

$$\{ a^i b^j c^k \mid i=j \text{ or } i=k \}$$

$$\{ a^n b^n c^n \mid n \geq 0 \}$$

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$$\{ ww^R \mid w \in \{a,b\}^* \}$$

$$\{ ww \mid w \in \{a,b\}^* \}$$

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$\forall$  CFL  $A \exists p$  st  $\forall w \in A$

if  $|w| \geq p$  then  $\exists u, v, x, y, z \in \Sigma^*$

st

(i)  $w = u \cdot v \cdot x \cdot y \cdot z$

(ii)  $\forall i \geq 0 \quad u v^i x y^i z \in A$

(iii)  $|v y| \geq 0$

(iv)  $|v x y| \leq p$

Lemma: a b-ary tree of height  $h$  has  $\leq b^h$  leaves



Conversely, more than  $b^h$  leaves implies height  $> h$

# Proof idea

$G$ : a CFG for  $A$

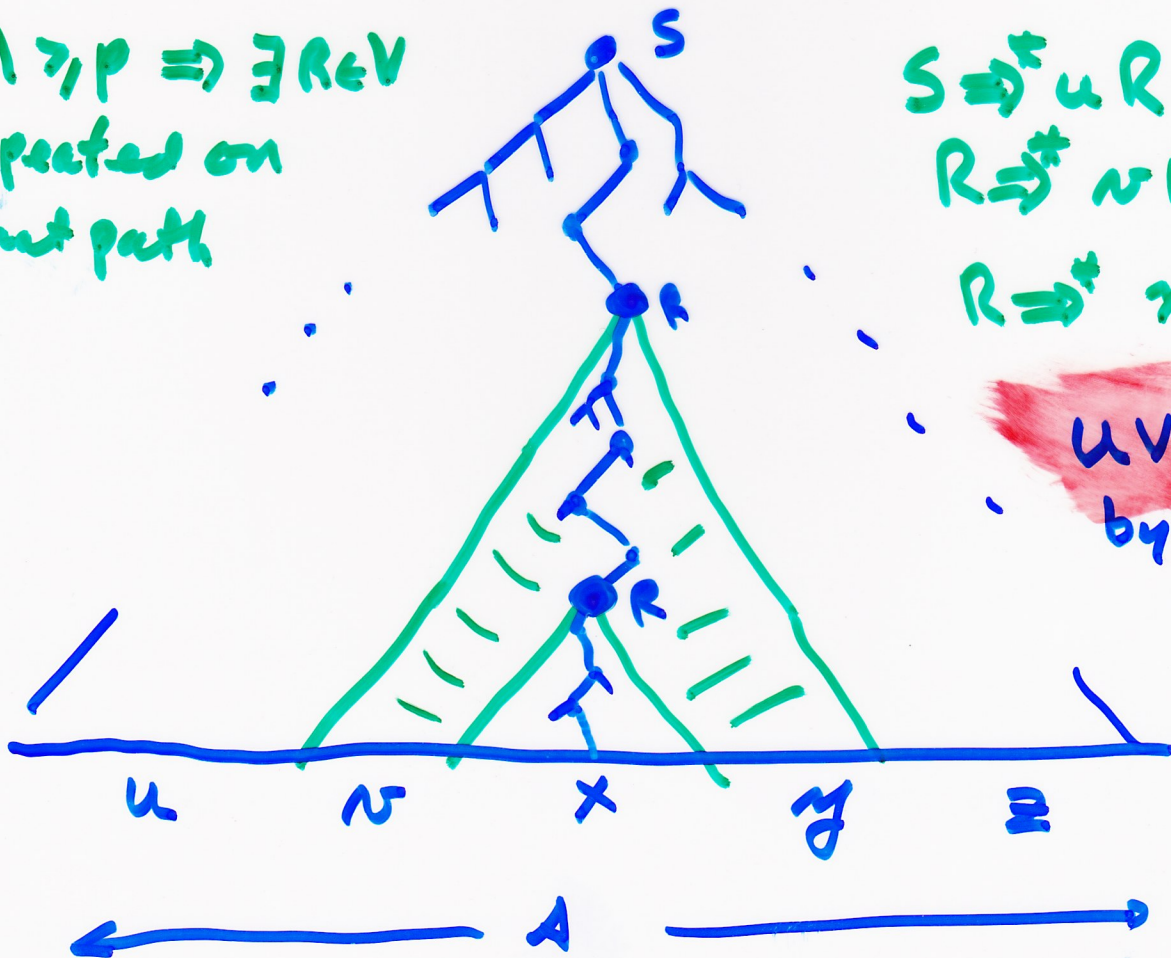
$b$  = length of longest r.h.s of a rule in  $G$

