

Closure under \cup

even parity $\leq \{0,1\}^*$

or 3rd from right is a $\leq \{a,b\}^*$

... - - - - - 1 $\leq \{a,1\}^*$

$$M_i = (Q_i, \Sigma, \delta_i, q_{0i}, F_i)$$

$$M = (Q_1 \times Q_2, \Sigma, \delta, (q_{01}, q_{02}), F)$$

$$\forall q_1 \in Q_1, q_2 \in Q_2, a \in \Sigma$$

$$\delta((q_1, q_2), a) = (\delta_1(q_1, a), \delta_2(q_2, a))$$

$$F = F_1 \times Q_2 \cup Q_1 \times F_2 - \text{union}$$

$$F' = F_1 \times F_2 - \text{intersection}$$

(*) To show by induction on length of w $\forall q_1 \in Q_1, q_2 \in Q_2. \forall w \in \Sigma^*$
 M is in state (q_1, q_2) after reading $w \iff M_1$ is in state q_1 and M_2 is in state q_2
 \uparrow after reading w \uparrow after reading w

Cor

M accepts $L(M_1) \cup L(M_2)$

* Note that there is a key difference between statement (*) and the definition of M : the latter, via defn of δ , says something about the finite set of ~~com~~ letters in Σ ; (*) talks about the infinite set of strings in Σ^* .

$$X, Y \subseteq \Sigma^*$$

$$X \cdot Y = \{x \cdot y \mid x \in X \& y \in Y\}$$

Examples

$$L_{\text{odd parity}} \cdot L_{\text{odd parity}} = L_{\text{even}} - \{0\}^*$$

$$L_{\text{odd parity}} \cdot L_{\text{even}} = L_{\text{odd}}$$

$$\begin{array}{ccc}
 A \cdot B & \stackrel{?}{=} & C \\
 \uparrow \quad \swarrow & & \uparrow \\
 \text{infinite} & \text{finite} & \text{finite}
 \end{array}
 \left. \begin{array}{l} \\ \\ \end{array} \right\} \begin{array}{l} \text{possible?} \\ \\ \text{yes} \end{array}$$

$$\Sigma^* \cdot \emptyset = \emptyset$$

$$\begin{array}{l}
 X \cdot Y \stackrel{?}{=} Y \cdot X \\
 \{0\} \cdot \{1\} \neq \{1\} \cdot \{0\}
 \end{array}$$

$\left. \begin{array}{l} \\ \\ \end{array} \right\} \begin{array}{l} \text{always} \\ \text{true?} \\ \\ \text{no} \end{array}$



$X, Y \subseteq \mathbb{R}^n, \mathbb{R}^n \setminus \{0\}$, $\$ \notin \mathbb{R}^n$

$X \cap Y \neq \emptyset$ \leftarrow "marked construction"