Database Management Systems CSEP 544

Lecture 2: SQL

CSEP 544 - Fall 2017

Announcements

- HW1 due tonight (11:59pm)
- PA2 & HW2 released
- Fill out HW3 email account form by tonight!
- Final information posted on piazza
- Check website for up to date OH info



- Data models
 - Instance
 - Schema
 - Language
- Relational data model
 - Relations are flat
 - Tuples are not ordered
- Logical and physical data independence

Reading Assignment 1

Selections in SQL

SELECT *

FROM Product

WHERE price > 100.0

Projections in SQL

SELECT CName FROM Product

Joins in SQL

]
pname	price	category	manufacturer	cname
MultiTouch	199.99	gadget	Canon	Gizmo
SingleTouch	49.99	photography	Canon	Canon
Gizom	50	gadget	GizmoWorks	Hitach
SuperGizmo	250.00	gadget	GizmoWorks] [

cname	country	
GizmoWorks	USA	
Canon	Japan	
Hitachi	Japan	

Retrieve all Japanese products that cost < \$150

Joins in SQL

				Ι.	
pname	price	category	manufacturer		cname
MultiTouch	199.99	gadget	Canon		Gizmo'
SingleTouch	49.99	photography	Canon		Canon
Gizom	50	gadget	GizmoWorks		Hitachi
SuperGizmo	250.00	gadget	GizmoWorks		

cnamecountryGizmoWorksUSACanonJapanHitachiJapan

Retrieve all Japanese products that cost < \$150

```
SELECT pname, price
FROM Product, Company
WHERE ...
```

Joins in SQL

				Ι.		
pname	price	category	manufacturer	,	cname	country
MultiTouch	199.99	gadget	Canon		GizmoWorks	USA
SingleTouch	49.99	photography	Canon	-	Canon	Japan
Gizom	50	gadget	GizmoWorks	-	Hitachi	Japan
SuperGizmo	250.00	gadget	GizmoWorks			-

Retrieve all Japanese products that cost < \$150

SELECT	pname, price				
FROM	Product, Company				
WHERE	manufacturer=cname AND				
	country='Japan' AND price < 150				

Joins in SQL

				Ι.		
pname	price	category	manufacturer		cname	country
MultiTouch	199.99	gadget	Canon		GizmoWorks	USA
SingleTouch	49.99	photography	Canon		Canon	Japan
Gizom	50	gadget	GizmoWorks		Hitachi	Japan
SuperGizmo	250.00	gadget	GizmoWorks			

Retrieve all Japanese products that cost < \$150

SELECT P.pname, P.price
FROM Product as P, Company as C
WHERE P.manufacturer=C.cname AND
C.country='Japan' AND C.price < 150</pre>

Joins in SQL

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

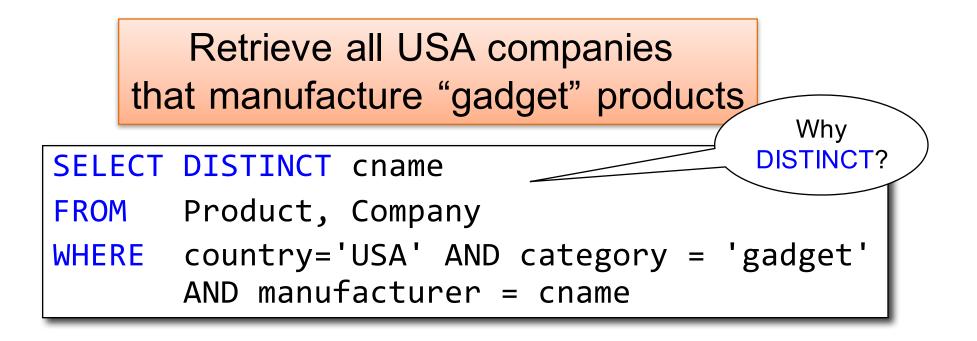
cname	country	
GizmoWorks	USA	
Canon	Japan	
Hitachi	Japan	

Retrieve all USA companies that manufacture "gadget" products

Joins in SQL

pname	price	category	manufacturer	cr
MultiTouch	199.99	gadget	Canon	G
SingleTouch	49.99	photography	Canon	С
Gizom	50	gadget	GizmoWorks	Н
SuperGizmo	250.00	gadget	GizmoWorks	

cname	country	
GizmoWorks	USA	
Canon	Japan	
Hitachi	Japan	



Joins in SQL

- The standard join in SQL is sometimes called an inner join
 - Each row in the result must come from
 both tables in the join
- Sometimes we want to include rows from only one of the two table: outer join

Joins and Aggregates

Product(<u>pname</u>, price, category, manufacturer)
Company(<u>cname</u>, country)

-- manufacturer is foreign key to Company

SELECT	DISTINCT cname
FROM	Product, Company
WHERE	<pre>country='USA' AND category = 'gadget' AND manufacturer = cname</pre>

SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget' AND manufacturer = cname

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
Camera	Photo	Hitachi
OneClick	Photo	Hitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget' AND manufacturer = cname

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
Camera	Photo	Hitachi
OneClick	Photo	Hitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

SELECT DISTINCT cname

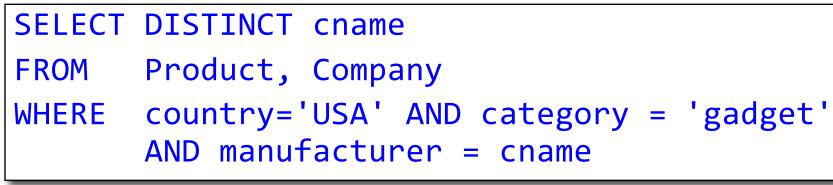
FROM Product, Company

WHERE country='USA' AND category = 'gadget' AND manufacturer = cname

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
Camera	Photo	Hitachi
OneClick	Photo	Hitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan



•Product

	pname	category	ma	nufacturer			CNa	ame	CO	untry
	Gizmo	gadget	Gi	GizmoWorks			Gizmo	oWorks	L	JSA
	Camera	Photo		Hitachi			Ca	non	Ja	apan
(OneClick	Photo	Hitachi				Hitachi		Ja	apan
	pname	category	/	manufactu	rer	cn	ame	cour	ntry	
	Gizmo	gadget		GizmoWo	ks	Gizm	oWorks	US	A	25

SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget' AND manufacturer = cname

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
Camera	Photo	Hitachi
OneClick	Photo	Hitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget' AND manufacturer = cname

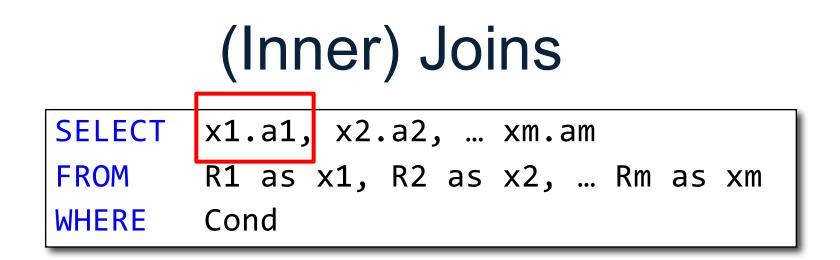
Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
Camera	Photo	Hitachi
OneClick	Photo	Hitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

SELECT	DISTINCT cname
FROM	Product, Company
WHERE	country='USA' AND category = 'gadget' AND manufacturer = cname

SELECT	DISTINCT cname
FROM	<pre>Product JOIN Company ON country = 'USA' AND category = 'gadget' AND manufacturer = cname</pre>



```
for x1 in R1:
for x2 in R2:
```

...
for xm in Rm:
 if Cond(x1, x2...):
 output(x1.a1, x2.a2, ... xm.am)
This is called nested loop semantics since we are
 interpreting what a join means using a nested loop

Another example

Product(pname, price, category, manufacturer)
Company(cname, country)

-- manufacturer is foreign key to Company

Retrieve all Japanese companies that manufacture products in both 'gadget' and 'photography' categories

Another example

Product(pname, price, category, manufacturer)
Company(cname, country)

-- manufacturer is foreign key to Company

Retrieve all Japanese companies that manufacture products in both 'gadget' and 'photography' categories

```
SELECT DISTINCT cname
FROM Product P1, Product P2, Company
WHERE country = 'Japan' AND P1.category = 'gadget'
AND P2.category = 'photography'
AND P1.manufacturer = cname
AND P2.manufacturer = cname;
```

Self-Joins and Tuple Variables

- Find all companies that manufacture both products in the 'gadgets' and 'photo' category
- Joining Product with Company is insufficient: need to join Product, with Product, and with Company
- When a relation occurs twice in the FROM clause we call it a *self-join*
 - in that case we must use tuple variables (why?)