$p = b^{n+1}$  where  $n \geq |V|$ , # of vars in  $G$

$\Delta \in L(G)$  with  $|\Delta| \geq p$

Pick a smallest parse tree for  $\Delta$   
and a longest path in that tree

$|\Delta| \geq p \Rightarrow \exists R \in V$   
repeated on  
that path

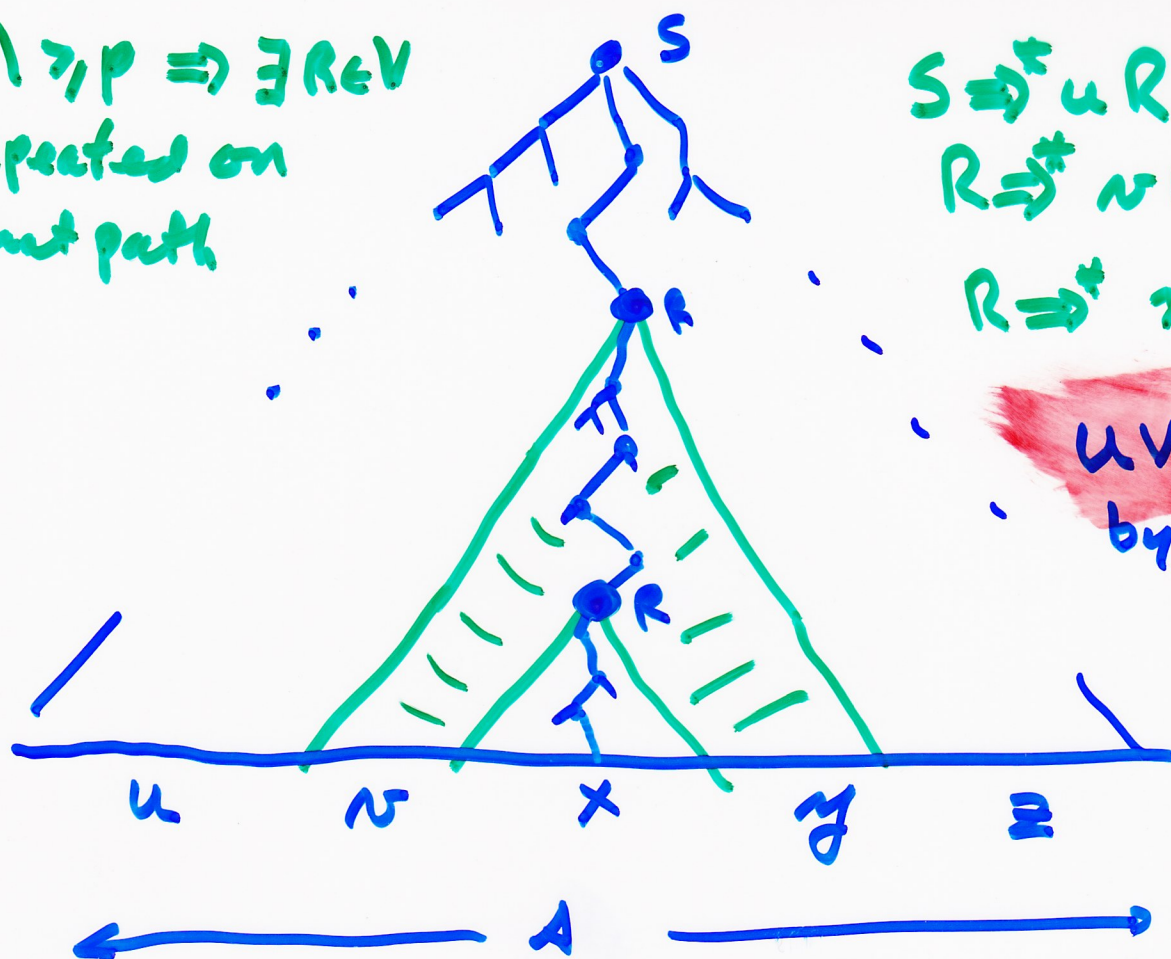


$S \Rightarrow^* u R z$   
 $R \Rightarrow^* v R y$   
 $R \Rightarrow^* x$

$u v^i x y^i z$   
by repeating  
 $i$  times.



$|A| \geq p \Rightarrow \exists R \in V$   
 repeated on  
 that path



$S \Rightarrow^+ u R z$   
 $R \Rightarrow^+ v R y$   
 $R \Rightarrow^+ x$

$uv^i xy^i z$   
 by repeating  
 $i$  times.

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Why a repeat?

$> b^{|V|}$  leaves  $\Rightarrow > |V|$  path length  
 $\Rightarrow$  some variable  $R$  repeated.

Why  $uv \neq \epsilon$ ?

because it was smallest tree

Why  $|uxy| \leq p$ ?

