Database Management Systems CSEP 544

Lecture 2: SQL

Announcements

- HW1 due tonight (11:59pm)
- PA2 & HW2 released
- Fill out HW3 email account form by tonight!
- Final information posted on piazza
- Check website for up to date OH info

Review

- Data models
 - Instance
 - Schema
 - Language
- Relational data model
 - Relations are flat
 - Tuples are not ordered
- Logical and physical data independence

Reading Assignment 1

Selections in SQL

```
SELECT *
FROM Product
WHERE price > 100.0
```

Projections in SQL

```
SELECT CName
FROM Product
```

Joins in SQL

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

```
Product(<u>pname</u>, price, category, manufacturer)
Company(<u>cname</u>, country)
```

Joins in SQL

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

Joins in SQL

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

cname	country
GizmoWorks/	USA
Canon	Japan
Hitachi	Japan

```
SELECT pname, price

FROM Product, Company

WHERE manufacturer=cname AND

country='Japan' AND price < 150

CSEP 544 - Fall 2017
```

Joins in SQL

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

```
FROM Product as P, Company as C
WHERE P.manufacturer=C.cname AND
C.country='Japan' AND C.price < 150
```

Joins in SQL

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

Retrieve all USA companies that manufacture "gadget" products

Joins in SQL /

pname	price	category	manufacturer
MultiTouch	199.99	gadget	Canon
SingleTouch	49.99	photography	Canon
Gizom	50	gadget	GizmoWorks
SuperGizmo	250.00	gadget	GizmoWorks

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

Why

DISTINCT?

Retrieve all USA companies that manufacture "gadget" products

SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname

Joins in SQL

- The standard join in SQL is sometimes called an inner join
 - Each row in the result must come from both tables in the join
- Sometimes we want to include rows from only one of the two table: outer join

Joins and Aggregates

```
Product(<u>pname</u>, price, category, manufacturer)
Company(<u>cname</u>, country)
-- manufacturer is foreign key to Company
```

```
SELECT DISTINCT cname
FROM Product, Company
WHERE country='USA' AND category = 'gadget'
          AND manufacturer = cname
```

```
SELECT DISTINCT cname
```

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
Camera	Photo	Hitachi
OneClick	Photo	Hitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

```
SELECT DISTINCT cname
```

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname

Product

_	pname	category	manufacturer	
	Gizmo	gadget	GizmoWorks	
	Camera	Photo	Hitachi	
	OneClick	Photo	Hitachi	

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname

Product

_	pname	category	manufacturer	
	Gizmo	gadget	GizmoWorks	
	Camera	Photo	Hitachi	
	OneClick	Photo	Hitachi	

	cname	country
G	izmoWorks	USA
	Canon	Japan
	Hitachi	Japan

```
SELECT DISTINCT cname
FROM Product, Company
WHERE country='USA' AND category = 'gadget'
AND manufacturer = cname
```

Product

pnamecategorymanufacturerGizmogadgetGizmoWorksCameraPhotoHitachiOneClickPhotoHitachi

Company

	cname	country
	GizmoWorks	USA
	Canon	Japan
<u> </u>	Hitachi	Japan

25

pname	category	manufacturer	cname	country
Gizmo	gadget	GizmoWorks	GizmoWorks	USA

```
SELECT DISTINCT cname
```

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname

Product

pnamecategorymanufacturerGizmogadgetGizmoWorksCameraPhotoHitachiOneClickPhotoHitachi

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

```
SELECT DISTINCT cname
```

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname

Product

_	pname	category	manufacturer	
	Gizmo	gadget	GizmoWorks	
	Camera	Photo	Hitachi	
	OneClick	Photo	Hitachi	

cname	country
GizmoWorks	USA
Canon	Japan
Hitachi	Japan

```
SELECT DISTINCT cname

FROM Product, Company

WHERE country='USA' AND category = 'gadget'

AND manufacturer = cname
```

```
FROM Product JOIN Company ON country = 'USA' AND category = 'gadget'
AND manufacturer = cname
```

where

```
SELECT x1.a1, x2.a2, ... xm.am

FROM R1 as x1, R2 as x2, ... Rm as xm
WHERE Cond
```

