

CSE 341 - Programming Languages

Final Exam - Autumn 2005

Your Name:

Open book and notes. No laptop computers, PDAs, or similar devices. (Calculators are OK, although you won't need one.) Please answer the problems on the exam paper — if you need extra space use the back of a page.

110 minutes, 100 points total

1. (14 points)

Suppose the following ML functions and exception have been defined:

```
exception UnequalLengthLists;

fun zip2 ([], []) = []
| zip2 ((x::xs), (y::ys)) = (x,y) :: zip2 (xs, ys)
| zip2 _ = raise UnequalLengthLists;

fun zip2curried [] [] = []
| zip2curried (x::xs) (y::ys) = (x,y) :: zip2curried xs ys
| zip2curried _ _ = raise UnequalLengthLists;

fun average (x,y) = (x+y)/2.0;
```

What is the *value* of each of the following expressions? If evaluating it raises an exception, say so. (Hint: this code uses the built-in function `map` in ML, which is curried.)

- (a) `zip2([1.0, 2.0, 3.0], [3.0, 4.0, 5.0])`
- (b) `map average (zip2([1.0, 2.0, 3.0], [3.0, 4.0, 5.0]))`
- (c) `zip2([1.0, 2.0, 3.0], [3.0, 4.0])`

What is the *type* of each of the following expressions? Some of them may give type errors — if so, say that.

- (a) `zip2`
- (b) `zip2curried`
- (c) `zip2 average`
- (d) `map average`

2. (8 points - 2 points each for parts a and b; 4 points for part c) What is the result of evaluating each of the following Scheme expressions? If there is more than one expression, just give the result of evaluating the final expression.

(a)

```
(define a 3)
(define b 4)
(define c 5)
(let ((b (+ 2 4))
      (c (+ b 10)))
  (+ a b c))
```

(b)

```
(define a 3)
(define b 4)
(define c 5)
(let* ((b (+ 2 4))
       (c (+ b 10)))
  (+ a b c))
```

(c)

```
(define z 'y)
(define y 'x)
(define x 42)
(define w 'w)
(list
  z
  (eval z)
  (eval (eval z))
  (eval (eval (eval z))))
  w
  (eval w)
  (eval (eval w))
  (eval (eval (eval w))))
```

3. (6 points)

(a) What is the result of evaluating the following expressions in Scheme? (Just give the value of the final expression.)

```
(define y 3)
(define (f x) (+ x y))
(let ((x 10)
      (y 20))
  (f 100))
```

(b) Suppose Scheme used dynamic scoping rather than lexical scoping. In that case, what would be the value of the final expression?

4. (12 points) Write a Scheme definition for my-cond, a macro for the built-in Scheme “cond”. Use define-syntax. The definition for cond is as follows:

```
(cond <clause1> <clause2> ...)
```

Syntax: Each <clause> should be of the form

```
(<test> <expression1> ...)
```

where <test> is any expression. The last <clause> may be an “else clause,” which has the form

```
(else <expression1> <expression2> ...).
```

Semantics: A cond expression is evaluated by evaluating the <test> expressions of successive <clause>s in order until one of them evaluates to a true value. When a <test> evaluates to a true value, then the remaining <expression>s in its <clause> are evaluated in order, and the result of the last <expression> in the <clause> is returned as the result of the entire cond expression. If the selected <clause> contains only the <test> and no <expression>s, then the value of the <test> is returned as the result. If there are no true <clause>s, cond should return nothing. In your macro, you can get this result with the expression (void).

Macros: As a refresher on macro syntax, here are a couple examples we saw in class:

```
(define-syntax my-if                ; macro name
  (syntax-rules (then else)       ; literals it uses, if any
    ((my-if e1 then e2 else e3)   ; pattern
     (if e1 e2 e3))))           ; template
```

```
(define-syntax my-or
  (syntax-rules ()
    ((my-or) #f)
    ((my-or e) e)
    ((my-or e1 e2 ...)
     (let ((temp e1))
       (if temp
           temp
           (my-or e2 ...))))))
```

5. (6 points)

(a) Suppose that you have a global variable `k` declared in Scheme:

```
(define k #f)
```

What is the value returned by evaluating the following Scheme expression?

```
(+ 10
  (call/cc
   (lambda (c)
     (set! k c)
     (* 5 6))))
20)
```

(b) After the above expression has been evaluated and `k` reassigned, what is the value returned by this expression?

```
(k 100)
```

(c) And this one?

```
(* 3 (k 200))
```

6. (12 points) Tacky but easy-to-grade true/false questions!

