

CSE 341: Programming Languages

Autumn 2005

Lecture 5 — Type synonyms, more pattern-matching, accumulators

Goals

- Contrast type synonyms with new types
- See pattern-matching for built-in “one of” types (not really a concept, but important for ML programming) and “each of” types
- Investigate why accumulator-style recursion can be more efficient

Type synonyms

You can bind a *type name* to a type. Example:

```
type intpair = int * int
```

(We call something else a *type variable*.)

In ML, this creates a *synonym*, also known as a *transparent* type definition. Recursion not allowed.

So a type name is *equivalent* to its definition.

To contrast, the type a datatype binding introduces is not equivalent to any other type (until possibly a later type binding).

Review: datatypes and pattern-matching

Evaluation rules for datatype bindings and case expressions:

datatype $t = C_1 \text{ of } t_1 \mid C_2 \text{ of } t_2 \mid \dots \mid C_n \text{ of } t_n$

Adds constructors C_i where $C_i \ v$ is a value (and C_i has type $t_i \rightarrow t$).

case e of $p_1 \Rightarrow e_1 \mid p_2 \Rightarrow e_2 \mid \dots \mid p_n \Rightarrow e_n$

- Evaluate e to v
- If p_i is the first pattern to *match* v , then result is evaluation of e_i in environment extended by the match.
- If C is a constructor of type $t_1 * \dots * t_n \rightarrow t$, then $C(x_1, \dots, x_n)$ is a pattern that matches $C(v_1, \dots, v_n)$ and the match extends the environment with x_1 to $v_1 \dots x_n$ to v_n .
- Coming soon: many more pattern forms.

Why patterns?

Even without more pattern forms, this design has advantages over functions for “testing and destructing” (e.g., `null`, `hd`, and `tl`):

- easier to check for missing and redundant cases
- more concise syntax by combining “test, destruct, and bind”
- you can easily define testing and destructing in terms of pattern-matching

In fact, case expressions are the preferred way to test variants and extract values from all ML’s “one-of” types, including predefined ones (`[]` and `::` just funny syntax).

So: Do *not* use functions `hd`, `tl`, `null`, `isSome`, `valOf`

Teaser: These functions are useful for *passing as values*

Tuple/record patterns

You can also use patterns to extract fields from tuples and records:

pattern `{f1=x1, ..., fn=xn}` (or `(x1, ..., xn)`) matches `{f1=v1, ..., fn=vn}` (or `(v1, ..., vn)`).

For record-patterns, field-order does not matter.

This is better style than `#1` and `#foo`, and it means you do not (ever) need to write function-argument types.

Instead of a case with one pattern, better style is a pattern directly in a `val` binding.

Next time: “deep” (i.e., nested) patterns.

Recursion

You should now have the hang of recursion:

- It's no harder than using a loop (whatever that is)
- It's much easier when you have multiple recursive calls (e.g., with functions over ropes or trees)

But there are idioms you should learn for *elegance*, *efficiency*, and *understandability*.

Today: using an *accumulator*.

Accumulator lessons

- Accumulators can avoid data-structure copying
- Accumulators can reduce the depth of recursive calls that are not *tail calls*
- Key idioms:
 - Non-accumulator: compute recursive results and combine
 - Accumulator: use recursive result as new accumulator
 - The base case becomes the initial accumulator

You will use recursion in non-functional languages—this lesson still applies.

Let's investigate the evaluation of `to_list_1` and `to_list_2`.