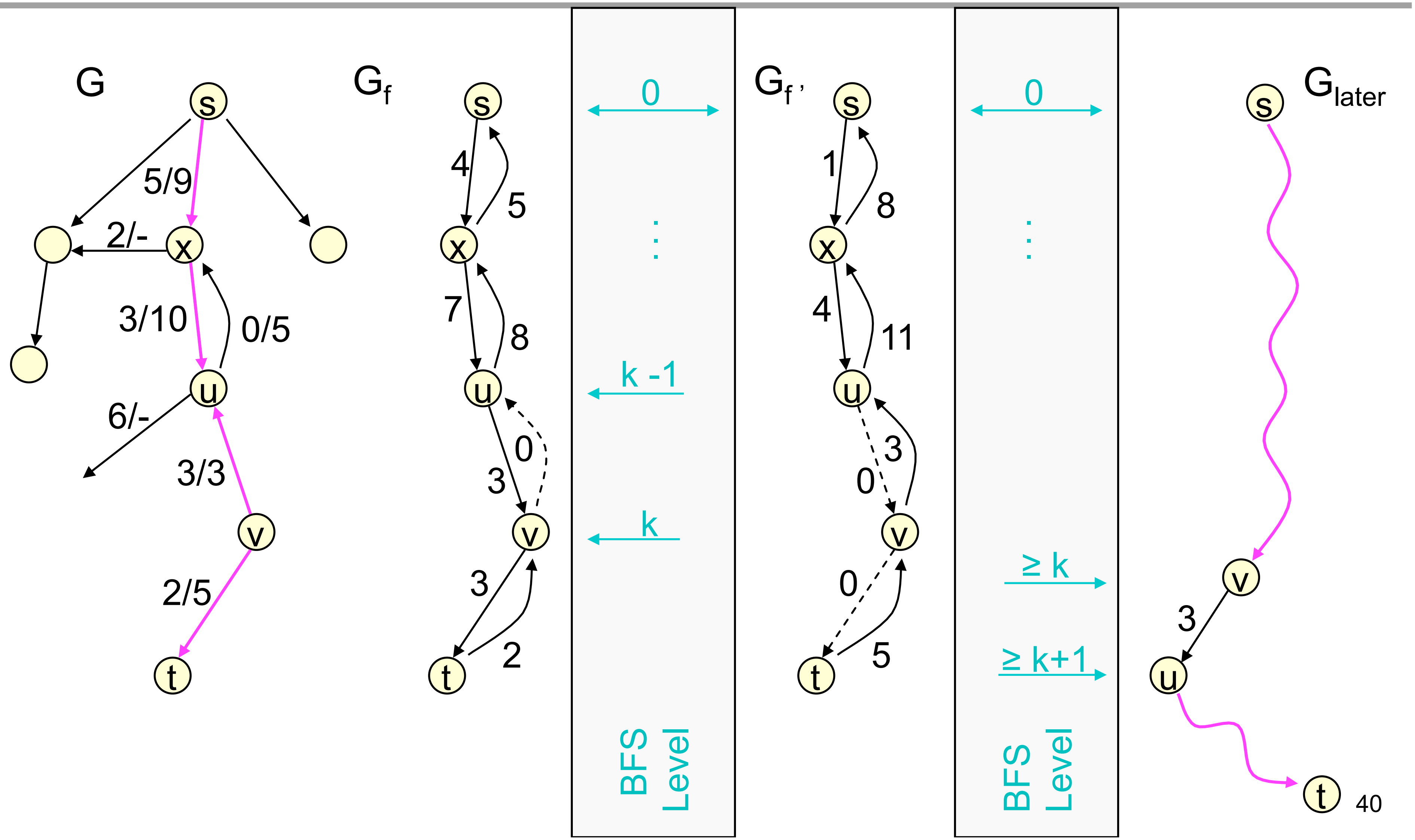
Proof:

{*u*,*v*} is critical on augmenting path *p* if it's closest to *s* having min residual capacity.

Won't be critical again until farther from *s*.

So each edge critical at most *n* times.

