## CSE 332 Data Abstractions, Winter 2013 Homework 1

Due: Friday, Jan 18, 2013 at the BEGINNING of lecture. Your work should be readable as well as correct. You should refer to the written homework guidelines on the course website. This assignment has 6, count them SIX fabulous questions! Have fun!

## Problem 1. Some Important Sums

Much of algorithm analysis involves working with summations. Fortunately, finding the closedform solutions to these sums often involve the simple algebraic manipulation. In the following two problems, you will compute two sums and prove a third to be true. They are the same in the 2nd and 3rd ed. of Weiss (p. 27 in 3rd ed., p. 26 in 2nd ed.).
(a) Weiss 1.8 a
(b) Weiss 1.8 b
(c) Weiss 1.12a

## Problem 2. Recurrence Relations

Consider the following recurrence relation: $T(1)=6$, and for $n>1, T(n)=1+2 T(\lfloor n / 2\rfloor)$
Note. $\lfloor x\rfloor$ is the the floor function. It rounds $x$ down to the nearest integer.
(a) Determine the value for $T(n)$ for integers $n$ from 1 to 8 .
(b) Expand the recurrence relation to get the closed form. Show your work; do not just show the final equation. For arithmetic simplicity, you may assume $n$ is a sufficiently large power of 2 such that the floor function does not lead to rounding issues.

## Problem 3. Budgeting Time

For each function $f(n)$ and time $t$ in the following table, determine the largest size $n$ of a problem that can be solved in time $t$, assuming that the algorithm to solve the problem takes $f(n)$ microseconds ( 1 second equals 1 million microseconds). For large entries (say, those that warrant scientific notation), an estimate is sufficient. Note that for one of the rows, you will not be able to solve it analytically, and will need a calculator, spreadsheet, or small program.

| $\mathbf{f}(\mathbf{n})$ | 1 second | 1 minute | 1 hour $\quad$ 1 day | 1 month | 1 year |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $500 \log _{2} n$ |  |  |  |  |  |
| $1000 n$ |  |  |  |  |  |
| $100 n \log _{2} n$ |  |  |  |  |  |
| $10 n^{2}$ |  |  |  |  |  |
| $2 n^{3}$ |  |  |  |  |  |
| $\frac{1}{20} 2^{n}$ |  |  |  |  |  |

This problem gives an orthogonal view of comparative running times from that given in lecture. Be sure to look at the patterns in your table when you have completed it.

## Problem 4. Super Crazy Fun with Induction

The following statement is clearly not true. Can you spot the error in the inductive "proof" below? Specify which of the 5 numbered lines are wrong, and clearly describe the error.

## All umbrellas are the same color

"Proof": The proof is by induction on $n$, the number of umbrellas in a set.
Base case ( $n=1$ ):
(1) If there is only one umbrella in the set, then the statement trivially holds.

Induction Step $(n=k+1)$ :
Assume the statement holds for $n=k$. Now suppose you have $k+1$ umbrellas.
(2) Set the first umbrella aside. The remaining $k$ must be the same color (lets say blue).
(3) All we have to do now is show that the first one is also blue.
(4) To do this, remove a second umbrella and put the first umbrella back in to form a new set of size $k$. By the inductive hypothesis, all the umbrellas in this new set are also the same color.
(5) Since this set contains $k-1$ umbrellas that we already know are blue, it follows that they are all blue (including the first).

## Problem 5. Big- $O$, Big- $\Theta$

For each of the following statements, use our formal denitions of $\operatorname{Big}-O$, $\operatorname{Big}-\Theta$, and $\operatorname{Big}-\Omega$ either to prove the statement is true or to explain why it is false.
(a) If we have an algorithm that runs in $O(n)$ time and make some changes that cause it to run 10 times slower for all inputs, it will still run in $O(n)$ time.
(b) If $f(n)=O(g(n))$ and $h(n)=O(k(n))$, then $f(n) h(n)=O(g(n) k(n))$.
(c) If $f(n)=O(g(n))$ and $h(n)=O(k(n))$, then $f(n)+h(n)=O(g(n)+k(n))$.
(d) $2^{n+3}=\Theta\left(2^{n}\right)$
(e) $\left(2^{n}\right)^{1 / 3}=\Theta\left(2^{n}\right)$

## Problem 6. Algorithm Analysis

These problems are the same in both the 2nd and 3rd edition of Weiss.
(a) Weiss 2.7a (Give the best big- $O$ bound you can for the six program fragments. You do not have to explain why.)
(b) Weiss 2.11