- (a) An advantage of static typing over dynamic typing is that it is more flexible — more programs can execute correctly.
- (b) An advantage of static typing over dynamic typing is that more errors can be caught at compile time rather than run time.
- (c) An advantage of static typing over dynamic typing is that type declarations provide machine-checkable documentation.
- (d) After a continuation has been resumed in Scheme, it is used up — it can't be resumed again.
- (e) In a pure functional language, any function will return the same answers if call-by-name is used instead of lazy evaluation.
- (f) In a pure functional language, any function will return the same answers if call-by-name is used instead of call-by-value.

7. (6 points) Discuss two reasons **inner classes** are useful for defining iterators in Java.

8. (8 points) Consider the following Smalltalk class definitions.

```
Object subclass: #Animal
  classVariableNames: ''
  poolDictionaries: ''

describe
  Transcript show: 'An animal'.
  self additionalDescription.

additionalDescription
  Transcript show: ' - no further information available'.
```

```
Animal subclass: #SeaCreature
  instanceVariableNames: ''
  classVariableNames: ''
  poolDictionaries: ''

additionalDescription
  Transcript show: ' that lives in the sea'.
  Transcript cr.
  self yetMoreDescription.

yetMoreDescription
  Transcript show: 'Also a source of 341 variable names'.
```

```
SeaCreature subclass: #Octopus
  instanceVariableNames: ''
  classVariableNames: ''
  poolDictionaries: ''

additionalDescription
  Transcript show: ' that has tentacles and'.
  super additionalDescription.
```

```
Octopus subclass: #GiantOctopus
  instanceVariableNames: ''
  classVariableNames: ''
  poolDictionaries: ''

yetMoreDescription
  Transcript show: 'It can grow up to 300 pounds'.
```

What is printed when each of the following expressions is evaluated?

- `Animal new describe`

- `SeaCreature new describe`

- `Octopus new describe`

- `GiantOctopus new describe`

9. (16 points)

Consider the following Java code fragments. In each case, does the code compile correctly? If so, does it execute without error, or is there an exception? (Hints: the `List` interface includes an `add` method. The class `ArrayList` implements the `List` interface.)

- (a)

```
List<String> s = new ArrayList<String>();  
s.add("squid");
```

- (b)

```
ArrayList<Object> s = new ArrayList<String>();  
s.add("squid");
```

- (c)

```
ArrayList<String> s = new ArrayList<Object>();  
s.add("squid");
```

- (d)

```
String [] s = new String [10];  
s[0] = "squid";
```

- (e)

```
Object [] s = new String[10];  
s[0] = "squid";
```

- (f)

```
String [] s = new Object[10];  
s[0] = "squid";
```

- (g)

```
String [] s = new String [10];  
s[0] = "squid";  
s[1] = new Point(10,20);
```

- (h)

```
Object [] s = new String[10];  
s[0] = "squid";  
s[1] = new Point(10,20);
```

10. (12 points) Suppose that we have two definitions of a static method `shapes` in Java:

```
public static boolean shapes1
    (ArrayList<? extends RectangularShape> s) {
    ....
}
```

```
public static boolean shapes2
    (ArrayList<RectangularShape> s) {
    ....
}
```

Suppose we also have some variables declared:

```
ArrayList<RectangularShape> a;
ArrayList<Rectangle2D> b;
ArrayList<Ellipse2D> c;
ArrayList<Object> d;
```

(`Rectangle2D` and `Ellipse2D` are both subclasses of `RectangularShape`.)

(a) Which of the following method calls are legal? Circle the legal ones:

`shapes1(a);`

`shapes1(b);`

`shapes1(c);`

`shapes1(d);`

`shapes2(a);`

`shapes2(b);`

`shapes2(c);`

`shapes2(d);`

(b) Within the method `shapes1`, would it be legal to add a new `Rectangle2D` to the array list `s`?

(c) Within the method `shapes2`, would it be legal to add a new `Rectangle2D` to the array list `s`?