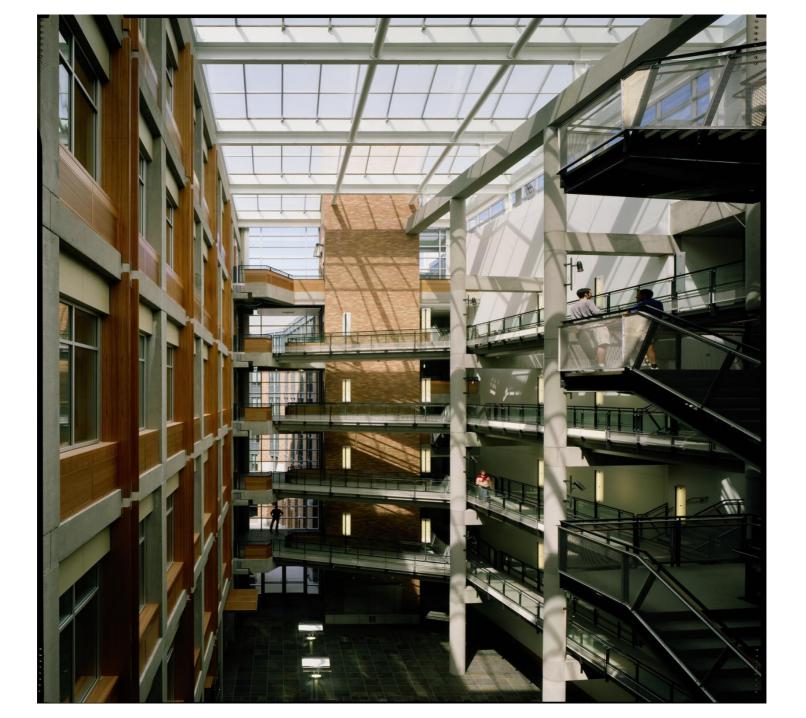
Content Delivery and File Sharing in the Modern Internet

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Work with:

- Steve Gribble
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Outline and Goals

- Explore and characterize content delivery in today's Internet
 - Web, Akamai CDN, Kazaa & Gnutella P2P
 - » What is the impact of these systems on the Internet?
 - » What are the characteristics of the new delivery systems?
- 2. Understand what drives P2P file sharing systems
 - Dynamics of the web are understood
 - » Driven by changes to documents, Zipf popularity
 - Dynamics of P2P are unknown

Outline

1. Characterize content delivery in today's Internet

2. Understand what drives P2P file sharing systems

Surprise: The WWW is changing even more rapidly now!

- Thirst for data (+ new types of data) continues to increase.
- People are using new means to provide and obtain that data.
- The result -- the web is now seeing a mixture of new and old content-delivery mechanisms:
 - Conventional web clients and web servers
 - Global-scale content-delivery networks (e.g., Akamai)
 - Self-organizing peer-to-peer file-sharing systems (e.g., Gnutella, Kazaa)

Quick overview: Peer-to-peer (P2P) systems

- Peers are individually owned computers, most on modems or broadband
- Peers collaborate to exchange content among themselves
- Each peer is both a client and server
- Peers issue or broadcast text queries to the peer network to find content
- Example -- Kazaa:
 - No centralized components
 - Two-level structure some peers are "supernodes"
 - Supernode indexes content from the peers underneath it
 - Supernodes can communicate with each other to find content
 - Files transferred in segments from multiple peers simultaneously
 - The protocol is proprietary

Peer-to-peer systems

- These systems are technically interesting
- They are autonomous, totally distributed, self-organizing
- There is a *huge* amount of research on P2P
- There is almost no data on P2P

- Questions:
 - What are the characteristics of the new P2P systems?
 - What is the impact of these P2P systems on the web?

Methodology

- Data is based on a 9-day trace collected at UW from May 28, 2002 through June 6, 2002.
- We use passive network monitoring.
- Our trace machine sees every packet going in and out.
- We classify traffic based on port number and other information in the message headers.
- We anonymize all sensitive information before writing to disk.

• Trace machine is a dual-CPU Dell with 2.0 GHz Xeon processors and a gigabit network card, running FreeBSD.

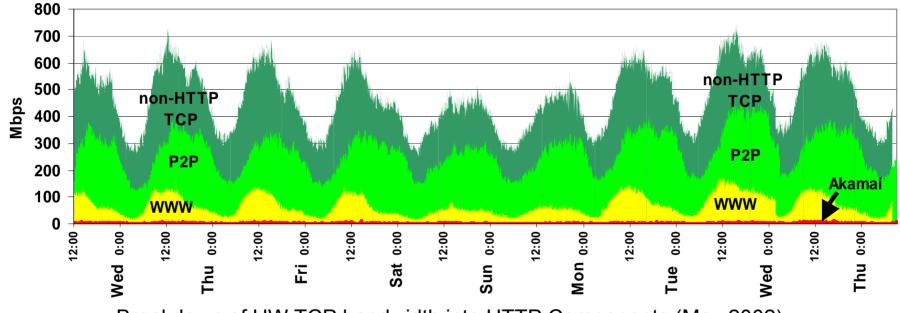
Limitations

- Only studied one population (UW)
- Finite trace period
- Could see data transfers, but not encrypted control traffic
- Cannot see UW-internal traffic

Question 1:

• What is the bandwidth impact of new P2P and contentdelivery systems?

Where has all the bandwidth gone?

