

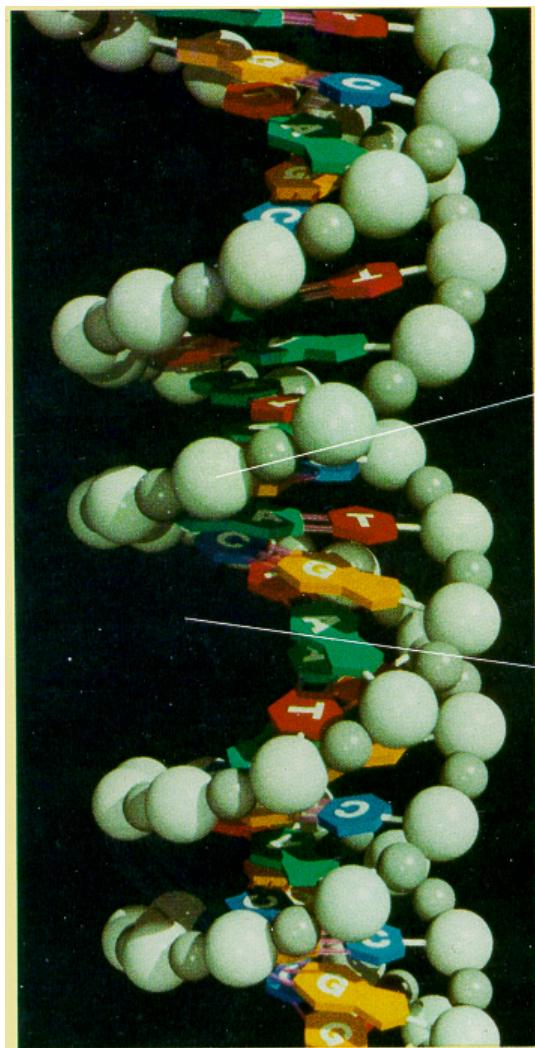
CSE 417: Algorithms and Computational Complexity

Winter 2009

W. L. Ruzzo

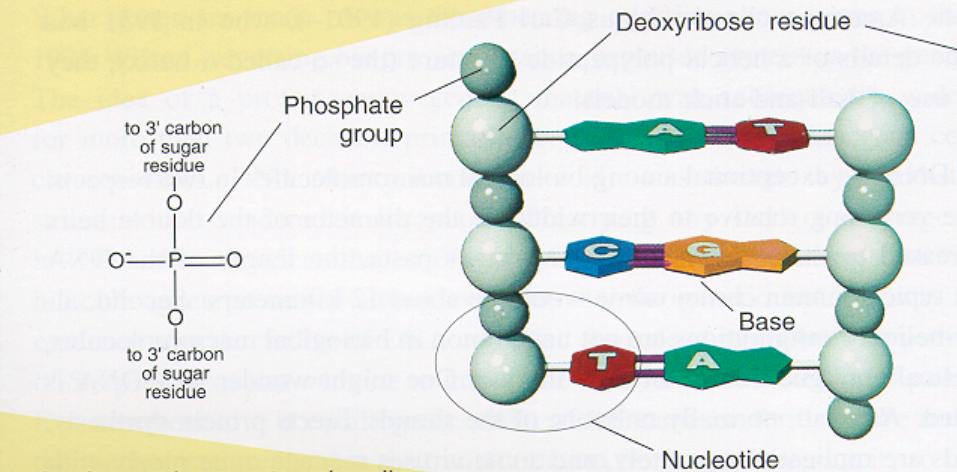
Dynamic Programming, II
RNA Folding

The Double Helix



(a) Computer-generated Image of DNA
(by Mel Prueitt)

(b) Uncoiled DNA Fragment



As shown, the two strands coil about each other in a fashion such that all the bases project inward toward the helix axis. The two strands are held together by hydrogen bonds (pink rods) linking each base projecting from one backbone to its so-called complementary base projecting from the other backbone. The base A always bonds to T (A and T are comple-

Shown in (b) is an uncoiled fragment of (a three complementary base pair. In a chemist's viewpoint, each strand is a polymer made up of four repeating units called deoxyribonucleotides.

Los Alamos Science

NATURE VOL. 227 AUGUST 8 1970

Central Dogma of Molecular Biology

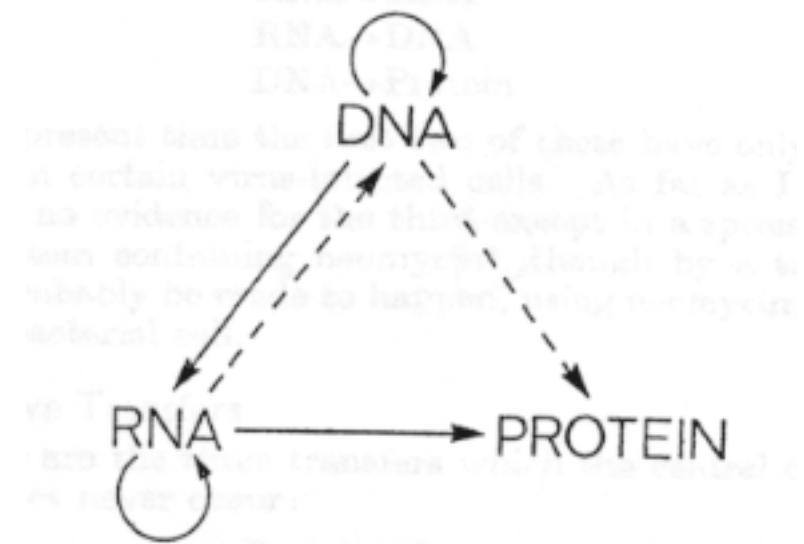
by

FRANCIS CRICK
MRC Laboratory
Hills Road,
Cambridge CB2 2QH

The central dogma of molecular biology deals with the detailed residue-by-residue transfer of sequential information. It states that such information cannot be transferred from protein to either protein or nucleic acid.

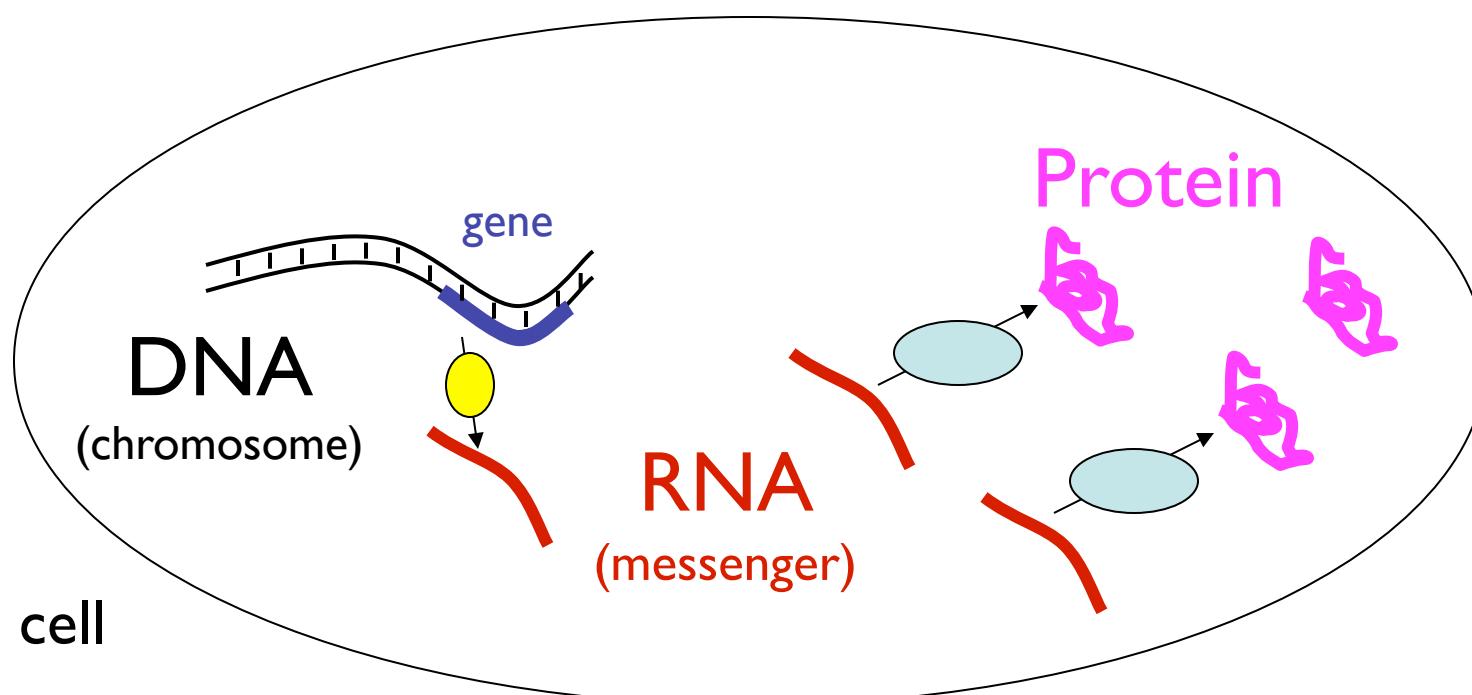
"The central dogma, enunciated by Crick in 1958 and the keystone of molecular biology ever since, is likely to prove a considerable over-simplification."

Fig. 2. The arrows show the situation as it seemed in 1958. Solid arrows represent probable transfers, dotted arrows possible transfers. The absent arrows (compare Fig. 1) represent the impossible transfers postulated by the central dogma. They are the three possible arrows starting from protein.



The “Central Dogma” of Molecular Biology

DNA → RNA → Protein



Non-coding RNA

Messenger RNA - codes for proteins

Non-coding RNA - all the rest

Before, say, mid 1990's, 1-2 dozen known
(critically important, but narrow roles)

Since mid 90's dramatic discoveries

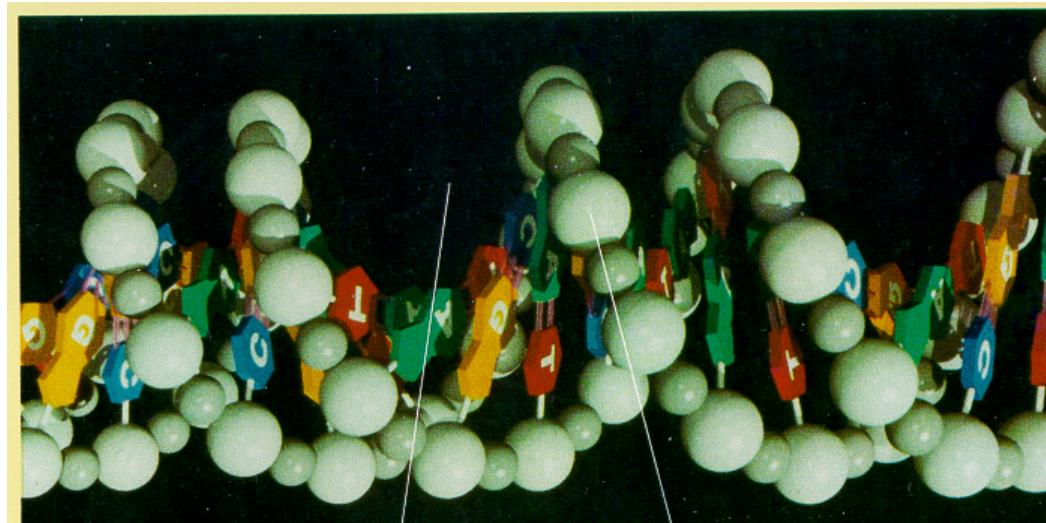
Regulation, transport, stability/degradation

