



CSE341: Programming Languages

Lecture 3 Local Bindings; Options; Benefits of No Mutation

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Review

Huge progress already on the core pieces of ML:

- Types: `int bool unit t1*...*tn t list t1*...*tn->t`
 - Types “nest” (each `t` above can be itself a compound type)
- Variables, environments, and basic expressions
- Functions
 - Build: `fun x0 (x1:t1, ..., xn:tn) = e`
 - Use: `e0 (e1, ..., en)`
- Tuples
 - Build: `(e1, ..., en)`
 - Use: `#1 e, #2 e, ...`
- Lists
 - Build: `[] e1::e2`
 - Use: `null e hd e tl e`

Today

- The big thing we need: **local bindings**
 - For style and convenience
 - A big but natural idea: nested function bindings
 - For efficiency (**not** “just a little faster”)
- One last feature for last problem of homework 1: **options**
- Why **not having mutation** (assignment statements) is a valuable language feature
 - No need for you to keep track of sharing/aliasing, which Java programmers must obsess about

Let-expressions

3 questions:

- Syntax: `let b1 b2 ... bn in e end`
 - Each ***bi*** is any *binding* and ***e*** is any *expression*
- Type-checking: Type-check each ***bi*** and ***e*** in a static environment that includes the previous bindings.
Type of whole let-expression is the type of ***e***.
- Evaluation: Evaluate each ***bi*** and ***e*** in a dynamic environment that includes the previous bindings.
Result of whole let-expression is result of evaluating ***e***.

It is an expression

A let-expression is ***just an expression***, so we can use it ***anywhere*** an expression can go

Silly examples

```
fun silly1 (z : int) =  
  let val x = if z > 0 then z else 34  
      val y = x+z+9  
  in  
    if x > y then x*2 else y*y  
  end  
fun silly2 () =  
  let val x = 1  
  in  
    (let val x = 2 in x+1 end) +  
    (let val y = x+2 in y+1 end)  
  end
```

`silly2` is poor style but shows let-expressions are expressions

- Can also use them in function-call arguments, if branches, etc.
- Also notice shadowing

What's new

- What's new is **scope**: where a binding is in the environment
 - *In* later bindings and body of the let-expression
 - (Unless a later or nested binding shadows it)
 - *Only in* later bindings and body of the let-expression
- *Nothing else is new*:
 - Can put any binding we want, even function bindings
 - Type-check and evaluate just like at “top-level”

Any binding

According to our rules for let-expressions, we can define functions inside any let-expression

```
let b1 b2 ... bn in e end
```

This is a natural idea, and often good style

(Inferior) Example

```
fun countup_from1 (x : int) =  
  let fun count (from : int, to : int) =  
        if from = to  
        then to :: []  
        else from :: count(from+1, to)  
      in  
        count (1, x)  
      end
```

- This shows how to use a local function binding, but:
 - Better version on next slide
 - `count` might be useful elsewhere

Better:

```
fun countup_from1_better (x : int) =  
  let fun count (from : int) =  
        if from = x  
        then x :: []  
        else from :: count(from+1)  
      in  
        count 1  
      end
```

- Functions can use bindings in the environment where they are defined:
 - Bindings from “outer” environments
 - Such as parameters to the outer function
 - Earlier bindings in the let-expression
- Unnecessary parameters are usually bad style
 - Like `to` in previous example

Nested functions: style

- Good style to define helper functions inside the functions they help if they are:
 - Unlikely to be useful elsewhere
 - Likely to be misused if available elsewhere
 - Likely to be changed or removed later
- A fundamental trade-off in code design: reusing code saves effort and avoids bugs, but makes the reused code harder to change later