This is called nested loop semantics since we are interpreting what a join means using a nested loop

Another example

```
Product(pname, price, category, manufacturer)
Company(cname, country)
-- manufacturer is foreign key to Company
```

Retrieve all Japanese companies that manufacture products in both 'gadget' and 'photography' categories

Another example

```
Product(pname, price, category, manufacturer)
Company(cname, country)
-- manufacturer is foreign key to Company
```

Retrieve all Japanese companies that manufacture products in both 'gadget' and 'photography' categories

```
SELECT DISTINCT cname
FROM Product P1, Product P2, Company
WHERE country = 'Japan' AND P1.category = 'gadget'
        AND P2.category = 'photography'
        AND P1.manufacturer = cname
        AND P2.manufacturer = cname;
```

Self-Joins and Tuple Variables

- Find all companies that manufacture both products in the 'gadgets' and 'photo' category
- Joining Product with Company is insufficient: need to join Product, with Product, and with Company
- When a relation occurs twice in the FROM clause we call it a self-join
 - in that case we must use tuple variables (why?)

```
FROM Product x Product y Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

Product

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
SingleTouch	photo	Hitachi
MultiTouch	Photo	GizmoWorks

cname	country
GizmoWorks	USA
Hitachi	Japan

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

Product

V	
Л	

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
SingleTouch	photo	Hitachi
MultiTouch	Photo	GizmoWorks

cname	country
GizmoWorks	USA
Hitachi	Japan

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

Product

X	

\/
~/
y

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
SingleTouch	photo	Hitachi
MultiTouch	Photo	GizmoWorks

cname	country	
GizmoWorks	USA	
Hitachi	Japan	

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

	Flouuci		
X	pname	category	manufacturer
y <	Gizmo	gadget	GizmoWorks
•	SingleTouch	photo	Hitachi
	MultiTouch	Photo	GizmoWorks

Droduct

Company		_
cname	country	
GizmoWorks	USA	
Hitachi	Japan	

7

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

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IU	u	u	Оl

X

pname	category	manufacturer	
Gizmo	gadget	GizmoWorks	
SingleTouch	photo	Hitachi	

Company

cname	country
GizmoWorks	USA
Hitachi	Japan

Z

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

Product

X

pname	category	manufacturer
Gizmo	gadget	GizmoWorks
SingleTouch	photo	Hitachi
MultiTouch	Photo	GizmoWorks

Company

Z

cname	country
GizmoWorks	USA
Hitachi	Japan

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'

AND x.category = 'gadget'

AND y.category = 'photo'

AND x.manufacturer = cname

AND y.manufacturer = cname;
```

Droduct

7 chame

	Product						Company	,		_	•
X	pname	cate	gory	manu	facturer	1	cname		countr	У	
Y-5	Gizmo	gad	get	et GizmoWorks			GizmoWorks	S	USA		
<u> </u>	SingleTouc	ch pho	oto	o Hitachi		_	Hitachi		Japan		•
<u>y</u> ->	MultiTouc	h Pho	Photo GizmoWorks		o GizmoWorks						
x.pnam	ne x.category	x.manufacture	r y.pı	name	y.category	/	y.manufacturer		z.cname	Z.COU	ıntry
Gizmo	gadget	GizmoWorks	Mult	iTouch	Photo		GizmoWorks		GizmoWorks	US	Α

```
FROM Product x, Product y, Company z

WHERE z.country = 'USA'
AND x.category = 'gadget'
AND y.category = 'photo'
AND x.manufacturer = cname
AND y.manufacturer = cname;
```

Z **Product** Company X manufacturer category pname country cname Gizmo gadget **GizmoWorks GizmoWorks** USA SingleTouch photo Hitachi Japan Hitachi MultiTouch GizmoWorks **Photo**

x.pname	x.category	x.manufacturer	y.pname y.category		y.manufacturer	z.cname	z.country
Gizmo	gadget	GizmoWorks	MultiTouch	Photo	GizmoWorks	GizmoWorks	USA

Outer joins

```
Product(name, category)
Purchase(prodName, store)
```

-- prodName is foreign key

```
SELECT Product.name, Purchase.store
FROM Product, Purchase
WHERE Product.name = Purchase.prodName
```

We want to include products that are never sold, but some are not listed! Why?