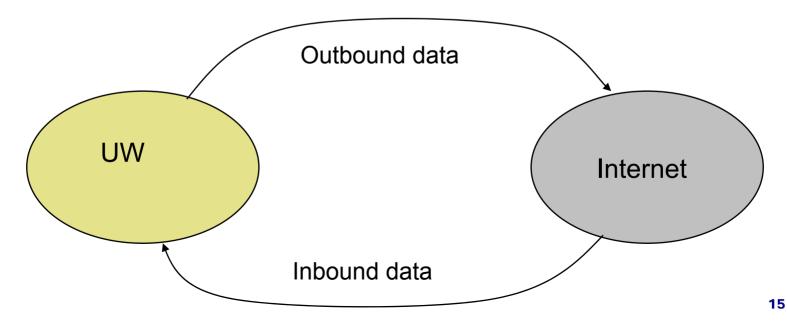


Breakdown of UW TCP bandwidth into HTTP Components (May 2002)

- WWW = 14% of TCP traffic; P2P = 43% of TCP traffic
- P2P dominates WWW in bandwidth consumed!!

Definition:

- <u>Inbound</u> traffic: data objects requested by UW clients, transmitted into the UW from an outside source.
- <u>Outbound</u> traffic: data objects requested by a source external to UW, transmitted by a UW server.



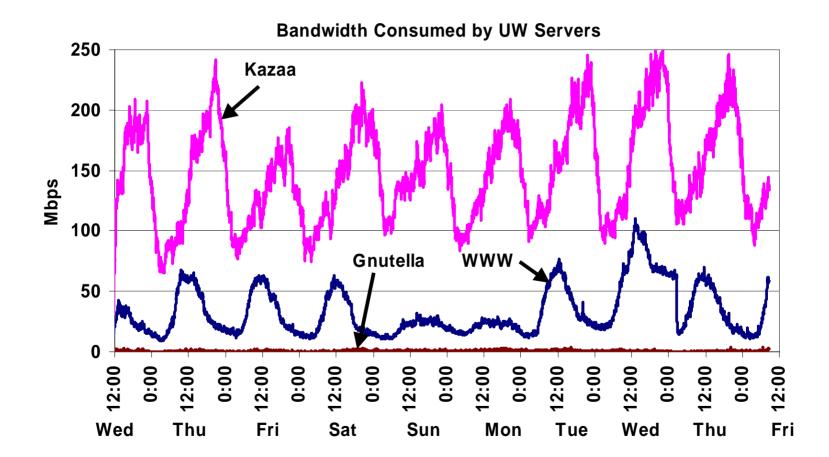
Detailed Trace Statistics

| | WWW | | |
|-----------------------|---------------------|------------|--|
| | inbound | outbound | |
| HTTP Xactions | 323,072,253 | 73,001,891 | |
| Unique objects | 72,818,997 3,412,64 | | |
| Clients | 39,285 | 1,231,308 | |
| Servers | 403,087 | 9,821 | |
| Bytes Xferred | 1.51 TB | 3.02 TB | |
| Median Object Size | 1,976 B | 4,646 B | |
| Mean Object Size | 22,491 | 334,944 | |

Detailed Trace Statistics

| | WV | VW | Kazaa | |
|-----------------------|-------------|------------|-----------------|------------|
| | inbound | outbound | inbound | outbound |
| HTTP Xactions | 323,072,253 | 73,001,891 | 11,140,861 | 19,190,902 |
| Unique objects | 72,818,997 | 3,412,647 | 111,437 | 166,442 |
| Clients | 39,285 | 1,231,308 | 4,644 | 611,005 |
| Servers | 403,087 | 9,821 | 281,026 | 3,888 |
| Bytes Xferred | 1.51 TB | 3.02 TB | 1.78 TB | 13.57 TB |
| Median Object Size | 1,976 B | 4,646 B | 3.75 MB 3.67 ME | |
| Mean Object Size | 22,491 | 334,944 | 29.4MB | 26.1MB |

Bandwidth consumed by UW servers (outbound traffic)





October 22, 2002

UNIVERSITY OF WASHINGTON

Bandwidth restrictions save almost \$1 million by Steven Friederich 10/22/2002

Those using peer-to-peer software on campus, such as the file-sharing program KaZaA, may notice their network connection has been acting slow lately. New technical restrictions placed within the campus networks have provided a limit to the amount of bandwidth users may access for Web sites and servers.

Question 2:

• What are all the the bytes carrying?

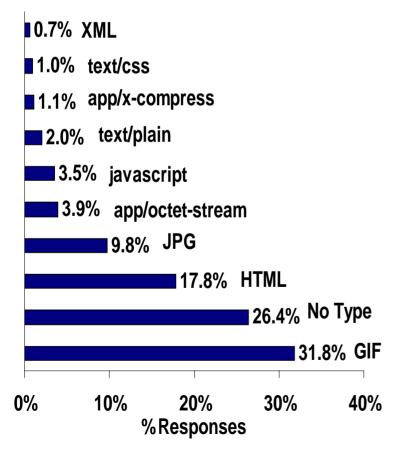
What data types are being downloaded (HTTP)?