SELECT DISTINCT z.cname
FROM Product x, Product y, Company z
WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname

AND y.manufacturer = cname;

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
SingleTouch	photo	Hitachi
MultiTouch	Photo	GizmoWorks

cname	country
GizmoWorks	USA
Hitachi	Japan

SELECT DISTINCT z.cname
FROM Product x, Product y, Company z
WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname
AND y.manufacturer = cname;

Product

X

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
SingleTouch	photo	Hitachi
MultiTouch	Photo	GizmoWorks

cname	country
GizmoWorks	USA
Hitachi	Japan

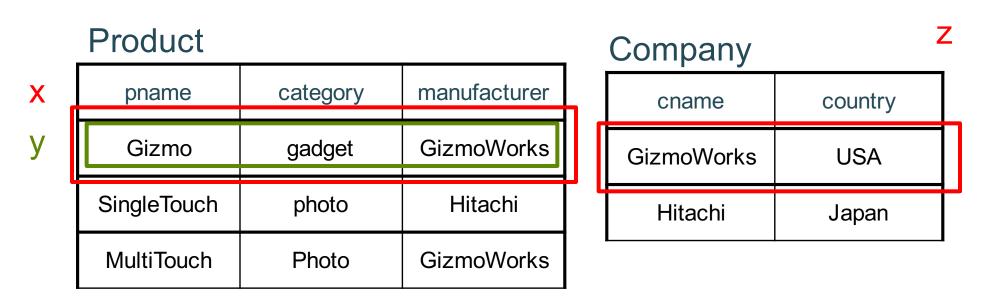
SELECT DISTINCT z.cname
FROM Product x, Product y, Company z
WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname
AND y.manufacturer = cname;

Product

x	pname	category	manufacturer
y	Gizmo	gadget	GizmoWorks
	SingleTouch	photo	Hitachi
	MultiTouch	Photo	GizmoWorks

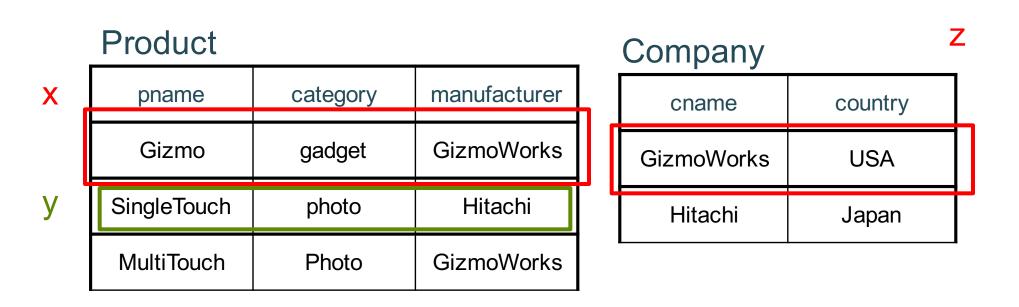
cname	country
GizmoWorks	USA
Hitachi	Japan

SELECT DISTINCT z.cname
FROM Product x, Product y, Company z
WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname
AND y.manufacturer = cname;



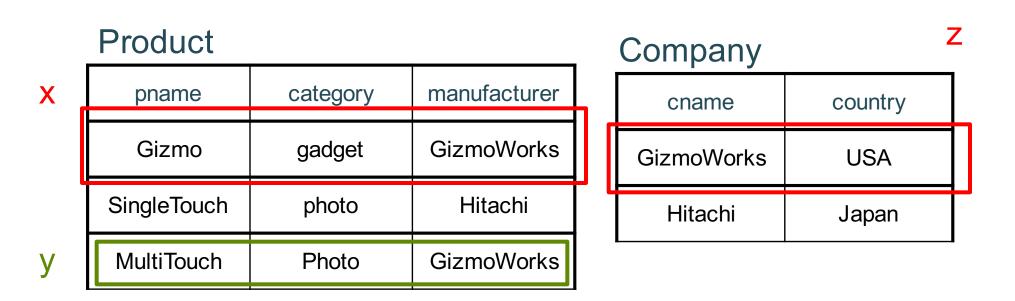
SELECT DISTINCT z.cname
FROM Product x, Product y, Company z
WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname

AND y.manufacturer = cname;



SELECT DISTINCT z.cname
FROM Product x, Product y,

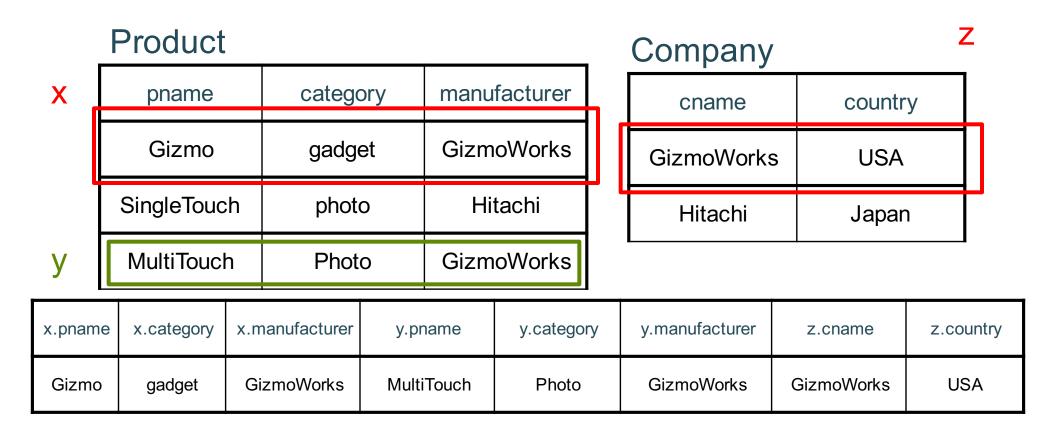
FROM Product x, Product y, Company z
WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname
AND y.manufacturer = cname;



SELECT DISTINCT z.cname

FROM Product x, Product y, Company z

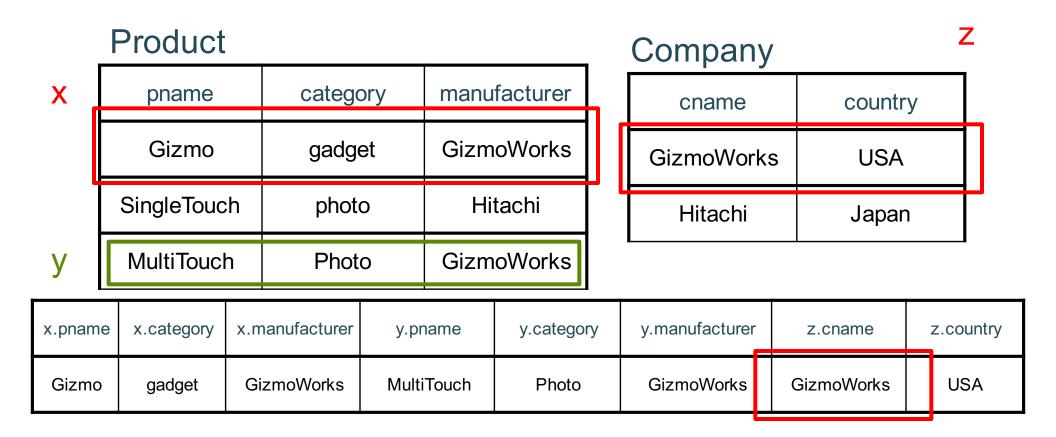
WHERE z.country = 'USA' AND x.category = 'gadget' AND y.category = 'photo' AND x.manufacturer = cname AND y.manufacturer = cname;



SELECT DISTINCT z.cname

FROM Product x, Product y, Company z

WHERE z.country = 'USA' AND x.category = 'gadget' AND y.category = 'photo' AND x.manufacturer = cname AND y.manufacturer = cname;



Outer joins

Product(<u>name</u>, category)
Purchase(prodName, store)

-- prodName is foreign key

SELECT	Product.name, Purchase.store
FROM	Product, Purchase
WHERE	<pre>Product.name = Purchase.prodName</pre>

We want to include products that are never sold, but some are not listed! Why?