E.g. "microRNA": 100s in humans => 50% of genes

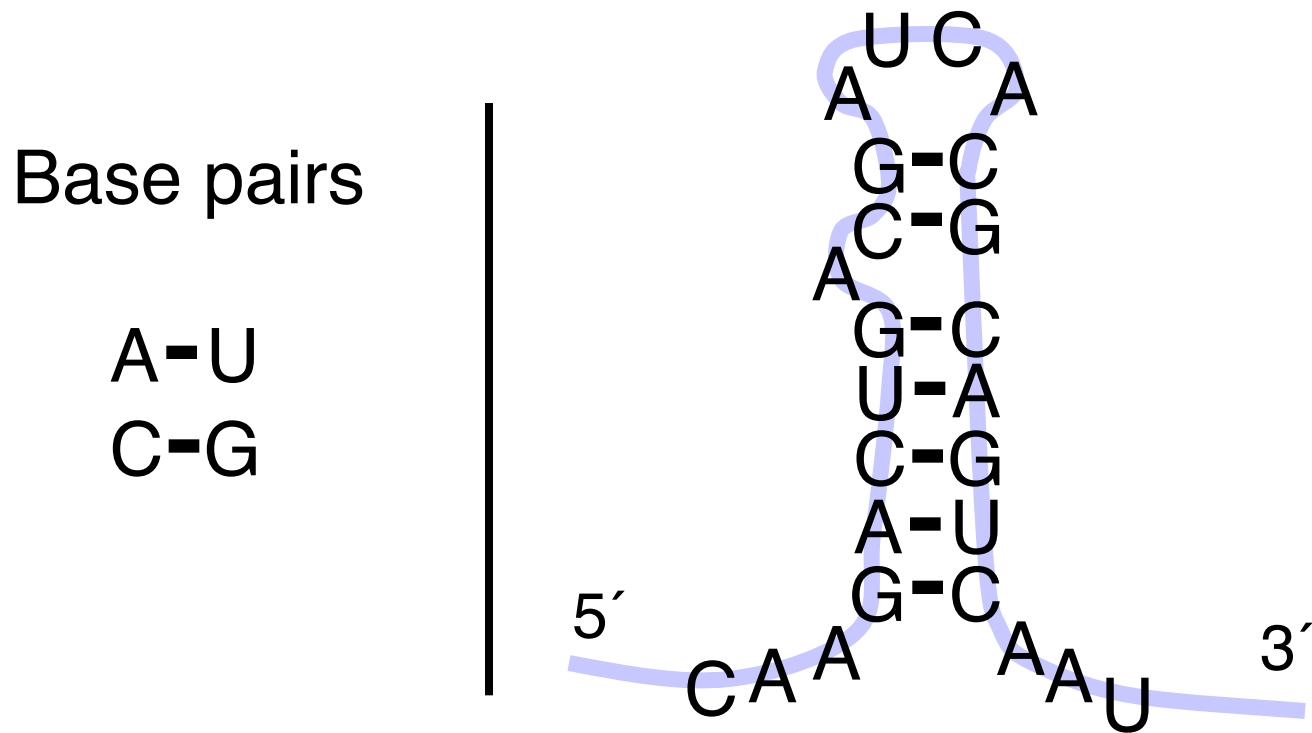
E.g. "riboswitches": 1000s in bacteria

DNA structure: dull

5' ...ACCGCTAGATG... 3'
 |||||
3' ...TGGCGATCTAC... 5'

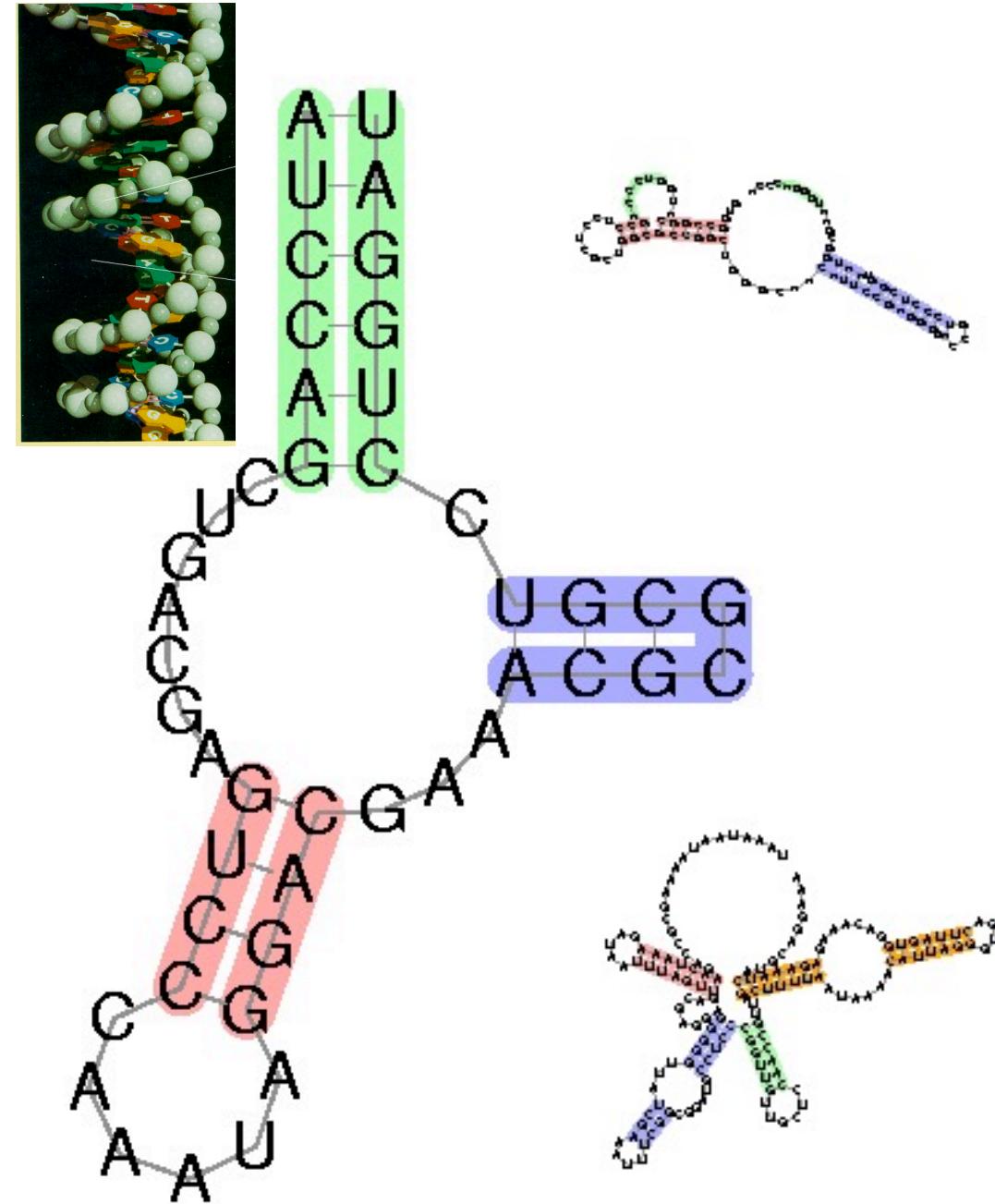


RNA Secondary Structure: RNA makes helices too

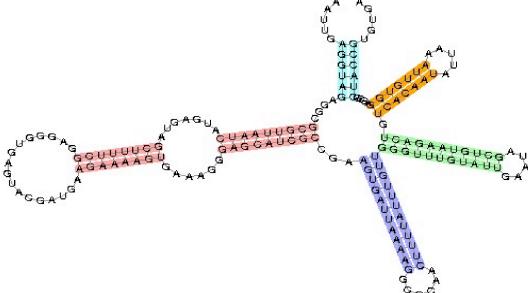


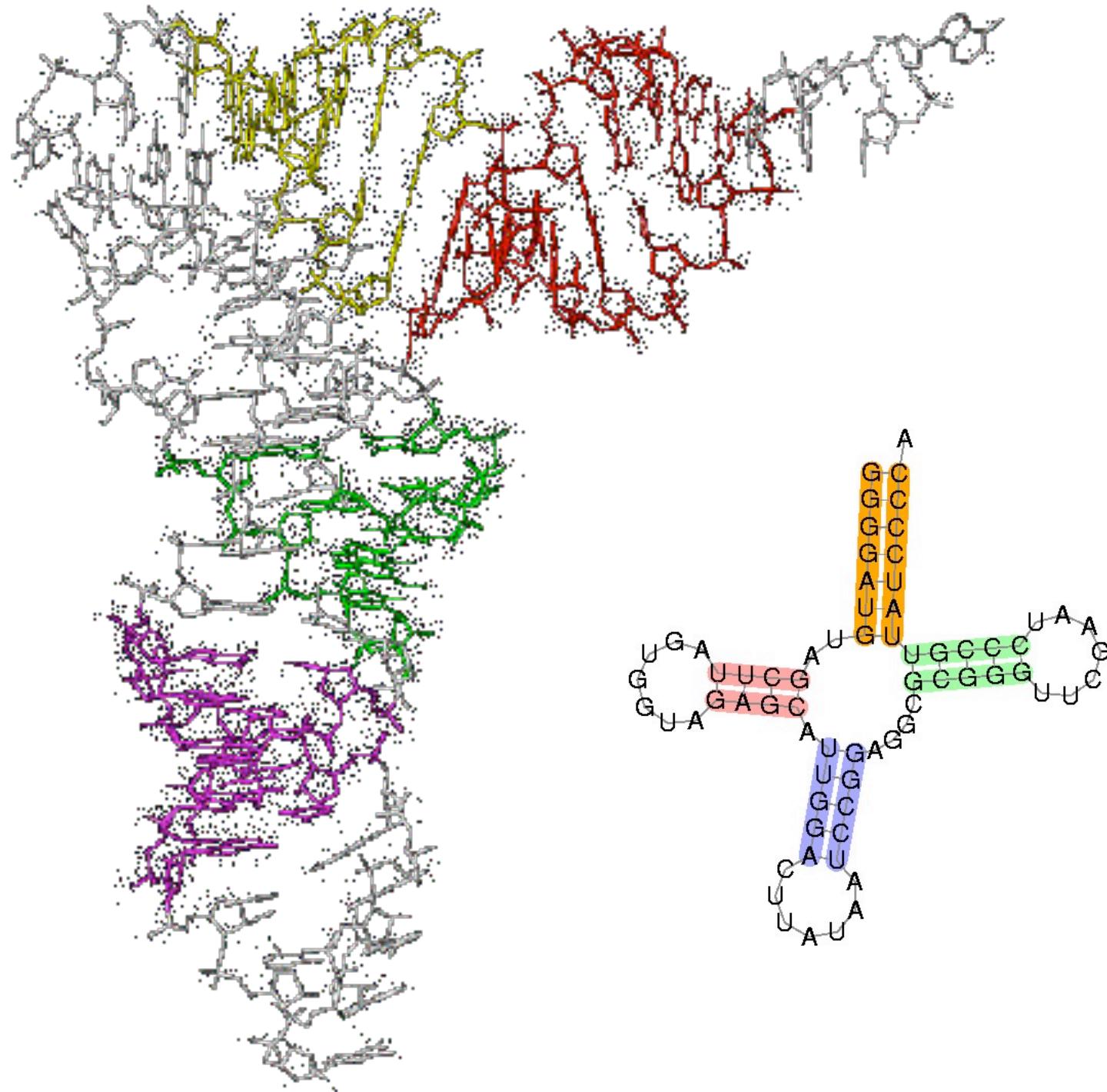
Usually *single* stranded

RNA Structure: Rich



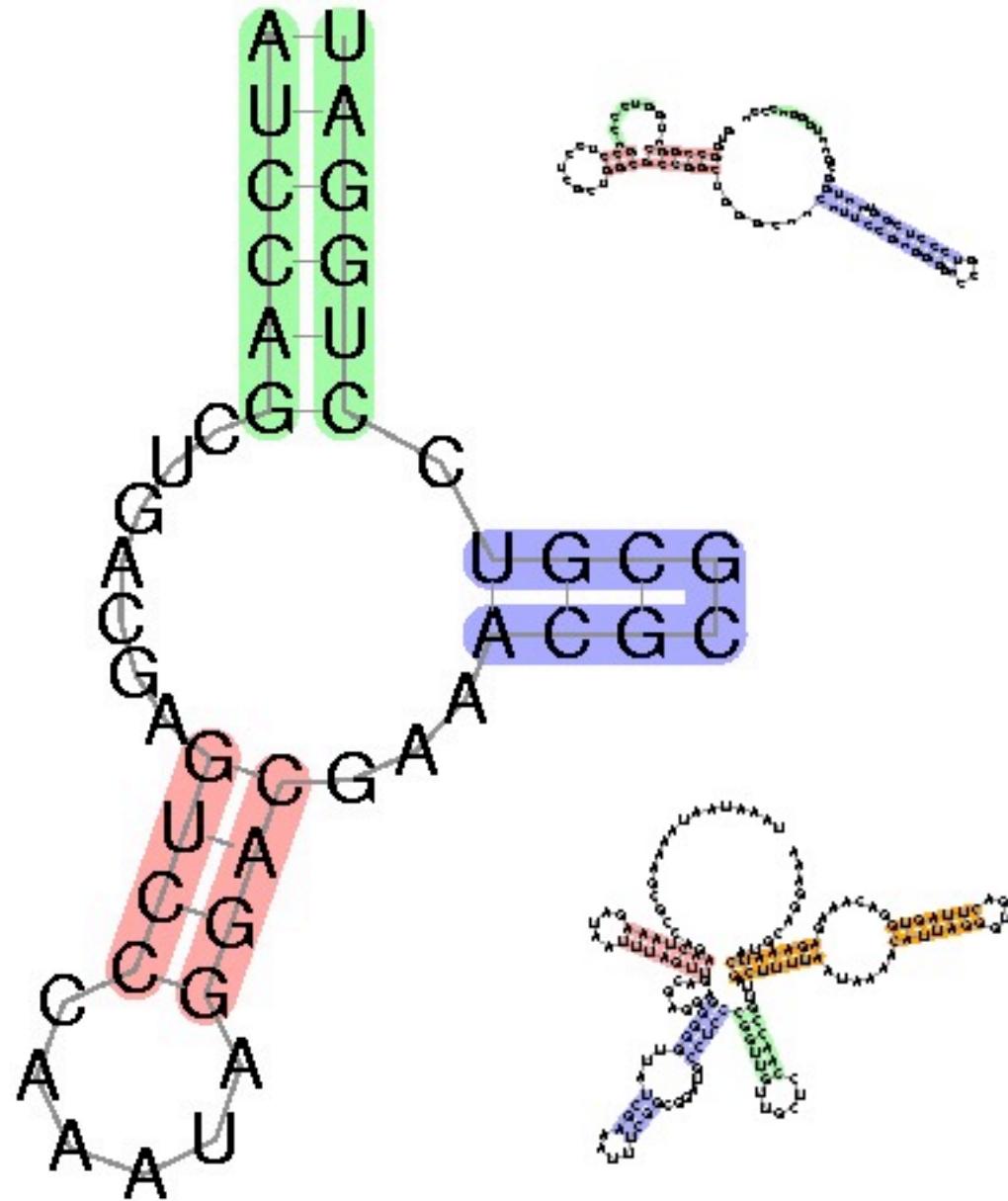
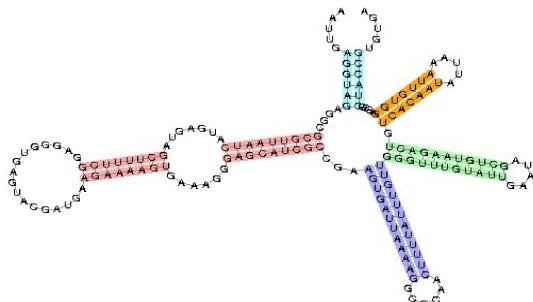
- RNA's fold, and function
- Nature uses what works