Avoid repeated recursion

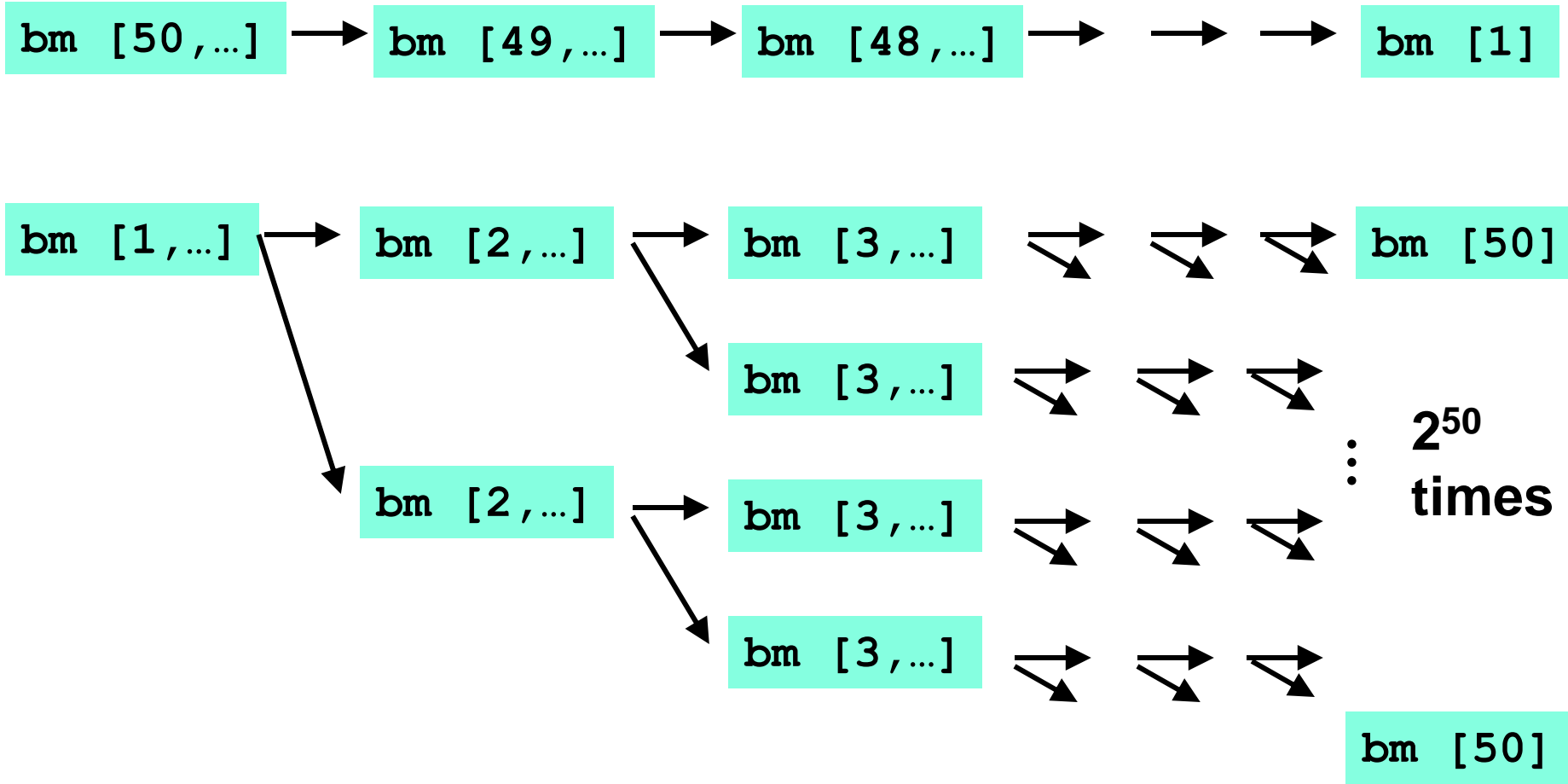
Consider this code and the recursive calls it makes

- Don't worry about calls to `null`, `hd`, and `tl` because they do a small constant amount of work

```
fun bad_max (xs : int list) =  
  if null xs  
  then 0 (* horrible style; fix later *)  
  else if null (tl xs)  
  then hd xs  
  else if hd xs > bad_max (tl xs)  
  then hd xs  
  else bad_max (tl xs)  
  
let x = bad_max [50,49,...,1]  
let y = bad_max [1,2,...,50]
```

Fast vs. unusable

```
if hd xs > bad_max (tl xs)
then hd xs
else bad_max (tl xs)
```



Math never lies

Suppose one `bad_max` call's if-then-else logic and calls to `hd`, `null`, `t1` take 10^{-7} seconds

- Then `bad_max [50, 49, ..., 1]` takes 50×10^{-7} seconds
- And `bad_max [1, 2, ..., 50]` takes 1.12×10^8 seconds
 - (over 3.5 years)
 - `bad_max [1, 2, ..., 55]` takes over 1 century
 - Buying a faster computer won't help much 😊

The key is not to do repeated work that might do repeated work that might do...

- Saving recursive results in local bindings is essential...

Efficient max

```
fun good_max (xs : int list) =  
  if null xs  
  then 0 (* horrible style; fix later *)  
  else if null (tl xs)  
  then hd xs  
  else  
    let val tl_ans = good_max(tl xs)  
    in  
      if hd xs > tl_ans  
      then hd xs  
      else tl_ans  
    end  
end
```

Fast vs. fast

```
let val tl_ans = good_max(tl xs)
in
  if hd xs > tl_ans
  then hd xs
  else tl_ans
end
```

gm [50, ...] → gm [49, ...] → gm [48, ...] → → → gm [1]

gm [1, ...] → gm [2, ...] → gm [3, ...] → → → gm [50]

Options

- `t option` is a type for any type `t`
 - (much like `t list`, but a different type, not a list)

Building:

- `NONE` has type `'a option` (much like `[]` has type `'a list`)
- `SOME e` has type `t option` if `e` has type `t` (much like `e :: []`)

Accessing:

- `isSome` has type `'a option -> bool`
- `valOf` has type `'a option -> 'a` (exception if given `NONE`)

Example

```
fun better_max (xs : int list) =
  if null xs
  then NONE
  else
    let val tl_ans = better_max(tl xs)
    in
      if isSome tl_ans
      andalso valOf tl_ans > hd xs
      then tl_ans
      else SOME (hd xs)
    end
end
```

```
val better_max = fn : int list -> int option
```

- Nothing wrong with this, but as a matter of style might prefer not to do so much useless “`valOf`” in the recursion