Outer joins

Product(<u>name</u>, category)
Purchase(prodName, store)

-- prodName is foreign key

SELECT Product.name, Purchase.store
FROM Product LEFT OUTER JOIN Purchase ON
Product.name = Purchase.prodName

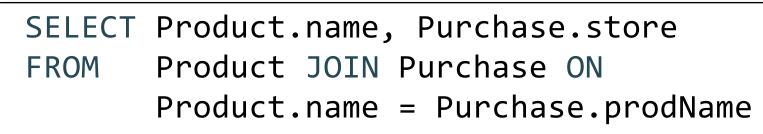
SELECT Product.name, Purchase.store FROM Product JOIN Purchase ON Product.name = Purchase.prodName

Product

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

Purchase

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

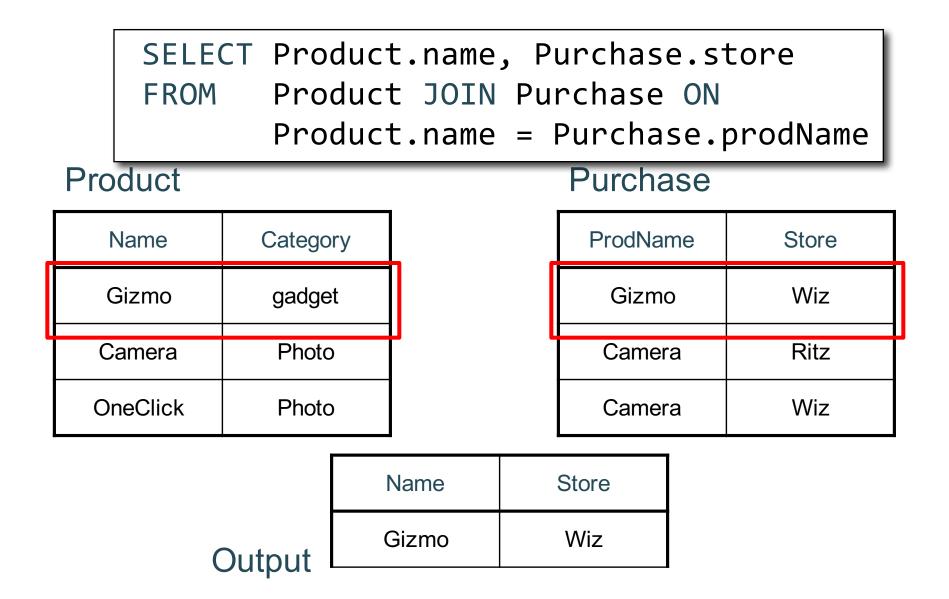


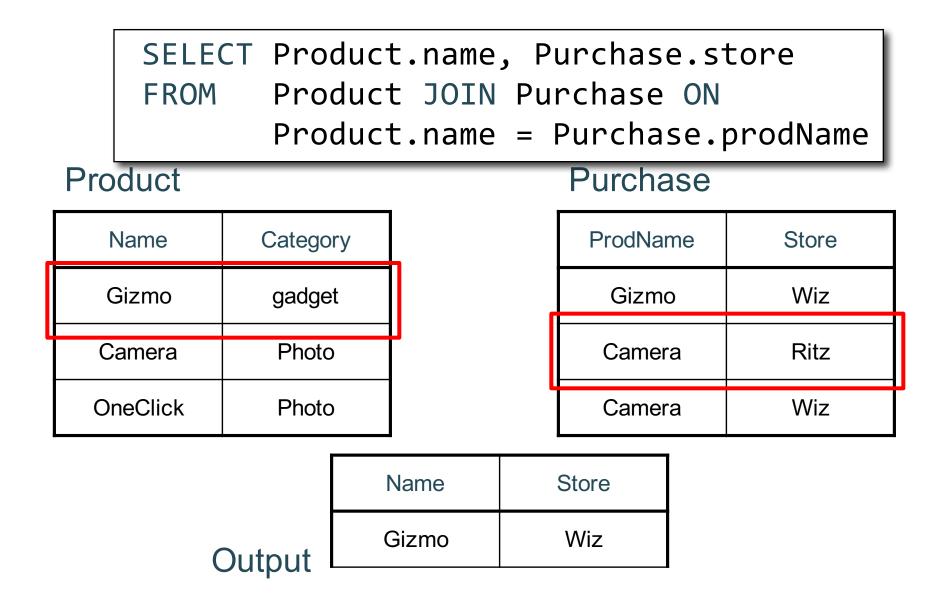
Product

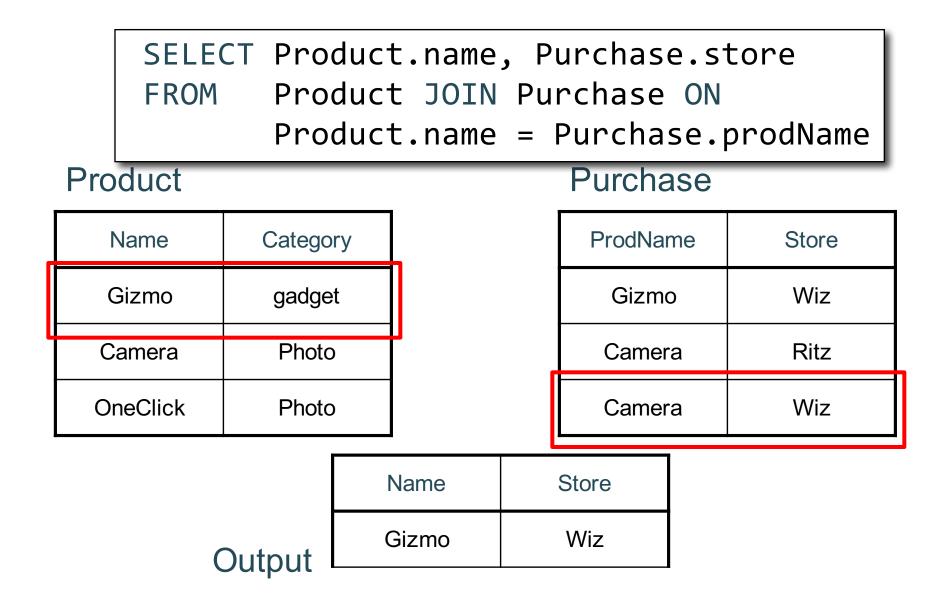
Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