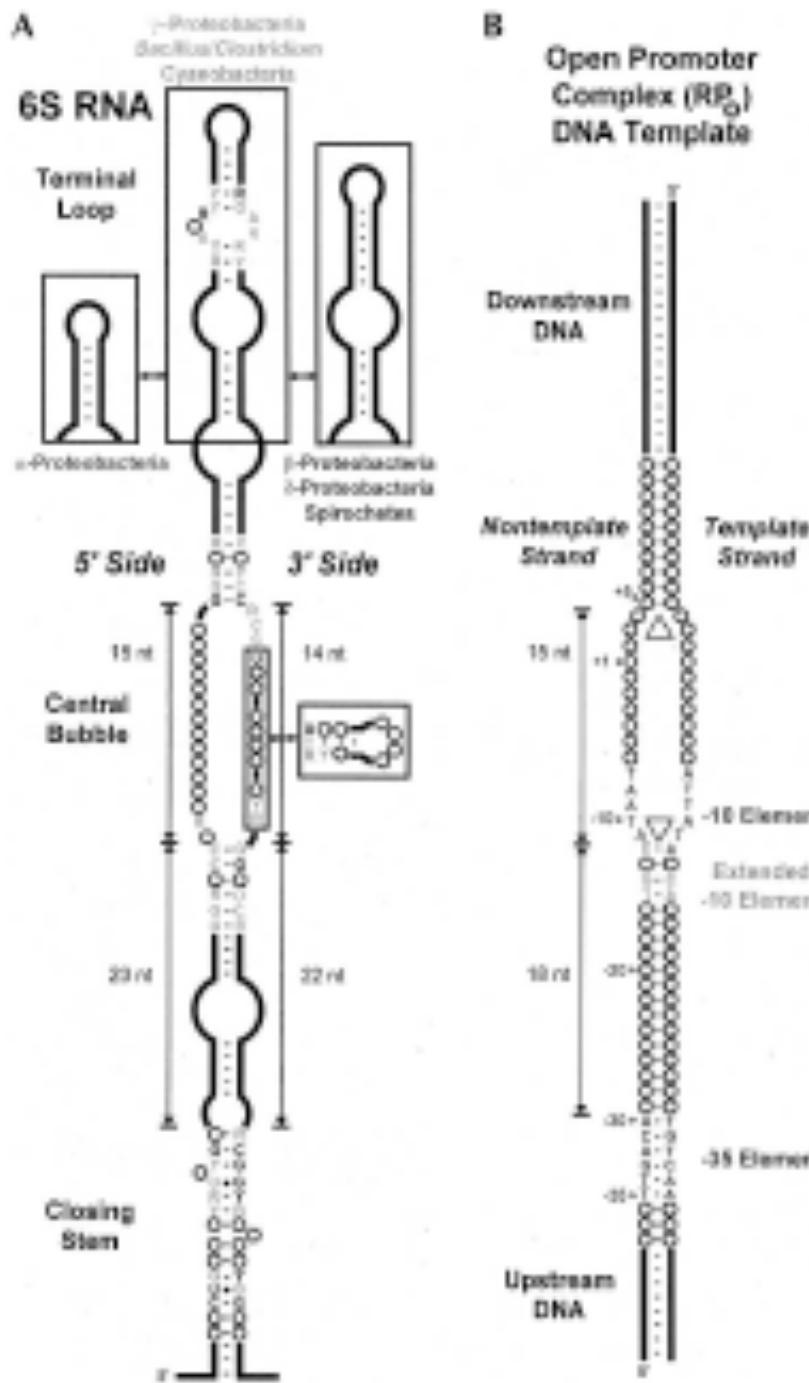
RNA Secondary Structure:

Not everything,
but important,
easier than 3d

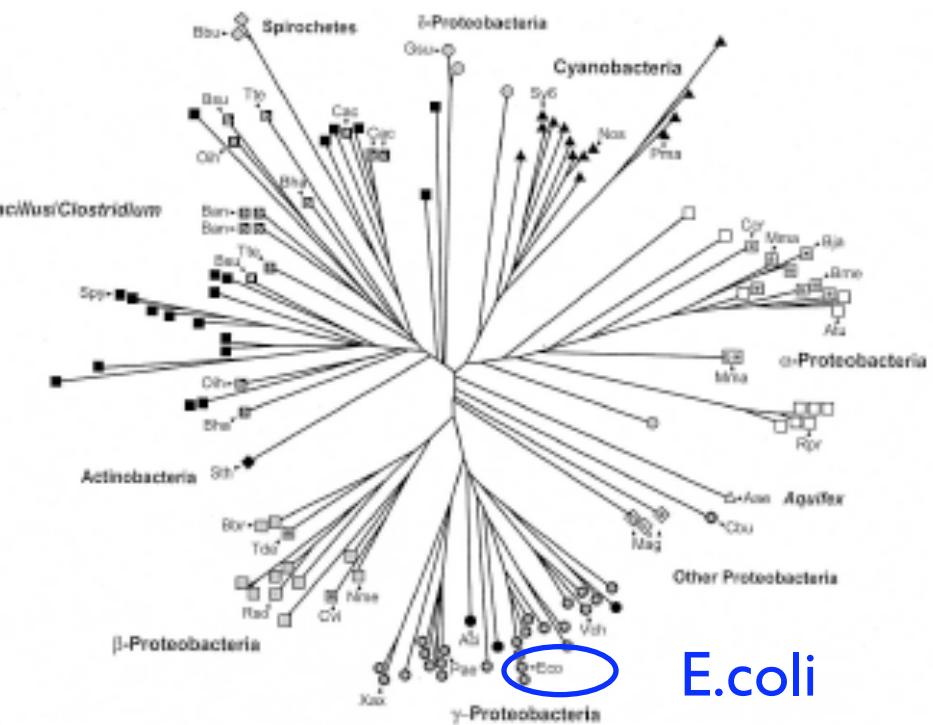


Why is structure important?

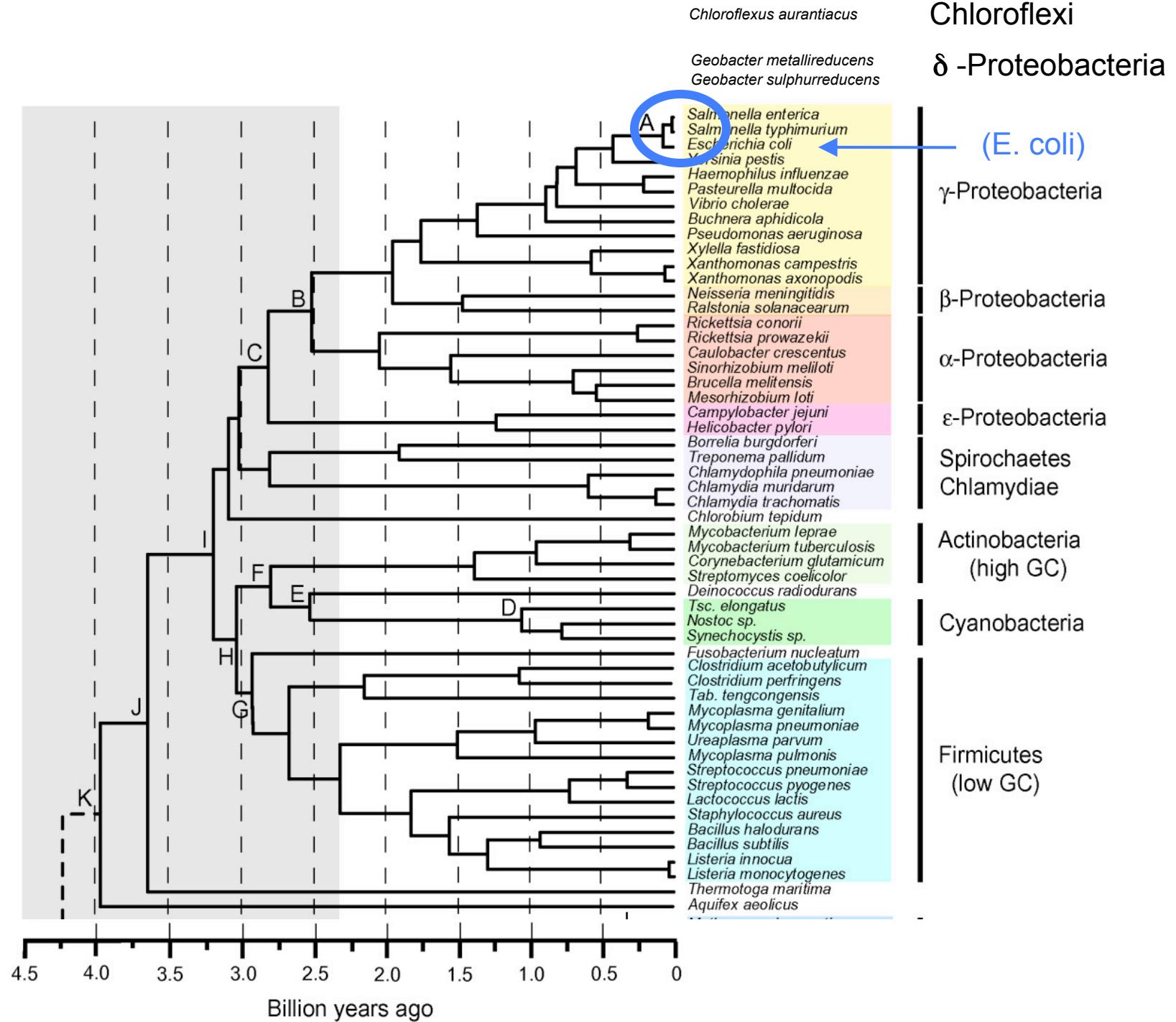
- For protein-coding, similarity in sequence is a powerful tool for finding related sequences
 - e.g. “hemoglobin” is easily recognized in all vertebrates
- For non-coding RNA, many different sequences have the same structure, and structure is most important for function.
 - So, using structure plus sequence, can find related sequences at much greater evolutionary distances