Content Types Ordered by Number of Downloads

| 0.7% X | ML | | | |
|-------------|----------|-----------------|---------|--------|
| 1.0% | text/css | | | |
| 1.1% | app/x-co | ompress | | |
| 2.0% | text/pla | in | | |
| 3.5% | 6 javas | cript | | |
| 3.99 | % app/c | octet-strea | m | |
| | 9.8% | JPG | | |
| | | 17.8% | HTML | |
| | | | 26.4% N | о Туре |
| | | | 31. | 8% GIF |
| 0% | 10% % | 20% Response | | 40% |

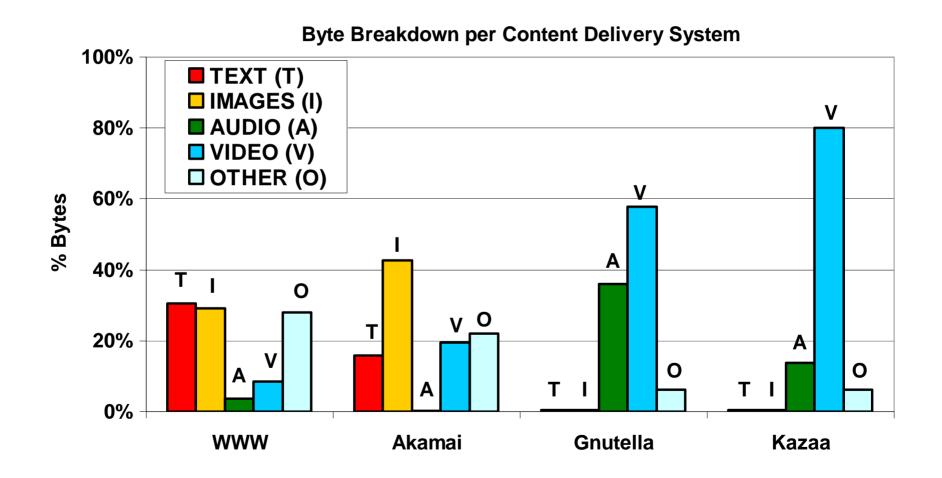
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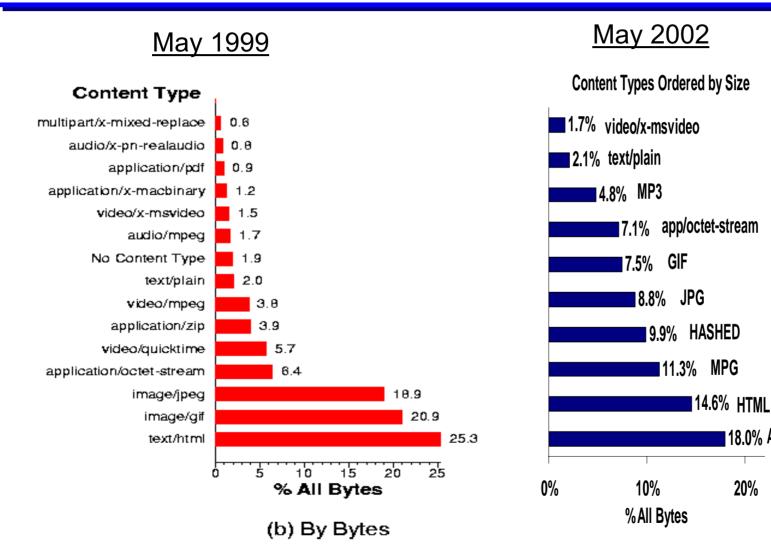


| Con | tent Types Ordered | d by Size |
|------|--------------------|------------|
| 1.7% | video/x-msvideo | |
| 2.1% | % text/plain | |
| | 14.8% MP3 | |
| | 7.1% app/oc | tet-stream |
| | 7.5% GIF | |
| | 8.8% JPC | 3 |
| | 9.9% H | ASHED |
| | 11.3% | MPG |
| | 1 | 4.6% HTML |
| | | 18.0% AVI |
| 0% | 10% | 20% |
| | %All Bytes | |

Object type for different systems



HTTP download content type: now vs. then



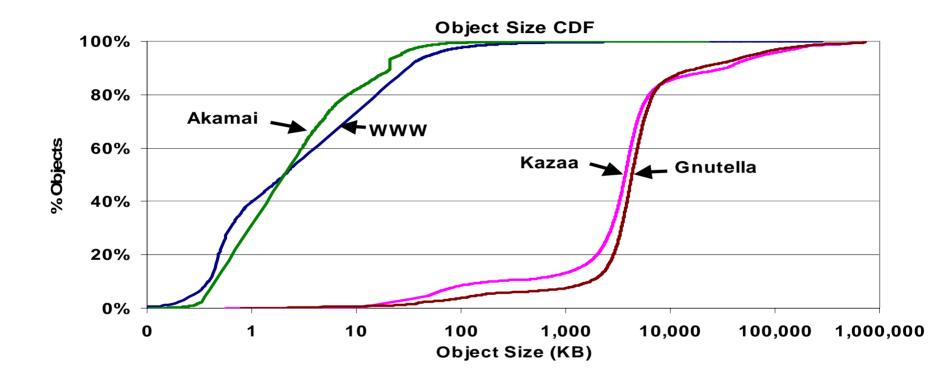
18.0% AVI

20%

Question 3:

 How do workload characteristics differ in P2P systems, compared to www?

Object size

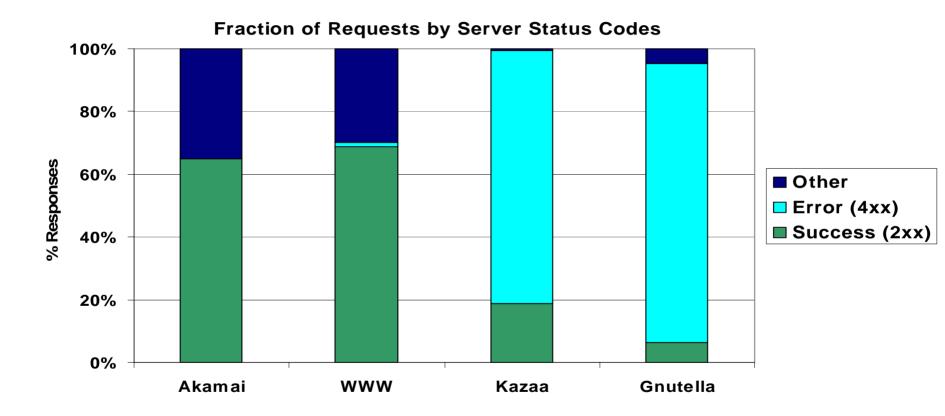


The median Kazaa object is **1000 times larger** than the median WWW object!