Outer joins

```
Product(name, category)
Purchase(prodName, store)
```

-- prodName is foreign key

```
SELECT Product.name, Purchase.store
FROM Product LEFT OUTER JOIN Purchase ON
Product.name = Purchase.prodName
```


Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

SELECT Product.name, Purchase.store FROM Product JOIN Purchase ON

Product.name = Purchase.prodName

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Output

Name Store
Gizmo Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Output

Name Store
Gizmo Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz
Camera	Ritz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Product Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Product Purchase

Name	Category	
Gizmo	gadget	
Camera	Photo	_
OneClick	Photo	4

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

Name	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz
OneClick	NULL

Store

Wiz

Ritz

Wiz

NULL

Foo

Product

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

Name

Gizmo

Camera

Camera

OneClick

(NULI

Oı	ıtr	ı ıt
U	սւբ	ut

Purchase

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz
Phone	Foo

Outer Joins

tableA (LEFT/RIGHT/FULL) OUTER JOIN tableB ON p

- Left outer join:
 - Include tuples from tableA even if no match
- Right outer join:
 - Include tuples from tableB even if no match
- Full outer join:
 - Include tuples from both even if no match
- In all cases:
 - Patch tuples without matches using NULL

Comment about SQLite

- Cannot load NULL values such that they are actually loaded as null values
- So we need to use two steps:
 - Load null values using some type of special value
 - Update the special values to actual null values

```
update Purchase
  set price = null
  where price = 'null'
```

Simple Aggregations

Five basic aggregate operations in SQL

```
select count(*) from Purchase
select sum(quantity) from Purchase
select avg(price) from Purchase
select max(quantity) from Purchase
select min(quantity) from Purchase
```

Except count, all aggregations apply to a single attribute

Aggregates and NULL Values

```
Null values are not used in aggregates
   insert into Purchase
   values(12, 'gadget', NULL, NULL, 'april')
Try the following at home:
    select count(*) from Purchase
    select count(quantity) from Purchase
    select sum(quantity) from Purchase 
   select count(*)
    from Purchase
    where quantity is not null;
```

Counting Duplicates

COUNT applies to duplicates, unless otherwise stated:

```
SELECT count(product)
FROM Purchase
WHERE price > 4.99
```

same as count(*) if no nulls

We probably want:

```
SELECT count(DISTINCT product)
FROM Purchase
WHERE price > 4.99
```

More Examples

SELECT Sum(price * quantity) FROM Purchase

SELECT Sum(price * quantity)
FROM Purchase
WHERE product = 'bagel'

What do they mean?

Grouping and Query Evaluation

Grouping and Aggregation

Purchase(product, price, quantity)

Find total quantities for all sales over \$1, by product.

Grouping and Aggregation

Product	Price	Quantity
Bagel	3	20
Bagel	1.50	20
Banana	0.5	50
Banana	2	10
Banana	4	10

Product	TotalSales
Bagel	40
Banana	20

```
SELECT product, Sum(quantity) AS TotalSales
FROM Purchase
WHERE price > 1
GROUP BY product
64
```

Other Examples

Compare these two queries:

```
SELECT product, count(*)
FROM Purchase
GROUP BY product
```

GROUP BY product

```
SELECT month, count(*)
FROM Purchase
GROUP BY month
```

```
SELECT product,
    sum(quantity) AS SumQuantity,
    max(price) AS MaxPrice
FROM Purchase
```

What does it mean?

Need to be Careful...

```
SELECT product,

max(quantity)

FROM Purchase

GROUP BY product
```

SELECT product, quantity
FROM Purchase
GROUP BY product

Product	Price	Quantity
Bagel	3	20
Bagel	1.50	20
Banana	0.5	50
Banana	2	10
Banana	4	10

SELECT product, quantity
FROM Purchase
GROUP BY product

Need to be Careful...

