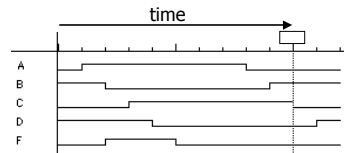
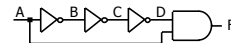


Overview

- ◆ Last lecture
 - Summary of 2-level combinational logic
 - Class example: A 4-bit prime-number detector
 - Quiz
- ◆ Today
 - Timing diagrams
 - Multilevel logic
 - ☛ Multilevel NAND/NOR conversion
 - ☛ AOI and OAI gates
 - Hazards

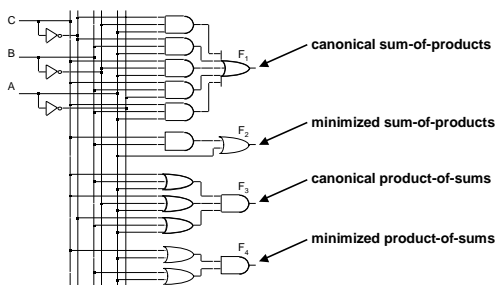
Timing diagrams

- ◆ "Sideways" truth tables
 - Show time-response of circuits
 - ☛ Real gates have real delays
 - Example: $A' \cdot A = 0$
 - ☛ Delays cause transient $F=1$



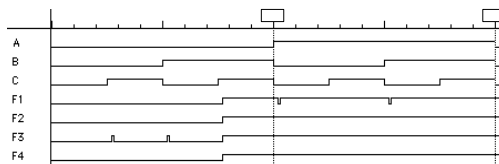
Some texts call timing diagrams "waveforms"

Example: $F=A+BC$ in 2-level logic



Timing diagram for $F = A + BC$

- ◆ Time waveforms for $F_1 - F_4$ are identical
 - Except for timing hazards (glitches)
 - More on this shortly...

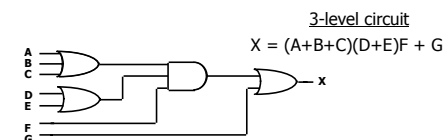


Multilevel logic

- ◆ Basic idea: Simplify logic using >2 gate levels
 - Time-space (speed versus gate count) tradeoff
- ◆ Two-level logic *usually*
 - Has smaller delays (faster circuits)
 - ☛ But more gates and more wires (more circuit area)
 - ☛ Sometimes has large fan-ins (slow)
 - Easier to eliminate hazards
- ◆ Multilevel logic *usually*
 - Has less gates (smaller circuits)
 - ☛ But can be slower (more gate delays)
 - Harder to eliminate hazards

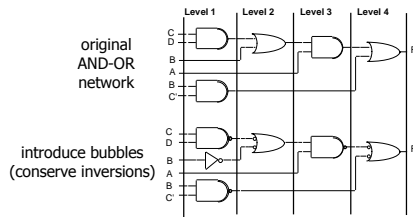
Multilevel logic example

- ◆ Function X
 - SOP: $X = ADF + AEF + BDF + BEF + CDF + CEF + G$
 - ☛ X is minimized!
 - ☛ Six 3-input ANDs; one 7-input OR; 25 wires
 - Multilevel: $X = (A+B+C)(D+E)F + G$
 - ☛ Factored form
 - ☛ One 3-input OR, two 2-input OR's, one 3-input AND; 10 wires



Multilevel NAND/NAND conversion

$$F = A(B+CD) + BC'$$

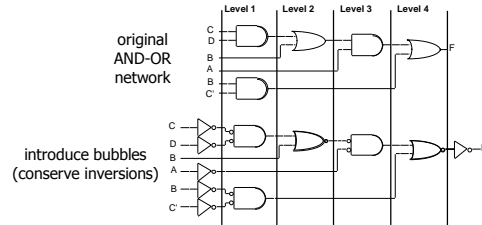


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Multilevel NOR/NOR conversion

$$F = A(B+CD) + BC'$$

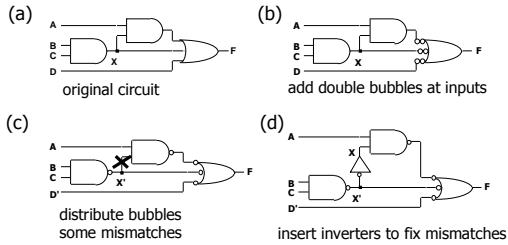


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Generic multilevel conversion

$$F = ABC + BC + D = AX + X + D$$

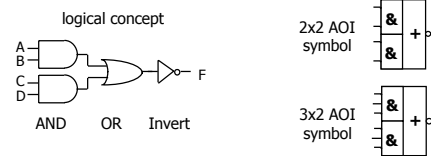


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AND-OR-Invert and OR-AND-Invert blocks

- ◆ AOI and OAI are dense 3-level building blocks
 - 2x2 AOI uses only 8 transistors
 - Minimal delay