Top BW-consuming Kazaa objects

| | Kazaa (inbound) | | | Kazaa (outbound) | | | | |
|---|------------------------|--------------------|---------|------------------|---------------------|--------------------|---------|--------|
| | object size (MB) | GB consum ed | clients | srvers | object size (MB) | GB consum ed | clients | srvers |
| 1 | 694 | 8.1 | 20 | 164 | 696 | 119.0 | 397 | 1 |
| 2 | 702 | 6.4 | 14 | 91 | 699 | 110.5 | 1000 | 4 |
| 3 | 690 | 6.1 | 22 | 83 | 699 | 78.7 | 390 | 10 |
| 4 | 775 | 5.6 | 16 | 105 | 700 | 73.3 | 558 | 2 |
| 5 | 698 | 4.7 | 14 | 74 | 634 | 64.9 | 540 | 1 |

Availability

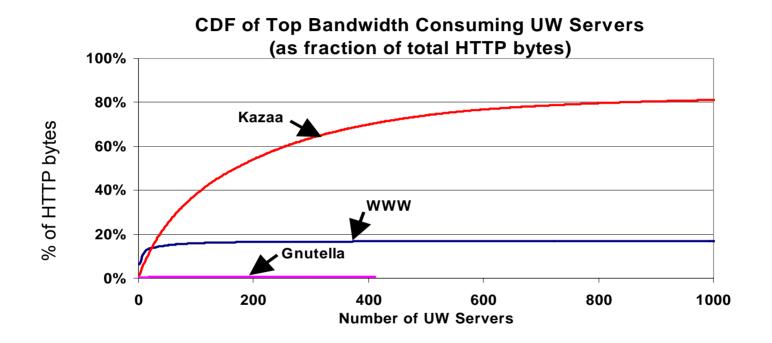


It's mostly unavailable, but it's free (and as long as you get it eventually, you don't care)

Question 4

 How is bandwidth use distributed among Kazaa clients and servers?

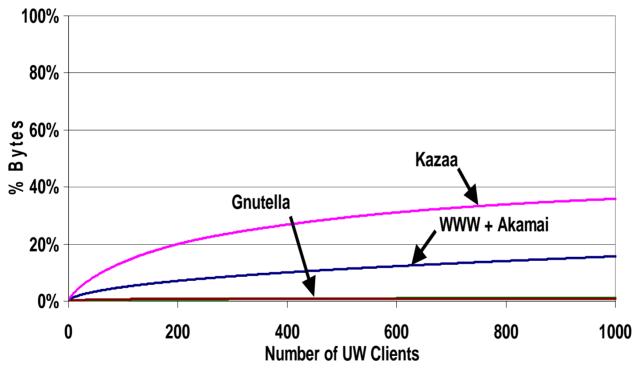
Top bandwidth producing UW servers



- 20 WWW servers supply 20% of HTTP bytes served
- 400 Kazaa peers generate 70% of all HTTP outgoing traffic!

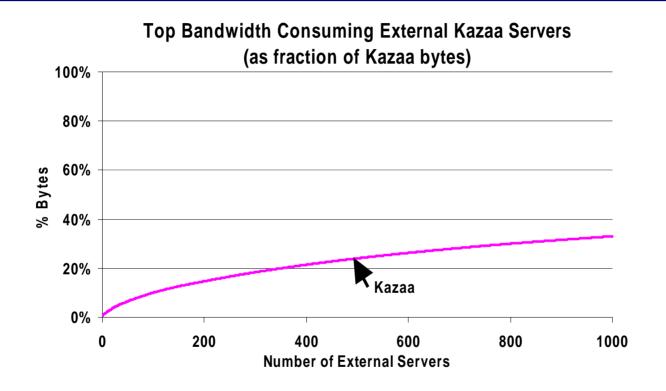
Top bandwidth consuming UW clients

Top Bandwidth Consuming UW Clients (as fraction of total HTTP)



 Top 200 Kazaa clients are responsible for 20% of all HTTP bytes downloaded

Top bandwidth producing external Kazaa servers



- 600 external Kazaa peers (out of 281K) serve 26% of the Kazaa bytes to UW.
- Surprising -- given the "scalability based design" of P2P systems and Kazaa.

Summary of CDN/P2P study

- Internet has undergone a huge qualitative change in only a few years.
- We've moved from interactive transfer of small files (10s of KBs) to batch transfer of enormous files (100s of MBs).
- P2P now accounts for the majority of HTTP bytes, exceeding WWW traffic by nearly 3X at UW
- P2P documents avg. 3 orders of magnitude larger than WWW docs
- A small number of huge objects are responsible for an enormous fraction of transfers (300 Kazaa objects used 5.6TB of BW!)
- A small number of P2P clients are causing much of the traffic
- On average, a single UW Kazaa peer uses <u>90 times</u> the bandwidth of a single WWW client!