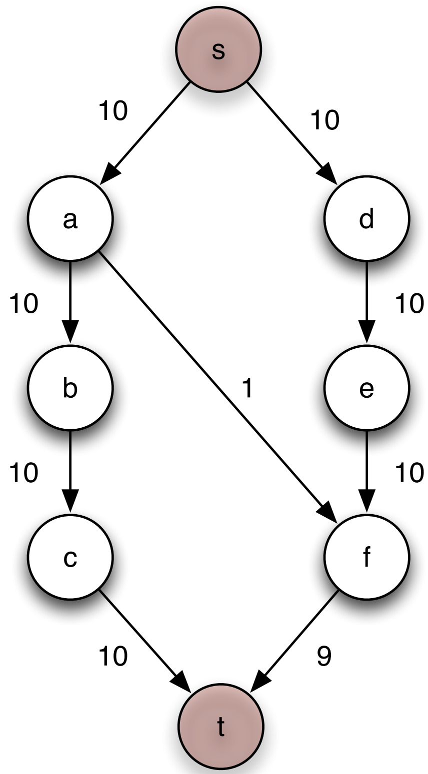
Augmentation vs BFS Level



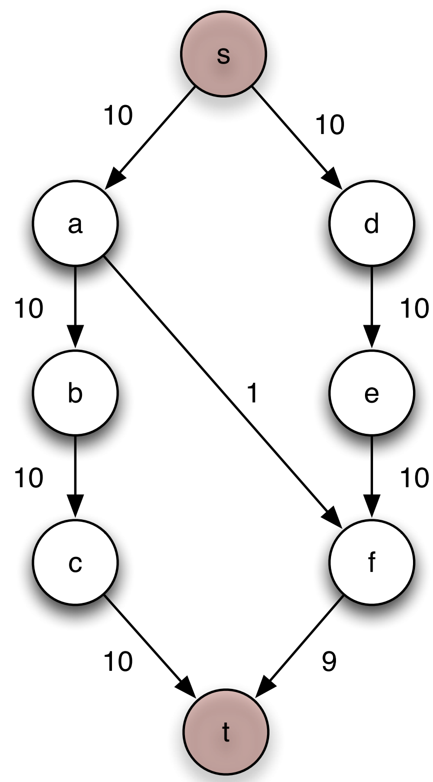
Corollary

Edmonds-Karp-Dinitz runs in $O(nm^2)$

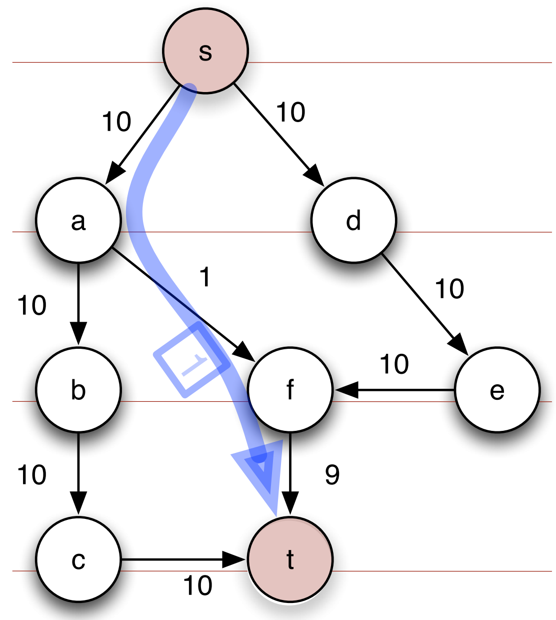
Example



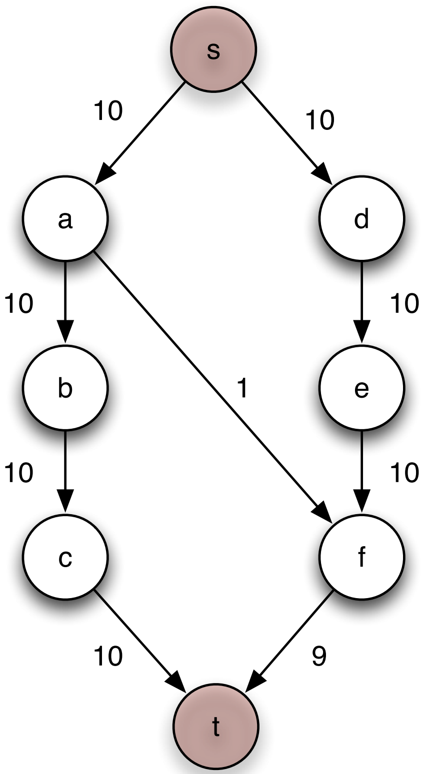
G_0: the flow problem



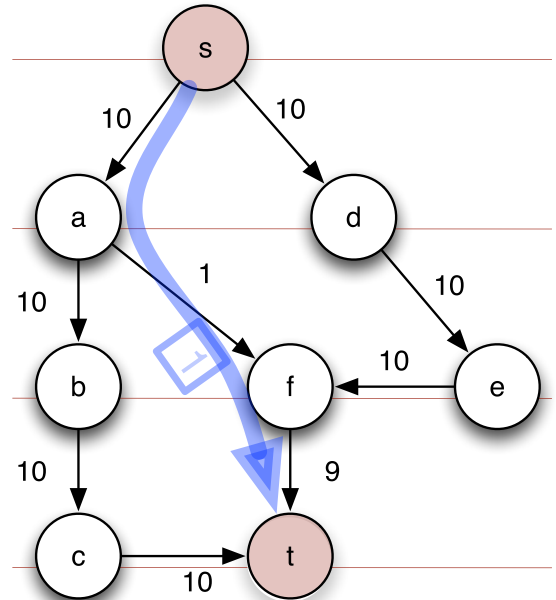
G_0: the flow problem



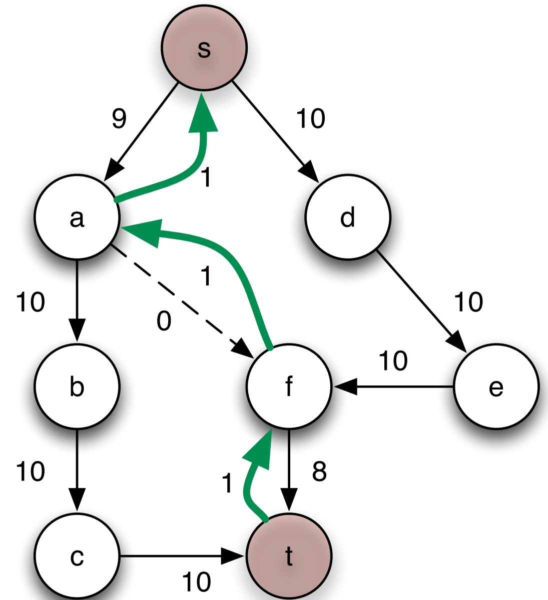
G_0: BFS layering + Aug Path



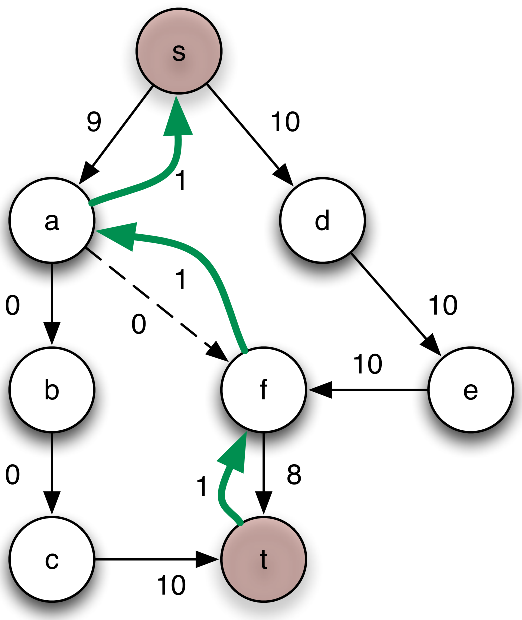
G_0: the flow problem



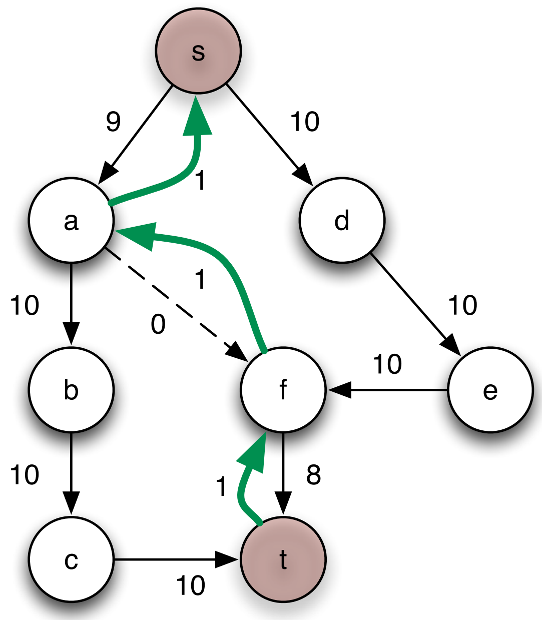
G_0: BFS layering + Aug Path



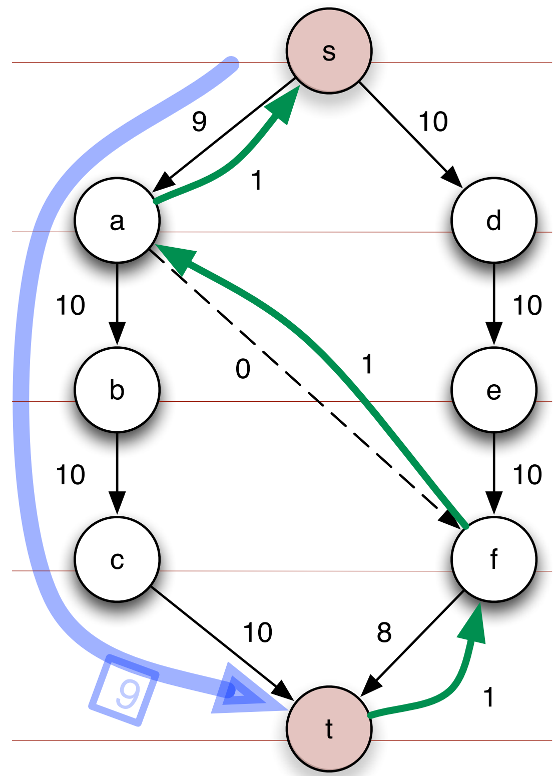
G_1: 1st Residual Graph



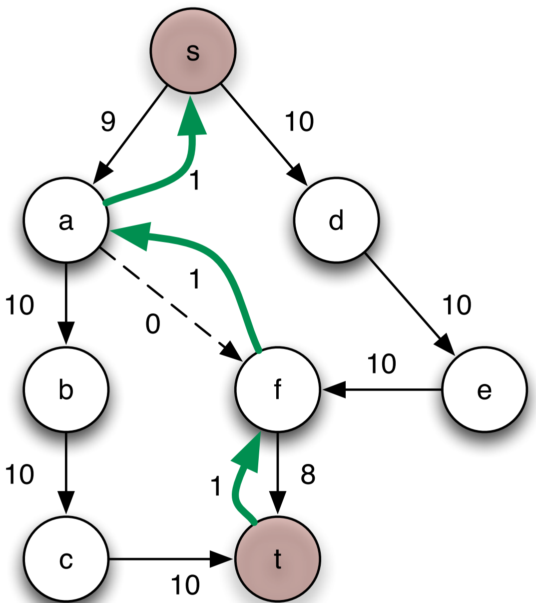
G$_1$: 1st Residual Graph



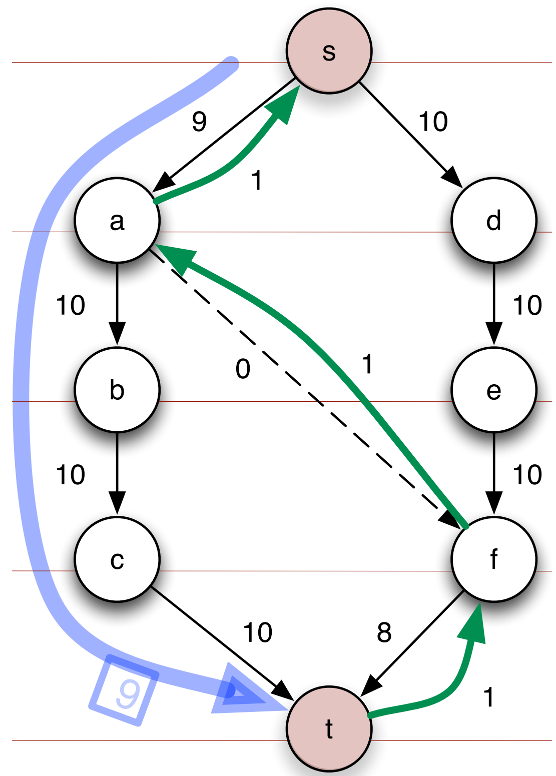
G_1: 1st Residual Graph



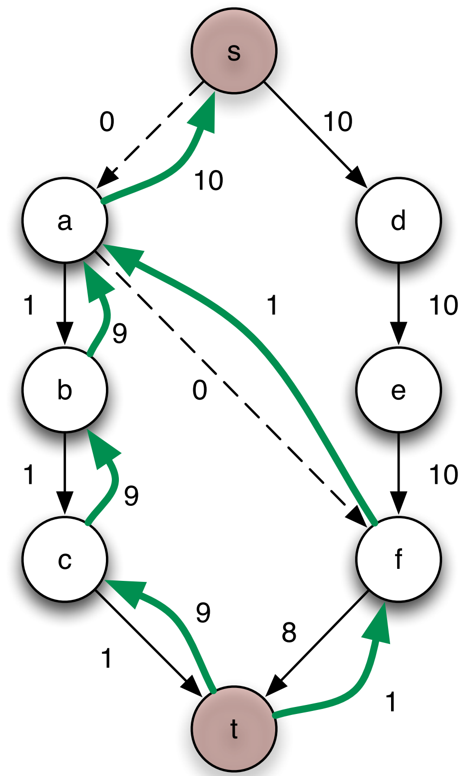
G_1: BFS layering + Aug Path



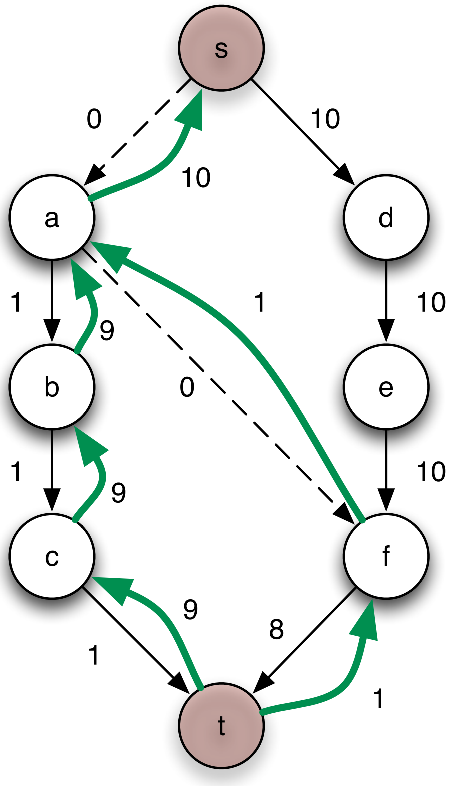
G_1: 1st Residual Graph



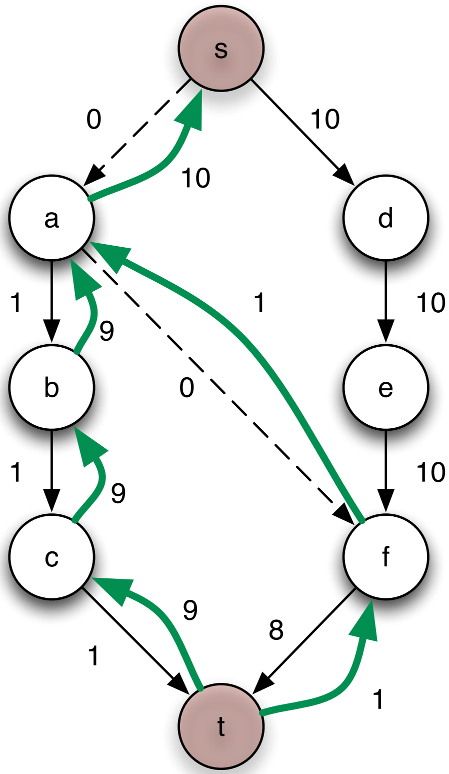
G_1: BFS layering + Aug Path



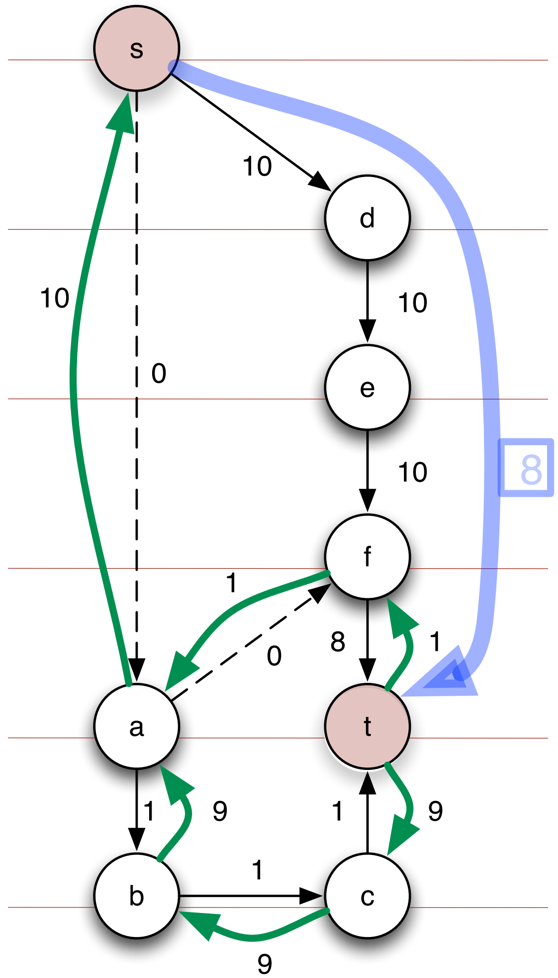
G_2: 2nd Residual Graph



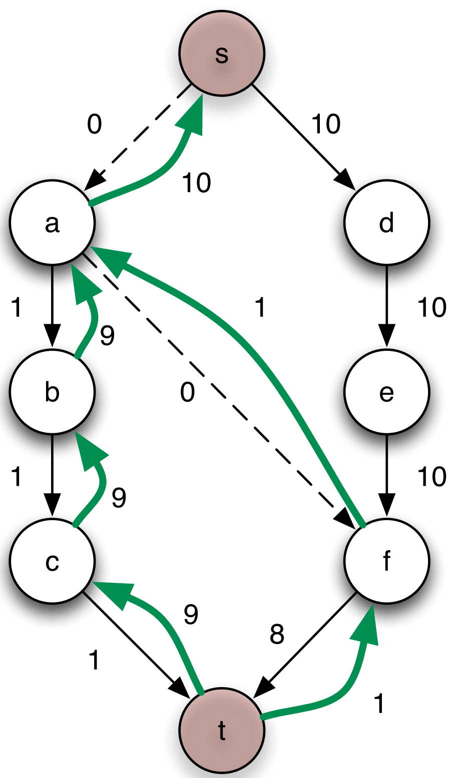
G_2: 2nd Residual Graph



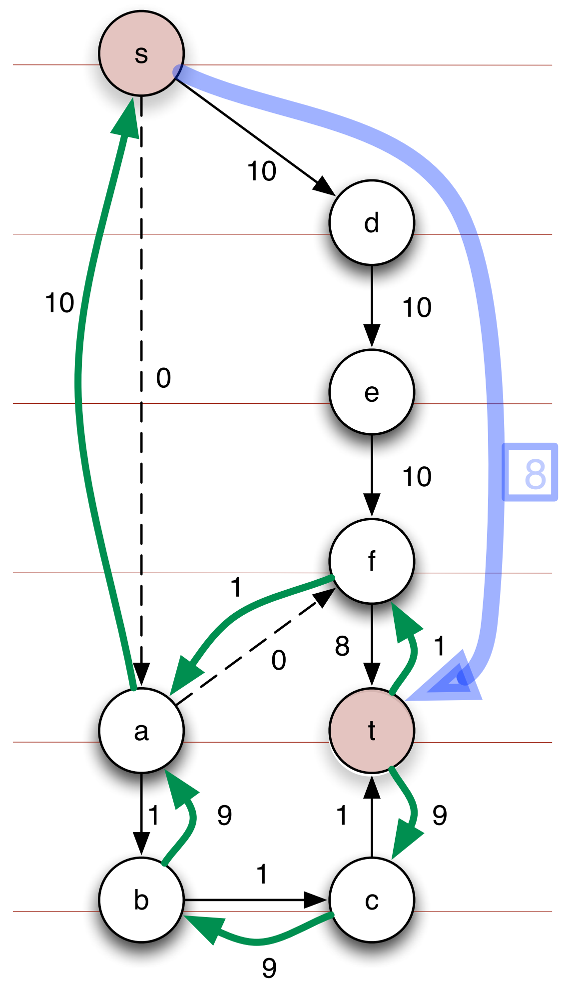
G$_2$: 2nd Residual Graph



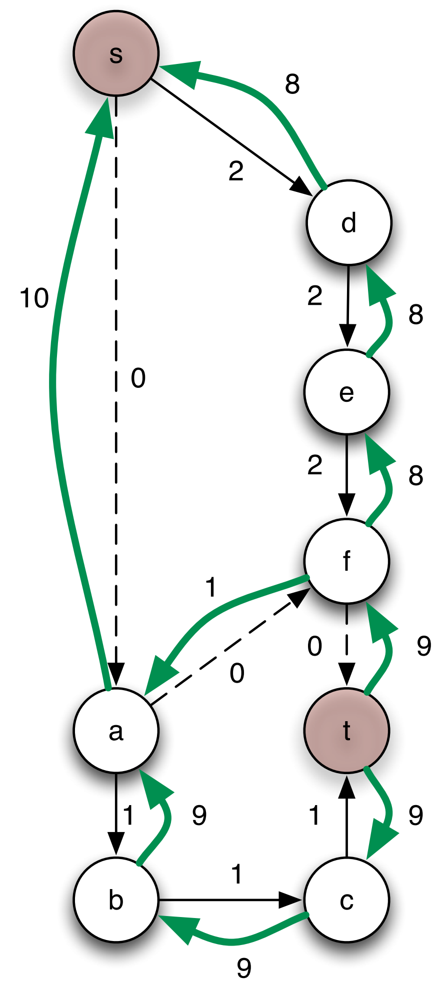
G$_2$: BFS layering + Aug Path



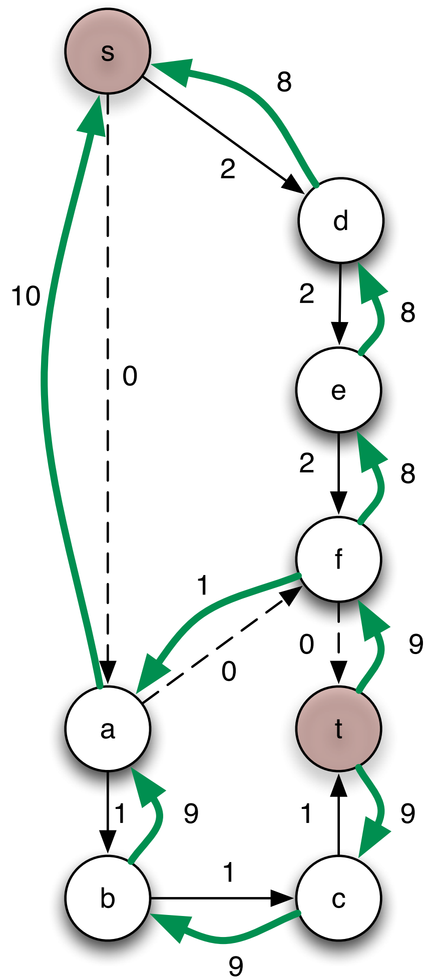