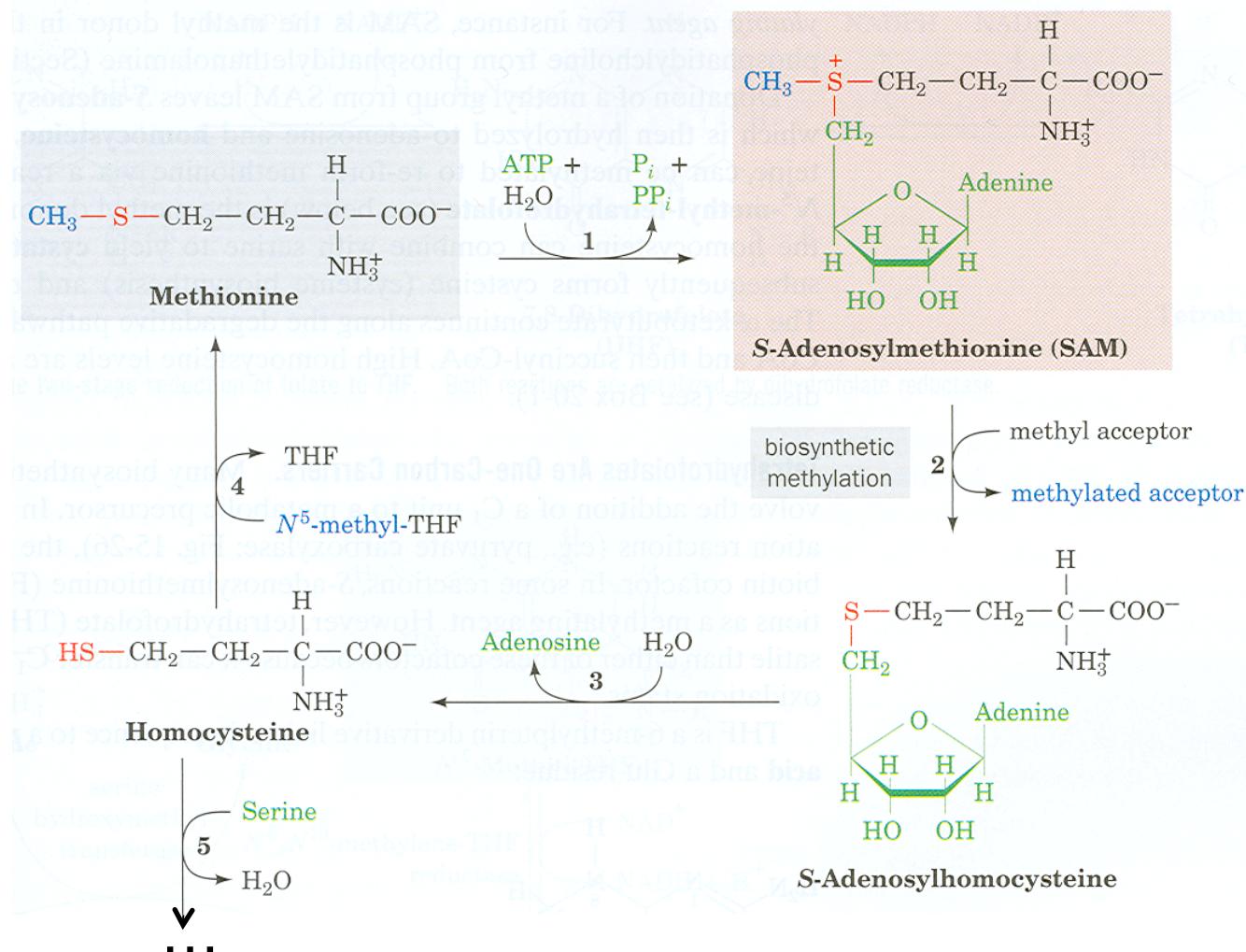
6S mimics an open promoter



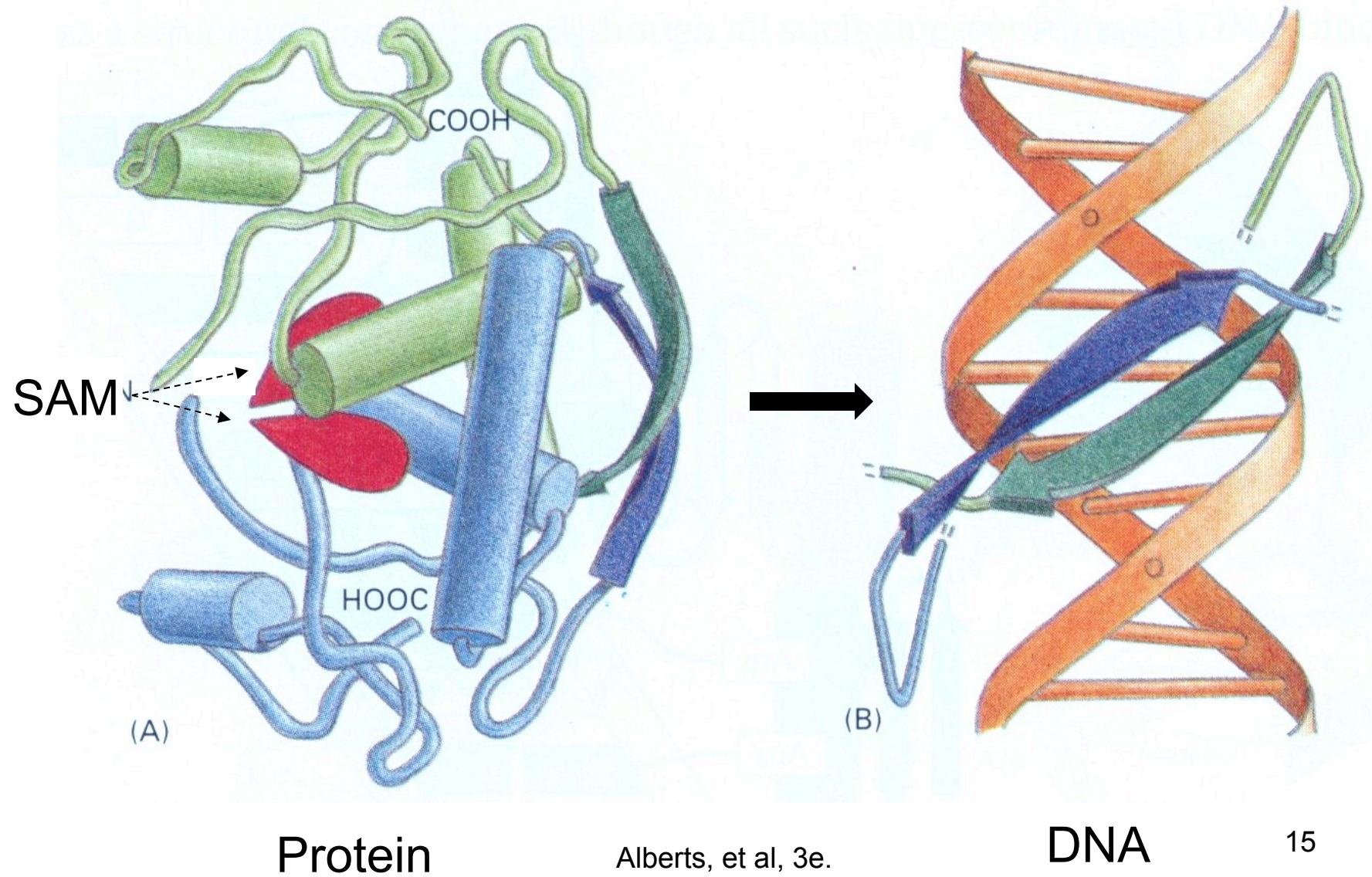
Barrick et al. RNA 2005
Trotochaud et al. NSMB 2005
Willkomm et al. NAR 2005

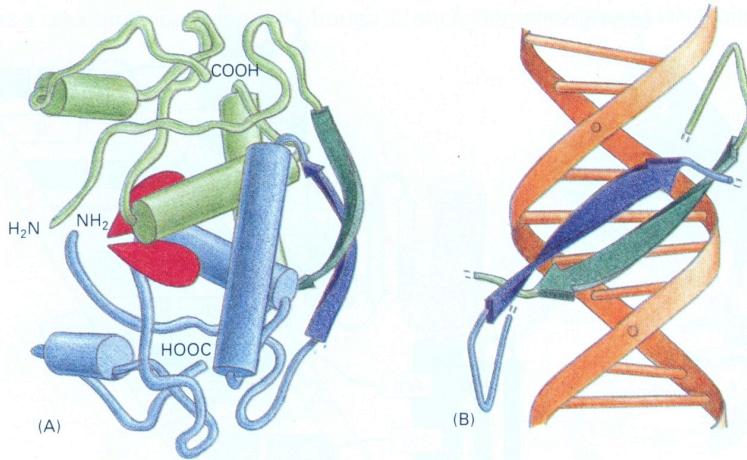


In Bacteria: A typical biosynthetic cycle around a critical metabolite (“SAM”)



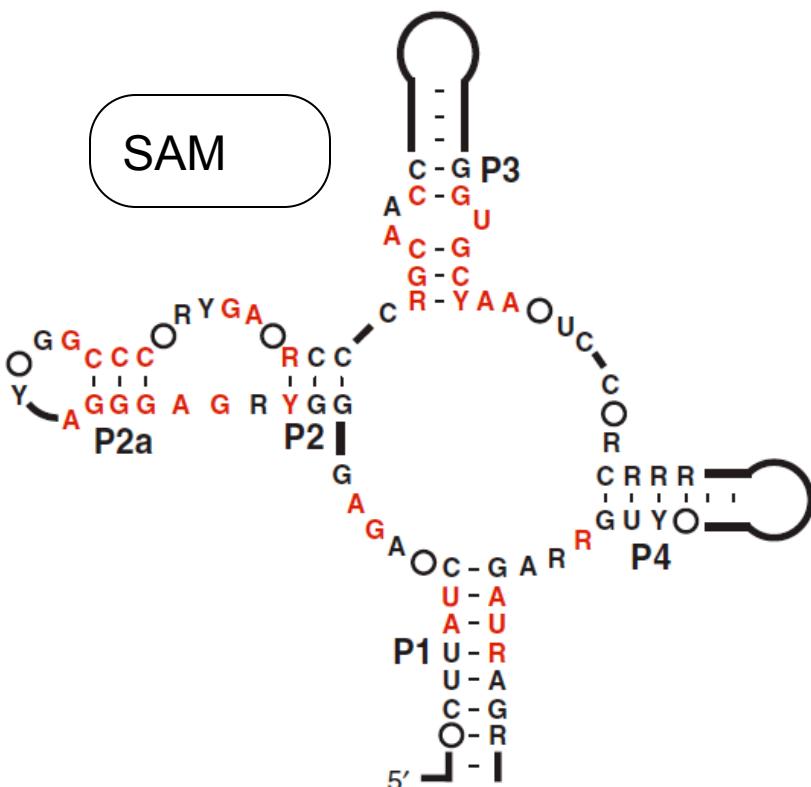
Gene Regulation: The MET Repressor





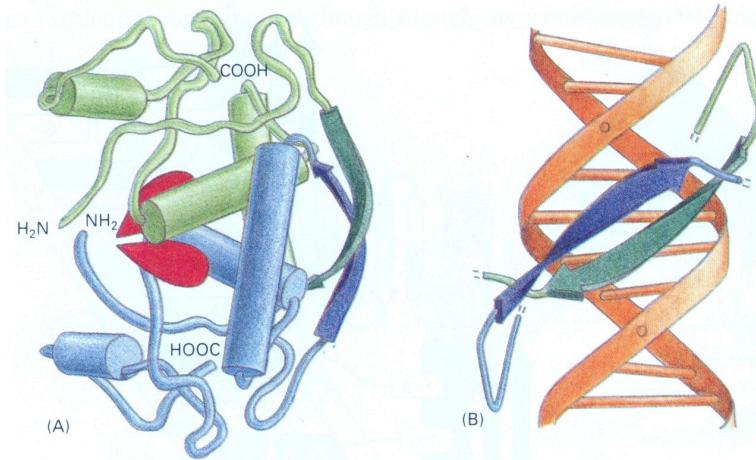
The
protein
way

Riboswitch
alternative



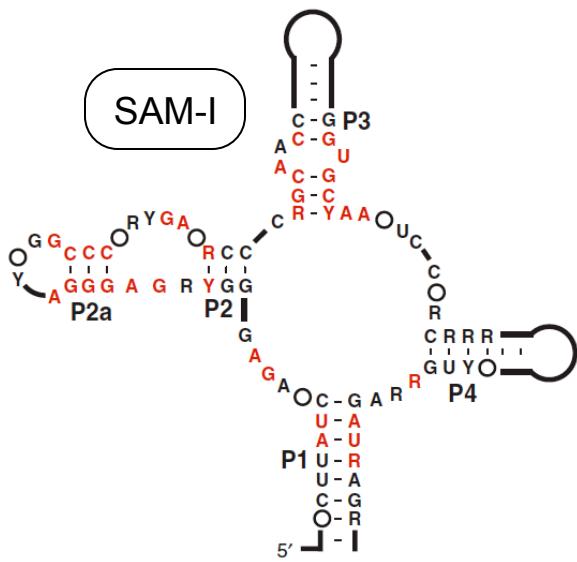
Grundy & Henkin, Mol. Microbiol 1998
Epshtein, et al., PNAS 2003
Winkler et al., Nat. Struct. Biol. 2003

Alberts, et al., 3e.