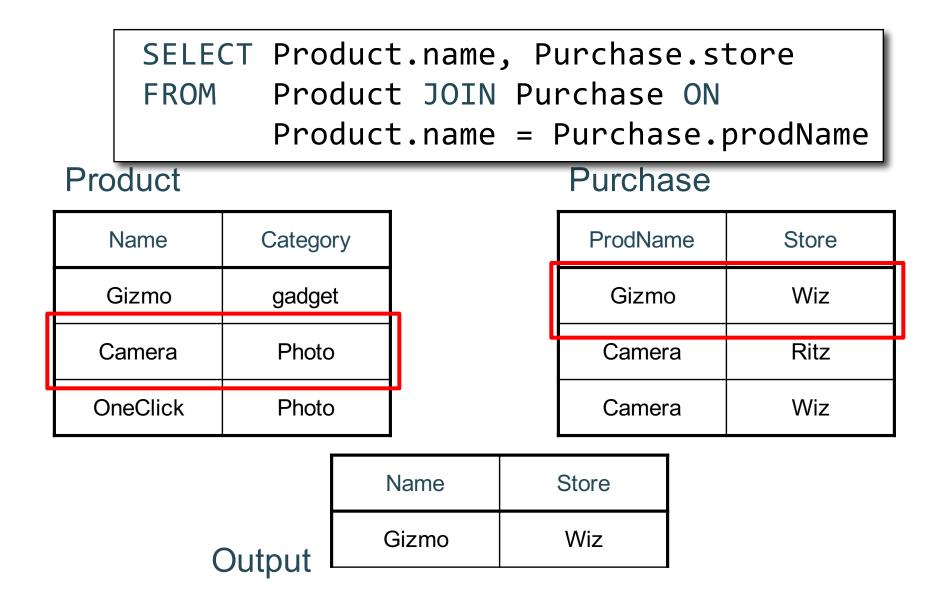
Purchase

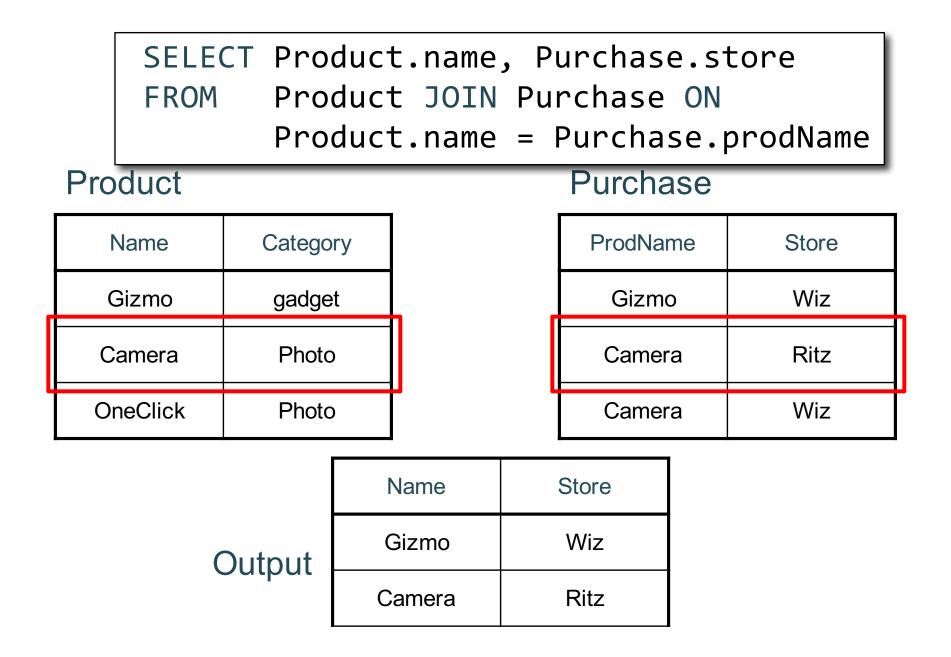
ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