G_2: 2nd Residual Graph



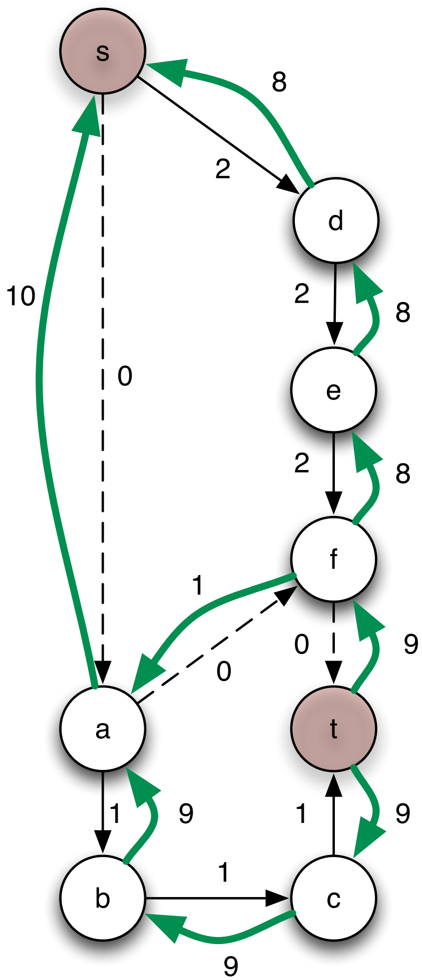
G_2: BFS layering + Aug Path



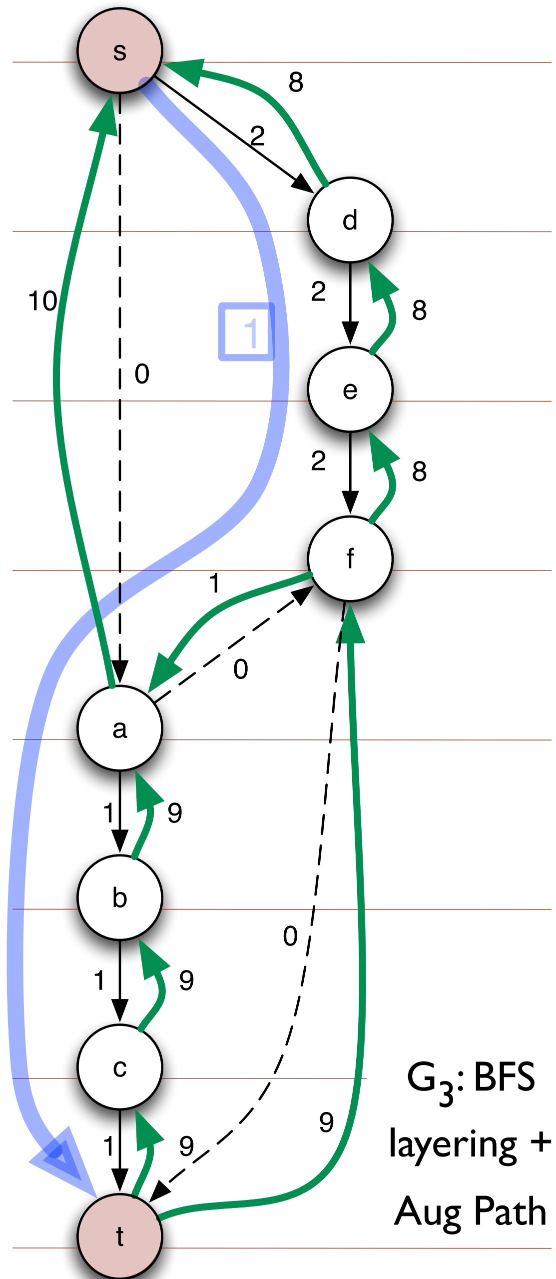
G_3: 3rd Residual Graph



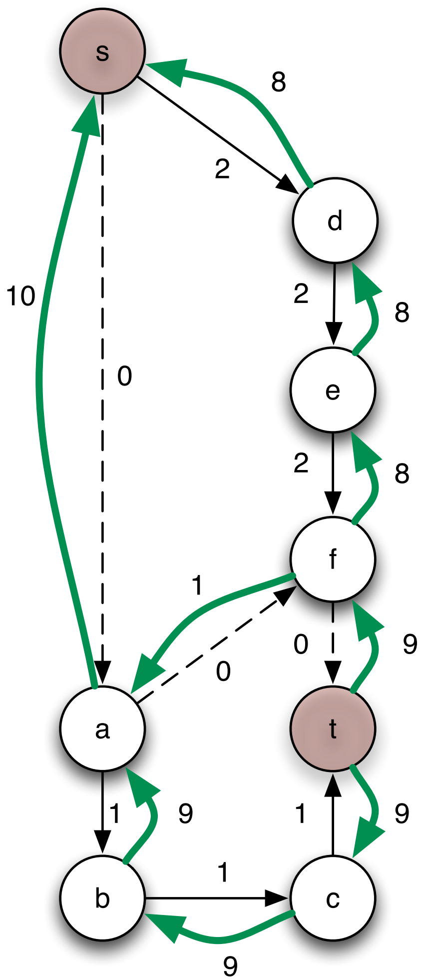
G$_3$: 3rd Residual Graph



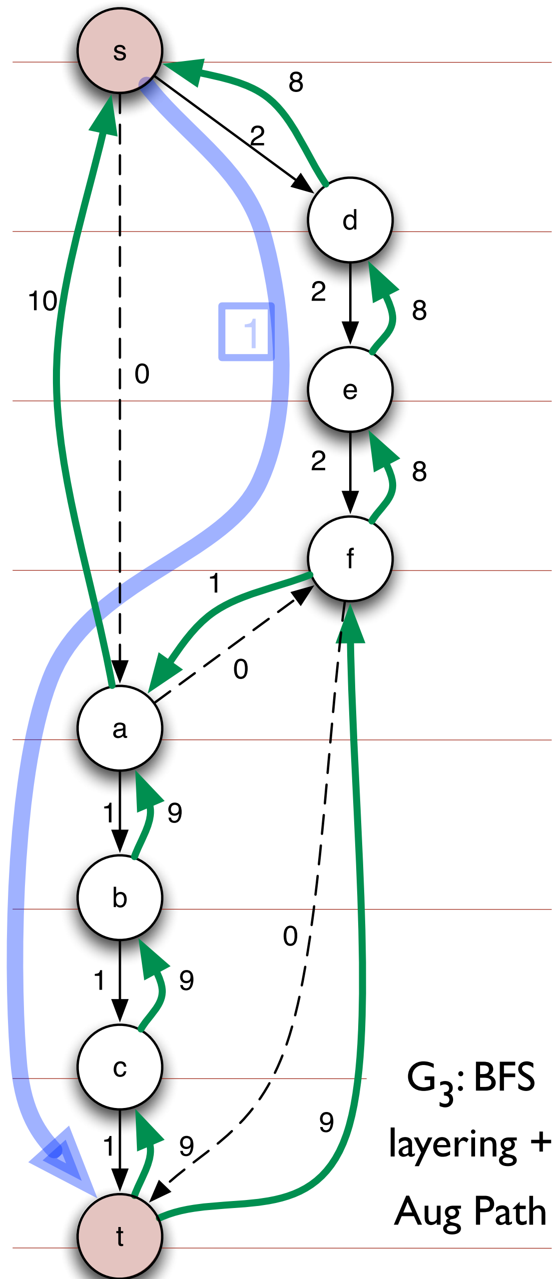
G_3: 3rd Residual Graph



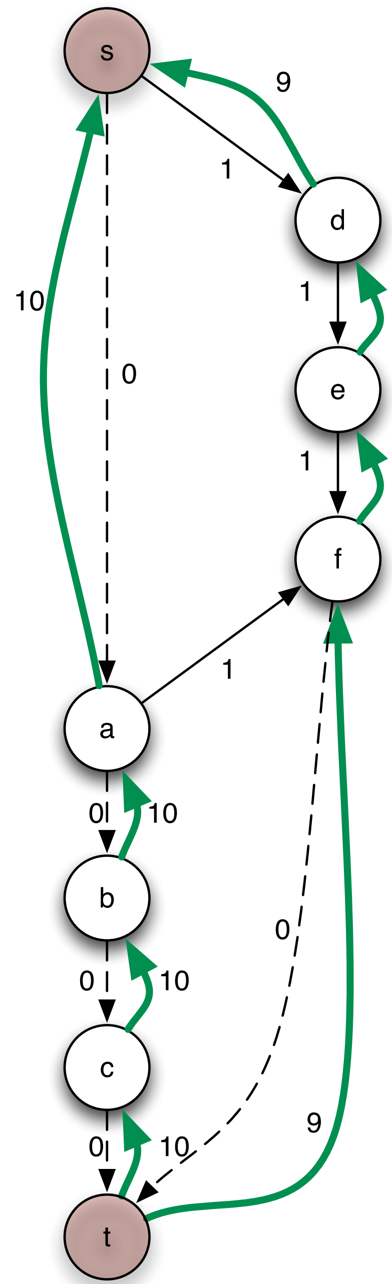
G_3: BFS layering + Aug Path



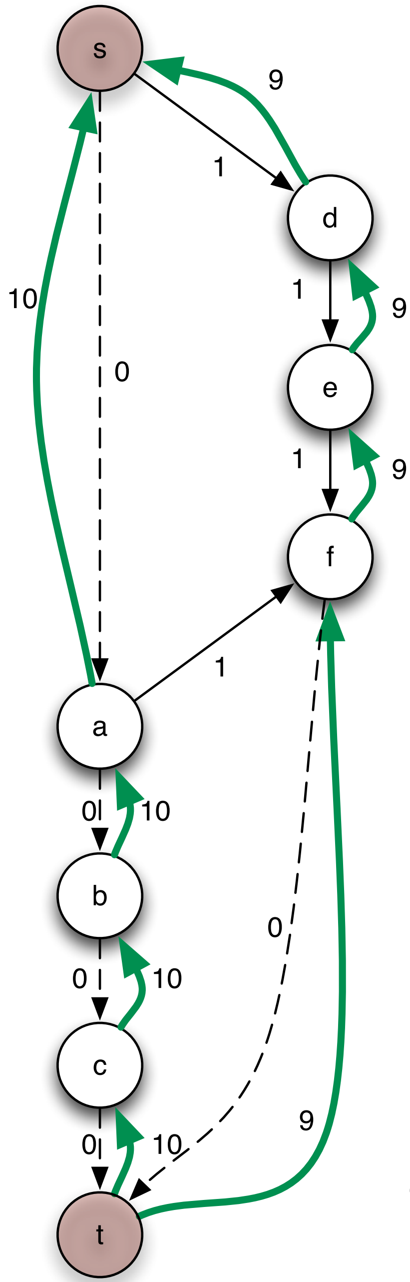
G_3: 3rd Residual Graph



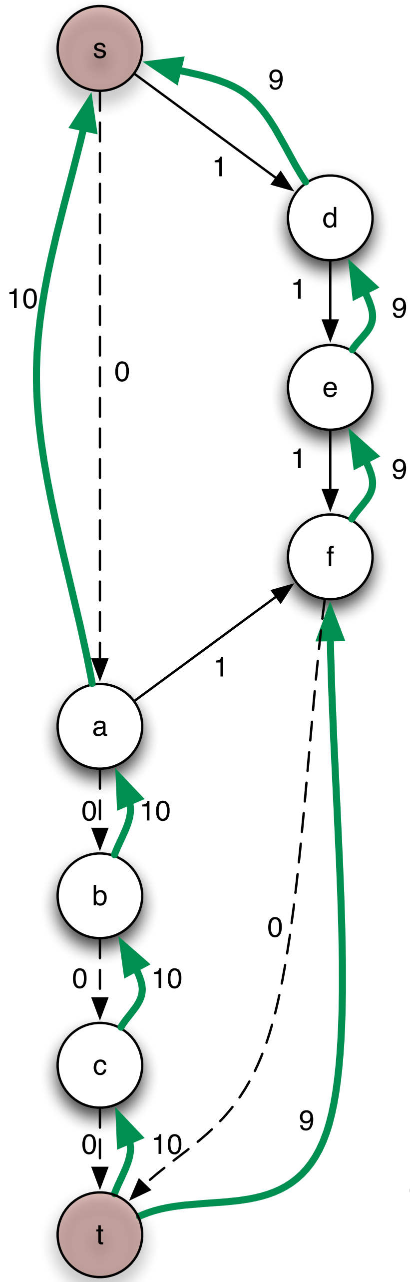
G_3: BFS layering + Aug Path



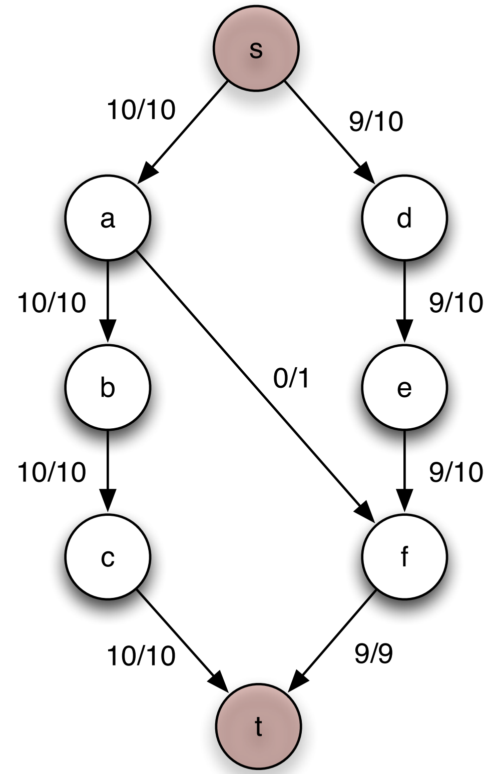
G_4: 4th residual graph



G_4: 4th residual graph



G_4: 4^{th} residual graph



G_5:The Max Flow (19)

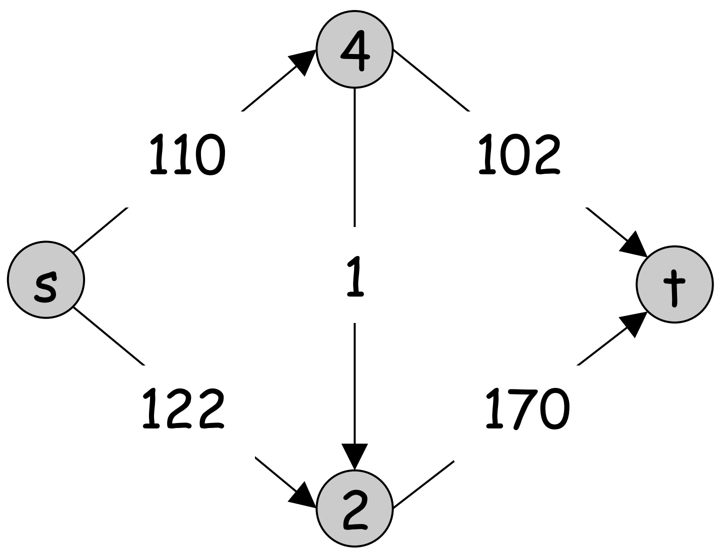
Edmonds-Karp-Dinitz '72

- A natural greedy idea – choose *max capacity* augmenting path first
- Finding *the* max is slow; finding *a large* capacity augmenting path is quicker, and gives the $O(m^2 \log C)$ algorithm "capacity scaling" algorithm below

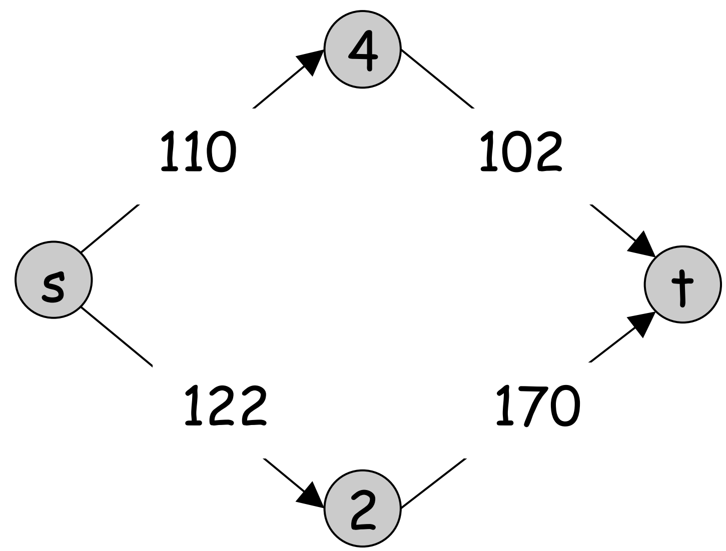
Capacity Scaling

Intuition. Choosing path with highest bottleneck capacity increases flow by max possible amount.

- Don't worry about finding exact highest bottleneck path.
- Maintain scaling parameter Δ.
- Let $G_f(\Delta)$ be the subgraph of the residual graph consisting of only arcs with capacity at least Δ.



