

Boolean algebra

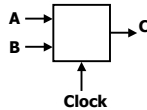
- ◆ Last lecture
 - Binary numbers
 - Base conversion
 - Number systems
 - ☞ Twos-complement
 - A/D and D/A conversion
- ◆ Today's lecture
 - Boolean algebra
 - ☞ Axioms
 - ☞ Useful laws and theorems
 - ☞ Simplifying Boolean expressions

Major topic: Combinational logic

- ◆ Axioms and theorems of Boolean algebra
- ◆ Logic functions and truth tables
 - AND, OR, Buffer, NAND, NOR, NOT, XOR, XNOR
- ◆ Gate logic
 - Networks of Boolean functions
- ◆ Canonical forms
 - Sum of products and product of sums
- ◆ Simplification
 - Boolean cubes and Karnaugh maps
 - Two-level simplification

Combinational versus sequential

- ◆ Combinational: Memoryless
 - Apply fixed inputs A, B
 - Wait for clock edge
 - Observe C
 - Wait for another clock edge
 - Observe C again: C will stay the same
- ◆ Sequential: With Memory
 - Apply fixed inputs A, B
 - Wait for clock edge
 - Observe C
 - Wait for another clock edge
 - Observe C again: C may be different



Boolean algebra

- ◆ A Boolean algebra comprises...
 - A set of elements B
 - Binary operators $\{+, \cdot\}$
 - A unary operation $\{\prime\}$
- ◆ ...and the following axioms
 1. The set B contains at least two elements $\{a, b\}$ with $a \neq b$
 2. Closure: $a+b$ is in B $a \cdot b$ is in B
 3. Commutative: $a+b = b+a$ $a \cdot b = b \cdot a$
 4. Associative: $a+(b+c) = (a+b)+c$ $a \cdot (b \cdot c) = (a \cdot b) \cdot c$
 5. Identity: $a+0 = a$ $a \cdot 1 = a$
 6. Distributive: $a+(b \cdot c) = (a+b) \cdot (a+c)$ $a \cdot (b+c) = (a \cdot b) + (a \cdot c)$
 7. Complementarity: $a+a' = 1$ $a \cdot a' = 0$

Digital (binary) logic is a Boolean algebra

- ◆ Substitute
 - $\{0, 1\}$ for B
 - AND for \cdot
 - OR for $+$
 - NOT for \prime
- ◆ All the axioms hold for binary logic
- ◆ Definitions
 - Boolean function
 - ☞ Maps inputs from the set $\{0,1\}$ to the set $\{0,1\}$
 - Boolean expression
 - ☞ An algebraic statement of Boolean variables and operators

AND, OR, Not

◆ AND	$X \cdot Y$	XY		<table border="1"> <thead> <tr><th>X</th><th>Y</th><th>Z</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	X	Y	Z	0	0	0	0	1	0	1	0	0	1	1	1
X	Y	Z																	
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◆ OR	$X+Y$			<table border="1"> <thead> <tr><th>X</th><th>Y</th><th>Z</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	X	Y	Z	0	0	0	0	1	1	1	0	1	1	1	1
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◆ NOT	\bar{X}	X'		<table border="1"> <thead> <tr><th>X</th><th>Y</th></tr> </thead> <tbody> <tr><td>0</td><td>1</td></tr> <tr><td>1</td><td>0</td></tr> </tbody> </table>	X	Y	0	1	1	0									
X	Y																		
0	1																		
1	0																		

Logic functions and Boolean algebra

- Any logic function that is expressible as a truth table can be written in Boolean algebra using +, •, and '.

X	Y	Z	Z = X•Y	X	Y	X'	Z	Z = X'•Y
0	0	0		0	0	1	0	
0	1	0		0	1	1	1	
1	0	0		1	0	0	0	
1	1	1		1	1	0	0	

X	Y	X'	Y'	X•Y	X'•Y'	Z	Z = (X•Y)+(X'•Y')
0	0	1	1	0	1	1	
0	1	1	0	0	0	0	
1	0	0	1	0	0	0	
1	1	0	0	1	0	1	

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Two key concepts

- Duality** (a meta-theorem— *a theorem about theorems*)
 - All Boolean expressions have logical duals
 - Any theorem that can be proved is also proved for its dual
 - Replace: • with +, + with •, 0 with 1, and 1 with 0
 - Leave the variables unchanged
- de Morgan's Theorem**
 - Procedure for complementing Boolean functions
 - Replace: • with +, + with •, 0 with 1, and 1 with 0
 - Replace all variables with their complements

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Useful laws and theorems

Identity:	$X + 0 = X$	Dual: $X \cdot 1 = X$
Null:	$X + 1 = 1$	Dual: $X \cdot 0 = 0$
Idempotent:	$X + X = X$	Dual: $X \cdot X = X$
Involution:	$(X')' = X$	
Complementarity:	$X + X' = 1$	Dual: $X \cdot X' = 0$
Commutative:	$X + Y = Y + X$	Dual: $X \cdot Y = Y \cdot X$
Associative:	$(X+Y)+Z=X+(Y+Z)$	Dual: $(X•Y)•Z=X•(Y•Z)$
Distributive:	$X•(Y+Z)=(X•Y)+(X•Z)$	Dual: $X+(Y•Z)=(X+Y)•(X+Z)$
Uniting:	$X•Y+X•Y'=X$	Dual: $(X+Y)•(X+Y')=X$

