

## Bridging

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Bridging makes networks:

- Bigger
- Smaller

(and it's not a multiple choice question)

## Bridges, Spanning Trees & Switches

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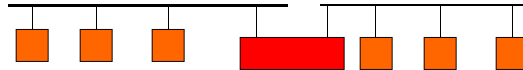
- Focus:
  - What to do when one shared LAN isn't big enough?
- Interconnecting LANs
  - Hubs
  - LAN bridges/switches
  - A preview of the Network layer

Application
Presentation
Session
Transport
Network
Data Link
Physical

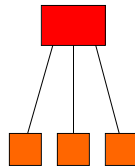
## Warning: confusing terminology / pictures

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Original Ethernet (repeater)



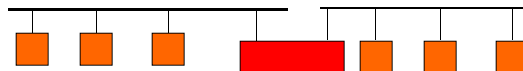
Modern Ethernet  
(Hub)



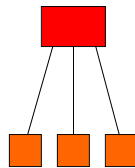
Multi-port "repeaters" –  
Hubs

## That is not all....

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Bridges – well this  
is historical too ...



## Why do we need a bridge/switch?

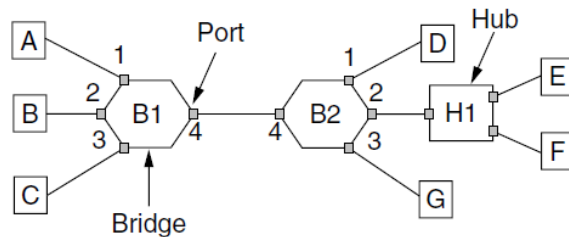
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- One shared LAN can limit us in terms of:
  - Distance (*why?*)
  - Number of nodes (*why?*)
- How do we scale to a larger, more efficient networks?
  - We must be able to interconnect LANs

## Bridges

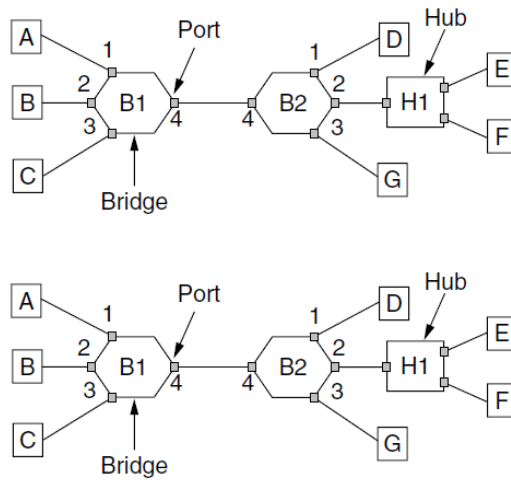
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- Connected to different LANs over *ports*
- Operate in “promiscuous mode”
  - receives packet on one port and forwards it to the outgoing port
    - Example. From A to D
  - is “never” a communication endpoint itself
- What should each bridge do with an incoming frame?



## Bridges: What to do with incoming frame

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## How are bridges configured?

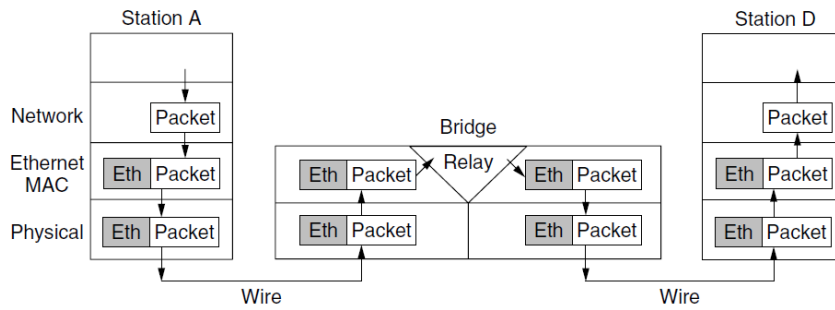
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- Backward learning algorithm
  1. At first, bridge doesn't know where any address is. Send incoming frame "everywhere"
  2. But, the incoming frame carries a source address. Remember it.
  3. If an incoming frame has a destination address I've remembered, send it only there
- Once we learn the ports, can we stop the backward learning algorithm?

