CSE 331: Software Design & Engineering — Spring 2023

Midterm: Section A — 5/19/2023 10:30am

Name: ______________________

UW Email: ____________`@uw.edu`

This exam contains 9 pages (including this cover page) and 4 problems. Check to see if any pages are missing. Enter all requested information on the top of this page.

Instructions:

- Closed book, closed notes, no cell phones, no calculators.
- You have **50 minutes** to complete the exam.
- Answer all problems on the exam paper.
- If you need extra space use the back of a page.
- Problems are not of equal difficulty; if you get stuck on a problem, move on.
- It may be to your advantage to read all the problems before beginning the exam.

Problem	Points	Score
1	18	
2	18	
3	28	
4	16	
Total:	80	

The following function `findIndex` searches for a string in an array of strings that is promised to be sorted in **decreasing** order. In other words, we are promised that $A[0] \geq A[1] \geq \cdots \geq A[n-1]$, where the ordering of strings is according to >= in TypeScript, (reverse) alphabetical ordering.

```
/**
 * Finds the index where x appears in the given sorted array or where, if
 * it is not in the array, it could be inserted to maintain sorted order.
 * @param A Array of strings in *decreasing* order
 * @param x String to look for in a.
 * @returns an integer k such that A[j] > x for any 0 <= j < k and
 *    x >= A[j] for any k <= j < A.length
 */
function findIndex(A: string[], x: string): number
```

Suppose that the function returns k. If x is in the array, then we must have $A[k] = x$. If x is not in the array, then we must have ($k = n$ or $k \geq 0$) and $A[k] \neq x$.

For example, suppose that `A` is the array `["mouse", "dog", "dog", "cat"]`. Then, the specification above tells us that

- A call to `findIndex(A, "zebra")` would return 0.
- A call to `findIndex(A, "dog")` would return 1 (not 2).
- A call to `findIndex(A, "cat")` would return 3.
- A call to `findIndex(A, "bat")` would return 4.
- A call to `findIndex(A, "kangaroo")` would return 1.

1. (18 points) **Loop, There It Is**

Consider the following code, which claims to implement `findIndex` from the prior page.

The precondition is that $A[j] \geq A[j+1]$ for any $0 \leq j < n-1$, where n is `A.length`.

```
let k: number = A.length;
```
$\{\{\, P_1 : k = n \,\}\}$

$\{\{$ **Inv:** $x \geq A[j]$ for any $k \leq j < n$ and $k \geq 0 \,\}\}$

```
while (k !== 0 && x >= A[k - 1]) {
```
$\{\{\, P_2 : x \geq A[j]$ for any $k \leq j < n$ and $k \geq 0$ and $k \neq 0$ and $x \geq A[k-1] \,\}\}$

$\{\{\, Q_2 : x \geq A[j]$ for any $k-1 \leq j < n$ and $k-1 \geq 0 \,\}\}$

```
    k = k - 1;
```
$\{\{\, x \geq A[j]$ for any $k \leq j < n$ and $k \geq 0 \,\}\}$

```
}
```
$\{\{\, P_3 : x \geq A[j]$ for any $k \leq j < n$ and $k \geq 0$ and $(k = 0$ or $A[k-1] > x) \,\}\}$

$\{\{\, Q_3 : A[j] > x$ for any $0 \leq j < k$ and $x \geq A[j]$ for any $k \leq j < n \,\}\}$

```
return k;
```

(a) Use reasoning to fill in all blank assertions above. The 'P_i's should be filled in with forward reasoning and the 'Q_i's should be filled in with backward reasoning.

(b) Prove that P_1 implies Inv.

Solution: Since $k = n$, Inv says that "$x \geq A[j]$ for any $n \leq j < n$". This is vacuously true since there are no such numbers j. We can also see that $k = n \geq 0$.

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(c) Prove that P_2 implies Q_2.

Solution: $k \geq 0$ and $k \neq 0$ imply that $k \geq 1$, which i the second part.
All the facts of the first part are included in P_2's first part except $x \geq A[k-1]$, which is the last part, so all the facts of Q_2 are actually included.

(d) Prove that P_3 implies Q_3.

Solution: The second part of Q_3 is incluedin P_3.

For the first part, we argue by cases.

If $k = 0$, then the first part says "$A[j] > x$ for any $0 \leq j < 0$", which is vacuously true because there are no such j's.

If $A[k-1] > x$, then for any $0 \leq j < k$, we have $A[j] \geq A[k-1] > x$ since A is sorted.

