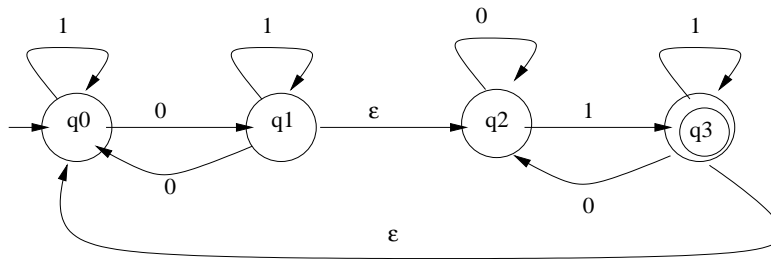


CSE 322
 Winter Quarter 2001
 Assignment 3
 Due Friday, January 26

All solutions should be neatly written or type set. All major steps in proofs and algorithms must be justified.

- (10 points) For this problem you will practice converting a NFA to a DFA. Convert the following NFA to a DFA. Show only the reachable states of the DFA. The transition function should be given in a table.



- (10 points) For this problem you will have practice in showing that regular languages are closed under more operations using finite automata constructions. We define the *interleaving* of two languages A and B over Σ by

$$A \parallel B = \{x_1y_1 \cdots x_ny_n : x_i, y_i \in \Sigma^*, x_1x_2 \cdots x_n \in A, \text{ and } y_1y_2 \cdots y_n \in B\}.$$

For example if $A = \{a, ab\}$ and $B = \{01\}$ then $A \parallel B = \{a01, 0a1, 01a, ab01, a0b1, a01b, 0ab1, 0a1b, 01ab\}$. Show that if A and B are regular then so is $A \parallel B$. Start with DFA's M_1 and M_2 that accept A and B , respectively. Then construct an NFA that accepts $A \parallel B$.

- (10 points) For this problem you should design algorithms in the style of problem 3 of the first assignment to decide properties of deterministic finite automata. In particular given a DFA $M = (Q, \Sigma, \delta, q_0, F)$ there is a natural directed graph G_M that models the transitions. The set of vertices of G_M is Q and (q, p) is an edge in G_M if $\delta(q, \sigma) = p$ for some $\sigma \in \Sigma$. The graph G_M is the transition diagram of M with the labels on the edges removed. You should use G_M in your algorithms.
 - Design an algorithm to decide whether a DFA accepts any strings at all. That is, the algorithm given a finite automaton M determines if $L(M) = \phi$.
 - Design an algorithm to decide whether a DFA accepts infinitely many strings.