

Student Activity

```
Mystery(int array a[]) {
  for (int p = 1; p < length; p++) {
    int tmp = a[p];
    for (int j = p; j > 0 && tmp < a[j-1]; j--)
      a[j] = a[j-1];
    a[j] = tmp;
  }
}
```

What sort is this? Insertion

What is its running time?
Best?
Avg?
Worst?

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Merge Sort: Complexity

Base case: $T(1) = c$
 $T(n) = 2 T(n/2) + n$
 ...
 $T(n) = O(n \log n)$
 (best, worst)

We Want:
 $n/2^k = 1$
 $n = 2^k$
 $\log n = k$

Base case: $T(1) = c$
 $T(n) = 2 T(n/2) + n$
 $= 2 (2T(n/4) + n/2) + n$
 $= 4T(n/4) + n + n$
 $= 4T(n/4) + 2n$
 $= 4 (2T(n/8) + n/4) + 2n$
 $= 8T(n/8) + n + 2n$
 $= 8T(n/8) + 3n$
 $= 2^k T(n/2^k) + kn$
 $= nT(1) + n \log n$
 $= n + n \log n$

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QuickSort: Best case complexity

$T(n) = 2T(n/2) + n$
 ...
 $T(n) = O(n \log n)$

Same as Mergesort

What is best case? Always chooses a pivot that splits array in half at each step

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QuickSort: Worst case complexity

$T(1) = c$
 $T(n) = n + T(n-1)$
 $T(n) = n + (n-1) + T(n-2)$
 $T(n) = n + (n-1) + (n-2) + T(n-3)$
 $T(n) = 1 + 2 + 3 + \dots + N$
 ...
 $T(n) = O(n^2)$

$T(n) = n + T(n-1)$
 ...
 $T(n) = O(n^2)$

Always chooses WORST pivot – so that one array is empty at each step

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