G_f



$G_f(100)$

Capacity Scaling

```
Scaling-Max-Flow(G, s, t, c) {
   foreach e ∈ E  f(e) ← 0
   C ← max capacity of any edge
   Δ ← least power of 2 ≥ C

   while (Δ ≥ 1) {
      Gf(Δ) ← Δ-residual graph
      while (∃ augmenting path P in Gf(Δ)) {
         f ← augment(f, c, P)
         update Gf(Δ)
      }
      Δ ← Δ / 2
   }
   return f
}
```

Capacity Scaling: Correctness

Assumption. All edge capacities are integers between 1 and C.

Integrality invariant. All flow and residual capacity values are integral.

Correctness. If the algorithm terminates, then f is a max flow.
Pf.

- By integrality invariant, when $\Delta = 1 \Rightarrow G_f(\Delta) = G_f$.
- Upon termination of $\Delta = 1$ phase, there are no augmenting paths. ▪

Capacity Scaling: Running Time

Lemma 1. The outer while loop repeats $1 + \lceil \log_2 C \rceil$ times.
Pf. Initially $C \leq \Delta < 2C$. Δ decreases by a factor of 2 each iteration. ▪

Lemma 2. Let f be the flow at the end of a Δ-scaling phase. Then the value of the maximum flow is at most $v(f) + m\,\Delta$. ← proof on next slide

Lemma 3. There are at most 2m augmentations per scaling phase.

- Let f be the flow at the end of the previous scaling phase.
- Lemma 2 $\Rightarrow$ $v(f^*) \leq v(f) + m\,(2\Delta)$.