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Using AOI and OAI blocks

- ◆ Approach
 - Start by finding F'
 - If using AOI, find F' in minimized SOP form
 - If using OAI, find F' in minimized POS form
 - Form AOI as (F')
- ◆ Example: $F = A'B + AB'$
 - SOP form: $F' = A'B' + AB$
 - AOI form: $F = (A'B' + AB)'$



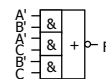
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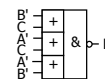
Example: AOI and OAI

	A			
	0	1	1	1
C	0	0	1	0
	B			

AOI form – Use SOP
 $F' = A'B' + A'C + B'C$
 $F = (A'B' + A'C + B'C)'$



OAI form – Use POS
 $F' = (B'+C)(A'+C)(A'+B')$
 $F = [(B'+C)(A'+C)(A'+B')]'$



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Issues with multilevel design

- ◆ No global definition of "optimal" multilevel circuit
 - Optimality depends on user-defined goals
 - Synthesize an implementation that meets design goals
- ◆ Synthesis requires CAD-tool help
 - No simple hand methods like K-maps
 - CAD tools manipulate Boolean expressions
 - ▣ Factoring, decomposition, etc.
 - Covered in more detail in CSE467

Multilevel logic summary

- ◆ Advantages over 2-level logic
 - Smaller circuits
 - Reduced fan-in
 - Less wires
- ◆ Disadvantages w.r.t 2-level logic
 - More difficult design
 - Less powerful optimizing tools
 - Dynamic hazards
- ◆ What you should know for CSE370
 - The basic multilevel idea
 - Multilevel NAND/NAND and NOR/NOR conversion
 - AOI gates

Hazards/glitches

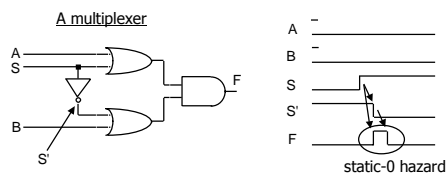
- ◆ Hazards/glitches: Undesired output switching
 - Occurs when different pathways have different delays
 - Wastes power; causes circuit noise
 - Dangerous if logic makes a decision while output is unstable
 - Dangerous if using asynchronous circuits
- ◆ Solutions
 - Design hazard-free circuits
 - ▣ Difficult when logic is multilevel
 - Wait until signals are stable
 - Use synchronous circuits

Types of hazards

- ◆ Static 1-hazard
 - Output should stay logic 1
 - Gate delays cause brief glitch to logic 0
- ◆ Static 0-hazard
 - Output should stay logic 0
 - Gate delays cause brief glitch to logic 1
- ◆ Dynamic hazards
 - Output should toggle cleanly
 - Gate delays cause multiple transitions

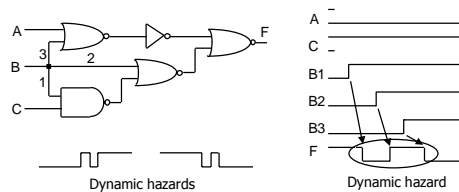
Static hazards

- ◆ Occur when a literal and its complement momentarily assume the same value
 - Through different paths with different delays
 - Causes an (ideally) static output to *glitch*



Dynamic hazards

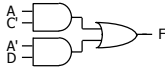
- ◆ Occur when a literal assumes multiple values
 - Through different paths with different delays
 - Causes an output to toggle multiple times



Eliminating static hazards

- ◆ In 2-level logic circuits
 - Assuming single-bit changes
- ◆ Key idea: Glitches happen when a changing input spans separate k-map encirclements
 - Example: 1101 to 0101 change can cause a static-1 glitch

$$F = AC' + A'D$$



	AB		A		
	00	01	11	10	
CD	00	0	0	1	1
	01	1	1	1	1
C	11	1	1	0	0
	10	0	0	0	0
		B			

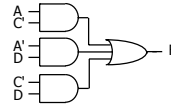
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Eliminating static hazards (con't)

- ◆ Solution: Add redundant k-map encirclements
 - Ensure that all single-bit changes are covered
 - First eliminate static-1 hazards: Use SOP form

$$F = AC' + A'D + C'D$$



	AB		A		
	00	01	11	10	
CD	00	0	0	1	1
	01	1	1	1	1
C	11	1	1	0	0
	10	0	0	0	0
		B			

- Still need to eliminate static-0 hazards: Use POS form

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Summary of hazards

- ◆ We can eliminate static hazards in 2-level logic
 - For single-bit changes
 - Eliminating static hazards also eliminates dynamic hazards
- ◆ Hazards are a difficult problem
 - Multiple-bit changes in 2-level logic are hard
 - Static hazards in multilevel logic are harder
 - Dynamic hazards in multilevel logic are harder yet
- ◆ CAD tools and simulation/testing are indispensable
 - Test vectors probe a design for hazards

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