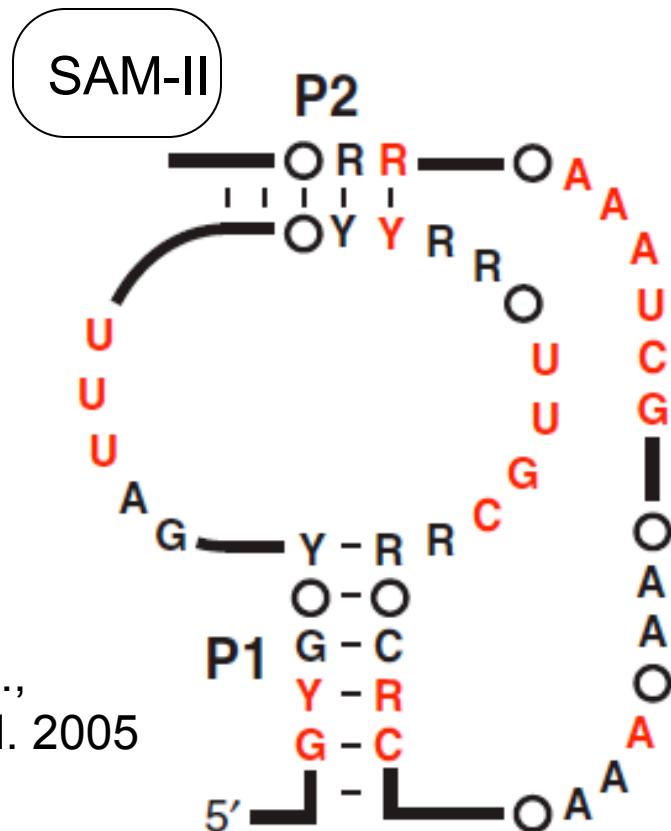
The
protein
way

Riboswitch
alternatives

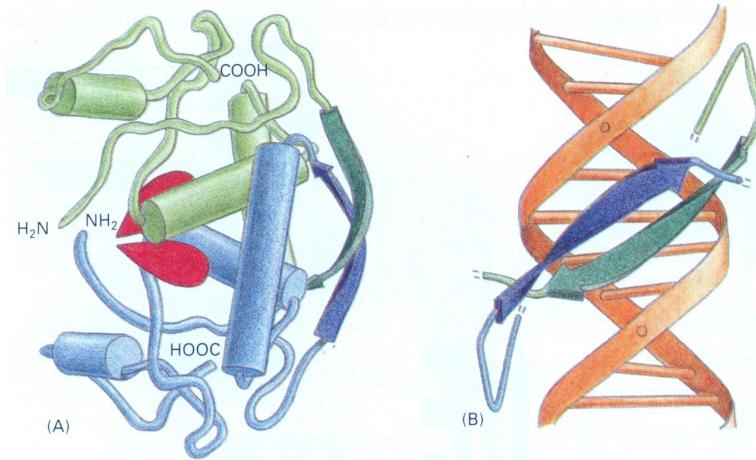


Grundy, Epshteyn, Winkler
et al., 1998, 2003

Corbino et al.,
Genome Biol. 2005



Alberts, et al, 3e.

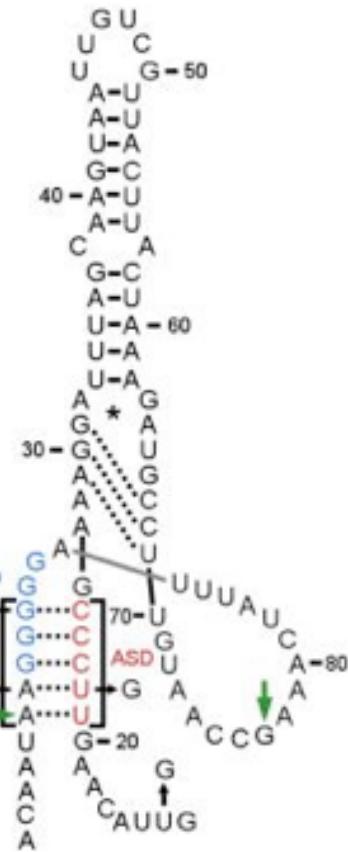


The
protein
way

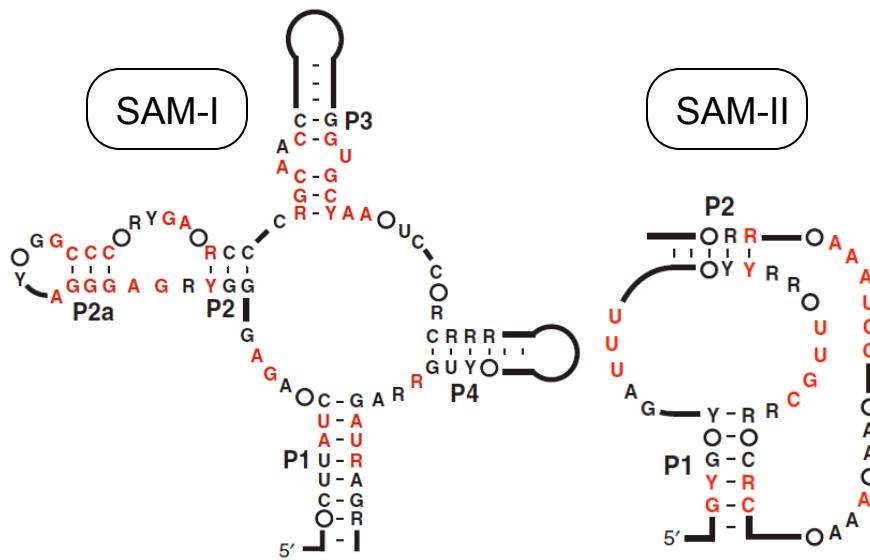
Riboswitch
alternatives



SAM-III

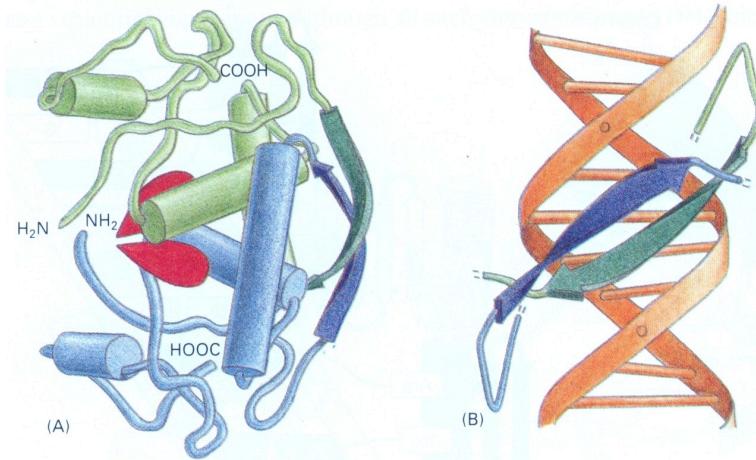


Fuchs et al.,
NSMB 2006



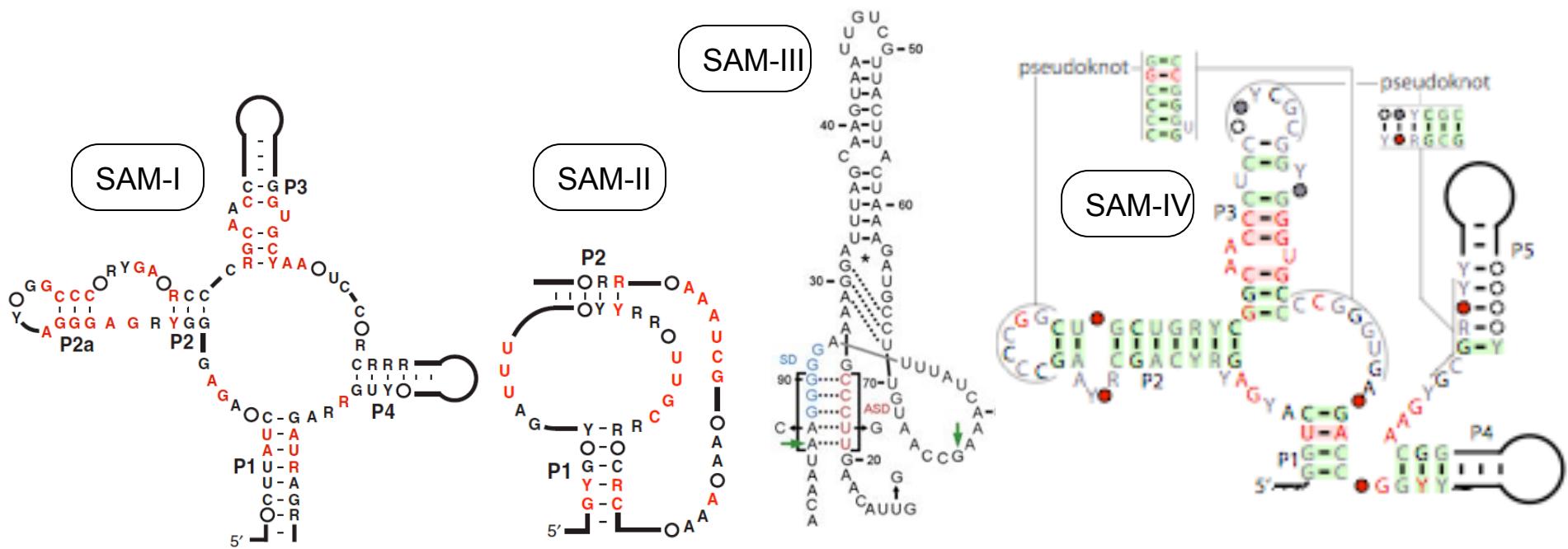
Grundy, Epshtain, Winkler
et al., 1998, 2003

Corbino et al.,
Genome Biol. 2005



The
protein
way

Riboswitch
alternatives

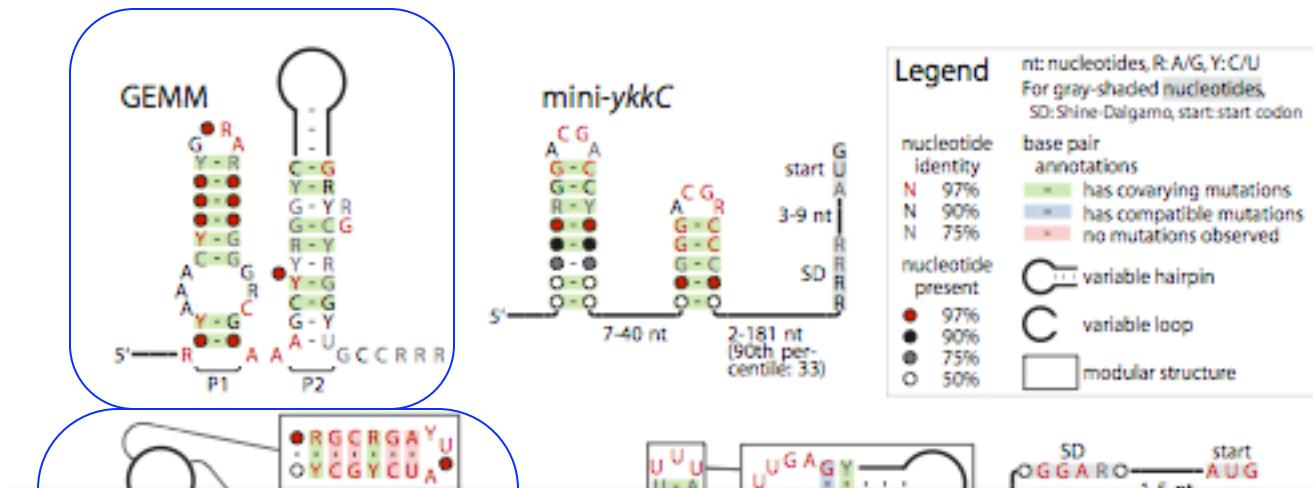


Grundy, Epshtain, Winkler
et al., 1998, 2003

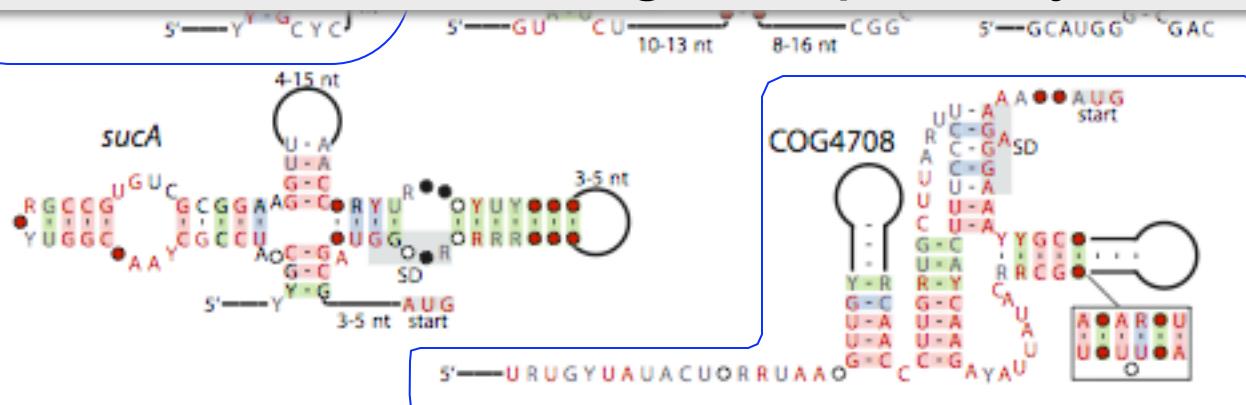
Corbino et al.,
Genome Biol. 2005

Fuchs et al.,
NSMB 2006

Weinberg et al.,¹⁹
RNA 2008



And many other examples. Widespread, deeply conserved, structurally sophisticated, functionally diverse, biologically important uses for ncRNA throughout prokaryotic world.

