

Reading Assignment: 0.1-0.4 (review) and 1.1-1.2

Problems:

1. Sipser's book, Exercise 1.3 (same in both editions.) Make sure you include everything that a state diagram should include!
2. The rule for valid names for variables in C programs is the following. Variables must begin with a character (that is, a letter in the English alphabet) or underscore and may be followed by any combination of characters, underscores, or the digits 0 – 9. Design a DFA that accepts strings that are valid variable names (For simplicity assume that $\Sigma = \{ \langle c \rangle, \langle d \rangle, \langle u \rangle, \# \}$ where $\langle c \rangle$ denotes a character, $\langle d \rangle$ denotes a digit, and $\langle u \rangle$ denotes and underscore, and $\#$ denotes any other possible ASCII character.
3. Give state diagrams of DFAs recognizing the following languages. In each parts assume that the alphabet is $\Sigma = \{0, 1\}$. As documentation for your DFA, for each state, give a description of the strings which will *end* at that given state. This means that for each state you should give a description of the set of strings which, if they were inputed to the DFA, would end at that state.
 - (a) $L_1 = \{ w \mid w \text{ contains at least three 1s and at least three 0s.} \}$
 - (b) $L_2 = \{ w \mid w \text{ has length at least 2 and its second symbol is a 0.} \}$
 - (c) $L_3 = \{ w \mid w \text{ has an even number of 0s and an odd number of 1s.} \}$
 - (d) $L_4 = \{ w \mid w \text{ begins with a 1, and which, interpreted as the binary representation of a positive integer, is divisible by 4} \}$.
For this last part assume that the DFA starts reading the string from its most significant bit. For example if $w = 1000$, then w is the binary representation of the (decimal) number 8 (and thus, is in the language). and the DFA starts by reading the bit 1.
4. The reversal of a string w denoted by w^R , is the string when you look at it backwards: for example, $\text{homer}^R = \text{remoh}$. Here is the formal inductive definition (where the alphabet is Σ):

Definition of reversal of a string

Base case. If $w = \epsilon$, then $w^R = \epsilon$.

Inductive step. If $w = va$ for $v \in \Sigma^*$ and $a \in \Sigma$, then $w^R = av^R$.

Note in this definition that a is a single element of the alphabet.

Prove by induction (on the length of y) that for all strings $x, y \in \Sigma^*$,

$$(xy)^R = y^R x^R$$

Note that you will be doing an inductive proof of the above fact using the inductive definition.

5. **Extra Credit** (minimal points, do it for the glory!) Sipser's book, Problem 1.37 in second edition (Problem 1.30 in the first edition.)