

# **CSE 417: Algorithms and Computational Complexity**

## **4: Dynamic Programming, I Fibonacci**

Winter 2006

Lecture 12

W. L. Ruzzo

# Some Algorithm Design Techniques, I

- General overall idea
  - Reduce solving a problem to a smaller problem or problems of the same type
- Greedy algorithms
  - Used when one needs to build something a piece at a time
  - Repeatedly make the **greedy** choice - the one that looks the best right away
    - e.g. closest pair in TSP search
  - Usually fast if they work (but often don't)

# Some Algorithm Design Techniques, II

- Divide & Conquer
  - Reduce problem to one or more sub-problems of the same type
  - Typically, each sub-problem is at most a constant fraction of the size of the original problem
    - e.g. Mergesort, Binary Search, Strassen's Algorithm, Quicksort (kind of)

# Some Algorithm Design Techniques, III

- Dynamic Programming
  - Give a solution of a problem using smaller sub-problems, e.g. a recursive solution
  - Useful when the same sub-problems show up again and again in the solution

# “Dynamic Programming”

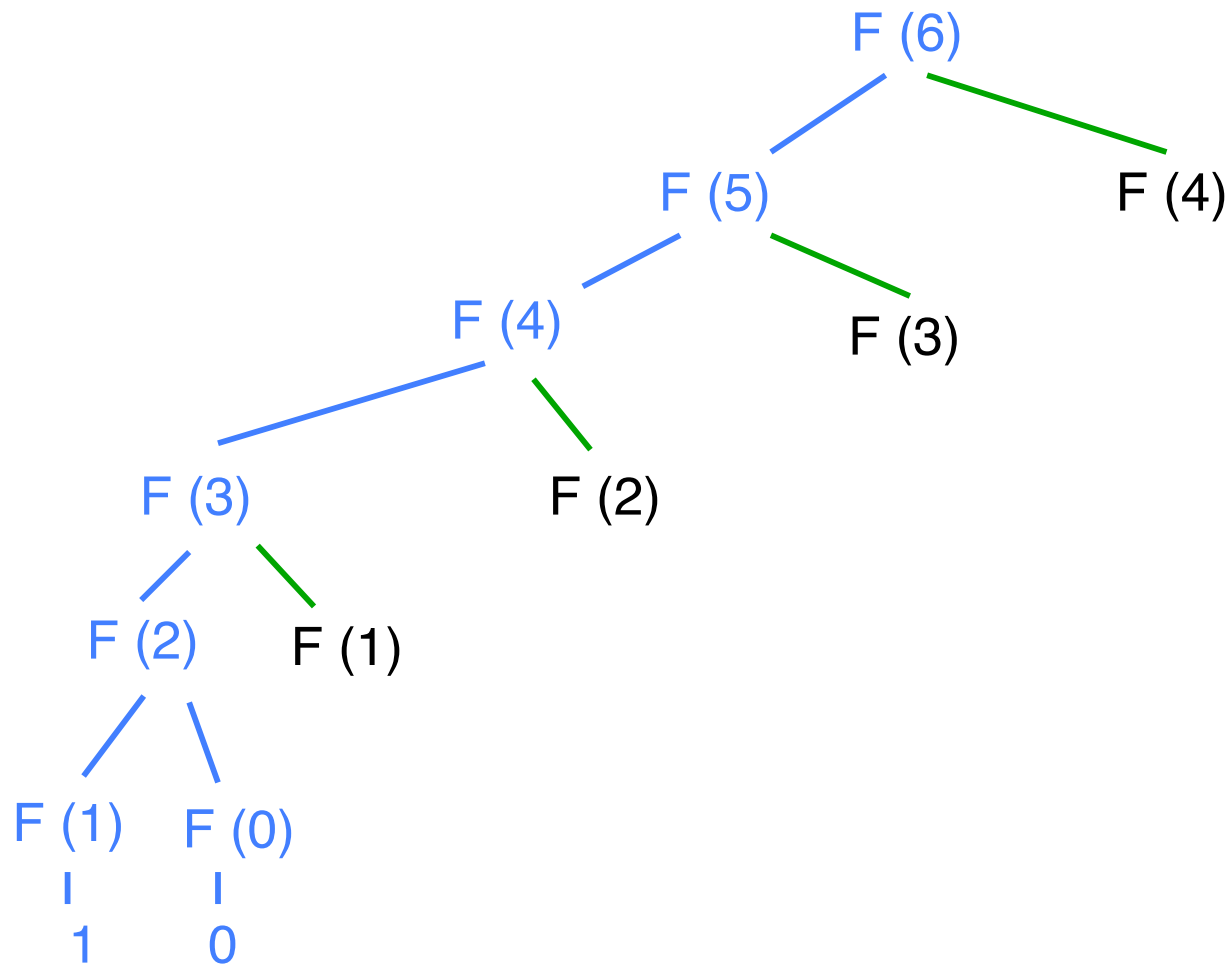
Program — A plan or procedure for dealing with some matter

– Webster’s New World Dictionary

# A simple case: Computing Fibonacci Numbers

- Recall  $F_n = F_{n-1} + F_{n-2}$  and  $F_0 = 0, F_1 = 1$
- Recursive algorithm:
  - `Fibo(n)`
    - if** `n=0` **then return**(0)
    - else if** `n=1` **then return**(1)
    - else return**(`Fibo(n-1)+Fibo(n-2)`)

# Call tree - start







# Memo-ization (Caching)

- Remember all values from previous recursive calls
- Before recursive call, test to see if value has already been computed
- Dynamic Programming
  - Convert memo-ized algorithm from a recursive one to an iterative one (top-down → bottom-up)

# Fibonacci - Memo-ized Version

initialize:  $F[i] \leftarrow$  undefined for all  $i$

$F[0] \leftarrow 0$

$F[1] \leftarrow 1$

FiboMemo( $n$ ):

**if**( $F[n]$  undefined) {

$F[n] \leftarrow$  FiboMemo( $n-2$ )+FiboMemo( $n-1$ )

}

**return**( $F[n]$ )

# Fibonacci - Dynamic Programming Version

FiboDP(n):

$F[0] \leftarrow 0$

$F[1] \leftarrow 1$

**for**  $i=2$  **to**  $n$  **do**

$F[i] \leftarrow F[i-1] + F[i-2]$

**endfor**

**return**( $F[n]$ )

# Dynamic Programming

- Useful when
  - same recursive sub-problems occur repeatedly
  - Can anticipate the parameters of these recursive calls
  - The solution to whole problem can be figured out **without** knowing the internal details of how the sub-problems are solved
    - principle of optimality