SELECT product
FROM Purchase
GROUP BY product

SELECT quantity FROM Purchase

		Quantity
Product		20
Bagel	+	20

???

Banana

	Qualitity
	20
	20
,	50
	10
	10

Product	Price	Quantity
Bagel	3	20
Bagel	1.50	20
Banana	0.5	50
Banana	2	10
Banana	4	10

SELECT product, quantity

FROM Purchase GROUP BY product

Everything in SELECT must be either a GROUP-BY attribute, or an aggregate

Need to be Careful...

SELECT product,
 max(quantity)

FROM Purchase

GROUP BY product

SELECT product, quantity

FROM Purchase

GROUP BY product

Product	Price	Quantity	
Bagel	3	20	
Bagel	1.50	20	
Banana	0.5	50	
Banana	2	10	
Banana	4	10	

sqlite is WRONG on this query.

Advanced DBMS (e.g. SQL Server) gives an error

Grouping and Aggregation

Purchase(product, price, quantity)

Find total quantities for all sales over \$1, by product.

```
SELECT product, Sum(quantity) AS TotalSales
FROM Purchase
WHERE price > 1
GROUP BY product
```

How is this query processed?

Grouping and Aggregation

- 1. Compute the FROM and WHERE clauses.
- 2. Group by the attributes in the GROUPBY
- 3. Compute the SELECT clause: grouped attributes and aggregates.



1,2: From, Where

FWGS

Product	Price	Quantity	
Bagel	3	20	
Bagel	1.50	20	
Banana	0.5	50	
Banana	2	10	
Banana	4	10	

WHERE price > 1

SELECT product, Sum(quantity) AS TotalSales
FROM Purchase

WHERE price > 1

GROUP BY product

3,4. Grouping, Select

FWGS

Product	Price	Quantity			
Bagel	3	20		Product	TotalSales
Bagel	1.50	20		Bagel	40
Banana	0.5	50		Banana	20
Banana	2	10		18	
Banana	4	10			

```
product, Sum(quantity) AS TotalSales
SELECT
     Purchase
FROM
       price > 1
WHERE
GROUP BY product
                                          72
```

Ordering Results

```
SELECT product, sum(price*quantity) as rev FROM Purchase GROUP BY product ORDER BY rev desc
```

FWGOS

TM

Note: some SQL engines want you to say ORDER BY sum(price*quantity) desc

HAVING Clause

Same query as before, except that we consider only products that had at least 30 sales.

```
SELECT product, sum(price*quantity)
FROM Purchase
WHERE price > 1
GROUP BY product
HAVING sum(quantity) > 30
```

HAVING clause contains conditions on aggregates.

General form of Grouping and Aggregation

SELECT S

FROM $R_1, ..., R_n$

WHERE C1

GROUP BY $a_1,...,a_k$

HAVING C2

S = may contain attributes a₁,...,a_k and/or any aggregates but NO OTHER ATTRIBUTES

C1 = is any condition on the attributes in $R_1,...,R_n$

C2 = is any condition on aggregate expressions and on attributes $a_1,...,a_k$

Why?

Semantics of SQL With Group-By

SELECT S
FROM $R_1,...,R_n$ WHERE C1
GROUP BY $a_1,...,a_k$ HAVING C2

FWGHOS

Evaluation steps:

- 1. Evaluate FROM-WHERE using Nested Loop Semantics
- 2. Group by the attributes $a_1, ..., a_k$
- 3. Apply condition C2 to each group (may have aggregates)
- 4. Compute aggregates in S and return the result

Exercise

Compute the total income per month
Show only months with less than 10 items sold
Order by quantity sold and display as "TotalSold"

Exercise

Compute the total income per month
Show only months with less than 10 items sold
Order by quantity sold and display as "TotalSold"

FROM Purchase

Exercise

Compute the total income per month

Show only months with less than 10 items sold

Order by quantity sold and display as "TotalSold"

