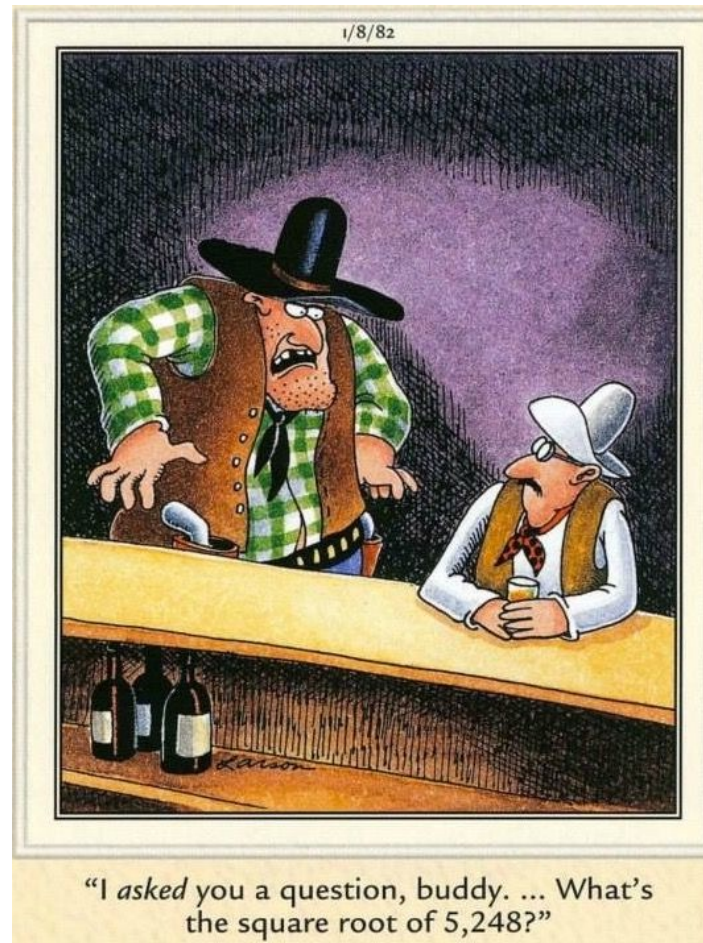


CSE 311: Foundations of Computing

Lecture 14: Modular Inverse, Exponentiation



Last time: Useful GCD Facts

If a and b are positive integers, then
$$\gcd(a, b) = \gcd(b, a \bmod b)$$

If a is a positive integer, $\gcd(a, 0) = a$.

Euclid's Algorithm

$$\text{gcd}(a, b) = \text{gcd}(b, a \bmod b) \qquad \text{gcd}(a, 0) = a$$

```
int gcd(int a, int b){ /* Assumes: a >= b, b >= 0 */
    if (b == 0) {
        return a;
    } else {
        return gcd(b, a % b);
    }
}
```

Note: $\text{gcd}(b, a) = \text{gcd}(a, b)$

Euclid's Algorithm

Repeatedly use $\gcd(a, b) = \gcd(b, a \bmod b)$ to reduce numbers until you get $\gcd(g, 0) = g$.

$$\gcd(660, 126) =$$

Euclid's Algorithm

Repeatedly use $\gcd(a, b) = \gcd(b, a \bmod b)$ to reduce numbers until you get $\gcd(g, 0) = g$.

$$\begin{aligned}\gcd(660, 126) &= \gcd(126, 660 \bmod 126) = \gcd(126, 30) \\ &= \gcd(30, 126 \bmod 30) = \gcd(30, 6) \\ &= \gcd(6, 30 \bmod 6) = \gcd(6, 0) \\ &= 6\end{aligned}$$

Euclid's Algorithm

Repeatedly use $\gcd(a, b) = \gcd(b, a \bmod b)$ to reduce numbers until you get $\gcd(g, 0) = g$.

Equations with recursive calls:

$$\begin{aligned}\gcd(660, 126) &= \gcd(126, 660 \bmod 126) = \gcd(126, 30) \\ &= \gcd(30, 126 \bmod 30) = \gcd(30, 6) \\ &= \gcd(6, 30 \bmod 6) = \gcd(6, 0) \\ &= 6\end{aligned}$$

Tableau form:

$$\begin{aligned}660 &= 5 * 126 + 30 \\ 126 &= 4 * 30 + \textcircled{6} \\ 30 &= 5 * 6 + 0\end{aligned}$$