## Learning Bridges

Bridges extend the Link layer:

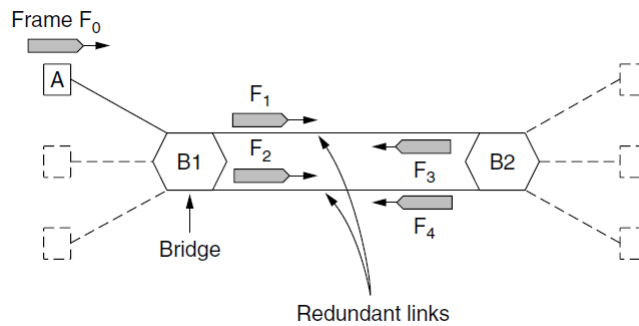
- Use but don't remove Ethernet header/addresses
- Don't inspect higher layer headers
- May need to store the packet



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## Adding (Un)Reliability

Redundant links



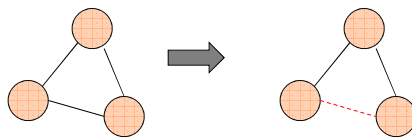
## Okay, just don't do that...

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## Spanning Tree Algorithm

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- Goal: prevent forwarding loops, even when there are physical loops, by having only a single bridge responsible for each LAN segment
  - “turn off” many bridge ports
  - Do this dynamically
- Results in a spanning tree for packet routing
  - Robust against failures, needs no organization

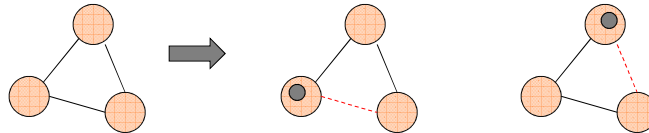


Llan.12

## Algorithm Overview

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1. Elect a root node for the tree
  - Remember that all bridges are doing this at once, and none know what the network topology is
2. Grow tree as shortest distances from the root
  - Remember that all bridges...
3. Add some details to actually make it work (!)



Llan.13

## Algorithm continued

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- Each bridge sends periodic messages to others containing:
  - [ its address,  
address of the bridge it believes is the root,  
distance (in hops) to root ]
- A bridge receiving a message may update its view of the “best” bridge on the LAN segment it received it from
  - First criterion: Smaller root address is preferred
  - Second criterion: Smaller distance to root is pre
  - Tie breaker: bridge with smaller address is better

djw // CSE 461, Fall 2009

Llan.14

## Algorithm continued

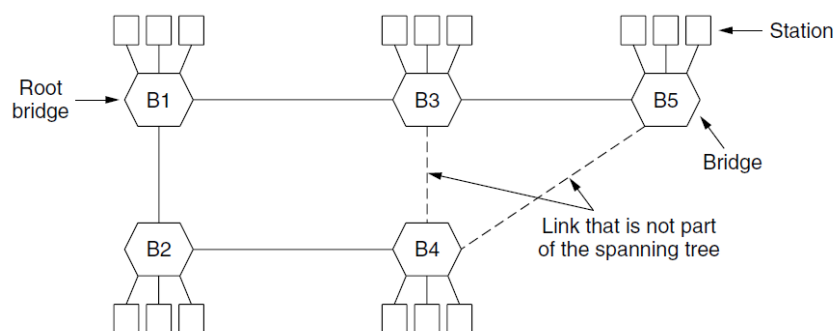
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- Initial condition: each bridge thinks it is the root
  - Sends configuration messages on all ports
    - [ I'm me,  
the root is me,  
I'm at distance 0 from the root ]
- Later, bridges send only “best” configs
  - Add 1 to distance, send configs on segments where it's still “best” (designated bridge)
  - Turn off forwarding on ports except those that send/receive “best”