FROM Purchase GROUP BY month

Exercise

Compute the total income per month
Show only months with less than 10 items sold
Order by quantity sold and display as "TotalSold"

```
FROM Purchase
GROUP BY month
HAVING sum(quantity) < 10
```

Purchase(pid, product, price, quantity, month)

Exercise

Compute the total income per month
Show only months with less than 10 items sold
Order by quantity sold and display as "TotalSold"

Purchase(pid, product, price, quantity, month)

Exercise

Compute the total income per month
Show only months with less than 10 items sold
Order by quantity sold and display as "TotalSold"

```
SELECT month, sum(price*quantity),
sum(quantity) as TotalSold
FROM Purchase
GROUP BY month
HAVING sum(quantity) < 10
ORDER BY sum(quantity)
```

WHERE vs HAVING

- WHERE condition is applied to individual rows
 - The rows may or may not contribute to the aggregate
 - No aggregates allowed here
- HAVING condition is applied to the entire group
 - Entire group is returned, or not at all
 - May use aggregate functions in the group

Purchase(pid, product, price, quantity, month)

Mystery Query

What do they compute?

```
SELECT month, sum(quantity), max(price)
FROM Purchase
GROUP BY month
```

```
SELECT month, sum(quantity)
FROM Purchase
GROUP BY month
```

SELECT month
FROM Purchase
GROUP BY month

Purchase(pid, product, price, quantity, month)

Mystery Query

What do they compute?

```
SELECT month, sum(quantity), max(price)
FROM Purchase
GROUP BY month
```

SELECT month, sum(quantity)
FROM Purchase
GROUP BY month

SELECT month
FROM Purchase
GROUP BY month

Lesson:
DISTINCT is
a special case
of GROUP BY

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: price is in Purchase, manufacturer is in Product...

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: price is in Purchase, manufacturer is in Product...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
```

manu facturer	 price	
Hitachi	150	
Canon	300	
Hitachi	180	

Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: price is in Purchase, manufacturer is in Product...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
  and y.price > 100
```

manu facturer	 price	
Hitachi	150	
Canon	300	
Hitachi	180	

```
-- step 2: do the group-by on the join
SELECT x.manufacturer, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
  and y.price > 100
GROUP BY x.manufacturer
```

manu facturer	count(*)	
Hitachi	2	
Canon	1	

Aggregate + Join

Variant:

For each manufacturer, compute how many products with price > \$100 they sold in each month

```
SELECT x.manufacturer, y.month, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
  and y.price > 100
GROUP BY x.manufacturer, y.month
```

manu facturer	month	count(*)
Hitachi	Jan	2
Hitachi	Feb	1
Canon	Jan	3



Including Empty Groups

In the result of a group by query, there
is one row per group in the result

Count(*) is never 0

```
SELECT x.manufacturer, count(*)
FROM Product x, Purchase y
WHERE x.pname = y.product
GROUP BY x.manufacturer
```

Including Empty Groups

```
SELECT x.manufacturer, count(y.pid)

FROM Product x LEFT OUTER JOIN Purchase y

ON x.pname = y.product

GROUP BY x.manufacturer

Count(pid) is 0

when all pid's in
the group are
NULL
```

Nested Queries

What have we learned so far

- Data models
- Relational data model
 - Instance: relations
 - Schema: table with attribute names
 - Language: SQL

What have we learned so far

- SQL features
 - Projections
 - Selections
 - Joins (inner and outer)
 - Aggregates
 - Group by
 - Inserts, updates, and deletes
- Make sure you read the textbook!

Subqueries

- A subquery is a SQL query nested inside a larger query
- Such inner-outer queries are called nested queries
- A subquery may occur in:
 - A SELECT clause
 - A FROM clause
 - A WHERE clause
- Rule of thumb: avoid writing nested queries when possible
 - But sometimes it's impossible, as we will see

Subqueries...