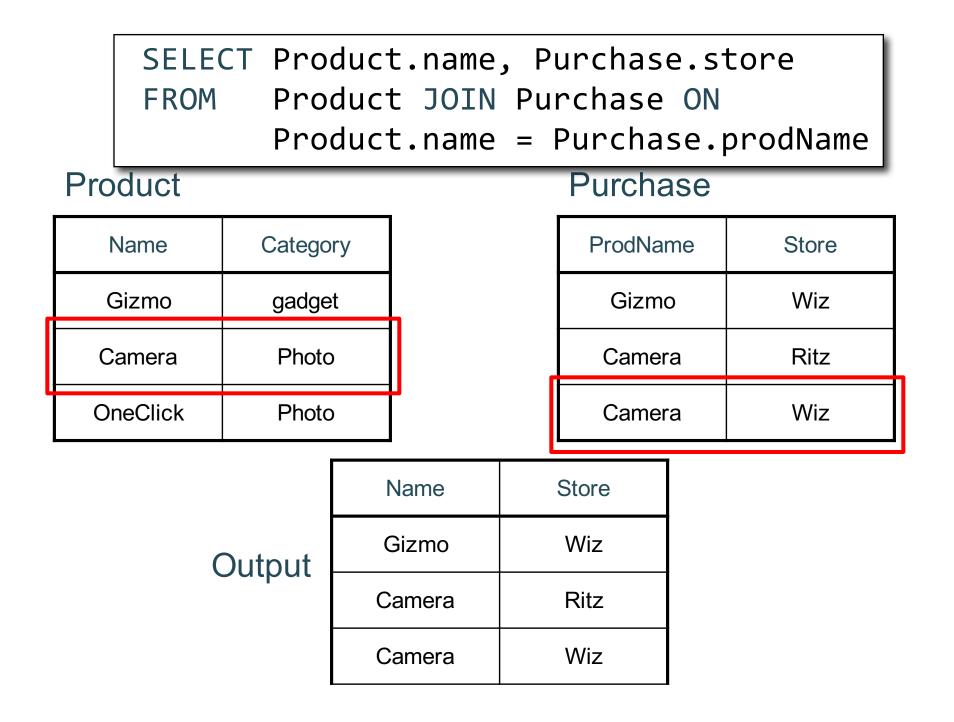


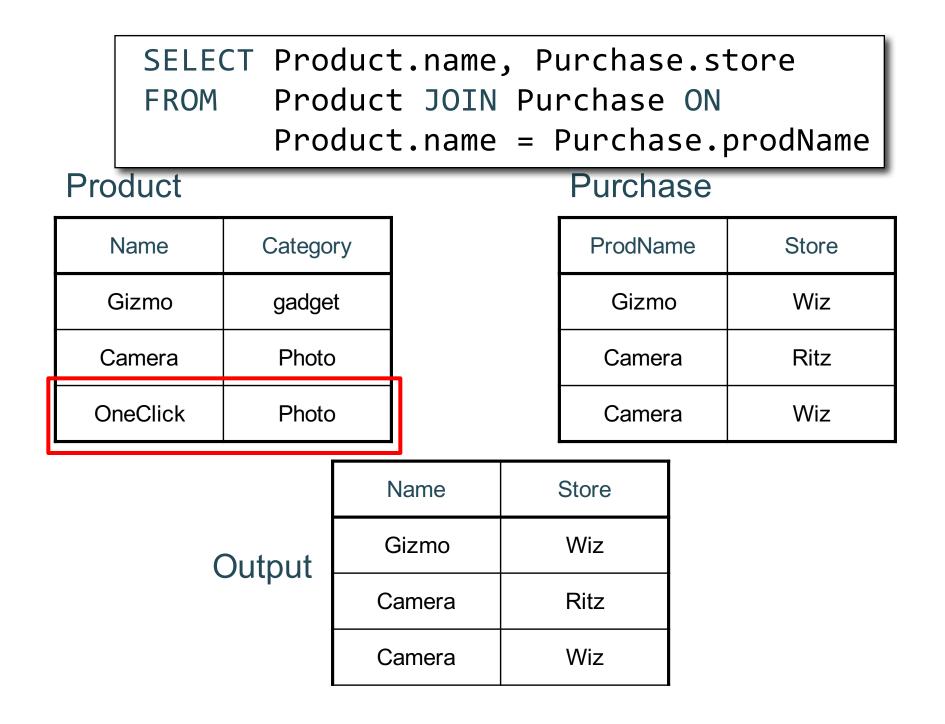


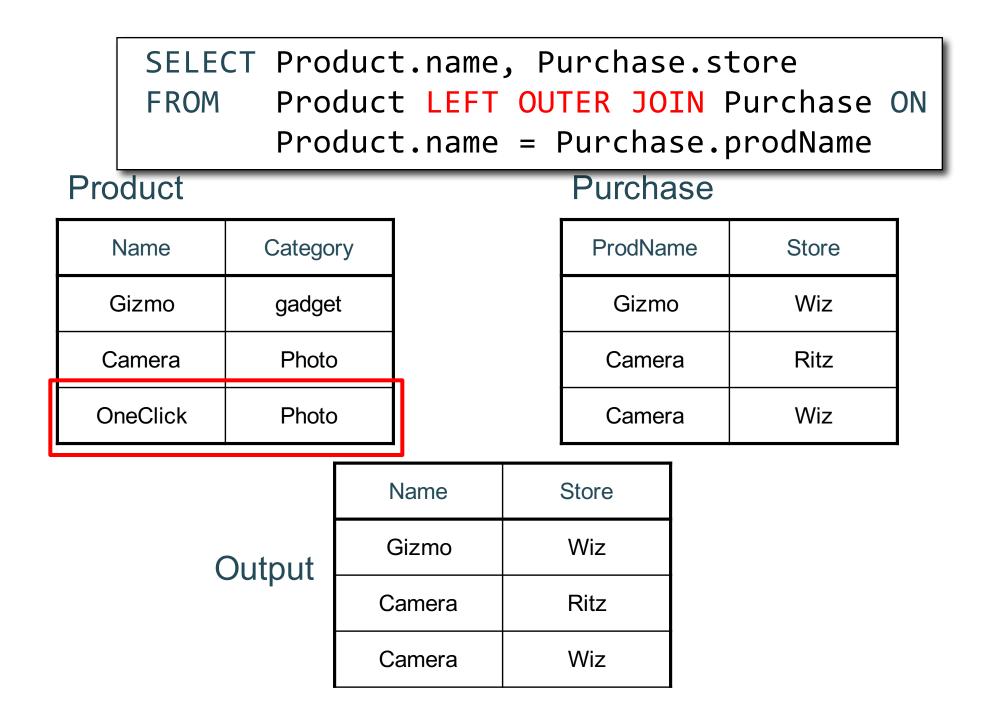


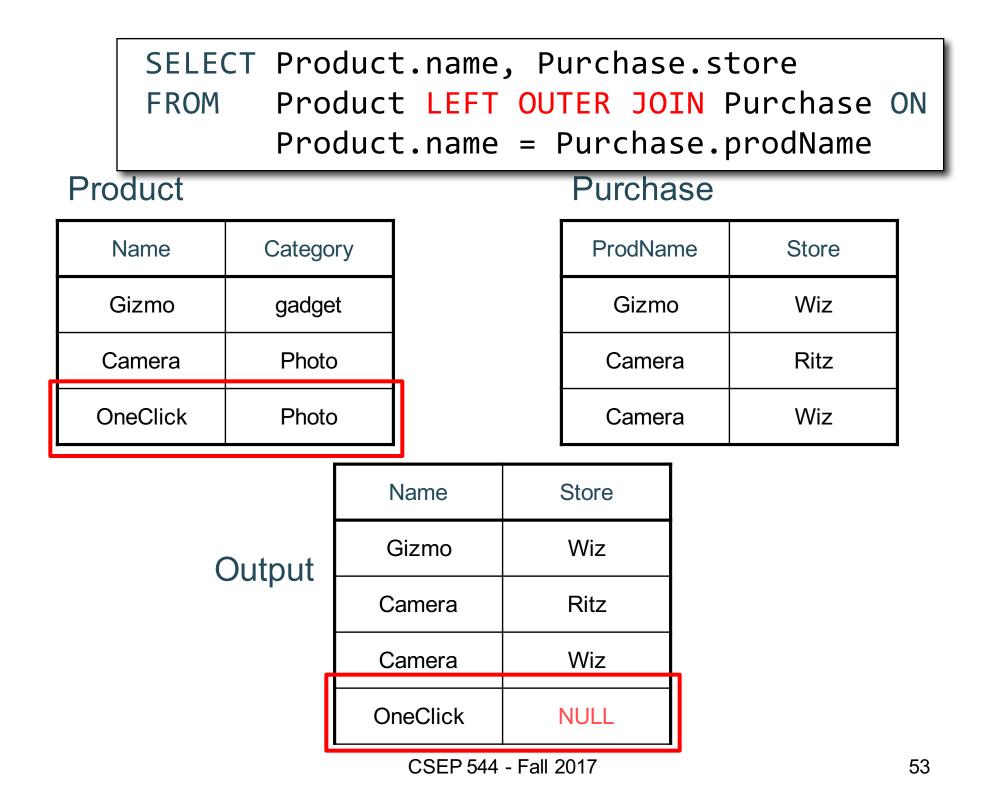












SELECT Product.name, Purchase.store FROM Product FULL OUTER JOIN Purchase ON Product.name = Purchase.prodName					ON
Product Purchase					
Name	Category		ProdName	Store]
Gizmo	gadget		Gizmo	Wiz	1
Camera	Photo		Camera	Ritz	1
OneClick	Photo		Camera	Wiz	1
	Name	Store	Phone	Foo	1
	Gizmo	Wiz			-
Output	Camera	Ritz			
	Camera	Wiz			
	OneClick	NULL			
	NULL	Foo			54

Outer Joins

tableA (LEFT/RIGHT/FULL) OUTER JOIN tableB ON p

- Left outer join:
 - Include tuples from tableA even if no match
- Right outer join:
 - Include tuples from tableB even if no match
- Full outer join:
 - Include tuples from both even if no match
- In all cases:
 - Patch tuples without matches using NULL

Comment about SQLite

- Cannot load NULL values such that they are actually loaded as null values
- So we need to use two steps:
 - Load null values using some type of special value
 - Update the special values to actual null values

```
update Purchase
  set price = null
  where price = 'null'
```

Simple Aggregations

Five basic aggregate operations in SQL

select count(*) from Purchase
select sum(quantity) from Purchase
select avg(price) from Purchase
select max(quantity) from Purchase
select min(quantity) from Purchase

Except count, all aggregations apply to a single attribute

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Aggregates and NULL Values

Null values are not used in aggregates

insert into Purchase
values(12, 'gadget', NULL, NULL, 'april')

Try the following at home:

select count(*) from Purchase
select count(quantity) from Purchase

```
select sum(quantity) from Purchase
```

```
select count(*)
from Purchase
where quantity is not null;
```

Counting Duplicates

COUNT applies to duplicates, unless otherwise stated:

SELECT	<pre>count(product)</pre>
FROM	
WHERE	price > 4.99

same as count(*) if no nulls

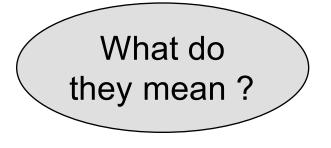
We probably want:

SELECT	<pre>count(DISTINCT product)</pre>
FROM	Purchase
WHERE	price > 4.99

More Examples

SELECTSum(price * quantity)FROMPurchase

SELECT Sum(price * quantity) FROM Purchase WHERE product = 'bagel'



Grouping and Query Evaluation

Grouping and Aggregation

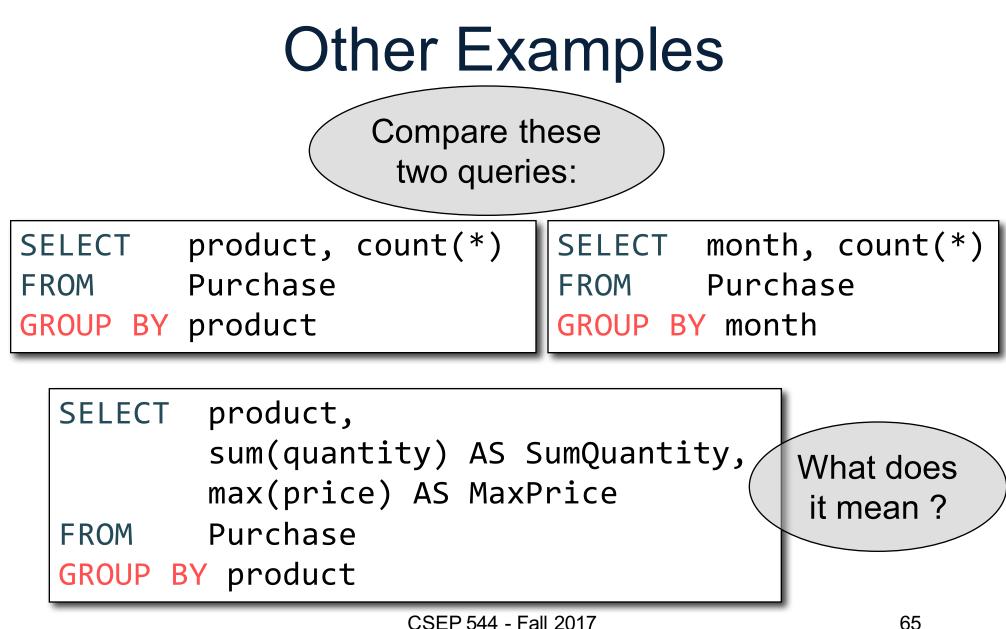
Purchase(product, price, quantity)

Find total quantities for all sales over \$1, by product.