One of these cases must occur because of the "or" in P_3, so Q_3 holds.

2. (18 points) **Give It Your Test Shot**

   Fill in the body of the following unit test for `findIndex`. Include comments explaining the test cases, as we did in the coding homework problems.

```
it('findIndex', function() {
 // 0 times through the loop
 assert.deepStrictEqual(
     findIndex([], "zebra"),
     0);

 // 0 times through the loop
 assert.deepStrictEqual(
     findIndex(["mouse"], "zebra"),
     0);

 // 1 time through the loop
 assert.deepStrictEqual(
     findIndex(["mouse"], "cat"),
     1);

 // 1 time through the loop
 assert.deepStrictEqual(
     findIndex(["mouse", "cat"], "dog"),
     1);

 // many times through the loop
 assert.deepStrictEqual(
     findIndex(["mouse", "dog", "cat", "bat"], "aardvark"),
     4);

 // many times through the loop
 assert.deepStrictEqual(
     findIndex(["mouse", "mouse", "dog", "cat"], "cat"),
     3);
}
```

The remaining problems involve the implementation of the following ADT:

```
/** An array of strings with no duplicates. */
interface StringSet {

  /**
   * Returns a set that includes all the current elements and x also
   * @param x a string to insert into the set (if not already present)
   * @returns obj if contains(obj, x) = T
   *          L   if contains(obj, x) = F
   *     where L = A ++ [x] ++ B with obj = A ++ B (i.e., L is an array
   *     containing the strings from obj with x inserted somewhere)
   */
   insert(x: string): StringSet;

  /**
   * Returns the largest string in the set
   * @requires obj.length > 0
   * @returns max(obj), where max is defined on non-empty lists by
   *             max([y]) := y
   *        max(A ++ [y]) := max(A)    if y < max(A)
   *        max(A ++ [y]) := y         if y >= max(A)
   */
  max(): string;
}
```

We will implement it with the following class, whose concrete representation is an array sorted in decreasing order.

```
class ArrayStringSet implements StringSet {

  // RI: elems[j] > elems[j+1] for any 0 <= j < elems.length - 1
  // AF: obj = this.elems
  readonly elems: readonly string[];

  // @requires elems is sorted in decreasing order, with no duplicates
  constructor(elems: readonly string[]) {
    this.elems = elems;
  }

  ...
}
```

3. (28 points) **Run Array! Run Array!**

Fill in the missing parts of the implementation of `insert`. Your code must be correct with the **provided invariants**. (You do not need to turn in a proof, but it must be correct.)

```
insert = (x: string): StringSet => {
  const k = findIndex(this.elems, x);

  if (k < this.elems.length && this.elems[k] === x) {
    return this;
  }

  // Create an array one longer than this.elems.
  const E: string[] = new Array(this.elems.length + 1);

  // Define A := this.elems[0 .. k-1]

  let i: number = 0;

  // Inv: E[0 .. i - 1] = A[0 .. i - 1]
  while (i !== k) {
    E[i] = this.elems[i];
    i = i + 1;
  }

  // Now have E[0 .. i - 1] = A and i = k

  E[i] = x;
  i = i + 1;

  // Now have E[0 .. i - 1] = A ++ [x] and i = k + 1
```

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```
  // Now have E[0 .. i - 1] = A ++ [x] and i = k + 1 (from previous page)

  // Define B := this.elems[k .. this.elems.length-1] as shorthand
  // With these definitions, we have this.elems = A ++ B.

  let j: number = 0;

  // Inv: E[0 .. i - 1] = A ++ [x] ++ B[0 .. j - 1] and i = k + 1 + j
  while (k + j !== this.elems.length) {
    E[i] = this.elems[k + j];
    i = i + 1;
    j = j + 1;
  }

  // Now have E[0 .. i - 1] = A ++ [x] ++ B and i = A.length + 1 + B.length,
  // so E = A ++ [x] ++ B
  return new ArrayStringSet(E);
};
```

4. (16 points) **Here Array, Gone Tomorrow**

(a) Fill in the implementation of `max` in `ArrayStringSet`.

```
max = (): string => {
  return this.elems[0];

};
```

(b) Explain in clear English (or prove formally, if you prefer) why your code above is correct.

Solution: The precondition, together with the AF, says that `this.elems.length > 0`, so this array access is legal.

The invariant says that the first element is larger than every later element, so this array element is the largest.