

CSE 332 Data Abstractions, Spring 2013

Homework 8

Due: **Friday, June 7, 2013** at the BEGINNING of lecture. Your work should be readable as well as correct. Homework Eight has two terrific questions!! Please write your **section** at the top of your homework.

Problem 1. Simple Concurrency with B-Trees

Note: Real databases and file systems use very fancy fine-grained synchronization for B-Trees such as “hand-over-hand locking” (which we did not discuss), but this problem considers some relatively simpler approaches.

Suppose we have a B Tree supporting operations insert and lookup (deletion is NOT a supported operation on this particular B Tree). A simple way to synchronize threads accessing the tree would be to have one lock for the entire tree that both operations acquire/release.

- (a) Suppose instead we have one lock per node in the tree. Each operation acquires the locks along the path to the leaf it needs and then at the end releases all these locks. Explain why this allegedly more fine-grained approach provides absolutely no benefit.
- (b) Now suppose we have one *readers/writer lock* per node and lookup acquires a read lock for each node it encounters whereas insert acquires a write lock for each node it encounters. Assume that the same policy from part (a) is followed in that a thread will only release all of its locks when it is done with its operation. How does this provide more concurrent access than the approach in part (a)? Is it any better than having one readers/writer lock for the whole tree (explain)?
- (c) Now suppose we modify the approach in part (b) so that insert acquires a write lock only for the leaf node (and read locks for other nodes it encounters). How would this approach increase concurrent access? When would this be incorrect? Explain how to fix this approach without changing the asymptotic complexity of insert by detecting when it is incorrect and in (only) those cases, starting the insert over using the approach in part (b) for that insert. Why would reverting to the approach in part (b) be fairly rare?

Problem 2. Minimum Spanning Trees

- (a) Weiss, problem 9.15(a). For Prim’s algorithm, start with vertex A, show the resulting table (see Figure 9.57 of the 3rd edition as an example, which is Figure 9.55 in the 2nd edition), ***and*** indicate the order in which vertices are added. For Kruskal’s algorithm, produce a table, similar to Figure 9.58 in the 3rd edition, which is Figure 9.56 in the 2nd edition. Ties may be broken arbitrarily.
- (b) Weiss, problem 9.15(b). Is this MST unique? Why? Be specific in your answer. If no, show another MST, if yes, explain why.