Grouping and Aggregation

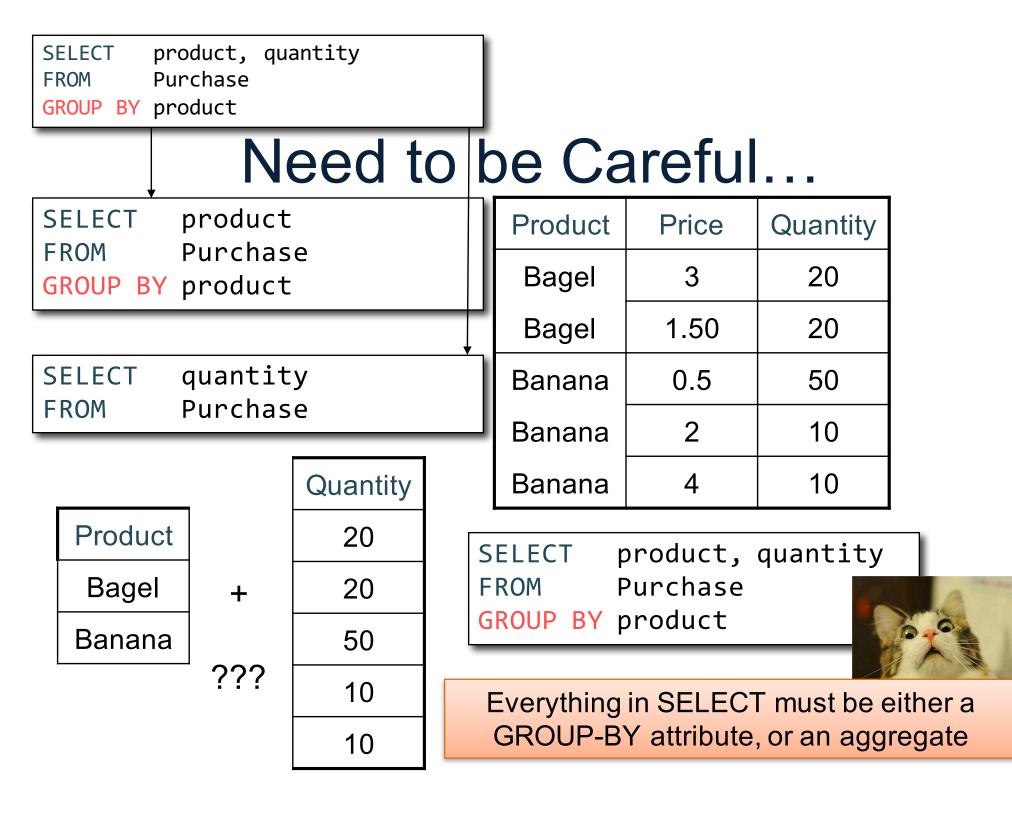
Product	Price	Quantity			
Bagel	3	20	N	Product	TotalSales
Bagel	1.50	20		Bagel	40
Banana	0.5	50		Banana	20
Banana	2	10			
Banana	4	10			

SELECT	<pre>product, Sum(quantity) AS TotalSales</pre>	5
FROM	Purchase	
WHERE	price > 1	
GROUP BY	product 64	



Need to be Careful...

SELECT product, max(quantity)		Product	Price	Quantity
FROM Purchase		Bagel	3	20
GROUP BY product		Bagel	1.50	20
SELECT product, quant	tity	Banana	0.5	50
FROM Purchase GROUP BY product		Banana	2	10
		Banana	4	10



Need to be Careful...

SELECT product, max(quantity)		Product	Price	Quantity
FROM Purchase		Bagel	3	20
GROUP BY product		Bagel	1.50	20
SELECT product, quar	ntity	Banana	0.5	50
FROMPurchaseGROUPBYproduct		Banana	2	10
		Banana	4	10
sqlite is WRONG on this query.		Advanced DE Server) giv	3MS (e.g. So ves an error	

Grouping and Aggregation

Purchase(product, price, quantity)

Find total quantities for all sales over \$1, by product.

SELECT	<pre>product, Sum(quantity) AS TotalSales</pre>
FROM	Purchase
WHERE	price > 1
GROUP BY	product

How is this query processed?

Grouping and Aggregation

- 1. Compute the FROM and WHERE clauses.
- 2. Group by the attributes in the GROUPBY
- 3. Compute the SELECT clause: grouped attributes and aggregates.

1,2: From, Where FWGS

Product	Price	Quantity	
Bagel	3	20	
Bagel	1.50	20	
Banana	0.5	50	
Banana	2	10	WHERE price > 1
Banana	4	10	WHERE price > 1

SELECT	<pre>product, Sum(quantity) AS TotalSa</pre>	les
FROM	Purchase	
WHERE	price > 1	
GROUP BY	product	71

3,4. Grouping, Select FWGS

Product	Price	Quantity		
Bagel	3	20	Product	TotalSales
Bagel	1.50	20	Bagel	40
Banana	0.5	50	Banana	20
Banana	2	10		
Banana	4	10		

SELECT	product, Sum(quantity) AS TotalSale	es :
FROM	Purchase	
WHERE	price > 1	
GROUP BY	product 72	

Ordering Results

SELECT product, sum(price*quantity) as rev
FROM Purchase
GROUP BY product
ORDER BY rev desc



Note: some SQL engines want you to say ORDER BY sum(price*quantity) desc

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HAVING Clause

Same query as before, except that we consider only products that had at least 30 sales.

SELECT	<pre>product, sum(price*quantity)</pre>	
FROM	Purchase	
WHERE	price > 1	
GROUP BY	product	
HAVING	<pre>sum(quantity) > 30</pre>	

HAVING clause contains conditions on aggregates.

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General form of Grouping and Aggregation

SELECTSFROM $R_1, ..., R_n$ WHEREC1GROUP BY $a_1, ..., a_k$ HAVINGC2

- S = may contain attributes a₁,...,a_k and/or any aggregates but NO OTHER ATTRIBUTES
- C1 = is any condition on the attributes in $R_1, ..., R_n$
- C2 = is any condition on aggregate expressions and on attributes a_1, \ldots, a_k

Why?

Semantics of SQL With Group-By

SELECT	S
FROM	R ₁ ,, R _n
WHERE	C1
GROUP BY	a ₁ ,,a _k
HAVING	C2

FWGHOS

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Evaluation steps:

- 1. Evaluate FROM-WHERE using Nested Loop Semantics
- 2. Group by the attributes a_1, \ldots, a_k
- 3. Apply condition C2 to each group (may have aggregates)
- 4. Compute aggregates in S and return the result CSEP 544 - Fall 2017

Exercise

Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

Exercise

Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

FROM Purchase

Exercise

Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

FROM		Purchase	
GROUP	BY	month	

Exercise

Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

FROM	Purchase
GROUP BY	month
HAVING	sum(quantity) < 10

Purchase(pid, product, price, quantity, month)

Exercise

Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

SELECT	<pre>month, sum(price*quantity),</pre>	
	<pre>sum(quantity) as TotalSold</pre>	
FROM	Purchase	
GROUP BY	month	
HAVING	sum(quantity) < 10	

Purchase(pid, product, price, quantity, month)

Exercise

Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

SELECT	<pre>month, sum(price*quantity), sum(quantity) as TotalSold</pre>
FROM	Purchase
GROUP BY	month
HAVING	sum(quantity) < 10
ORDER BY	<pre>sum(quantity)</pre>

WHERE vs HAVING

- WHERE condition is applied to individual rows
 - The rows may or may not contribute to the aggregate
 - No aggregates allowed here
- HAVING condition is applied to the entire group
 - Entire group is returned, or not at all
 - May use aggregate functions in the group

Purchase(pid, product, price, quantity, month)

Mystery Query

What do they compute?

SELECT	month, sum(quantity), max(price)
FROM	Purchase
GROUP BY	month

SELECTmonth, sum(quantity)FROMPurchaseGROUP BYmonth

SELECT	month
FROM	Purchase
GROUP BY	month

Purchase(pid, product, price, quantity, month)

Mystery Query

What do they compute?

