CSE 521: Design and Analysis of Algorithms

Assignment #3

due: Tuesday, Oct 25. 10:30AM

Each problem is worth 10 points. KT refers to Algorithm Design, First Edition, by Kleinberg and Tardos. "Give an algorithm" means pseudo-code, a high-level explanation and a proof of correctness. See the website for more grading guidelines.

1. Convergence of the gradient descent algorithm for solving Ax = b: Consider the following algorithm:

Algorithm 1 Gradient descent algorithm for solving linear systems of equations

- 1: $k \leftarrow 0$.
- 2: $x_0 \leftarrow 0$.
- 3: repeat

- $\begin{aligned} & \overset{\cdot}{r_k} \leftarrow b Ax_k \\ & \alpha_k \leftarrow \frac{r_k^T r_k}{r_k^T A r_k}. \\ & x_{k+1} \leftarrow x_k + \alpha_k r_k. \end{aligned}$
- $k \leftarrow k + 1$.
- 8: until $\|\alpha_k\| < 10^{-20}$
 - (a) Prove that $r_{k+1}^T r_k = 0$ for all k.
 - (b) Define $d_k = A^{-1}b x_k = A^{-1}r_k$. Define $\delta_k = d_k^T A d_k$ to be a measure of error. Prove that

$$\delta_{k+1} \le (1 - 1/\kappa)\delta_k$$
.

Here $\kappa := ||A|| \cdot ||A^{-1}||$ is the condition number of A, and $||M|| := \max_{z \neq 0} \frac{z^T M z}{z^T z}$.

- 2. KT, Chapter 5, Problem 1
- 3. KT, Chapter 5, Problem 4
- 4. KT, Chapter 5, Problem 5
- 5. For two sets X, Y of integers, the Minkowski sum X + Y is the set of all pairwise sums $\{x + y | x \in Y\}$ $X, y \in Y$. The goal of this problem is to compute |X + Y|; that is, the number of elements in X + Y. Let n = |X| = |Y| and assume that all elements of X, Y are between 0 and M. Further assume that M is small enough so that adds, multiplies, etc of $O(\log M)$ -bit numbers takes constant time.
 - (a) Describe an algorithm to compute |X + Y| in time $O(n^2 \log(n))$.
 - (b) Describe an algorithm to compute |X + Y| in time $O(M \log(M))$.

- (c) For k a positive integer, define $kX = X + X + \cdots + X$. Describe an algorithm to compute |kX|in time $O(kM \log(kM))$.
- (d) Extra credit: Let L = |kX|. Describe a randomized algorithm to compute |kX| with $\geq 2/3$ probability of success in time $O(L^2 \log(L))$.