Pick repeat nearest leaf

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$$L = \{ a^n b^n c^n \mid n \geq 0 \}$$

Let  $p$  be const from P.L.

$$\text{let } s = a^p b^p c^p$$

By PL  $\exists uvxyz \dots$

~~uvxy~~ does not have ~~all 3 letters~~ both a & c

Since ~~uvxy~~  $|s| \leq p$

if no c in ~~uvxy~~ then

$$uv^i xy^j = a^i b^j c^p$$

~~Then~~ for some  $i \leq p, j \leq p$

$$\underline{i+j} < 2p$$

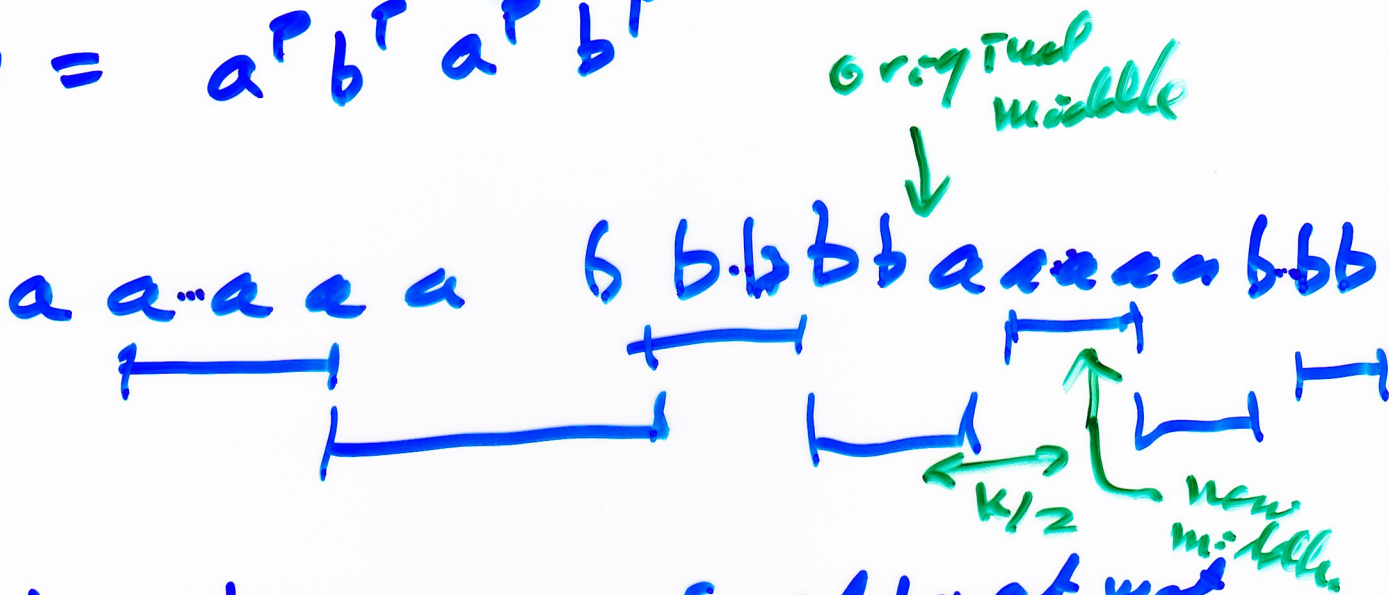
$$\text{Sim } |vxy| > 0$$

case 2 no a's  
Similar

$$L = \{ ww \mid w \in \{a, b\}^* \}$$

~~$$A = a^p b a^p b$$~~

$$A = a^p b^p a^p b^p$$



$|vxy| \leq p \therefore$  confined to at most 2 adjacent blocks of a's & b's.

case 1  $|uvxy| \leq 2p$

$uv^0xy^0z$  removed  $k$  letters from left half  $1 \leq k \leq p$

~~is~~ Last letter of (new) left half is  $a$ , but last of right half is  $b$ .

$\therefore \notin L$

↪



Case 2

$vxy$  in right half:  $i \geq p$

Case 3

$vxy$  straddles middle.

$$uv^i xy^j z = a^p b^i a^j b^p$$

for some  $i \leq p, j \leq p$

not both  $i=j=p$

$i < j$  too few  $b$ 's

$j < i$  ...  $a$ 's

$i=j < p$   $a^p b^i \neq a^i b^p$

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"Corollary"

$\{ww \mid w \in \{a,b\}^+\}$  not CFL  $\Rightarrow$  Java not CFL

This is representative programming languages ~~that~~ that require variables to be declared (1<sup>st</sup>  $w$ ) before use (2<sup>nd</sup>  $w$ ), none of which (C, Java, C++, ...) are CFL's at this level.