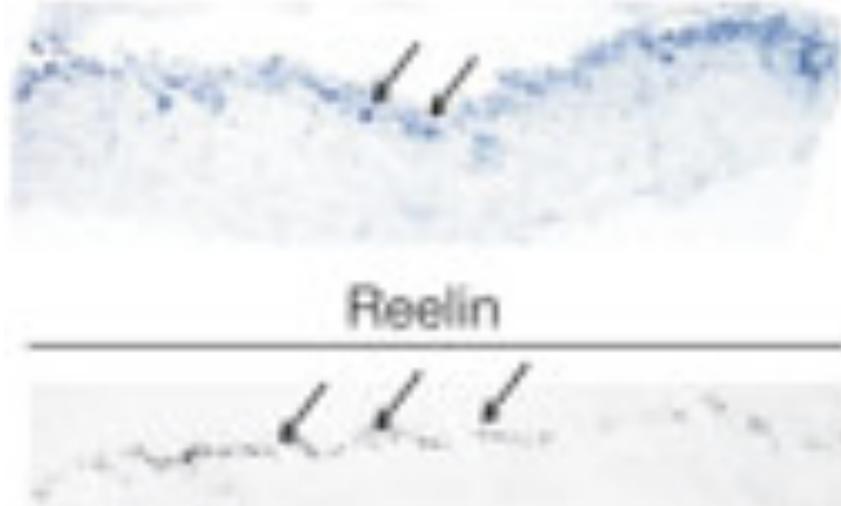


Vertebrates

- Bigger, more complex genomes
- <2% coding
- But >5% conserved in sequence?
- And 50-90% transcribed?
- And *structural* conservation, if any, invisible (without proper alignments, etc.)
 - What's going on?

Fastest Human Gene?

HAR1F AS



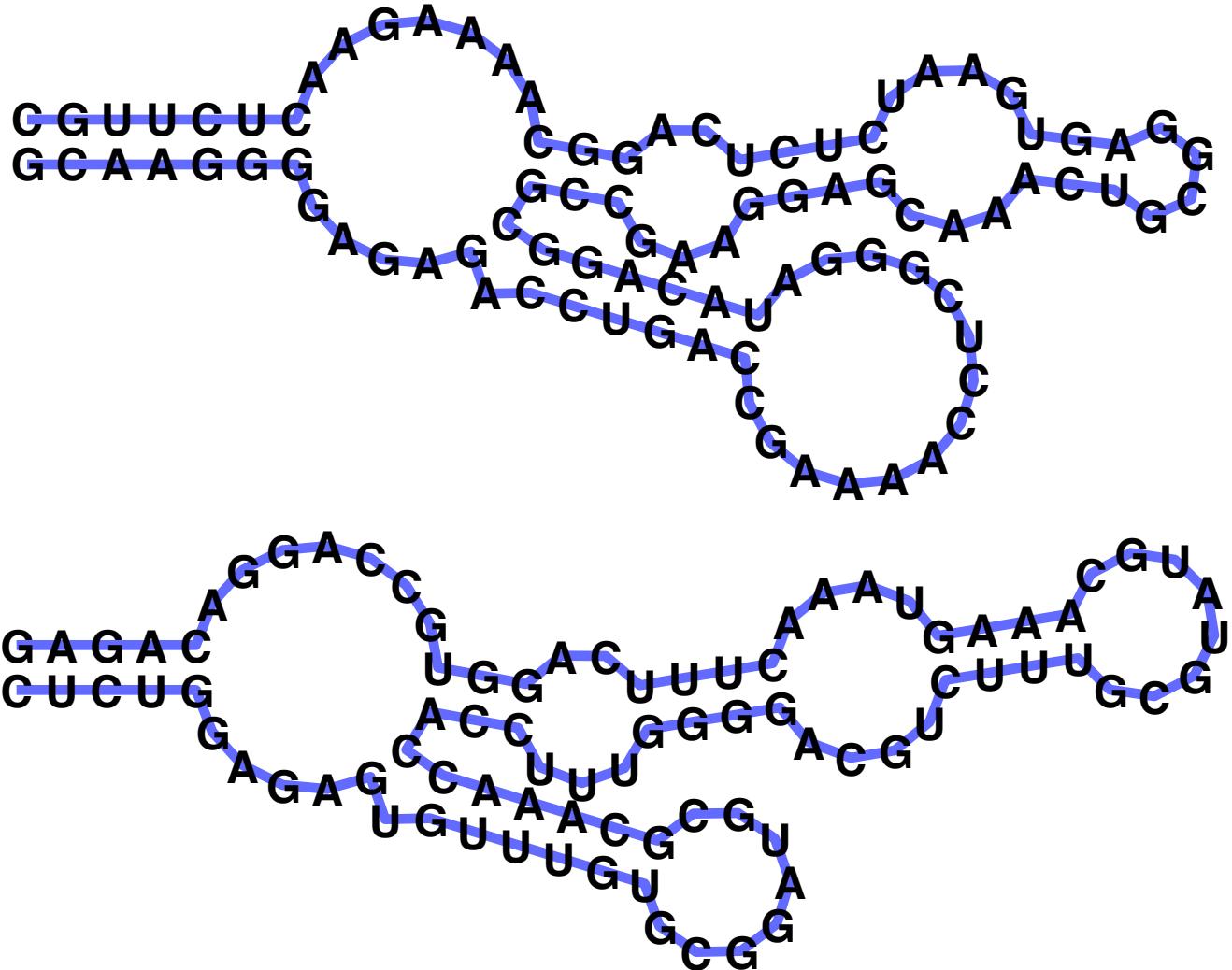
a

Position	20	30	40	50	
Human	AGACGTTACAGCAACCGTGTCA	GCTGAAATGATGGCGTAGACGCACGT			
Chimpanzee	AGAAAATTACAGCAATT	TATCAACTGAAATTATA	AGGTGTAGACACATGT		
Gorilla	AGAAAATTACAGCAATT	TATCAACTGAAATTATA	AGGTGTAGACACATGT		
Orang-utan	AGAAAATTACAGCAATT	TATCAACTGAAATTATA	AGGTGTAGACACATGT		
Macaque	AGAAAATTACAGCAATT	TATCA	GCTGAAATTATA	AGGTGTAGACACATGT	
Mouse	AGAAAATTACAGCAATT	TATCA	GCTGAAATTATA	AGGTGTAGACACATGT	
Dog	AGAAAATTACAGCAATT	TATCAACTGAAATTATA	AGGTGTAGACACATGT		
Cow	AGAAAATTACAGCAATT	CATCAG	CTGAAATTATA	AGGTGTAGACACATGT	
Platypus	ATAAAATTACAGCAATT	TATCAA	ATGAAATTATA	AGGTGTAGACACATGT	
Opossum	AGAAAATTACAGCAATT	TATCAACTGAAATTATA	AGGTGTAGACACATGT		
Chicken	AGAAAATTACAGCAATT	TATCAACTGAAATTATA	AGGTGTAGACACATGT		
Fold	(((((((.....))))....))	[[[[[+(((((.....))))+))]			
Pair symbol	lmnopqr	rqpon	ml	rstuvwxyz	xwvutsr

b



Q: What's so hard?



A: Structure often more important than sequence

Origin of Life?

Life needs

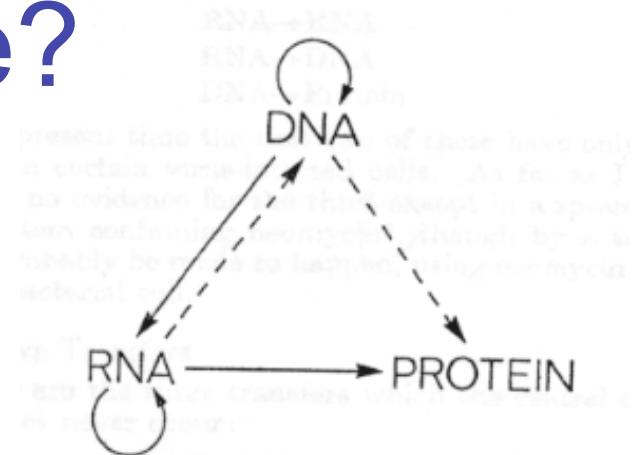
information carrier: DNA

molecular machines, like enzymes: Protein

making proteins needs DNA + RNA + proteins

making (duplicating) DNA needs proteins

Horrible circularities! How could it have arisen in an abiotic environment?



Origin of Life?

RNA can carry information, too

RNA double helix; RNA-directed RNA polymerase

RNA can form complex structures

RNA enzymes exist (ribozymes)

RNA can control, do logic (riboswitches)

The “RNA world” hypothesis:
1st life was RNA-based

6.5 RNA Secondary Structure

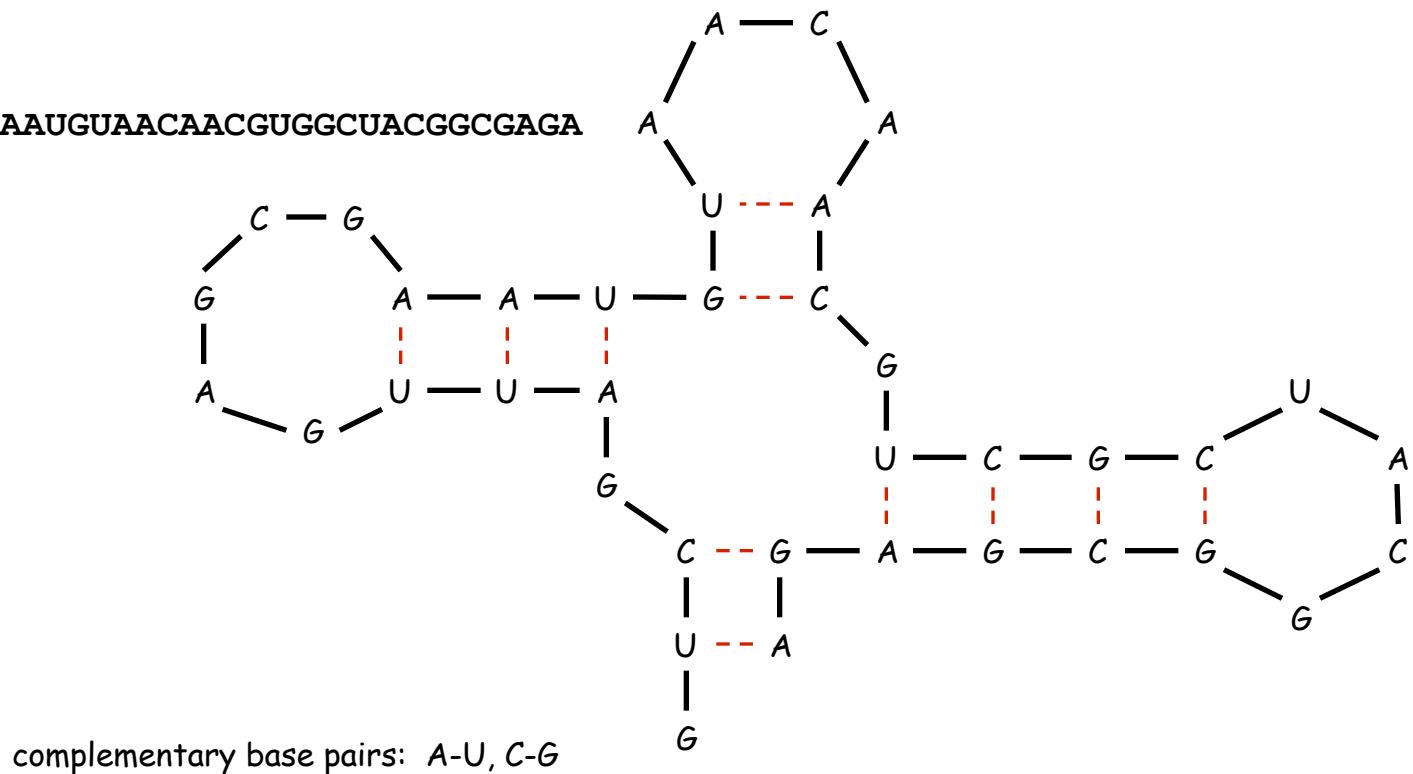
Nussinov's Algorithm – core technology
for RNA structure prediction

RNA Secondary Structure

RNA. String $B = b_1b_2\dots b_n$ over alphabet { A, C, G, U }.

Secondary structure. RNA is usually single-stranded, and tends to loop back and form base pairs with itself. This structure is essential for understanding behavior of molecule.

Ex: GUCAUUGAGCGAAUGUAACAAACGUGGUACGGCGAGA



RNA Secondary Structure (somewhat oversimplified)

Secondary structure. A set of pairs $S = \{ (b_i, b_j) \}$ that satisfy:

- [Watson-Crick.]
 - S is a *matching*, i.e. each base pairs with at most one other, and
 - each pair in S is a Watson-Crick pair: A-U, U-A, C-G, or G-C.
- [No sharp turns.] The ends of each pair are separated by at least 4 intervening bases. If $(b_i, b_j) \in S$, then $i < j - 4$.
- [Non-crossing.] If (b_i, b_j) and (b_k, b_l) are two pairs in S , then we cannot have $i < k < j < l$. (Violation of this is called a *pseudoknot*.)