SELECTmonth, sum(quantity), max(price)FROMPurchaseGROUP BYmonth

SELECTmonth, sum(quantity)FROMPurchaseGROUP BYmonth

SELECTmonthFROMPurchaseGROUP BYmonth

Lesson: DISTINCT is a special case of GROUP BY Product(pid,pname,manufacturer)
Purchase(id,product_id,price,month)

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Product(pid,pname,manufacturer)
Purchase(id,product_id,price,month)

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: price is in Purchase, manufacturer is in Product...

Product(<u>pid</u>,pname,manufacturer)
Purchase(<u>id</u>,product_id,price,month)

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: price is in Purchase, manufacturer is in Product...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
```

manu facturer	 price	
Hitachi	150	
Canon	300	
Hitachi	180	

Product(pid,pname,manufacturer)
Purchase(id,product_id,price,month)

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: price is in Purchase, manufacturer is in Product...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
```

manu facturer	 price	
Hitachi	150	
Canon	300	
Hitachi	180	

```
-- step 2: do the group-by on the join
SELECT x.manufacturer, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
GROUP BY x.manufacturer
```

manu facturer	count(*)	
Hitachi	2	
Canon	1	

Product(<u>pid</u>,pname,manufacturer)
Purchase(<u>id</u>,product_id,price,month)

Aggregate + Join

Variant:

For each manufacturer, compute how many products with price > \$100 they sold in each month

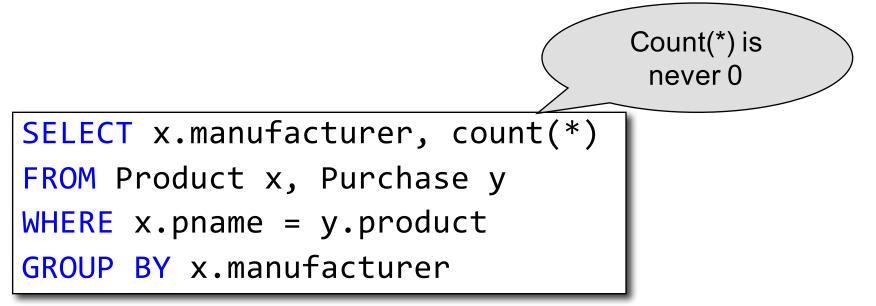
```
SELECT x.manufacturer, y.month, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
GROUP BY x.manufacturer, y.month
```

manu facturer	month	count(*)
Hitachi	Jan	2
Hitachi	Feb	1
Canon	Jan	3

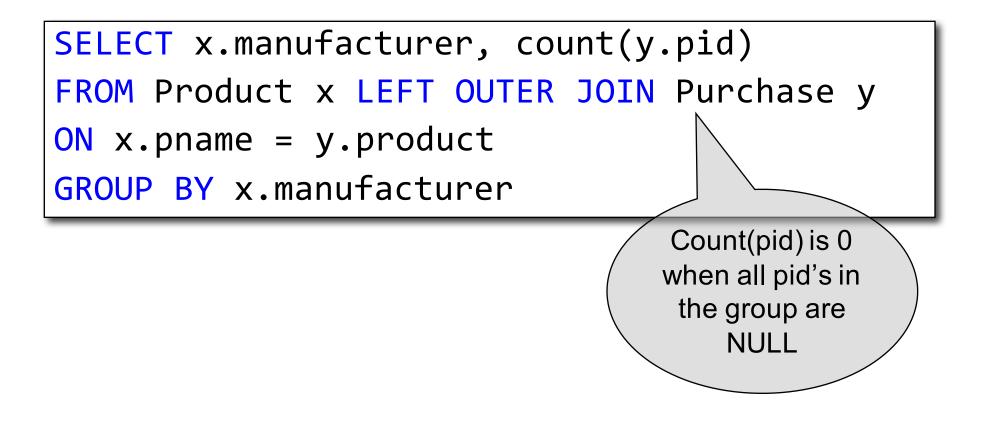


Including Empty Groups

• In the result of a group by query, there is one row per group in the result



Including Empty Groups



Nested Queries

What have we learned so far

- Data models
- Relational data model
 - Instance: relations
 - Schema: table with attribute names
 - Language: SQL

What have we learned so far

- SQL features
 - Projections
 - Selections
 - Joins (inner and outer)
 - Aggregates
 - Group by
 - Inserts, updates, and deletes
- Make sure you read the textbook!

Subqueries

- A subquery is a SQL query nested inside a larger query
- Such inner-outer queries are called nested queries
- A subquery may occur in:
 - A SELECT clause
 - A FROM clause
 - A WHERE clause
- Rule of thumb: avoid writing nested queries when possible
 - But sometimes it's impossible, as we will see

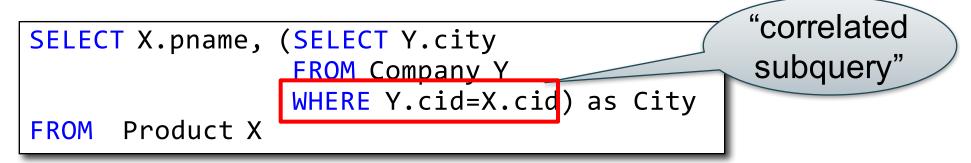
Subqueries...

- Can appear as computed values in a SELECT clause
- Can appear in FROM clauses and aliased using a tuple variable that represents the tuples in the result of the subquery
- Can return a single constant to be compared with another value in a WHERE clause
- Can return relations to be used in WHERE clauses

1. Subqueries in SELECT

Product (pname, price, cid)
Company (cid, cname, city)

For each product return the city where it is manufactured



What happens if the subquery returns more than one city?

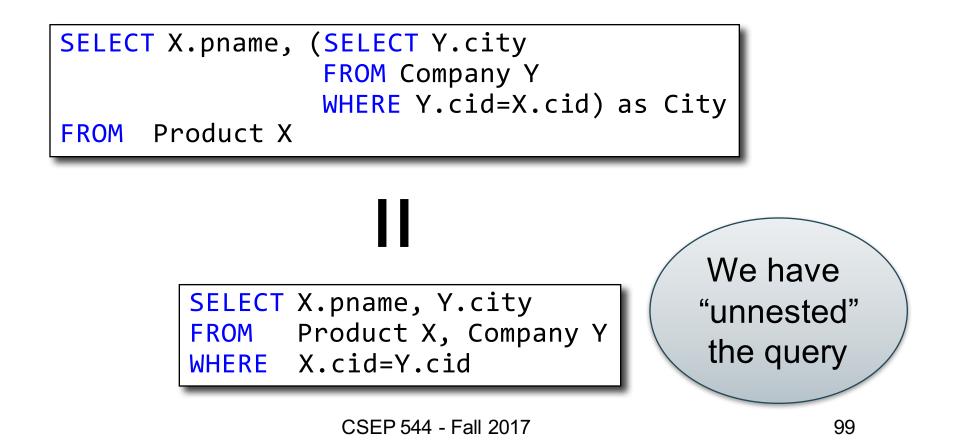
We get a runtime error

(and SQLite simply ignores the extra values...)