- Each augmentation in a Δ-phase increases $v(f)$ by at least Δ. ▪

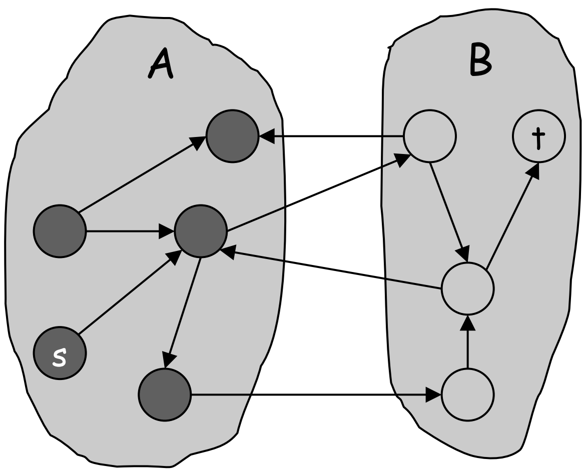
Theorem. The scaling algorithm finds a max flow in $O(m \log C)$ augmentations. It can be implemented to run in $O(m^2 \log C)$ time. ▪

Capacity Scaling: Running Time

Lemma 2. Let f be the flow at the end of a Δ-scaling phase. Then value of the maximum flow is at most $v(f) + m\,\Delta$.

Pf. (similar to proof of max-flow min-cut theorem)

- Let A be the set of nodes reachable from s in $G_f(\Delta)$; B = V-A
- By construction, $s \in A$, $t \in B$, so A, B is a cut
- We show that $cap(A, B) \leq v(f) + m\,\Delta$
- Key points: among edges crossing the cut
  - forward edges > Δ: saturated;
  - backward edges > Δ: empty



original network

$$\begin{aligned}
v(f) &= \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) \\
&\geq \sum_{e \text{ out of } A} (c(e) - \Delta) - \sum_{e \text{ in to } A} \Delta \\
&= \sum_{e \text{ out of } A} c(e) - \sum_{e \text{ out of } A} \Delta - \sum_{e \text{ in to } A} \Delta \\
&\geq cap(A, B) - m\Delta
\end{aligned}$$

Section 7.4: Preflow-Push

Goldberg/Goldberg-Tarjan 1986

Most (all?) prior methods based on augmenting paths

Preflow-Push is a fundamentally different idea

Read 7.4; just a sketch here!

