
CSE 331

Software Design & Implementation

Autumn 2021

Section 1 – Code Reasoning

Administrivia

- HW0 due tomorrow (Friday 10/1 by 5PM).
- Any questions before we dive in?
 - What are the most interesting/confusing/puzzling things so far in the course?

Agenda

- Introductions?
- Website tour
- Review and practice logical reasoning about code with Hoare Logic
- Review and practice logical strength of assertions (weaker vs. stronger)
- Onwards to forward reasoning

Introductions

Website Tour

Why reason about code?

- Prove that code is correct
- Understand *why* code is correct
- Diagnose why/how code is *not* correct
- Specify code behavior

From lecture:

Hoare Logic: First definitions

- **Program State:** Values of all related variables
- **Assertion:** True/False claim (proposition) about the program state at a certain point in execution
- An assertion **holds** for a program state if it is true at that point.
- **Precondition:** Assertion before the code
 - Assumptions about when the code is used
- **Postcondition:** Assertion after the code
 - What we want the result of the code to be



From lecture:

Hoare Logic

- **Hoare Triple:** Two assertions surrounding a piece of code
 - $\{ \{ P \} \} S \{ \{ Q \} \}$
 - P is the precondition, S is the code, Q is the postcondition
 - P,Q are **specifications**
- A Hoare triple $\{ \{ P \} \} S \{ \{ Q \} \}$ is **valid** if in any state that P holds, Q holds after running the code S.
 - If P is true, after running S we have that Q is true.
 - Otherwise the triple is **invalid**.

Let's practice!
(Q1, Q2)

A Note on Implication (\Rightarrow)

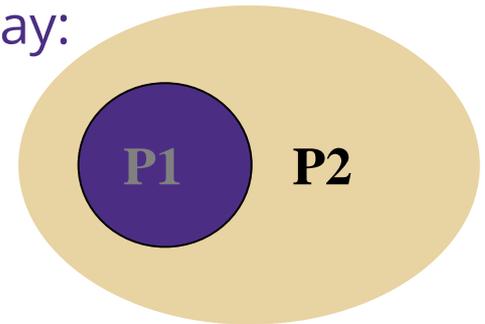
- Implication might be a bit new, but the basic idea is pretty simple. Implication $p \Rightarrow q$ is true as long as q is always true whenever p is

p	q	$p \Rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

From lecture:

Weaker / Stronger Assertions

- If P1 implies P2 (written $P1 \Rightarrow P2$) then we say:
 - P1 is **stronger** than P2
 - P2 is **weaker** than P1
- In other words:
 - P1 is “more difficult” to satisfy than P2
 - P1 puts more constraints on program states
 - P1 gives us more information about the program state



Let's practice!
(Q3)

Forward Reasoning

Forward Reasoning

- “What facts follow from initial assumptions about the code?”
- Precondition is **given**
- Fill in the **strongest** postcondition
 - For an assignment statement $x = y$
 - Add fact “ $x = y$ ” to what is known
 - important subtleties here (more later...)
 - Later: if statements and loops

From lecture:

Example of Forward Reasoning

$\{ \{ w > 0 \} \}$

$x = 17;$

$\{ \{ w > 0 \wedge x = 17 \} \}$

$y = 42;$

$\{ \{ w > 0 \wedge x = 17 \wedge y = 42 \} \}$

$z = w + x + y;$

$\{ \{ w > 0 \wedge x = 17 \wedge y = 42 \wedge z = w + 59 \} \}$



Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start...

```
{ { i >= 2 } }
```

```
x = 2 * i;
```

```
y = x;
```

```
z = (x + y) / 2;
```

Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start, what do we know about z at the end?

```
{ { i >= 2 } }
```

```
x = 2 * i;
```

```
{ { x = 2 * i ∧ i >= 2 } }
```

```
y = x;
```

```
z = (x + y) / 2;
```

Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start, what do we know about z at the end?

```
{{ i >= 2 }}
```

```
x = 2 * i;
```

```
{{ x = 2 * i & i >= 2 }}
```

```
y = x;
```

```
{{ y = x & x = 2 * i & i >= 2 }}
```

```
z = (x + y) / 2;
```

Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start, what do we know about z at the end?

```
{{ i >= 2 }}
```

```
x = 2 * i;
```

```
{{ x = 2 * i & i >= 2 }}
```

```
y = x;
```

```
{{ y = x & x = 2 * i & i >= 2 }}
```

```
z = (x + y) / 2;
```

```
{{ z = (x + y) / 2 & y = x & x = 2 * i & i >= 2  
}}
```

Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start, what do we know about z at the end?

$\{ \{ i \geq 2 \} \}$

$x = 2 * i;$

$\{ \{ x = 2 * i \wedge i \geq 2 \} \}$

$y = x;$

$\{ \{ y = x \wedge x = 2 * i \wedge i \geq 2 \} \}$

$z = (x + y) / 2;$

$\{ \{ z = (x + y) / 2 \wedge y = x \wedge x = 2 * i \wedge i \geq 2 \} \}$

$\Rightarrow \{ \{ z = (2 * i + 2 * i) / 2 \wedge i \geq 2 \} \}$

Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start, what do we know about z at the end?

$\{ \{ i \geq 2 \} \}$

$x = 2 * i;$

$\{ \{ x = 2 * i \wedge i \geq 2 \} \}$

$y = x;$

$\{ \{ y = x \wedge x = 2 * i \wedge i \geq 2 \} \}$

$z = (x + y) / 2;$

$\{ \{ z = (x + y) / 2 \wedge y = x \wedge x = 2 * i \wedge i \geq 2 \} \}$

$\Rightarrow \{ \{ z = (2 * i + 2 * i) / 2 \wedge i \geq 2 \} \}$

$\Rightarrow \{ \{ z = 2 * i \wedge i \geq 2 \} \}$

Let's Try One Together (Q4)

Suppose we know that $i \geq 2$ at the start, what do we know about z at the end?

$\{ \{ i \geq 2 \} \}$

$x = 2 * i;$

$\{ \{ x = 2 * i \wedge i \geq 2 \} \}$

$y = x;$

$\{ \{ y = x \wedge x = 2 * i \wedge i \geq 2 \} \}$

$z = (x + y) / 2;$

$\{ \{ z = (x + y) / 2 \wedge y = x \wedge x = 2 * i \wedge i \geq 2 \} \}$

$\Rightarrow \{ \{ z = (2 * i + 2 * i) / 2 \wedge i \geq 2 \} \}$

$\Rightarrow \{ \{ z = 2 * i \wedge i \geq 2 \} \}$

$\Rightarrow \{ \{ z \geq 4 \} \}$

Let's practice!
(Q5, Q6)

Questions?

- What is the most surprising thing about this?
- What is the most confusing thing?
- What will need a bit more thinking to digest?