

---

## Today

- (course motivation)
- Summarize what we've learned with a concise and well-known notation for recursively-defined language constructs
- Begin first-class functions

## CSE 341: Programming Languages

Autumn 2005

Lecture 7 — Motivation; BNF; First-Class Functions

CSE 341 Autumn 2005, Lecture 7

1

CSE 341 Autumn 2005, Lecture 7

2

---

## Why these 3 Languages?

*Functional programming* (ML, Scheme) encourages recursion, discourages mutation, provides elegant, lightweight support for first-class code. Support for extensibility complements OO.

- ML has a polymorphic type system (vindication imminent!) complementary to OO-style subtyping, a rich module system for abstract types, and rich pattern-matching.
- Scheme has dynamic typing, “good” macros, fascinating control operators, and a minimalist design.
- Smalltalk has classes but not types, an unconventional environment, and a complete commitment to OO.

CSE 341 Autumn 2005, Lecture 7

3

---

## Runners-Up

Runners-up: Miranda (laziness, simplicity and purity), Haskell (laziness, nonproprietary implementation, but difficult type system), Prolog (unification and backtracking), CPR(R) (everything that Prolog has plus constraints).

There are thousands of other languages as well ...

CSE 341 Autumn 2005, Lecture 7

4

## Why not some popular ones?

- Java: you've already studied it in 142/143. We'll look at some additional features at the end of course and compare it with the other languages we've been studying (e.g., interfaces, anonymous inner classes, container types)
- C: lots of "implementation-dependent" behavior (a bad property), and we have CSE303
- C++: an enormous language, and unsafe like C
- Perl: advantages (strings, files, ...) not foci of this course. Python or Ruby would be closer.

CSE 341 Autumn 2005, Lecture 7

5

CSE 341 Autumn 2005, Lecture 7

6

## Are these useful?

We focus on *interesting language concepts* in ML/Scheme/Smalltalk. "Real" programming needs file I/O, strings, floating-point, graphics libraries, project managers, unit testers, threads, foreign-function interfaces, ...

- These languages have all that and more!
- Just not course focus

## Summary and Some Notation

Learned the syntax, typing rules, and semantics for (a big) part of ML

Can summarize *abstract syntax* with (E)BNF. Informally:

```
t ::= int | bool | unit | dtype
    | t1 -> t2 | t1 * t2 | {x1=t1, ..., xn=tn}
e ::= 34 | x | (e1, e2) | if e1 then e2 else e3
    | let b1 ... bn in e end | e1 e2
    | case e of p1 => e1 | ... | pn => en
    | e1 + e2 | {x1=e1, ..., xn=en} | C e
b ::= val p = e | fun f p = e
    | datatype dtype = C1 of t1 | ... | Cn of tn
p ::= 34 | x | _ | C p | (p1, p2) | {x1=p1, ..., xn=pn}
```

Things left out of this *grammar*: *n*-tuples, field-accessors, floating-point, boolean constants, andalso/orelse, lists, ...

CSE 341 Autumn 2005, Lecture 7

7

CSE 341 Autumn 2005, Lecture 7

8

## First-Class Functions

- Functions are values. (Variables in the environment are bound to them.)
- We can pass functions to other functions.
  - *Factor* common parts and *abstract* different parts.
- We can return functions as values from other functions.

## Type Inference and Polymorphism

---

ML can infer function types based on function bodies. Possibilities:

- The argument/result must be one specific type.
- The argument/result can be *any* type, but may have to be the *same type* as other parts of argument/result.
- Some hand-waving about “equality types”

We will study this *parametric polymorphism* more later.

Without it, ML would be a pain (e.g., a different list library for every list-element type).

Curious fact: If  $f: \text{int} \rightarrow \text{int}$ , there are lots of values  $f$  could return. If  $f: 'a \rightarrow 'a$ , whenever  $f$  returns, it returns its argument!

CSE 341 Autumn 2005, Lecture 7

9

## Returning Functions

---

Syntax note:  $\rightarrow$  “associates to the right”

- $t1 \rightarrow t2 \rightarrow t3$  means  $t1 \rightarrow (t2 \rightarrow t3)$

Again, there is nothing new here.

The key question: What about *free variables* in a function value?

What *environment* do we use to *evaluate* them?

Are such free variables useful?

You must understand the answers to move beyond being a novice programmer.

CSE 341 Autumn 2005, Lecture 7

11

## Anonymous Functions

---

As usual, we can write functions anywhere we write expressions.

- We already could:  
(`let fun f x = e in f end`)
- Here is a more concise way (better style when possible):  
(`fn x => e`)
- Cannot do this for recursive functions (why?)

CSE 341 Autumn 2005, Lecture 7

10