Llan.15

## Spanning Tree Example: Final Tree

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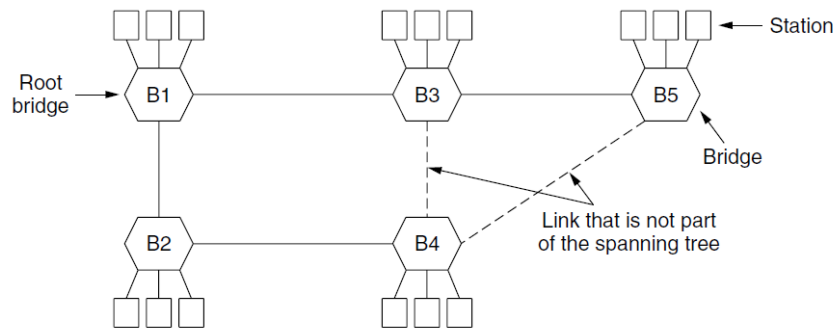


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## Spanning Tree Example: Transients

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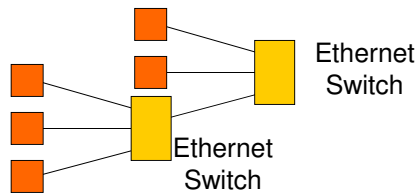
## General Design Principles

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- All bridges to run the same algorithm
- All bridges start with zero information and operate in parallel
- Bridges send periodic messages about their own state
- State that hasn't been refreshed in a while is deleted (*soft state*)
- The algorithm converges to a globally consistent state
  - If we all have the same inputs and have implemented it correctly and the inputs aren't changing too faster than we can converge

## LAN Switches – what is used today

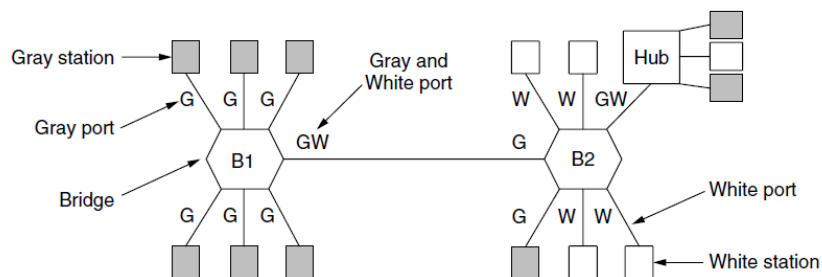
- LAN switches are multi-port bridges
  - Modern, high performance form of bridged LANs
  - Frames are switched, not shared like hubs
  - Every host on a separate port, can combine switches



## Virtual LANs

VLANs (Virtual LANs) split one physical LAN into multiple logical LANs, for management reasons

- Ports are “colored” according to their VLAN

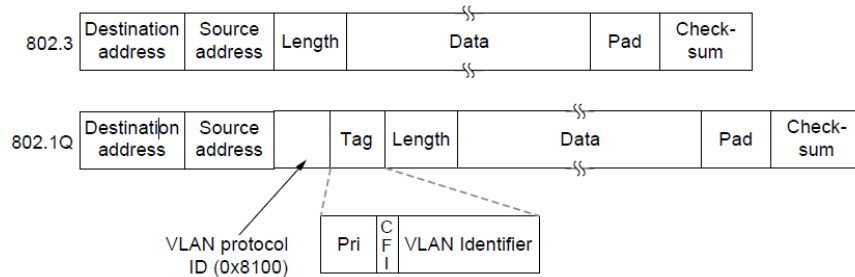


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## Virtual LANs -- IEEE 802.1Q

802.1Q frames carry a color tag (VLAN identifier)

- Length/Type value is 0x8100 for VLAN protocol



## VLAN Breaking News

The CSE wireless network used to be supported by CSE owned APs.

UW-IT has replaced our APs with theirs.

We'll still run an isolated CSE wireless network, and they'll run an isolated UW wireless network, both on their APs.

VLANs...

## Perspective on Bridges/Switches

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- LAN switches form an effective small-scale network
  - Designed to work transparently with no host changes
  - Plug and play for real!
- But they **don't scale well** to large networks
  - Little control over forwarding paths
  - Size of bridge forwarding tables grows with number of hosts
  - Broadcast traffic flows freely over whole extended LAN
  - Spanning tree algorithm limits reconfiguration speed
- We can layer on scalability by changing hosts too.
  - Next up: IP and routing.