Outline

- Characterize content delivery in today's Internet
- Understand what drives Kazaa P2P file-sharing system
 - 1. Some observations about Kazaa
 - 2. A model for studying P2P multimedia systems
 - 3. Locality-aware request distribution

Methodology: trace characteristics

• 6-month Kazaa trace gathered at UW's border router

| start date | May 28 th , 2002 |
|--------------------------------|----------------------------------|
| end date | December 17 th , 2002 |
| trace length | 203 days, 5 hours, 6 minutes |
| # of requests | 1,640,912 |
| # of transactions | 98,997,622 |
| # of unsuccessful transactions | 65,505,165 (66.2%) |
| # of clients | 24,578 |
| # of unique objects | 633,106 (totaling 8.85TB) |
| bytes transferred | 22.72TB |
| content demanded | 43.87TB |

Outline

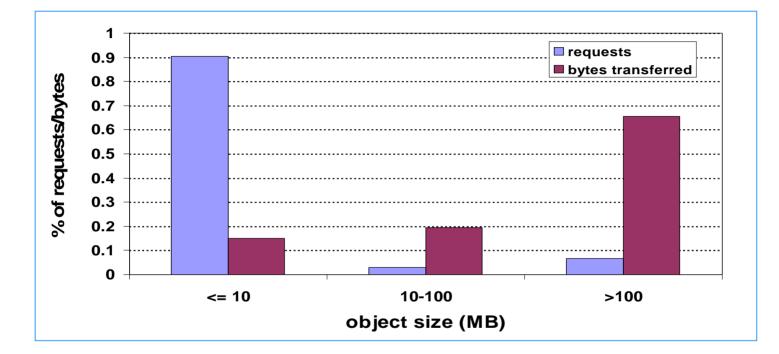
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1. Kazaa is really 2 phenomena

• Kazaa is *peer-to-peer*

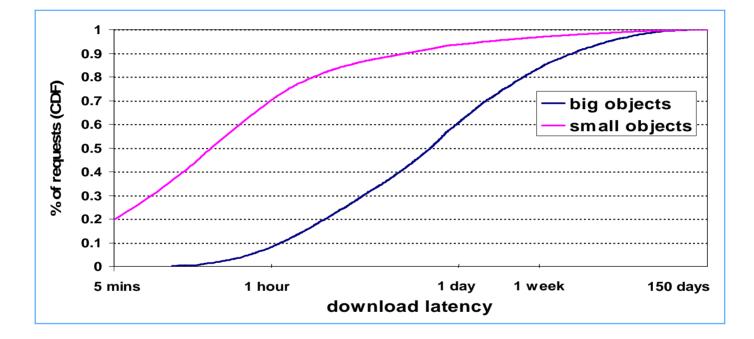
- *How* files are exchanged
- Content delivery system built out of volunteers
- Goal: understand behavior of the peers and implications
- Kazaa is a *multimedia* workload
 - What is exchanged
 - Broadband, MP3, huge disks sparked a new workload
 - Goal: understand implications of widespread multimedia on Internet, regardless of P2P vs. central, legality

2. Multimedia is really 2 workloads



- If you care about:
 - making users happy: make sure audio arrives quickly
 - making IT dept. happy: cache or rate limit video

3. Users are remarkably patient



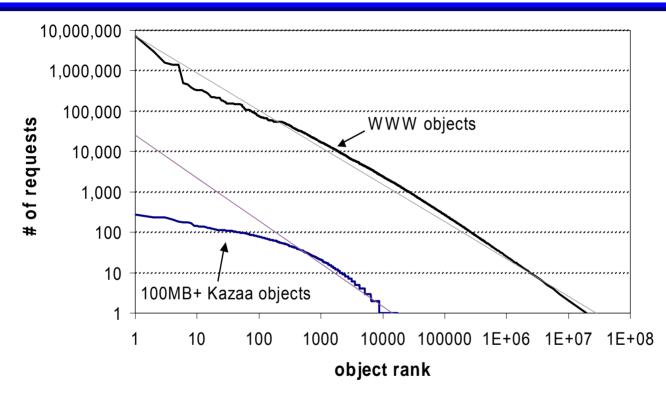
audio files take about an hour, video files a day

- but in either case, people will wait weeks!
- Web is an interactive system; Kazaa is a batch system!

4. Objects are immutable

- the Web is driven by object change
 - users revisit popular sites as their content changes
 - rate of change limits Web cache effectiveness [Wolman et al. 99]
- In contrast, Kazaa (multi-media) objects never change
 - users rarely re-download the same object
 - » 94% of the time, a user fetches an object at most once
 - » 99% of the time, a user fetches an object at most twice
 - Implication:
 - » popularity of most popular object is bounded by user population size

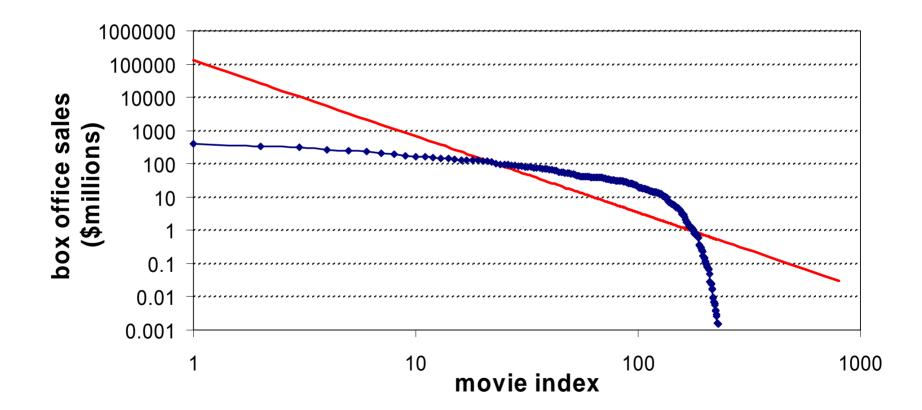
5. Kazaa does *not* obey Zipf's law



- Zipf: popularity(nth most popular object) ~ $1/n^{\alpha}$
- max # requests bounded by population size
- the most popular objects are much less popular than zipf would predict

Movie sales data (U.S.)