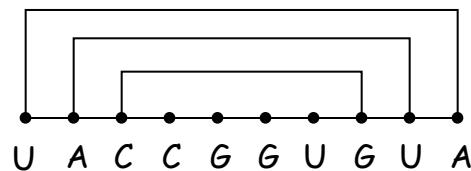
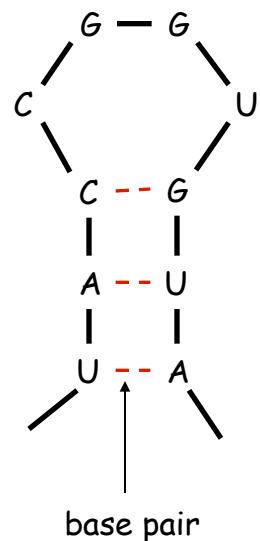
Free energy. Usual hypothesis is that an RNA molecule will form the secondary structure with the optimum total free energy.

approximate by number of base pairs

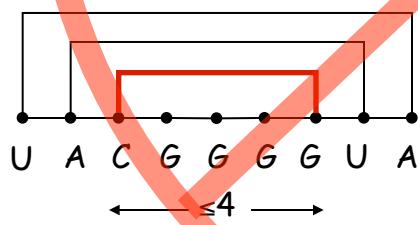
Goal. Given an RNA molecule $B = b_1 b_2 \dots b_n$, find a secondary structure S that maximizes the number of base pairs.

RNA Secondary Structure: Examples

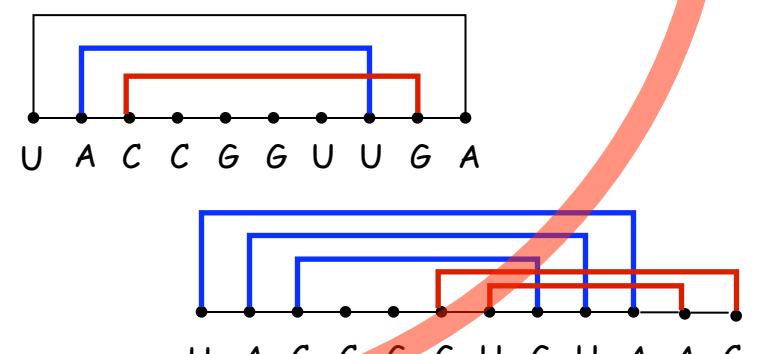
Examples.



ok



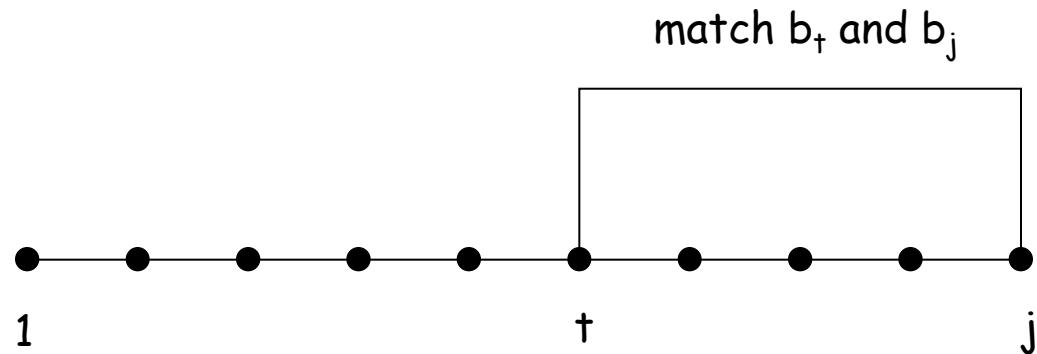
sharp turn



crossing

RNA Secondary Structure: Subproblems

First attempt. $\text{OPT}[j]$ = maximum number of base pairs in a secondary structure of the substring $b_1b_2\dots b_j$.



Difficulty. Results in two sub-problems.

- Finding secondary structure in: $b_1b_2\dots b_{t-1}$. $\leftarrow \text{OPT}(t-1)$
- Finding secondary structure in: $b_{t+1}b_{t+2}\dots b_{j-1}$. \leftarrow not OPT of anything;
need more sub-problems

Dynamic Programming Over Intervals: (R. Nussinov's algorithm)

Notation. $OPT[i, j]$ = maximum number of base pairs in a secondary structure of the substring $b_i b_{i+1} \dots b_j$.

- Case 1. If $i \geq j - 4$.
 - $OPT[i, j] = 0$ by no-sharp turns condition.
- Case 2. Base b_j is not involved in a pair.
 - $OPT[i, j] = OPT[i, j-1]$
- Case 3. Base b_j pairs with b_t for some $i \leq t < j - 4$.
 - non-crossing constraint decouples resulting sub-problems
 - $OPT[i, j] = 1 + \max_t \{ OPT[i, t-1] + OPT[t+1, j-1] \}$

↑
take max over t such that $i \leq t < j-4$ and
 b_t and b_j are Watson-Crick complements

Key point:
Either last base
is unpaired
(case 1,2) or
paired (case 3)

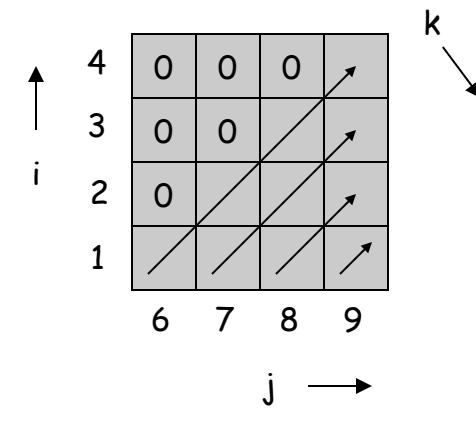
Remark. Same core idea in CKY algorithm to parse context-free grammars.

Bottom Up Dynamic Programming Over Intervals

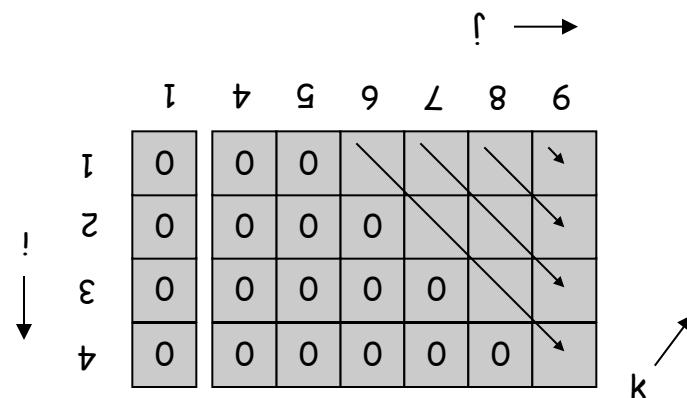
Q. What order to solve the sub-problems?

A. Do shortest intervals first.

```
RNA(b1, ..., bn) {
    for k = 5, 6, ..., n-1
        for i = 1, 2, ..., n-k
            j = i + k
            Compute OPT[i, j]
    return OPT[1, n] using recurrence
}
```



Running time. $O(n^3)$.



C U C C G G U U G C A A U G U C
((: (. . . .) .) . .) . .

$$n = 16$$

E.g.:
 $\text{OPT}[1,6] = 1:$

CUCCGG
(. . .)

E.g.:
 $\text{OPT}[6, 16] = 2$:

GUUGCAAUGUC
((.....)....)

0

Computing one cell: $\text{OPT}[2,18] = ?$

G G G A A A A C C C A A A A G G G G G U U U n= 20

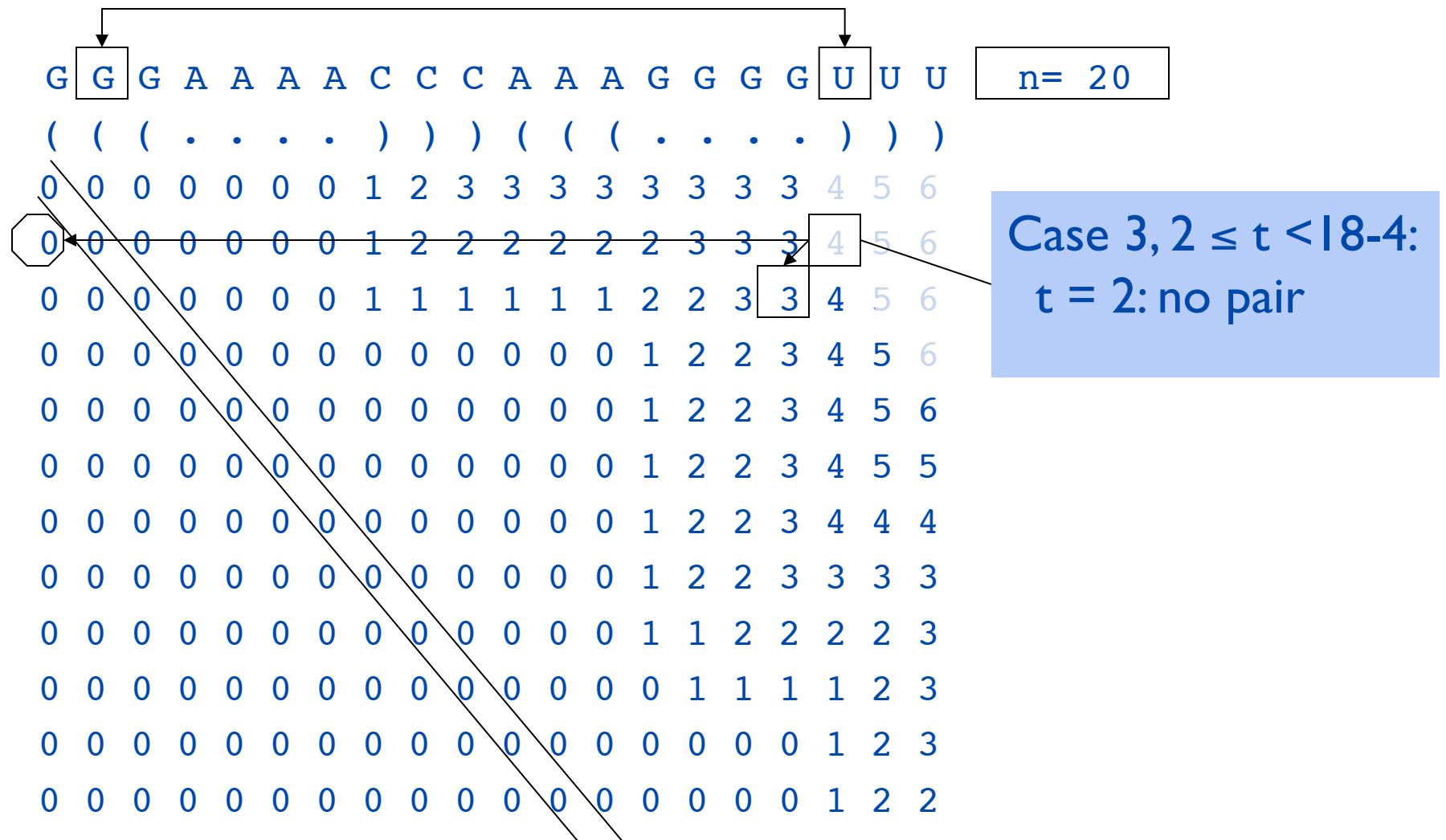
((((.	.	.	.)))	(((.	.	.	.)))				
0	0	0	0	0	0	0	1	2	3	3	3	3	3	3	3	3	3	4	5	6				
0	0	0	0	0	0	0	1	2	2	2	2	2	2	3	3	3	3	4	5	6				
0	0	0	0	0	0	0	1	1	1	1	1	1	1	2	2	3	3	4	5	6				
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6			
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6		
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	5	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	4	4	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	3	3	3	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2	2	2	3	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	2	3	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2

- Case 1:
 $2 \geq 18 - 4$? no.
- Case 2:
 B_{18} unpaired?
Always a possibility;
then $OPT[2, 18] \geq 3$

GGAAAACCCAAAGGGGU
((.....))(... . .) . . .