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Useful laws and theorems (con't)

Absorption:	$X+X•Y=X$	Dual: $X•(X+Y)=X$
Absorption (#2):	$(X+Y')•Y=X•Y$	Dual: $(X•Y')+Y=X+Y$
de Morgan's:	$(X+Y+...)'=X'•Y'•...$	Dual: $(X•Y•...)'=X'+Y'+...$
Duality:	$(X+Y+...)'•X•Y•...$	Dual: $(X•Y•...)'•X+Y+...$
Multiplying & factoring:	$(X+Y)•(X'+Z)=X•Z+X'•Y$	Dual: $X•Y+X'•Z=(X+Z)•(X'+Y)$
Consensus:	$(X•Y)+(Y•Z)+(X'•Z)=X•Y+X'•Z$	Dual: $(X+Y)•(Y+Z)•(X'+Z)=(X+Y)•(X'+Z)$

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Proving theorems

- Example 1: Prove the uniting theorem**

Distributive	$X•Y+X•Y' = X•(Y+Y')$
Complementarity	$= X•(1)$
Identity	$= X$
- Example 2: Prove the absorption theorem**

Identity	$X+X•Y = (X•1)+(X•Y)$
Distributive	$= X•(1+Y)$
Null	$= X•(1)$
Identity	$= X$

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Proving theorems

- Example 3: Prove the consensus theorem**

Complementarity	$XY+YZ+X'Z = XY+(X+X')YZ + X'Z$
Distributive	$= XYZ+XY+X'YZ+X'Z$
	<i>Use absorption {AB+A=A} with A=XY and B=Z</i>
	$= XY+X'YZ+X'Z$
Rearrange terms	$= XY+X'ZY+X'Z$
	<i>Use absorption {AB+A=A} with A=X'Z and B=Y</i>
	$XY+YZ+X'Z = XY+X'Z$

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de Morgan's Theorem

- ◆ Use de Morgan's Theorem to find complements
- ◆ Example: $F=(A+B) \cdot (A'+C)$, so $F'=(A \cdot B')+(A \cdot C')$

A	B	C	F	A	B	C	F'
0	0	0	0	0	0	0	1
0	0	1	0	0	0	1	1
0	1	0	1	0	1	0	0
0	1	1	1	0	1	1	0
1	0	0	0	1	0	0	1
1	0	1	1	1	0	1	0
1	1	0	0	1	1	0	1
1	1	1	1	1	1	1	0

Logic simplification

- ◆ Use the axioms to simplify logical expressions
 - Why? To use less hardware
- ◆ Example: A two-level logic expression

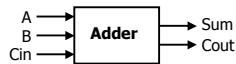
$$\begin{aligned}
 Z &= A'BC + AB'C' + AB'C + ABC + ABC \\
 &= AB'C + AB'C' + A'BC + ABC + ABC && \text{rearrange} \\
 &= AB'(C + C') + A'BC + AB(C + C) && \text{distributive} \\
 &= AB' + A'BC + AB && \text{comp.} \\
 &= AB' + AB + A'BC && \text{rearrange} \\
 &= A(B' + B) + A'BC && \text{distributive} \\
 &= A + A'BC && \text{comp.}
 \end{aligned}$$

Use absorption #2D $\{(X \cdot Y) + Y = X + Y\}$ with $X=BC$ and $Y=A$

$$Z = A + BC$$

Example: A full adder

- ◆ 1-bit binary adder
 - Inputs: A, B, Carry-in
 - Outputs: Sum, Carry-out



A	B	Cin	S	Cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

$$S = A'B'Cin + A'BCin' + AB'Cin' + ABCin$$

$$Cout = A'BCin + AB'Cin + ABCin' + ABCin$$

Simplifying the carry-out function

$$\begin{aligned}
 Cout &= A'BCin + AB'Cin + ABCin' + ABCin \\
 &= A'BCin + AB'Cin + ABCin' + ABCin + ABCin \\
 &= A'BCin + ABCin + AB'Cin + ABCin' + ABCin \\
 &= (A'+A)BCin + AB'Cin + ABCin' + ABCin \\
 &= (1)BCin + AB'Cin + ABCin' + ABCin && \text{associative} \\
 &= BCin + AB'Cin + ABCin' + ABCin \\
 &= BCin + AB'Cin + ABCin + ABCin' + ABCin \\
 &= BCin + A(B'+B)Cin + ABCin' + ABCin && \text{idempotent} \\
 &= BCin + A(1)Cin + ABCin' + ABCin \\
 &= BCin + ACin + AB'(Cin'+Cin) \\
 &= BCin + ACin + AB(1) \\
 &= BCin + ACin + AB
 \end{aligned}$$

Some notation

- ◆ Priorities: $\bar{A} \cdot B + C = ((\bar{A}) \cdot B) + C$
- ◆ Variables are sometimes called literals