Preflow-Push Sketch

Idea 1: ignore "flow conservation"; push as much flow as possible to neighbors of s, their neighbors, etc. (a "preflow")

excess(v) = flow_in(v) – flow_out(v); always ≥ 0

Idea 2: incrementally adjust flow to shed excess

Idea 3: nodes have *height*; push excess *downhill* (only)

h(s) = n; h(t) = 0; initially h(all others) = 0, but may rise

if excess(v) > 0 & downhill residual edge: push flow

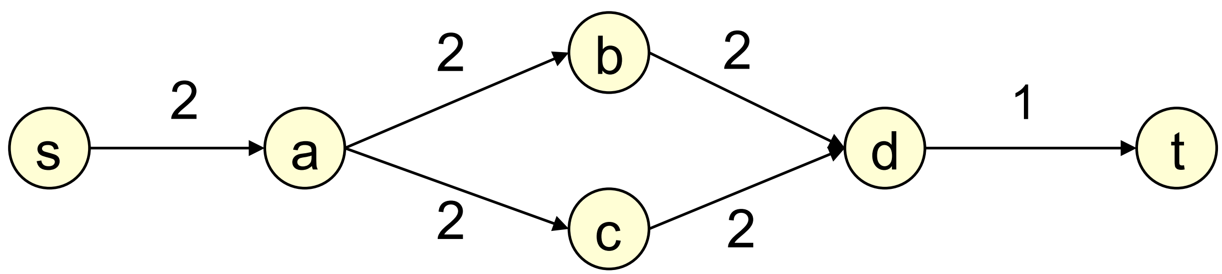
if excess(v) > 0 & no downhill residual edge: h(v)++

Magically, this stops, quickly, & you have a max flow

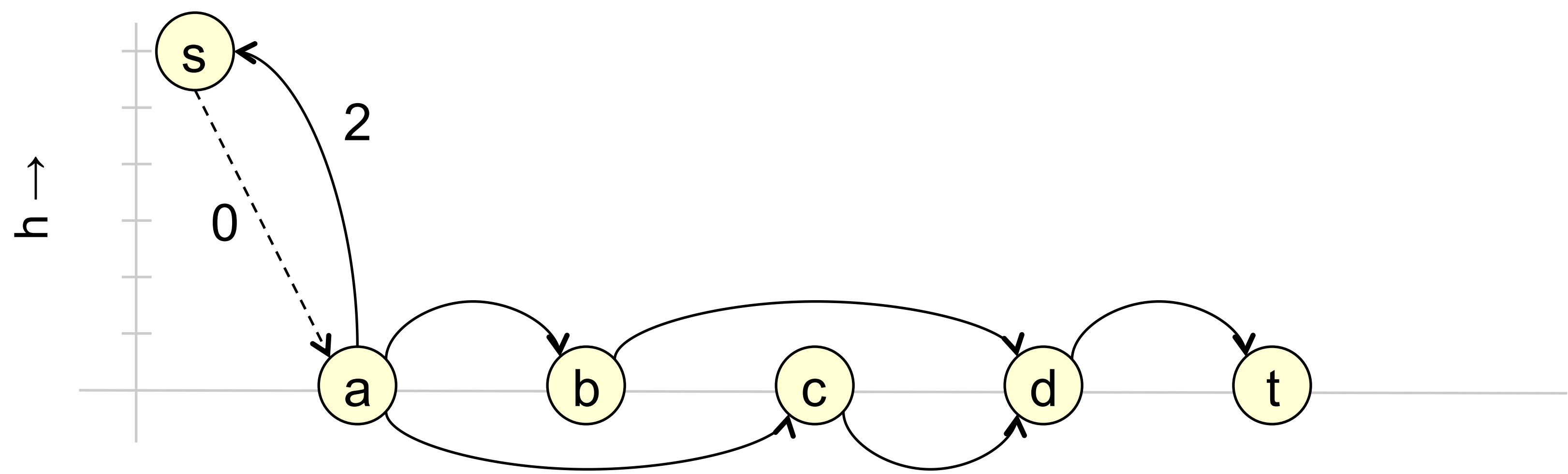
↳ $O(n^2m)$; $O(n^3)$ with "highest-first" rule

Example

graph:



initial preflow, heights, residual graph:



Flow Applications

Applications of Max Flow

Many!

Most look nothing like flow, at least superficially, but are deeply connected

Several interesting examples in 7.5-7.13

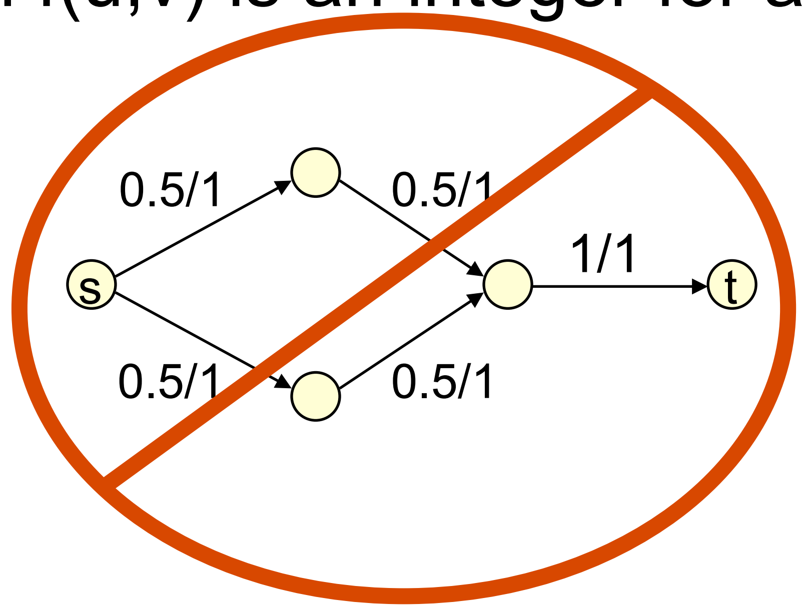
(7.8-7.11, 7.13 are optional, but interesting. Airline scheduling and image segmentation are especially recommended.)

A few more in following slides

Flow Integrality Theorem

Useful facts: If all capacities are integers

- Some max flow has an integer value
- Ford-Fulkerson method finds a max flow in which f(u,v) is an integer for all edges (u,v)



A valid flow,
but unnecessary

7.6: Disjoint Paths

Given a digraph with designated nodes s,t, are there k edge-disjoint paths from s to t?

You might try depth-first search; you might fail…

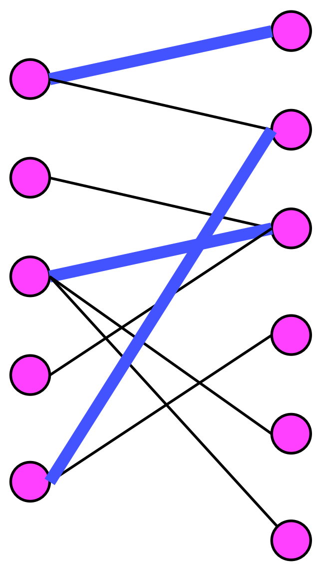
You might instead try "Is max flow $\geq$ k?" Success!

Max-flow/min-cut also implies max number of edge disjoint paths = min number of edges whose removal separates s from t.

Many variants: node-disjoint, undirected, …

See 7.6

7.5: Bipartite Maximum Matching



Bipartite Graphs:

- $G = (V,E)$
- $V = L \cup R \ \ (L \cap R = \emptyset)$
- $E \subseteq L \times R$

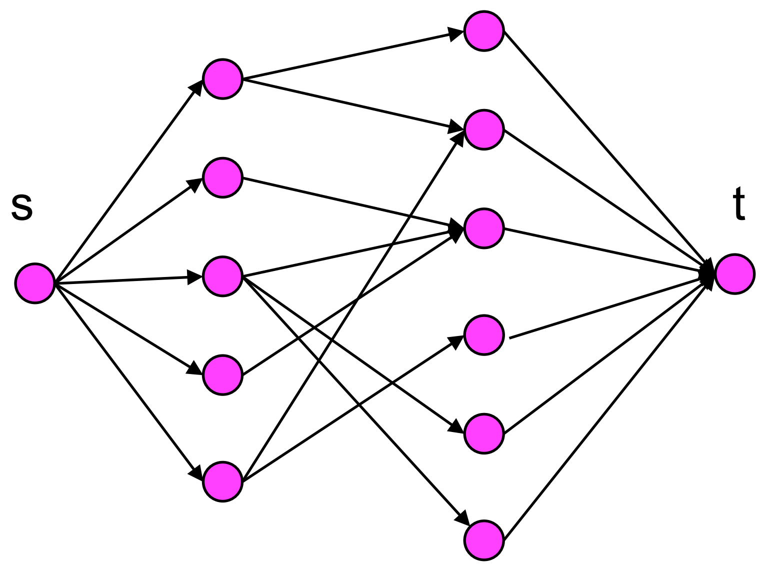
Matching:

- A set of edges $M \subseteq E$ such that no two edges touch a common vertex

Problem:

- Find a matching M of maximum size

Reducing Matching to Flow



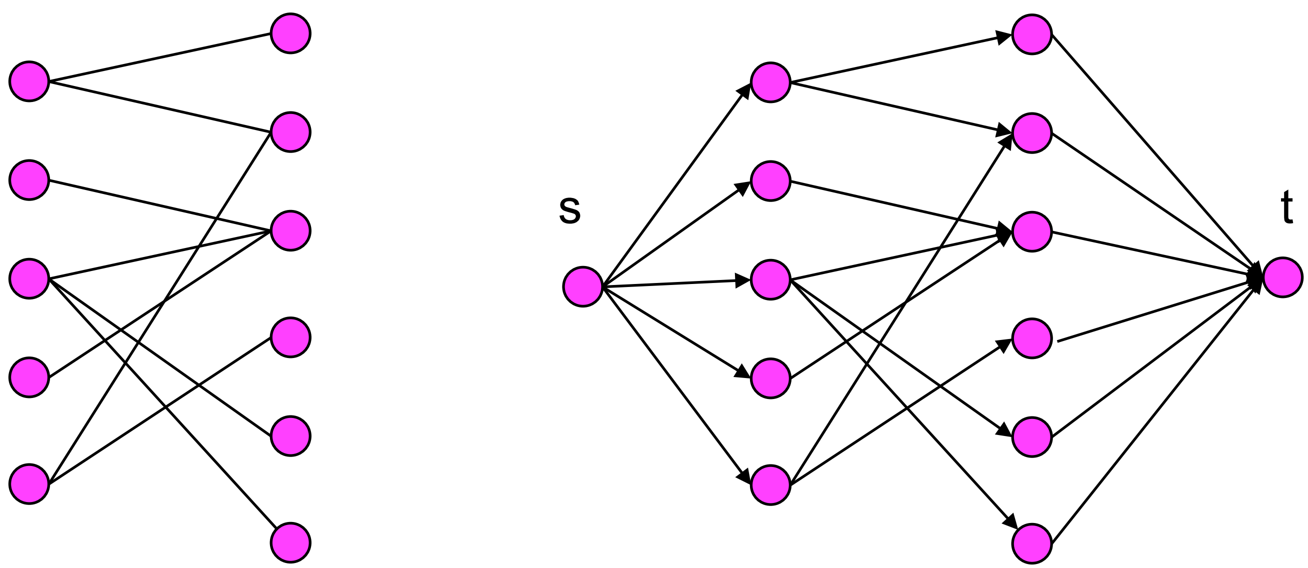
Given bipartite G, build flow network N as follows:

- Add source s, sink t
- Add edges s → L
- Add edges R→ t
- All edge capacities 1

Theorem:
Max flow iff
max matching

Reducing Matching to Flow

Theorem: Max matching size = max flow value



M → f? Easy – send flow only through M

f → M? Flow Integrality Thm, + cap constraints

Notes on Matching

- Max Flow Algorithm is probably overly general here
- But most direct matching algorithms use "augmenting path"-type ideas similar to that in max flow – See text (& homework?)
- Time $mn^{1/2}$ possible via Edmonds-Karp

7.12 Baseball Elimination

Some slides by Kevin Wayne

Baseball Elimination

Team i	Wins w_i	Losses l_i	To play g_i	Against = g_{ij}			
				Atl	Phi	NY	Mon
Atlanta	83	71	8	-	1	6	1
Philly	80	79	3	1	-	0	2
New York	78	78	6	6	0	-	0
Montreal	77	82	3	1	2	0	-

Which teams have a chance of finishing the season with most wins?

- Montreal eliminated since it can finish with at most 80 wins, but Atlanta already has 83.
- $w_i + g_i < w_j \Rightarrow$ team i eliminated.
- Only reason sports writers appear to be aware of.
- Sufficient, but not necessary!

Baseball Elimination

Team i	Wins w_i	Losses l_i	To play g_i	Against = g_{ij}			
				Atl	Phi	NY	Mon
Atlanta	83	71	8	-	1	6	1
Philly	80	79	3	1	-	0	2
New York	78	78	6	6	0	-	0
Montreal	77	82	3	1	2	0	-

Which teams have a chance of finishing the season with most wins?

- » Philly can win 83, but still eliminated . . .
- » If Atlanta loses a game, then some other team wins one.

Remark. Depends on *both* how many games already won and left to play, *and* on which opponents.

Baseball Elimination

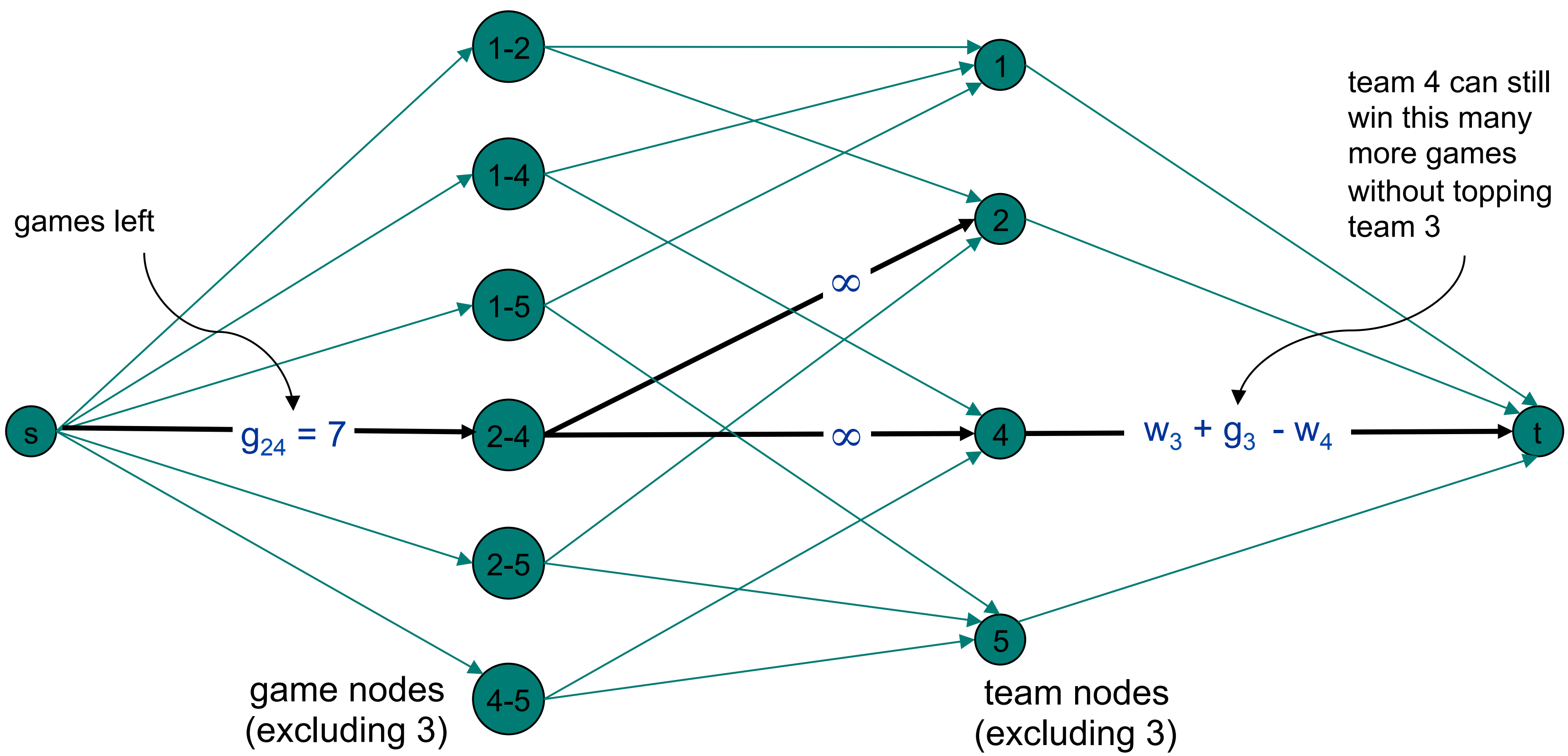
Baseball elimination problem.

- Set of teams S.
- Distinguished team $s \in S$.
- Team x has won w_x games already.
- Teams x and y play each other g_{xy} additional times.
- Is there any outcome of the remaining games in which team s finishes with the most (or tied for the most) wins?

Baseball Elimination: Max Flow Formulation

Can team 3 finish with most wins?

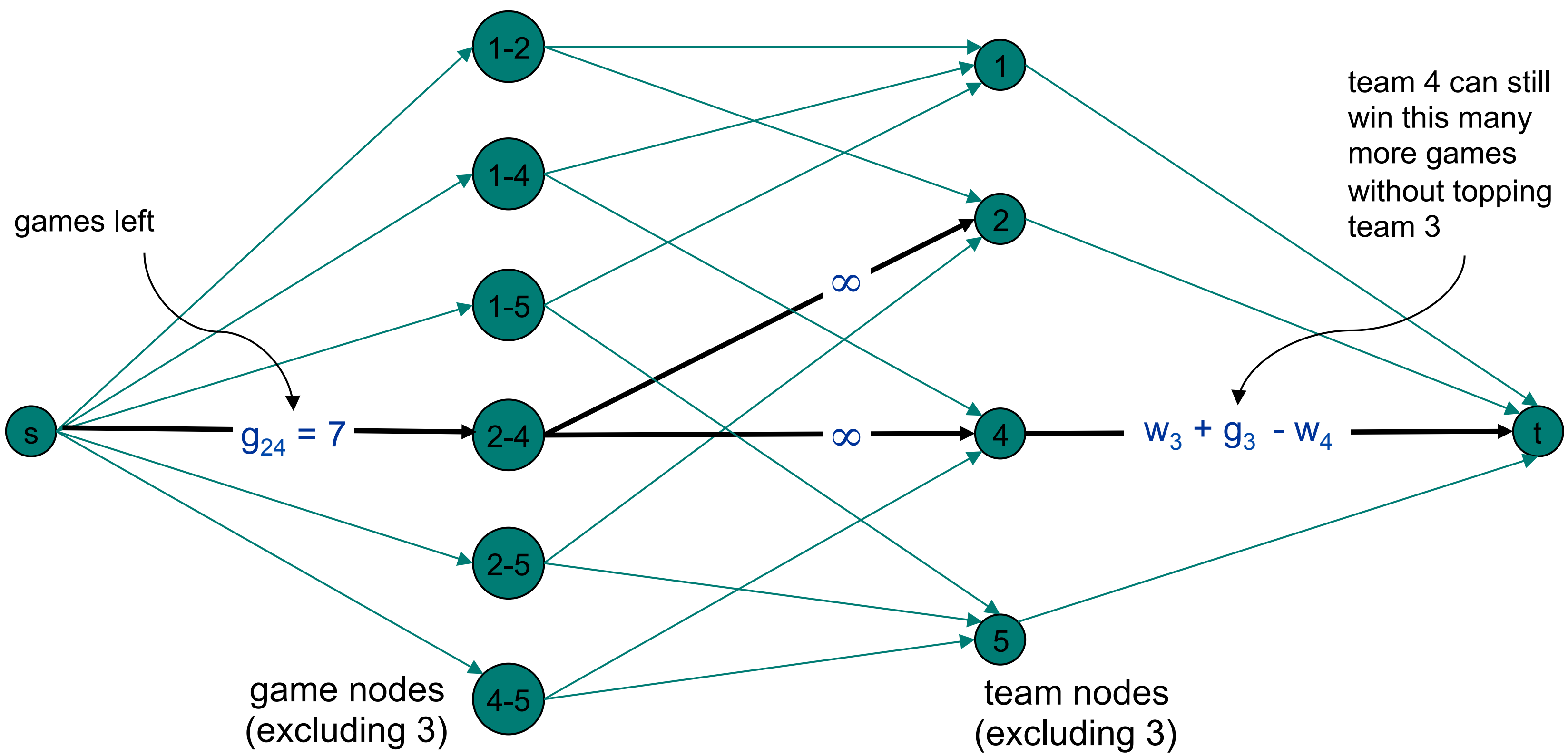
- Assume team 3 wins all remaining games $\Rightarrow$ $w_3 + g_3$ wins.
- Divvy remaining games so that all teams have $\leq$ $w_3 + g_3$ wins.



