

# Practical Aspects of Modern Cryptography

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# Side-Channel Attacks

Breaking a cryptosystem is a frontal attack, but there may be easier access through a side or back door – especially on embedded cryptographic devices such as SmartCards and RFIDs.

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- Cache Attacks
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- ... others?

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- Watch decryption times for  $z = E(m)$  where  $m < P$  and where  $m > P$ .
- If there is a minute difference,  $P$  can be determined with binary search.

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Decryption times may vary in a key-dependent manner based upon which lines have been flushed.

# Power Analysis

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Careful measurement and analysis of power consumption can be used to determine the key.

# Electromagnetic Emissions

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Careful analysis of the emissions may reveal a secret key.

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It can actually be possible to hear whether or not these conditional multiplications are performed.

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Careful manipulation of data may elicit responses which disclose information about a desired key or decryption value.

# Certificate Revocation

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- It is sometimes necessary to “revoke” a certificate before it expires.

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- Role Modification

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- Online Certificate Status Protocol (OCSP)

# Certificate Revocation Lists

- A CA revokes a certificate by placing the its identifying serial number on its Certificate Revocation List (CRL)
  - Every CA issues CRLs to cancel out issued certs
  - A CRL is like anti-matter – when it comes into contact with a certificate it lists it cancels out the certificate
  - Think “1970s-style credit-card blacklist”
- Relying parties are expected to check the most recent CRLs before they rely on a certificate
  - “The cert is valid unless you hear something telling you otherwise”

# The Problem with CRLs

Blacklists have numerous problems

- They can grow very large because certs cannot be removed until they expire.
- They are not issued frequently enough to be effective against a serious attack.
- Their size can make them expensive to distribute (especially on low-bandwidth channels).
- They are vulnerable to simple DOS attacks.  
(What do you do if you can't get the current CRL?)

# More Problems with CRLs

Poor CRL design has made the problem worse.

- CRLs can contain retroactive invalidity dates
  - A CRL issued today can say a cert was invalid as of last week.
    - Checking that something was valid at time  $t$  wasn't sufficient!
    - Back-dated CRLs can appear at any time in the future.
- CAs can even change the CRL rules retroactively.

# Yet More Problems with CRLs

- Revoking a cert used by a CA to issue other certs is even harder since this may invalidate an entire set of certs.
- “Self-signed” certificates are often used as a syntactic convenience. Is it meaningful for a cert to revoke itself?

# Even More Problems with CRLs

- CRLs can't be revoked.  
If a cert has been mistakenly revoked, the revocation can't be reversed.
- CRLs can't be updated.  
There's no mechanism to issue a new CRL to relying parties early – even if there's an urgent need to issue new revocations.

