

# Graphs

Chapter 9 in Weiss

5/28/2008

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# Today's Outline

- Announcements
  - HW #6-7
    - Partner Selection due Thurs May 29
    - Assignment due Thurs June 5th.
- **Graphs**
  - **Dijkstra's Algorithm**

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# Graph Traversals

- **Breadth-first search**
  - explore *all* adjacent nodes, then for each of *those* nodes explore *all* adjacent nodes
- **Depth-first search**
  - explore **first** child node, then its **first** child node, etc. until goal node is found or node has no children. Then backtrack, repeat with sibling.
- Both:
  - Work for arbitrary (directed or undirected) graphs
  - Must mark visited vertices so you do not go into an infinite loop!
- Either can be used to determine connectivity:
  - Is there a **path** between two given vertices?
  - Is the graph (weakly) connected?
- Which one:
  - Uses a queue?
  - Uses a stack?
  - Always finds the **shortest path** (for unweighted graphs)?

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# The Shortest Path Problem

Given a graph  $G$ , edge costs  $c_{i,j}$ , and vertices  $s$  and  $t$  in  $G$ , **find the shortest path from  $s$  to  $t$ .**

For a path  $p = v_0 v_1 v_2 \dots v_k$

– *unweighted length* of path  $p = k$  (a.k.a. *length*)

– *weighted length* of path  $p = \sum_{i=0, k-1} c_{i,i+1}$  (a.k.a. *cost*)

Path length equals path cost when ?

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# Single Source Shortest Paths (SSSP)

Given a graph  $G$ , edge costs  $c_{i,j}$ , and vertex  $s$ , **find the shortest paths from  $s$  to all vertices in  $G$ .**

# All Pairs Shortest Paths (APSP)

Given a graph  $G$  and edge costs  $c_{i,j}$ , **find the shortest paths between all pairs of vertices in  $G$ .**

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# Variations of SSSP

- Weighted vs. unweighted
- Directed vs undirected
- Cyclic vs. acyclic
- Positive weights only vs. negative weights allowed
- Shortest path vs. longest path
- ...

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## Applications

- Network routing
- Driving directions
- Cheap flight tickets
- Critical paths in project management (see textbook)
- ...

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## SSSP: Unweighted Version

*Ideas?*

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```
void Graph::unweighted (Vertex s){
    Queue q(NUM_VERTICES);
    Vertex v, w;
    q.enqueue(s);
    s.dist = 0;

    while (!q.isEmpty()){
        v = q.dequeue();
        for each w adjacent to v
            if (w.dist == INFINITY){
                w.dist = v.dist + 1;
                w.path = v;
                q.enqueue(w);
            }
    }
}
```

each edge examined at most once – if adjacency lists are used

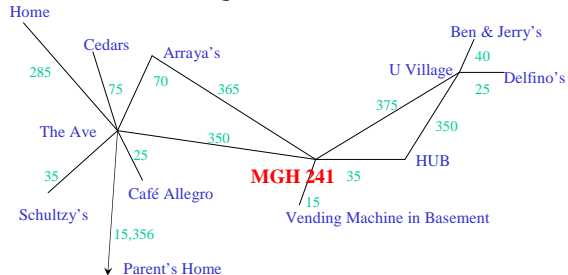
each vertex enqueued at most once

total running time:  $O(\quad)$

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## Weighted SSSP: The Quest For Food



*Can we calculate shortest distance to all nodes from MGH 241?*

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## Dijkstra, Edsger Wybe

Legendary figure in computer science; was a professor at University of Texas.

Supported teaching introductory computer courses without computers (pencil and paper programming)

Supposedly wouldn't (until very late in life) read his e-mail; so, his staff had to print out messages and put them in his box.



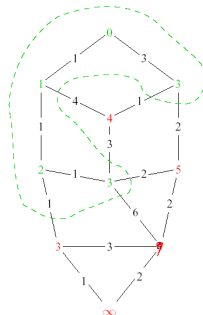
E.W. Dijkstra (1930-2002)

1972 Turing Award Winner, Programming Languages, semaphores, and ...

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## Dijkstra's Algorithm: Idea



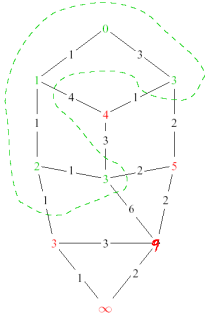
Adapt BFS to handle weighted graphs

Two kinds of vertices:

- Finished or **known** vertices
  - Shortest distance has been computed
- **Unknown** vertices
  - Have tentative distance

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## Dijkstra's Algorithm: Idea



At each step:

- 1) Pick closest **unknown** vertex
- 2) Add it to **known** vertices
- 3) Update distances

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## Dijkstra's Algorithm: Pseudocode

Initialize the cost of each node to  $\infty$

Initialize the cost of the source to 0

While there are **unknown** nodes left in the graph

Select an **unknown** node  $b$  with the lowest cost

Mark  $b$  as **known**

For each node  $a$  adjacent to  $b$

$a$ 's cost =  $\min(a$ 's old cost,  $b$ 's cost + cost of  $(b, a)$ )

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```
void Graph::dijkstra(Vertex s){
    Vertex v,w;

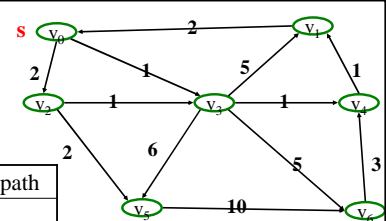
    Initialize s.dist = 0 and set dist of all other
    vertices to infinity

    while (there exist unknown vertices, find the
    one b with the smallest distance)
        b.known = true;

    for each a adjacent to b
        if (!a.known)
            if (b.dist + Cost_ba < a.dist){
                decrease(a.dist to= b.dist + Cost_ba);
                a.path = b;
            }
    }
}
```

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V	Known	Dist	path
v0			
v1			
v2			
v3			
v4			
v5			
v6			

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## Dijkstra's Algorithm: a Greedy Algorithm

**Greedy** algorithms always make choices that *currently* seem the best

- Short-sighted - no consideration of long-term or global issues
- Locally optimal - does not always mean globally optimal!!

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## Dijkstra's Algorithm: Summary

- Classic algorithm for solving SSSP in weighted graphs *without negative weights*
- A *greedy* algorithm (irrevocably makes decisions without considering future consequences)
- Intuition for correctness:
  - shortest path from source vertex to itself is 0
  - cost of going to adjacent nodes is at most edge weights
  - cheapest of these must be shortest path to that node
  - update paths for new node and continue picking cheapest path

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