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1. Subqueries in SELECT

Whenever possible, don't use nested queries:



1. Subqueries in SELECT

Compute the number of products made by each company

SELECT	DISTINCT	C.cname,	(SELECT	<pre>count(*)</pre>
			FROM Pr	roduct P
			WHERE F	<pre>P.cid=C.cid)</pre>
FROM	Company C			

1. Subqueries in SELECT

Compute the number of products made by each company

SELEC	T DISTINCT C.cname,	(SELECT count(*)
		FROM Product P
		WHERE P.cid=C.cid)
FROM	Company C	

Better: we can unnest using a GROUP BY SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname

1. Subqueries in SELECT

But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
```

FROM Company C

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

1. Subqueries in SELECT

But are these really equivalent?

SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
No! Different results if a
company has no products

```
SELECT C.cname, count(pname)
FROM Company C LEFT OUTER JOIN Product P
ON C.cid=P.cid
GROUP BY C.cname
```

2. Subqueries in FROM

Find all products whose prices is > 20 and < 500

SELECT X.pname
FROM (SELECT *
 FROM Product AS Y
 WHERE price > 20) as X
WHERE X.price < 500</pre>

Side note: This is not a correlated subquery. (why?)

Try unnest this query !

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2. Subqueries in FROM

At the end of the lecture we will see that sometimes we really need a subquery and one option will be to put it in the FROM clause. Product (<u>pname</u>, price, cid) Company (<u>cid</u>, cname, city)

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

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Find all companies that make <u>some</u> products with price < 200

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

```
Using EXISTS:

SELECT DISTINCT C.cname

FROM Company C

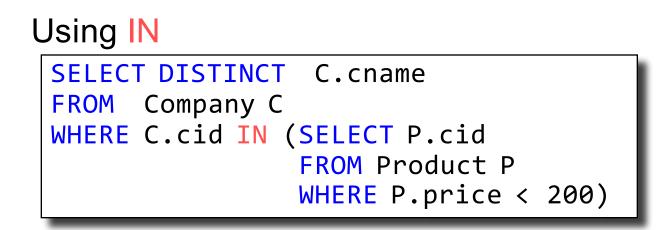
WHERE EXISTS (SELECT *

FROM Product P

WHERE C.cid = P.cid and P.price < 200)
```

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200



3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

```
Using ANY:

SELECT DISTINCT C.cname

FROM Company C

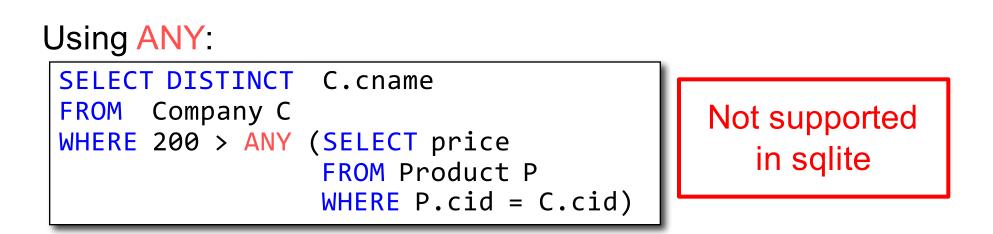
WHERE 200 > ANY (SELECT price

FROM Product P

WHERE P.cid = C.cid)
```

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200



3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Now let's unnest it:

SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200</pre>

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3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Now let's unnest it:

SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200</pre>

Existential quantifiers are easy! ③

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make <u>only</u> products with price < 200

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

same as:

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Universal quantifiers

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Find all companies s.t. <u>all</u> their products have price < 200

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Universal quantifiers

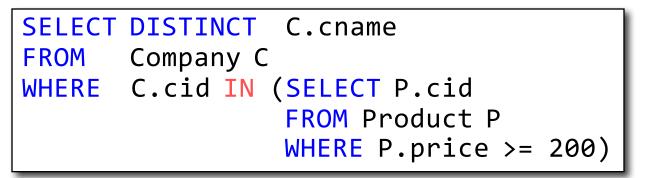
Universal quantifiers are hard! 🛞

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3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

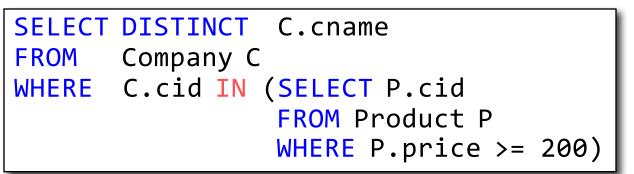
1. Find *the other* companies that make <u>some</u> product \geq 200



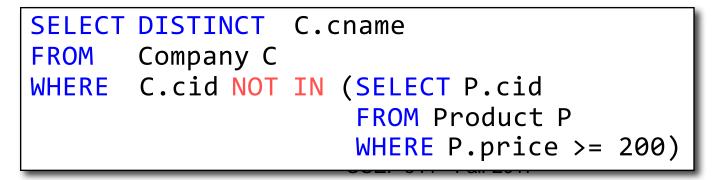
3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

1. Find *the other* companies that make <u>some</u> product \geq 200



2. Find all companies s.t. <u>all</u> their products have price < 200



3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using EXISTS:	
SELECT DISTINCT	C.cname
FROM Company C WHERE NOT EXISTS	(SELECT *
	FROM Product P
	WHERE P.cid = C.cid and P.price >= 200)

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

```
Using ALL:

SELECT DISTINCT C.cname

FROM Company C

WHERE 200 >= ALL (SELECT price

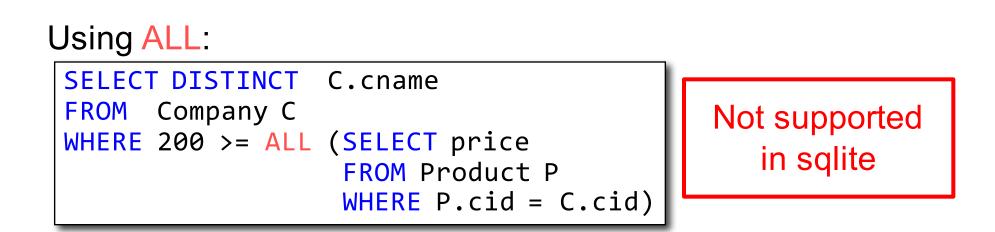
FROM Product P

WHERE P.cid = C.cid)
```

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers



Question for Database Theory Fans and their Friends

- Can we unnest the *universal quantifier* query?
- We need to first discuss the concept of *monotonicity*

Monotone Queries

- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Monotone Queries

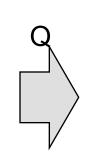
- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Product	

Company

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



pname	city
Gizmo	Lyon
Camera	Lodtz

Monotone Queries

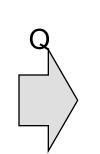
- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Product

Company

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003 Builder		Lodtz



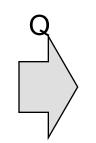
pname	city
Gizmo	Lyon
Camera	Lodtz

Product

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003
iPad	499.99	c001

Company

cname	city		
Sunworks	Bonn		
DB Inc.	Lyon		
Builder	Lodtz		
	Sunworks DB Inc.		



pname	city
Gizmo	Lyon
Camera	Lodtz
iPad	Lyon

Monotone Queries

• <u>Theorem</u>: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.

Monotone Queries

- <u>Theorem</u>: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.
- Proof. We use the nested loop semantics: if we insert a tuple in a relation R_i, this will not remove any tuples from the answer

SELECT
$$a_1$$
, a_2 , ..., a_k
FROM R_1 AS x_1 , R_2 AS x_2 , ..., R_n AS x_n
WHERE Conditions

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Monotone Queries

• The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

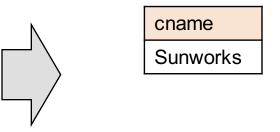
Monotone Queries

• The query:

Find all companies s.t. <u>all</u> their products have price < 200

is not monotone

	name city	cid	cid	price	pname
Gizmo 19.99 c001 c001 Sunworks Bo	Sunworks Bonn	c001	c001	19.99	Gizmo

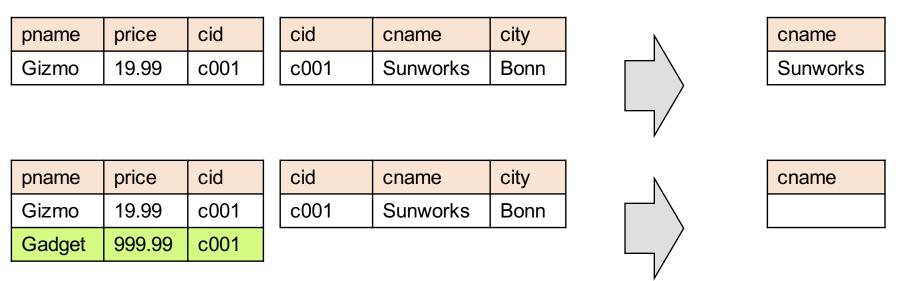


Monotone Queries

• The query:

Find all companies s.t. <u>all</u> their products have price < 200

is not monotone



 <u>Consequence</u>: If a query is not monotonic, then we cannot write it as a SELECT-FROM-WHERE query¹³⁰ without nested subqueries

Queries that must be nested

- Queries with universal quantifiers or with negation
- Queries that use aggregates in certain ways
 - sum(..) and count(*) are NOT monotone,
 because they do not satisfy set containment
 - select count(*) from R is not monotone!