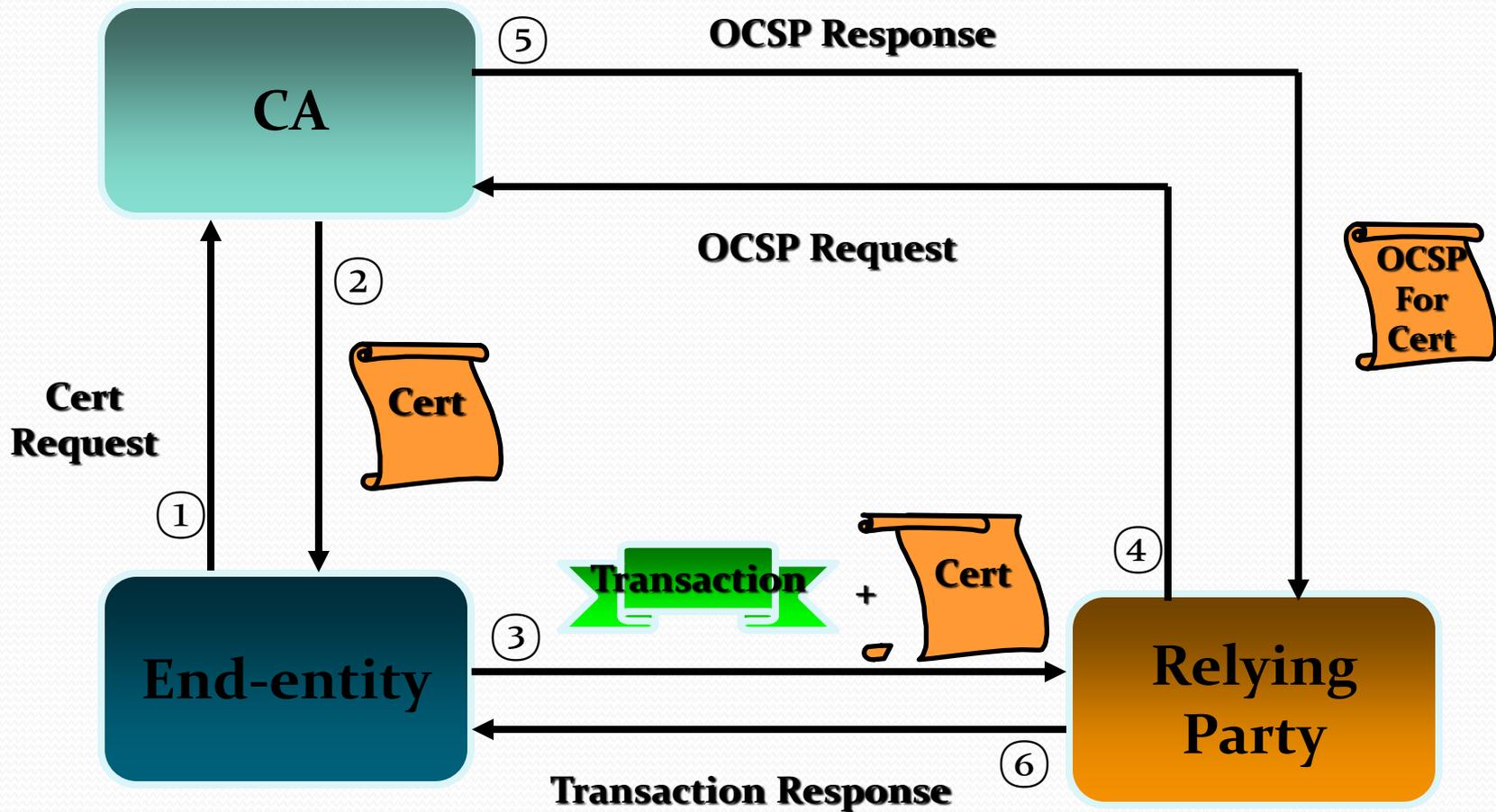
# Short-Lived Certificates

If you need to go to a CA to get a fresh CRL, why not just go to a CA to get a fresh cert?

# Online Status Checking

- OCSP: Online Certificate Status Protocol
  - A way to ask “is this certificate good right now?”
  - Get back a signed response from the OCSP server saying, “Yes, cert C is good at time t”
    - Response is like a “freshness certificate”
- OCSP response is like a selective CRL
  - Client indicates the certs for which he wants status information
  - OCSP responder dynamically creates a lightweight CRL-like response for those certs

# OCSP in Action



# Final thoughts on Revocation

- From a financial standpoint, it's the revocation data that is valuable, not the issued certificate itself.
  - For high-valued financial transactions, seller wants to know your cert is good right now.
  - This is similar to credit cards, where the merchant wants the card authorized “right now” at the point-of-sale.
- Card authorizations transfer risk from merchant to bank – thus they're worth \$\$\$.

# Design Charrette

How would you design a transit fare card system?

# Fare Card System Elements

- An RFID card for each rider
- Readers on each vehicle and/or transit station (Internet connected?)
- Card purchase/payment machines
- A web portal for riders to manage and/or enrich their cards