2002 U.S. box office ticket sales



6. Objects have quick turnover

- Popularity is short lived
 - only 5% of the top-100 audio objects stayed in the top-100 over our entire trace [video: 44%]
- Newly popular objects tend to be recently born
 - of audio objects that "broke into" the top-100, 79% were born a month before becoming popular [video: 84%]

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Modeling P2P Workloads

- Question: can we use a model to gain more insight into the forces driving P2P workloads?
- Objective: study three key issues identified by traces
 - 1. "fetch-at-most-once" behavior
 - 2. object birth rate
 - 3. client birth rate

Model basics

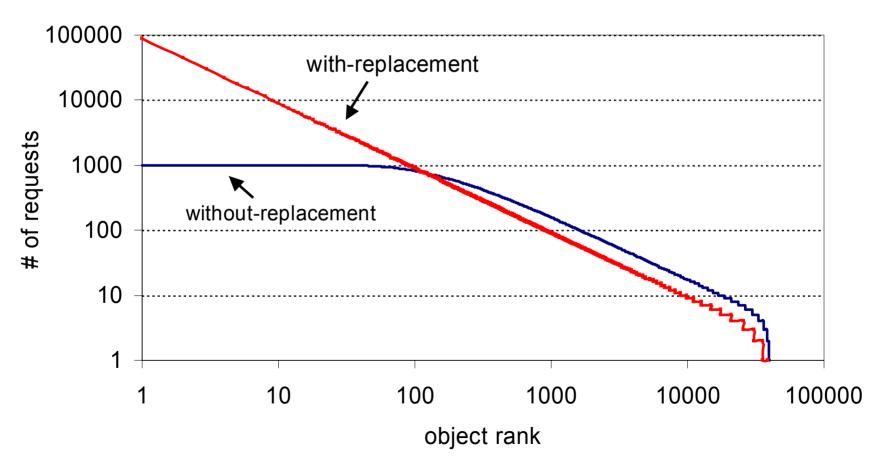
- Objects are chosen from an underlying Zipf curve, but assuming *"fetch-only-once"* behavior per client.
- Over time, users "coast" down the Zipf curve towards less popular objects
- New objects are inserted from the original Zipf (objects below pushed down)
- New users begin with a fresh Zipf curve

Model parameters

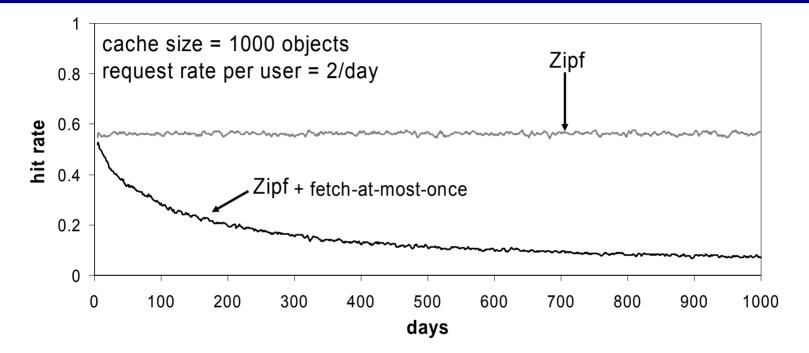
| С | # of clients | 1,000 |
|----------------|--|------------------------------------|
| 0 | # of objects | 40,000 |
| λ _R | client req. rate | 2 objects/day |
| α | Zipf param driving obj. popularity | 1.0 |
| P(x) | prob. of client req. object of pop rank x | Zipf (1.0) + fetch-at-most-once |
| A(x) | prob. of new object inserted at pop rank x | Zipf (1.0) |
| М | cache size (frac. of obj) | varies |
| λο | object arrival rate | varies |
| λ _c | client arrival rate | varies |

Fetching without replacement flattens Zipf curve head

modeled popularity results



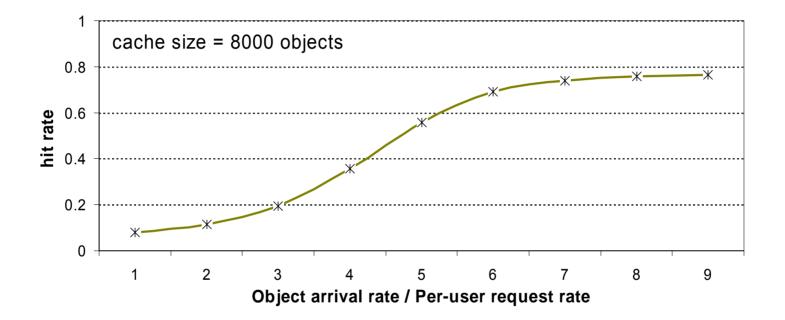
Caching implications



Given a fixed object and client population:

- with replacement \rightarrow hit rate stays flat over time
- without replacement \rightarrow hit rate degrades over time