Baseball Elimination: As Max Flow

Theorem. Team 3 is eliminated iff max flow ≠ games left.

- Integrality ⇒ each remaining x-y game added to # wins for x or y.
- Capacity on (x, t) edges ensure no team wins too many games.
- In max flow, unsaturated source edge = unplayed game; if played, (either) winner would push ahead of team 3



Baseball Elimination: Explanation for Sports Writers

Team i	Wins w_i	Losses l_i	To play g_i	Against $=g_{ij}$				
				NY	Bal	Bos	Tor	Det
NY	75	59	28	-	3	8	7	3
Baltimore	71	63	28	3	-	2	7	4
Boston	69	66	27	8	2	-	0	0
Toronto	63	72	27	7	7	0	-	-
Detroit	49	86	27	3	4	0	0	-

AL East: August 30, 1996

Which teams have a chance of finishing the season with most wins?

Detroit could finish season with 49 + 27 = 76 wins.

Baseball Elimination: Explanation for Sports Writers

Team i	Wins w_i	Losses l_i	To play g_i	Against $= g_{ij}$				
				NY	Bal	Bos	Tor	Det
NY	75	59	28	-	3	8	7	3
Baltimore	71	63	28	3	-	2	7	4
Boston	69	66	27	8	2	-	0	0
Toronto	63	72	27	7	7	0	-	-
Detroit	49	86	27	3	4	0	0	-

AL East: August 30, 1996

Which teams could finish the season with most wins?

- Detroit could finish season with 49 + 27 = 76 wins.

Certificate of elimination. R = {NY, Bal, Bos, Tor}

- Have already won w(R) = 278 games.
- Must win at least r(R) = 27 more.
- Average team in R wins at least 305/4 > 76 games.

Baseball Elimination: Explanation for Sports Writers

Certificate of elimination

$$T \subseteq S, \quad w(T) := \overbrace{\sum_{i \in T} w_i}^{\#\text{ wins}}, \quad g(T) := \overbrace{\sum_{\{x,y\} \subseteq T} g_{xy}}^{\#\text{ remaining games}},$$

If $\overbrace{\dfrac{w(T) + g(T)}{|T|}}^{\text{LB on avg \# games won}} > w_z + g_z$ then z eliminated (by subset T).

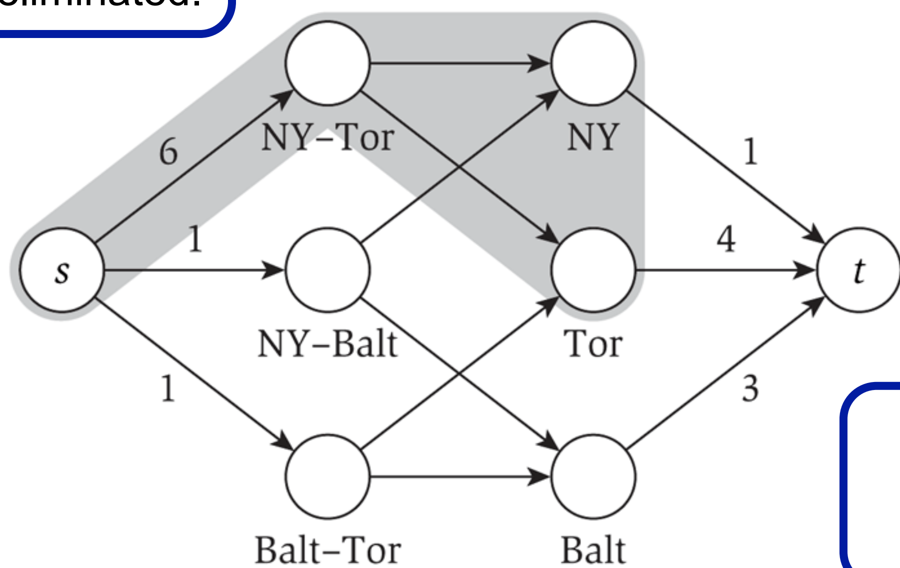
Theorem. [Hoffman-Rivlin 1967] Team z is eliminated iff there exists a subset T* that eliminates z.

Proof idea. Let T* = teams on source side of min cut.

	w	l	g	NY	Balt	Tor	Bos
NY	90		11	-	1	6	4
Baltimore	88		6	1	-	1	4
Toronto	87		10	6	1	-	4
Boston	79		12	4	4	4	-

g* = 1+6+1 = 8

(90 + 87 + 6) / 2 > 91,
so the set T = {NY, Tor}
proves Boston is eliminated.



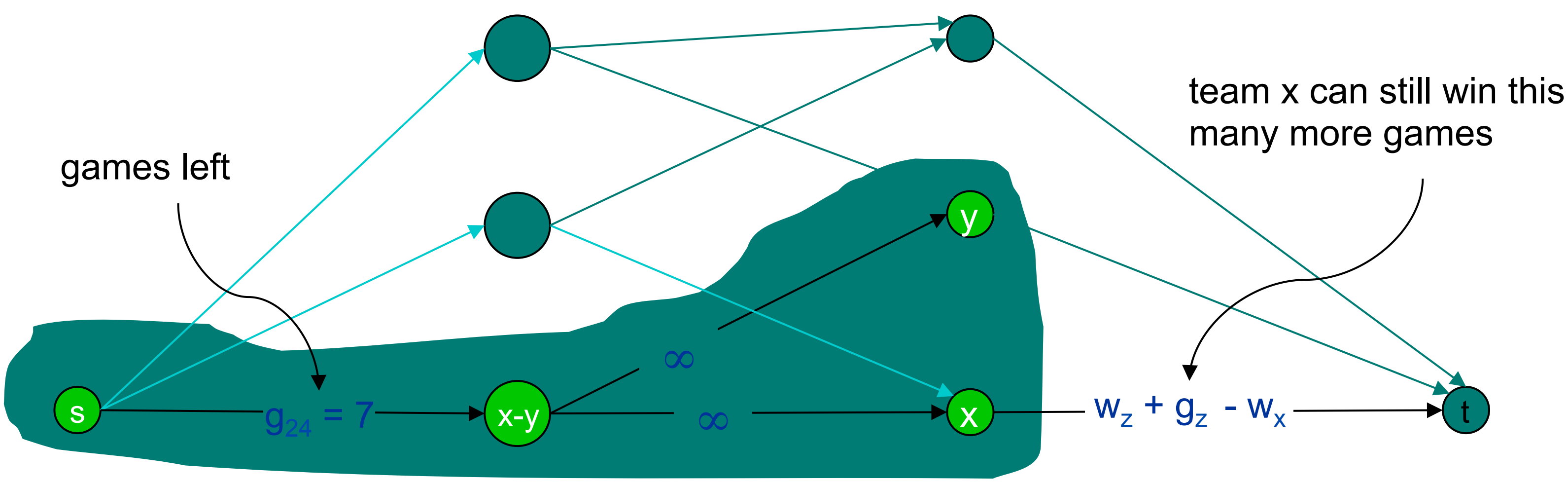
Note: T={NY,Tor, Balt} is NOT a certificate, since (90+88+87+8)/3 = 91

Fig 7.21 Min cut ⇒ no flow of value g*, so Boston eliminated.

Baseball Elimination: Explanation for Sports Writers

Pf of theorem.

- Use max flow formulation, and consider min cut (A, B).
- Define T* = team nodes on source side of min cut.
- Observe x-y $\in$ A iff both x $\in$ T* and y $\in$ T*.
  - infinite capacity edges ensure if x-y $\in$ A then x $\in$ A and y $\in$ A
  - if x $\in$ A and y $\in$ A but x-y $\notin$ T*, then adding x-y to A decreases capacity of cut



Baseball Elimination: Explanation for Sports Writers

Pf of theorem.

Use max flow formulation, and consider min cut (A, B).

Define T* = team nodes on source side of min cut.

Observe x-y ∈ A iff both x ∈ T* and y ∈ T*.

$$g(S-\{z\}) \quad > \quad cap(A,\, B)$$

$$= \quad \overbrace{g(S-\{z\}) - g(T^*)}^{\text{capacity of game edges leaving A}} \; + \; \overbrace{\sum_{x \in T^*} (w_z + g_z - w_x)}^{\text{capacity of team edges leaving A}}$$

$$= \quad g(S-\{z\}) - g(T^*) \; - \; w(T^*) \; + \; |T^*|(w_z + g_z)$$

Rearranging:

$$w_z + g_z \; < \; \frac{w(T^*) + g(T^*)}{|T^*|}$$

Matching & Baseball: Key Points

Can (sometimes) take problems that seemingly have *nothing* to do with flow & reduce them to a flow problem

How? Build a clever network; map allocation of stuff in original problem (match edges; wins) to allocation of flow in network. Clever edge capacities constrain solution to mimic original problem in some way. Integrality useful.

Matching & Baseball: Key Points

Furthermore, in the baseball example, min cut can be translated into a succinct *certificate* or *proof* of some property that is much more transparent than “see, I ran max-flow and it says flow must be less than g*”.

These examples suggest why max flow is so important – *it’s a very general tool used in many other algorithms*.