Example variation

```
fun better_max2 (xs : int list) =
  if null xs
  then NONE
  else let (* ok to assume xs nonempty b/c local *)
        fun max_nonempty (xs : int list) =
          if null (tl xs)
          then hd xs
          else
            let val tl_ans = max_nonempty (tl xs)
            in
              if hd xs > tl_ans
              then hd xs
              else tl_ans
            end
        in
          SOME (max_nonempty xs)
        end
```

Cannot tell if you copy

```
fun sort_pair (pr : int * int) =  
  if #1 pr < #2 pr  
  then pr  
  else (#2 pr, #1 pr)
```

```
fun sort_pair (pr : int * int) =  
  if #1 pr < #2 pr  
  then (#1 pr, #2 pr)  
  else (#2 pr, #1 pr)
```

In ML, these two implementations of `sort_pair` are **indistinguishable**

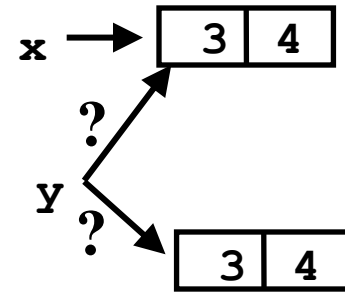
- But only because tuples are immutable
- The first is better style: simpler and avoids making a new pair in the then-branch
- In languages with mutable compound data, these are different!

Suppose we had mutation...

```
val x = (3,4)
val y = sort_pair x
```

somehow mutate #1 x to hold 5

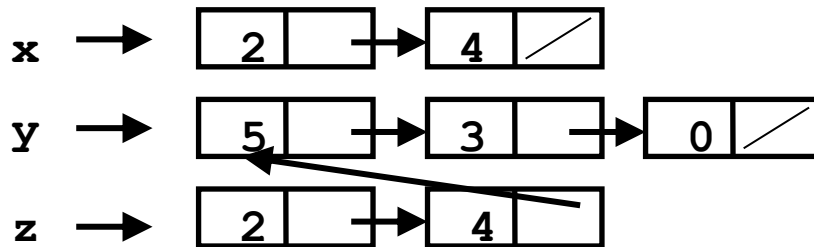
```
val z = #1 y
```



- What is **z**?
 - Would depend on how we implemented `sort_pair`
 - Would have to decide carefully and document `sort_pair`
 - But without mutation, we can implement “either way”
 - No code can ever distinguish aliasing vs. identical copies
 - No need to think about aliasing: focus on other things
 - Can use aliasing, which saves space, without danger

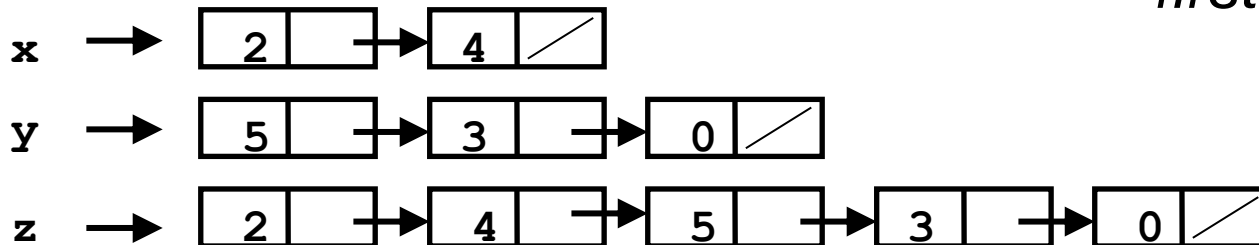
An even better example

```
fun append (xs : int list, ys : int list) =  
  if null xs  
  then ys  
  else hd (xs) :: append (tl(xs), ys)  
val x = [2,4]  
val y = [5,3,0]  
val z = append(x,y)
```



*(can't tell,
but it's the
first one)*

or



ML vs. Imperative Languages

- In ML, we create aliases all the time without thinking about it because it is *impossible* to tell where there is aliasing
 - Example: `tl` is constant time; does not copy rest of the list
 - So don't worry and focus on your algorithm
- In languages with mutable data (e.g., Java), programmers are *obsessed* with aliasing and object identity
 - They have to be (!) so that subsequent assignments affect the right parts of the program
 - Often crucial to make copies in just the right places
 - Consider a Java example...

Java security nightmare (bad code)

```
class ProtectedResource {
    private Resource theResource = ...;
    private String[] allowedUsers = ...;
    public String[] getAllowedUsers() {
        return allowedUsers;
    }
    public String currentUser() { ... }
    public void useTheResource() {
        for(int i=0; i < allowedUsers.length; i++) {
            if(currentUser().equals(allowedUsers[i])) {
                ... // access allowed: use it
                return;
            }
        }
        throw new IllegalAccessException();
    }
}
```


Have to make copies

The problem:

```
p.getAllowedUsers()[0] = p.currentUser();  
p.useTheResource();
```

The fix:

```
public String[] getAllowedUsers() {  
    ... return a copy of allowedUsers ...  
}
```

Reference (alias) vs. copy doesn't matter if code is immutable!