

CSE 341: Programming Languages

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Lecture 16— `define-struct`, `let/cc` for exceptions

Data in Scheme

Recall ML's approach to each-of, one-of, and self-referential types.

Pure Scheme's approach:

- There is One Big Datatype with built-in predicates.
- Use pairs (lists) for each-of types.
- Primitives implicitly raise errors for “wrong variant”
- Use helper functions like `caddr` and your own.

We'll discuss advantages/disadvantages next week.

define-struct

MzScheme extends Scheme with `define-struct`, e.g.:

```
(define-struct square (x y))  
(define-struct piece (squares))
```

Semantics:

- Binds constructors (`make-square`, `make-piece`) that take arguments and make values.
- Binds predicates (`square?`, `piece?`) that take one argument and return `#t` only for values built from the right constructor.
- Binds accessors (`square-x`, `square-y`, `piece-squares`) that take one argument, return the appropriate field, and call error for values not built from the right constructor.
- Binds mutators (`set-square-x!`, `set-square-y!`, `set-piece-squares!`).

define-struct is special

`define-struct` creates a new variant for the One Big Datatype.

Claim: `define-struct` is not a function.

Claim: `define-struct` is not a macro.

It could be a macro except for one key bit of its semantics: Values built from the constructor cause every *other* predicate (including all built-in ones) to return `#f`.

Advantage: abstraction

Disadvantage: Can't write "generic" code that has a case for every possible variant in every Scheme program.

Idiom for ML datatypes

Instead of a datatype with n constructors, you just use `define-struct` n times.

That “these n go together” is just convention.

Instead of `case`, you have a `cond` with n predicates and one “catch-all” error case.

Exceptions in Scheme

Recall exceptions in Java, ML: Transfer control to nearest *dynamically scoped* exception handler (i.e., nearest on “call stack”).

Transfer control: Forget what you’re doing. Result of entire program is now result of the handle (catch) in the “call stack” that existed when the handler was reached.

Scheme has a *more powerful* concept that can be a little less convenient for exceptions:

- You explicitly indicate what “handler” (*continuation*) to transfer control to.
- You do the transfer via a function application (that does not have function-application semantics)
- The continuation does not even have to be on the “call stack” when it’s transferred to!

Continuations for exceptions

Plan:

- Show how to use continuations for exceptions (needed for hw 5).
- Explain continuation-semantics “from scratch” (later)
- Hint at some advanced uses (later)

Syntax:

```
(let/cc k e) ; bind k to ‘‘current continuation’’  
(k e) ; ‘‘invoke’’ continuation bound to k
```

Exception idiom:

- Instead of handler, use `let/cc`
- Pass an appropriate function that invokes `k` to any function that needs to “raise”