Famous Algorithmic Problems

- **Primality Testing**

- Given an integer n , determine if n is prime

- **Factoring**

- Given an integer n , find an integer d (with $1 < d < n$) that divides n

- **Greatest Common Divisor**

- Given integers a and b , find the largest integer d that divides them both

Bézout's theorem

If a and b are positive integers, then there exist integers s and t such that

$$\gcd(a, b) = sa + tb.$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 1 (Compute GCD & Keep Tableau Information):

$$\begin{array}{cc} a & b \\ \gcd(35, 27) & = \gcd(27, 35 \bmod 27) = \gcd(27, 8) \end{array}$$

$$\begin{array}{l} a = q * b + r \\ 35 = 1 * 27 + 8 \end{array}$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 1 (Compute GCD & Keep Tableau Information):

$$\begin{array}{l} \begin{array}{cc} a & b \\ \gcd(35, 27) & = \gcd(27, 35 \bmod 27) & = \gcd(27, 8) \\ & = \gcd(8, 27 \bmod 8) & = \gcd(8, 3) \\ & = \gcd(3, 8 \bmod 3) & = \gcd(3, 2) \\ & = \gcd(2, 3 \bmod 2) & = \gcd(2, 1) \\ & = \gcd(1, 2 \bmod 1) & = \gcd(1, 0) \end{array} \end{array}$$

$$\begin{array}{l} a = q * b + r \\ 35 = 1 * 27 + 8 \\ 27 = 3 * 8 + 3 \\ 8 = 2 * 3 + 2 \\ 3 = 1 * 2 + \textcircled{1} \end{array}$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 2 (Solve the equations for r):

$$a = q * b + r$$

$$35 = 1 * 27 + 8$$

$$27 = 3 * 8 + 3$$

$$8 = 2 * 3 + 2$$

$$3 = 1 * 2 + 1$$

$$r = a - q * b$$

$$8 = 35 - 1 * 27$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 2 (Solve the equations for r):

$$a = q * b + r$$

$$35 = 1 * 27 + 8$$

$$27 = 3 * 8 + 3$$

$$8 = 2 * 3 + 2$$

$$3 = 1 * 2 + \textcircled{1}$$

$$r = a - q * b$$

$$8 = 35 - 1 * 27$$

$$3 = 27 - 3 * 8$$

$$2 = 8 - 2 * 3$$

$$\textcircled{1} = 3 - 1 * 2$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 3 (Backward Substitute Equations):

$$8 = 35 - 1 * 27$$

$$3 = 27 - 3 * 8$$

$$2 = 8 - 2 * 3$$

$$\textcircled{1} = 3 - 1 * 2$$

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 3 (Backward Substitute Equations):

$$8 = 35 - 1 * 27$$

$$3 = 27 - 3 * 8$$

$$2 = 8 - 2 * 3$$

$$1 = 3 - 1 * 2$$

$$1 = 3 - 1 * (8 - 2 * 3)$$

$$= 3 - 8 + 2 * 3$$

$$= (-1) * 8 + 3 * 3$$

Plug in the def of 2

Re-arrange into
3's and 8's



Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 3 (Backward Substitute Equations):

$$8 = 35 - 1 * 27$$

$$3 = 27 - 3 * 8$$

$$2 = 8 - 2 * 3$$

$$1 = 3 - 1 * 2$$

$$1 = 3 - 1 * (8 - 2 * 3)$$

$$= 3 - 8 + 2 * 3$$

$$= (-1) * 8 + 3 * 3$$

$$= (-1) * 8 + 3 * (27 - 3 * 8)$$

$$= (-1) * 8 + 3 * 27 + (-9) * 8$$

$$= 3 * 27 + (-10) * 8$$

Plug in the def of 2

Re-arrange into
3's and 8's

Plug in the def of 3

Re-arrange into
8's and 27's

Extended Euclidean algorithm

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

Step 3 (Backward Substitute Equations):

$$8 = 35 - 1 * 27$$

$$3 = 27 - 3 * 8$$

$$2 = 8 - 2 * 3$$

$$1 = 3 - 1 * 2$$

Re-arrange into
27's and 35's

$$1 = 3 - 1 * (8 - 2 * 3)$$

$$= 3 - 8 + 2 * 3$$

$$= (-1) * 8 + 3 * 3$$

Plug in the def of 2

Re-arrange into
3's and 8's

Plug in the def of 3

$$= (-1) * 8 + 3 * (27 - 3 * 8)$$

$$= (-1) * 8 + 3 * 27 + (-9) * 8$$

$$= 3 * 27 + (-10) * 8$$

Re-arrange into
8's and 27's

$$= 3 * 27 + (-10) * (35 - 1 * 27)$$

$$= 3 * 27 + (-10) * 35 + 10 * 27$$

$$= 13 * 27 + (-10) * 35$$

Multiplicative inverse mod m

Let $0 \leq a, b < m$. Then, b is the *multiplicative inverse of a (modulo m)* iff $ab \equiv_m 1$.