New objects help, not hurt



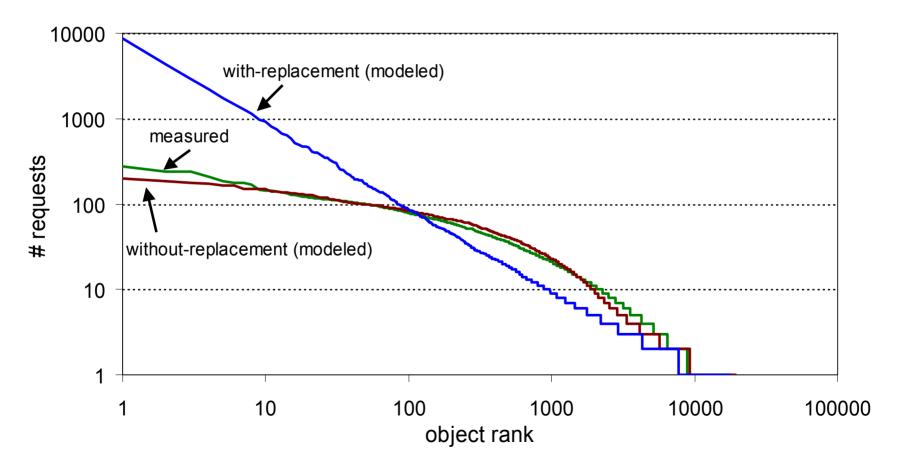
- New objects do cause cold misses
 - but they replenish the highly cacheable part of the Zipf curve
- A slow, constant arrival rate stabilizes performance
 - rate needed is proportional to avg. per-user request rate

New clients don't help

- they have some potential...
 - they have a "fresh" Zipf curve to draw from
 - therefore will have higher hit rate
- but new clients grow old too
 - ultimately they increase the size of the old population
 - to offset, must add clients at exponentially increasing rate
 - » not sustainable in long run...

Model validation

model parameterized by measured trace values



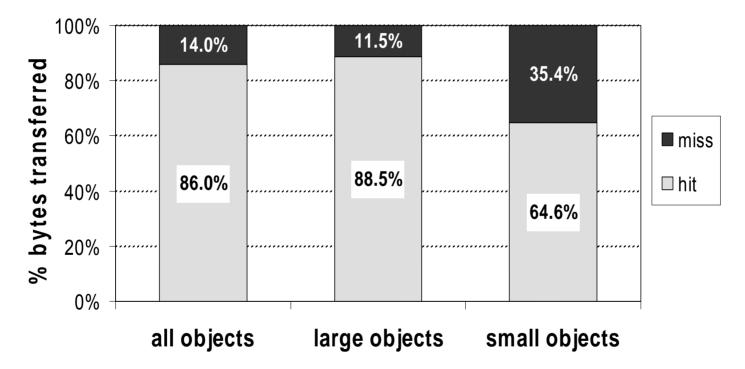
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Kazaa has untapped locality

simulated proxy cache hit rate for UW P2P environment

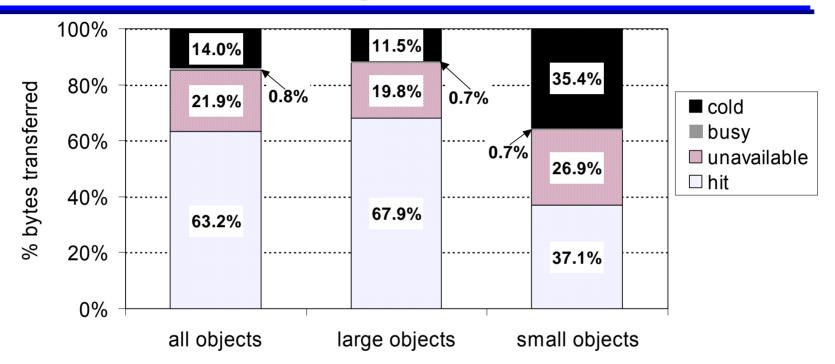


implication → 86% of bytes already exist within UW when they're downloaded externally!

Locality-aware request routing

- Alternative: make better use of local peers
- Scheme 1: use a *redirector* instead of a cache
 - Redirector sits at university border and watches traffic
 - It indexes content and *redirects* requests to local peers that can serve it
- Scheme 2: decentralized request distribution
 - Use location information in the P2P protocols
- We simulated locality-aware routing using our trace data
 - (note that both schemes are identical w.r.t. the simulation)

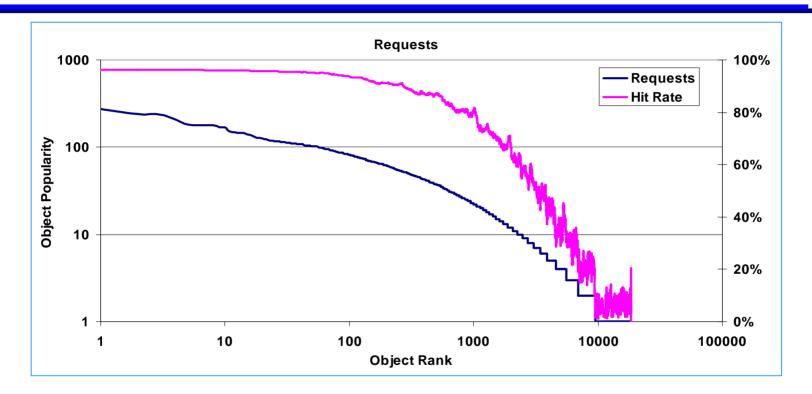
Basic locality-aware result



•Locality-aware schemes works!

- bars are a lower bound (assume worst availability from trace)
- •savings is huge (9.6 TB for large objects in our trace)
- •but there is still opportunity for improvement
- •need to figure out what causes unavailability misses

How can we eliminate availability misses?