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

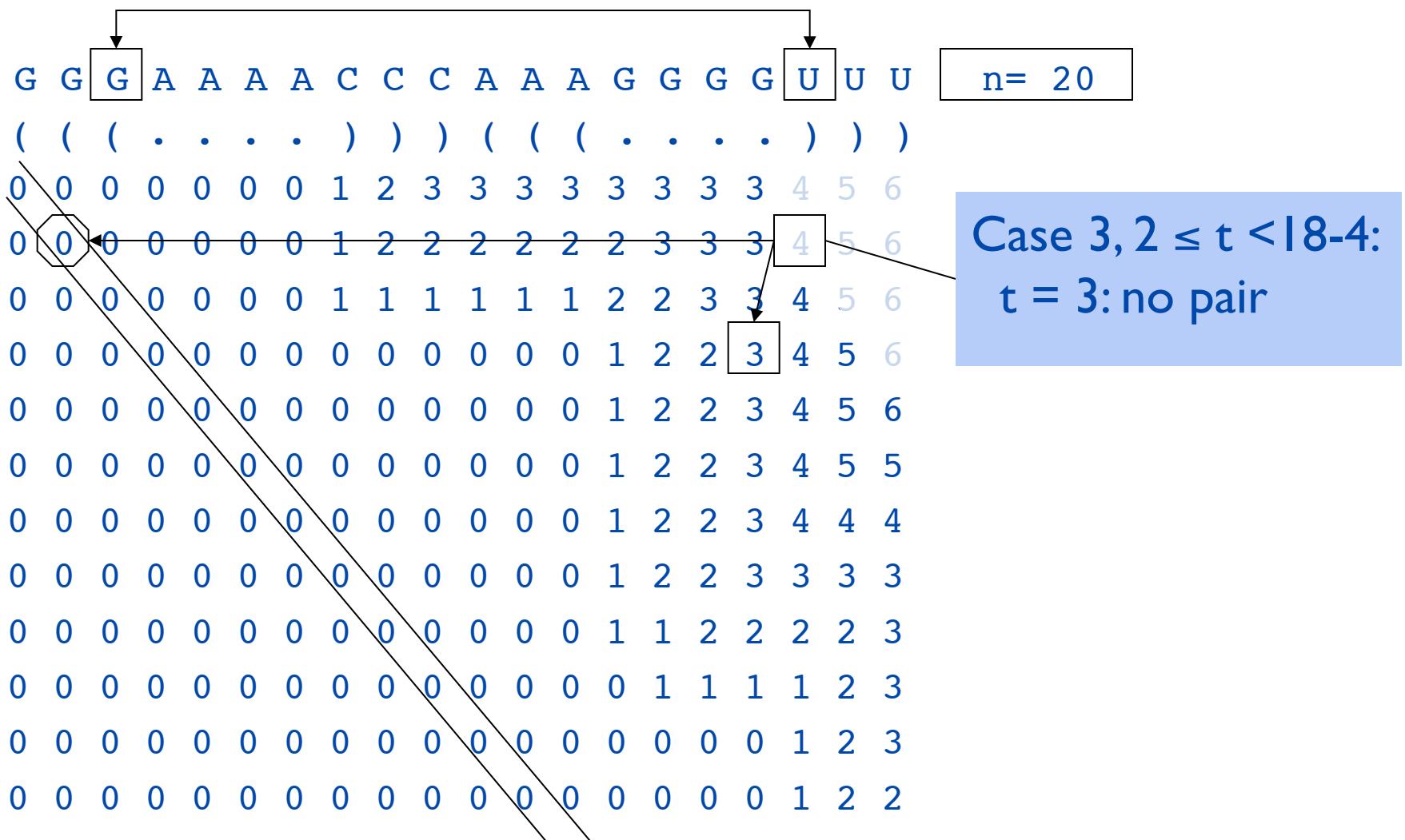
Computing one cell: $\text{OPT}[2,18] = ?$



$$\text{OPT}(i,j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

0
 0
 0

Computing one cell: $\text{OPT}[2,18] = ?$



$$\text{OPT}(i,j) = \begin{cases} 0 & \text{if } i \geq j-4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2,18] = ?$

G G G A A A A C C C A A A A G G G G U U U n= 20

Case 3, $2 \leq t < 18-4$:
 $t = 4$: yes pair
 $\text{OPT}[2, 18] \geq 1 + 0 + 3$

GGAAAACCCAAAGGGGU
..(....((...)))

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2,18] = ?$



Case 3, $2 \leq t < 18-4$:
 $t = 5$: yes pair
 $\text{OPT}[2,18] \geq 1 + 0 + 3$

GGAAAACCCAAAGGGGU
 $\dots(\dots((\dots)))$

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2,18] = ?$



Case 3, $2 \leq t < 18-4$:
 $t = 6$: yes pair
 $\text{OPT}[2,18] \geq 1 + 0 + 3$

GGAAAACCCAAAGGGGU
.....((....)))

$$\text{OPT}(i,j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: OPT[2,18] = ?

G G G A A A A C C C A A A A G G G G U U U n= 20

((((.	.	.	.)))	(((.	.	.	.)))				
0	0	0	0	0	0	0	1	2	3	3	3	3	3	3	3	3	4	5	6					
0	0	0	0	0	0	0	0	1	2	2	2	2	2	2	3	3	3	4	5	6				
0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	2	2	3	3	4	5	6			
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6			
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6		
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	5
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	4	4
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	3	3	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2	2	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	2	3	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3			
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3			

Case 3, $2 \leq t < 18 - 4$:
 $t = 7$: yes pair
 $\text{OPT}[2, 18] \geq 1 + 0 + 3$

GGAAAACCCAAAGGGGU
.....(((((.....))))

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: OPT[2,18] = ?

G G G A A A A C C C A A A A G G G G U U U n= 20

((((.))))	(((.)))				
0	0	0	0	0	0	0	0	1	2	3	3	3	3	3	3	3	3	3	3	4	5	6					
0	0	0	0	0	0	0	0	0	1	2	2	2	2	2	2	2	2	3	3	3	4	5	6				
0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	1	2	2	3	3	4	5	6			
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	5
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	5
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	4	4
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	3	3	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	3	3	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2	3	2	2
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2	3	2	2
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	2	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2

Case 3, $2 \leq t < 18-4$:
 $t = 8$: no pair

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: OPT[2,18] = ?

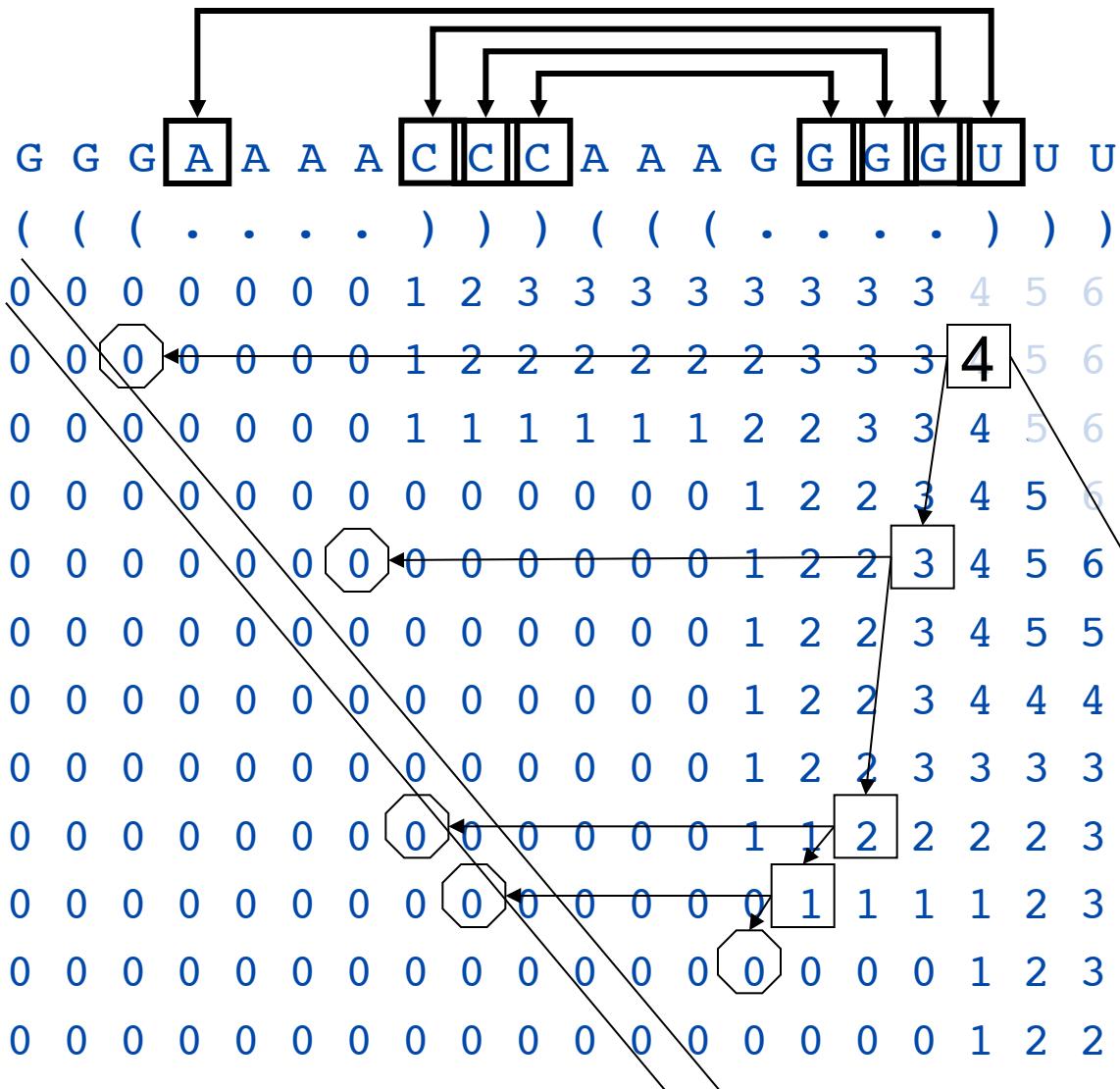
Case 3,
t = 11
OPT[20]

Case 3, $2 \leq t < 18 - 4$:
 $t = 11$: yes pair
 $\text{OPT}[2, 18] \geq 1 + 2 + 0$

GGAAAACCCAAAGGGGU
((.....)) (.....)

(not shown:
 $t=9, 10, 12, 13$)

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \right\} & \text{otherwise} \end{cases}$$

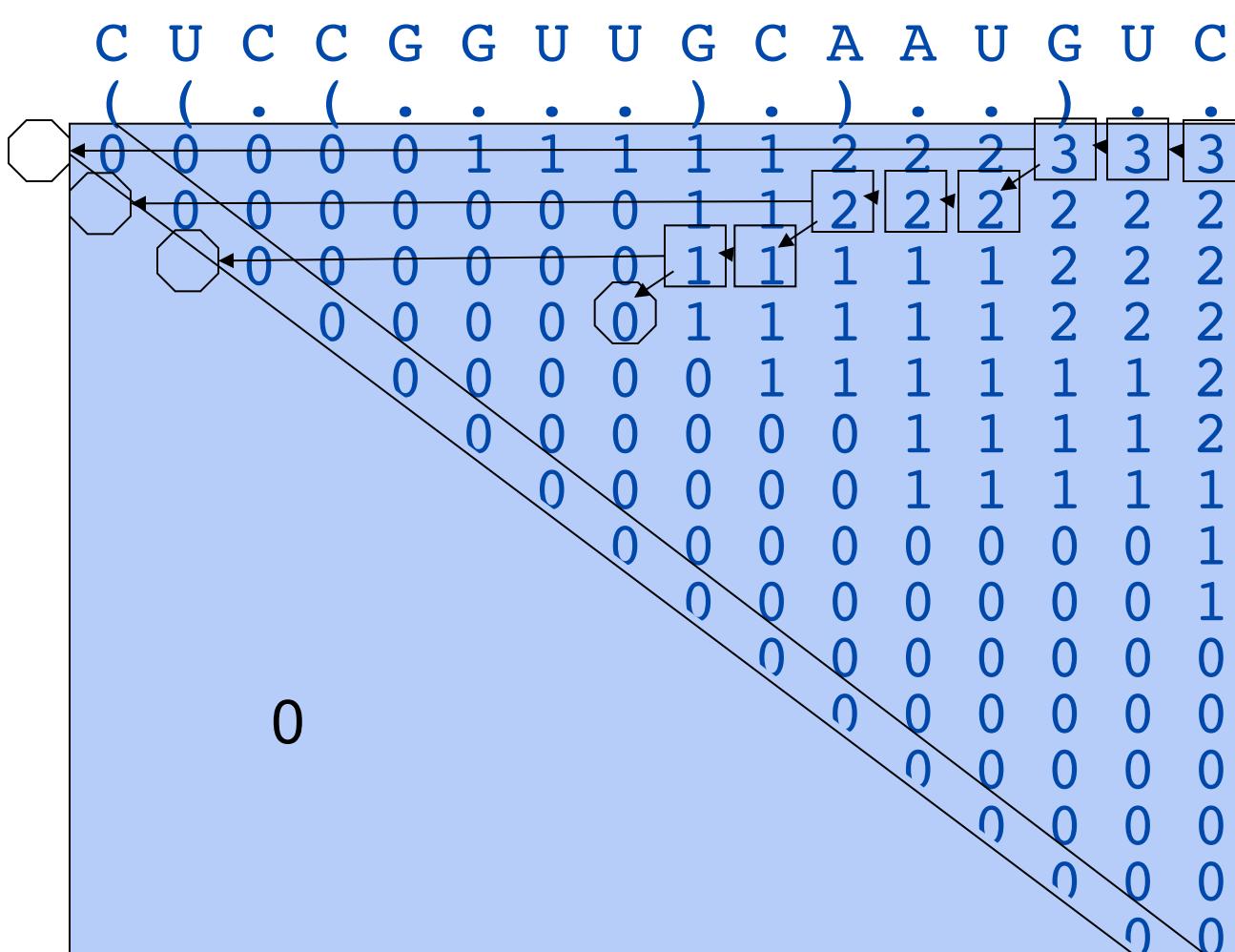


Overall, Max = 4
several ways, e.g.:
 GGAAAACCCAAAGGGGU
 ...((...((....))))

tree shows trace back:
 square = case 3
 octagon = case 1

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Another Trace Back Example



E.g.:
 $\text{OPT}[1, 16] = 3:$
CUCCGGUUGCAAUGUC
((...). . .) . . .)

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$