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

mod 7

x	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9
2	0	2	4	6	8	0	2	4	6	8
3	0	3	6	9	2	5	8	1	4	7
4	0	4	8	2	6	0	4	8	2	6
5	0	5	0	5	0	5	0	5	0	5
6	0	6	2	8	4	0	6	2	8	4
7	0	7	4	1	8	5	2	9	6	3
8	0	8	6	4	2	0	8	6	4	2
9	0	9	8	7	6	5	4	3	2	1

mod 10

Multiplicative inverse mod m

Suppose $\gcd(a, m) = 1$

By Bézout's Theorem, there exist integers s and t such that $sa + tm = 1$.

s is the multiplicative inverse of a (modulo m):

$$1 = sa + tm \equiv_m sa$$

So... we can compute multiplicative inverses with the extended Euclidean algorithm

These inverses let us solve modular equations...

Example

Solve: $7x \equiv_{26} 1$

Example

Solve: $7x \equiv_{26} 1$

$$\gcd(26, 7) = \gcd(7, 5) = \gcd(5, 2) = \gcd(2, 1) = 1$$

Example

Solve: $7x \equiv_{26} 1$

$$\gcd(26, 7) = \gcd(7, 5) = \gcd(5, 2) = \gcd(2, 1) = 1$$

$$26 = 3 * 7 + 5$$

$$7 = 1 * 5 + 2$$

$$5 = 2 * 2 + 1$$

Example

Solve: $7x \equiv_{26} 1$

$$\gcd(26, 7) = \gcd(7, 5) = \gcd(5, 2) = \gcd(2, 1) = 1$$

$$26 = 3 * 7 + 5 \qquad 5 = 26 - 3 * 7$$

$$7 = 1 * 5 + 2 \qquad 2 = 7 - 1 * 5$$

$$5 = 2 * 2 + 1 \qquad 1 = 5 - 2 * 2$$

Example

Solve: $7x \equiv_{26} 1$

$$\gcd(26, 7) = \gcd(7, 5) = \gcd(5, 2) = \gcd(2, 1) = 1$$

$$26 = 3 * 7 + 5 \quad 5 = 26 - 3 * 7$$

$$7 = 1 * 5 + 2 \quad 2 = 7 - 1 * 5$$

$$5 = 2 * 2 + 1 \quad 1 = 5 - 2 * 2$$

$$\begin{aligned} 1 &= 5 - 2 * (7 - 1 * 5) \\ &= (-2) * 7 + 3 * 5 \\ &= (-2) * 7 + 3 * (26 - 3 * 7) \\ &= (-11) * 7 + 3 * 26 \end{aligned}$$

Example

Solve: $7x \equiv_{26} 1$

$$\gcd(26, 7) = \gcd(7, 5) = \gcd(5, 2) = \gcd(2, 1) = 1$$

$$26 = 3 * 7 + 5 \qquad 5 = 26 - 3 * 7$$

$$7 = 1 * 5 + 2 \qquad 2 = 7 - 1 * 5$$

$$5 = 2 * 2 + 1 \qquad 1 = 5 - 2 * 2$$

$$\begin{aligned} 1 &= 5 - 2 * (7 - 1 * 5) \\ &= (-2) * 7 + 3 * 5 \\ &= (-2) * 7 + 3 * (26 - 3 * 7) \\ &= (-11) * 7 + 3 * 26 \end{aligned}$$

Multiplicative inverse of 7 modulo 26

Now $(-11) \bmod 26 = 15$. So, $x = 15 + 26k$ for $k \in \mathbb{Z}$.

Example of a more general equation

Now solve: $7y \equiv_{26} 3$

We already computed that **15** is the multiplicative inverse of **7** modulo **26**. That is, $7 \cdot 15 \equiv_{26} 1$

If y is a solution, then multiplying by **15** we have

$$15 \cdot 7 \cdot y \equiv_{26} 15 \cdot 3$$

Substituting $15 \cdot 7 \equiv_{26} 1$ into this on the left gives

$$y = 1 \cdot y \equiv_{26} 15 \cdot 3 \equiv_{26} 19$$

This shows that every solution y is congruent to **19**.

Example of a more general equation

Now solve: $7y \equiv_{26} 3$

Multiplying both sides of $y \equiv_{26} 19$ by 7 gives

$$7y \equiv_{26} 7 \cdot 19 \equiv_{26} 3$$

So, any $y \equiv_{26} 19$ is a solution.

Thus, the set of numbers of the form $y = 19 + 26k$, for any k , are exactly solutions of this equation.

Math mod a prime is especially nice

$\gcd(a, m) = 1$ if m is prime and $0 < a < m$ so
can always solve these equations mod a prime.

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

mod 7

Multiplicative Inverses and Algebra

Adding to both sides is an equivalence:

$$\begin{array}{ccc} -c & \rightarrow & x \equiv_m y \\ & & \searrow +c \\ & & x + c \equiv_m y + c \end{array}$$

The same is not true of multiplication...

unless we have a multiplicative inverse $cd \equiv_m 1$

$$\begin{array}{ccc} \times d & \rightarrow & x \equiv_m y \\ & & \searrow \times c \\ & & cx \equiv_m cy \end{array}$$