- a kind of "natural replication" is driven by popularity
 - this is descriptive but also predictive
 - focus on "middle" popularity objects when designing systems

Summary of Part 2

- Kazaa/P2P is driven by completely different forces than the Web
 - multimedia content dominates the workload
 - system usage is "batch," not interactive
 - object immutability leads to *fetch-at-most-once* behavior
 - workload is driven by object and client births
 - multimedia is not "Zipf"
- Current file sharing architectures miss opportunity
 - Kazaa makes poor or no use of locality
 - locality-aware architectures can save significant amounts of external bandwidth
- We have a model that captures behavior of this workload

Final note

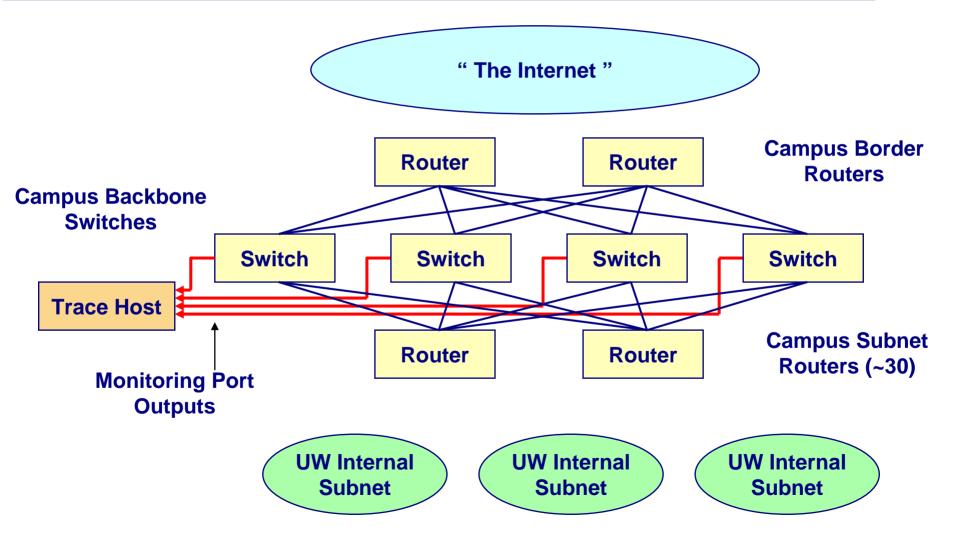
- The internet continues to reinvent itself every few years and can change dramatically at a moment's notice.
- Multi-media and peer-to-peer file sharing are going to be a huge component of the internet in the future -independent of current legal issues with P2P.

Pubs related to this talk

- An analysis of internet content delivery systems. *Proc.* 5th *Conf. on Operating Systems Design and Implementation* (OSDI), December 2002.
- Measurement, modelling, and analysis of a peer-to-peer file-sharing workload. Proc. 19th ACM Symp. on Operating Systems Principles (SOSP), October 2003.

http://www.cs.washington.edu/research/networking/websys

UW Trace Environment



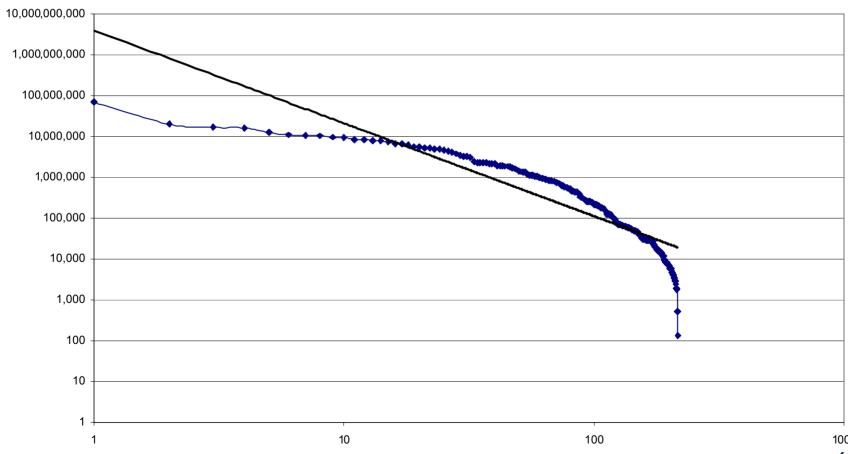
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Movie sales data (UK)

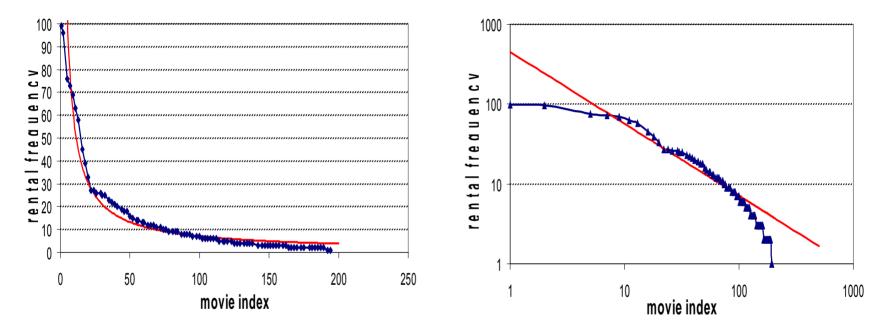
Revenue (in British Pounds)

UK Film Revenues in 1998



Movie Rank

Previously reported Zipf results



•Left side: graph from Dan, Sitaram, and Shahabuddin, "Scheduling policies for an on-demand video server with batching," *Proc. of ACM Multimedia 1994*, Oct. 1994, "demonstrating" Zipf fit of 1992 video rental data.

•Right side: same data plotted on a log/log scale.