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## Call-by: Best of both worlds?

call-by-value: eval every argument before call

call-by-name: eval arguments at every actual use (via thunk)

*call-by-need* ("lazy" evaluation): Evaluate every argument the first time it's used. Save answer for subsequent uses.

- Asymptotically it's the best
- But behind-the-scenes bookkeeping can be costly
- And it's hard to reason about with effects
  - Typically used in (sub)languages without effects
- Nonetheless, a key idiom with syntactic support in Scheme
  - And related to *memoization*

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## CSE 341: Programming Languages

Spring 2006

Lecture 18 — Delayed Evaluation, Memoization & Streams

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### Memoization

A "cache" of previous results is equivalent if results cannot change.

- Could be slower: cache too big or computation too cheap
- Could be faster: just a lookup
  - On homework: An example where it's a *lot* faster by preventing an exponential explosion.

An association list is not the fastest data structure for large memo tables, but works fine for 341.

Question: Why does `assoc` return the pair?

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### Streams

- A stream is an "infinite" list — you can ask for the rest of it as many times as you like and you'll never get null.
- The universe is finite, so a stream must really be an object that acts like an infinite list.
- The idea: use a function to describe what comes next.

Note: Deep connection to sequential feedback circuits

Note: Connection to UNIX pipes

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