- Can appear as computed values in a SELECT clause
- Can appear in FROM clauses and aliased using a tuple variable that represents the tuples in the result of the subquery
- Can return a single constant to be compared with another value in a WHERE clause
- Can return relations to be used in WHERE clauses

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

For each product return the city where it is manufactured

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cid=X.cid) as City
```

What happens if the subquery returns more than one city? We get a runtime error

(and SQLite simply ignores the extra values...)

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Whenever possible, don't use nested queries:

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cid=X.cid) as City
FROM Product X
```



```
SELECT X.pname, Y.city
FROM Product X, Company Y
WHERE X.cid=Y.cid
```

We have "unnested" the query

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Compute the number of products made by each company

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Compute the number of products made by each company

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

Better: we can unnest using a GROUP BY

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

No! Different results if a company has no products

```
SELECT C.cname, count(pname)
FROM Company C LEFT OUTER JOIN Product P
ON C.cid=P.cid
GROUP BY C.cname
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

2. Subqueries in FROM

Find all products whose prices is > 20 and < 500

```
SELECT X.pname
FROM (SELECT *
FROM Product AS Y
WHERE price > 20) as X
WHERE X.price < 500
```

Side note: This is not a correlated subquery. (why?)

Try unnest this query!

2. Subqueries in FROM

At the end of the lecture we will see that sometimes we really need a subquery and one option will be to put it in the FROM clause.

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using EXISTS:

```
SELECT DISTINCT C.cname

FROM Company C
WHERE EXISTS (SELECT*)
FROM Product P
WHERE C.cid = P.cid and P.price < 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using IN

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price < 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using ANY:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using ANY:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

Not supported in sqlite

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Now let's unnest it:

```
SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200</pre>
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Now let's unnest it:

```
SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200
```

Existential quantifiers are easy!

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make <u>only</u> products with price < 200

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

Universal quantifiers

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

Universal quantifiers

Universal quantifiers are hard!

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

1. Find *the other* companies that make <u>some</u> product ≥ 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. all their products have price < 200

1. Find *the other* companies that make <u>some</u> product ≥ 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

2. Find all companies s.t. <u>all</u> their products have price < 200

```
SELECT DISTINCT C.cname

FROM Company C

WHERE C.cid NOT IN (SELECT P.cid

FROM Product P

WHERE P.price >= 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using EXISTS:

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using ALL:

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using ALL:

Not supported in sqlite

Question for Database Theory Fans and their Friends

 Can we unnest the universal quantifier query?

 We need to first discuss the concept of monotonicity

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)

Monotone Queries

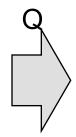
- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Product

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



pname	city
Gizmo	Lyon
Camera	Lodtz

Product (pname, price, cid)
Company (cid, cname, city)

Monotone Queries

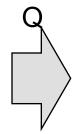
- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Product

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



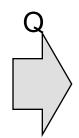
pname	city
Gizmo	Lyon
Camera	Lodtz

Product

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003
iPad	499.99	c001

Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



pname	city
Gizmo	Lyon
Camera	Lodtz
iPad	Lyon

• Theorem: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.

- Theorem: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.
- Proof. We use the nested loop semantics: if we insert a tuple in a relation R_i, this will not remove any tuples from the answer

```
for x<sub>1</sub> in R<sub>1</sub> do
  for x<sub>2</sub> in R<sub>2</sub> do
  ...
  for x<sub>n</sub> in R<sub>n</sub> do
    if Conditions
    output (a<sub>1</sub>,...,a<sub>k</sub>)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)

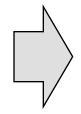
Monotone Queries

The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

pname	price	cid
Gizmo	19.99	c001

cid	cname	city
c001	Sunworks	Bonn



cname Sunworks Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)

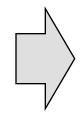
Monotone Queries

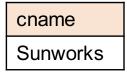
The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

pname	price	cid
Gizmo	19.99	c001

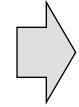
cid	cname	city
c001	Sunworks	Bonn

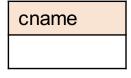




pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c001

cid	cname	city
c001	Sunworks	Bonn





 Consequence: If a query is not monotonic, then we cannot write it as a SELECT-FROM-WHERE query 130 without nested subqueries

Queries that must be nested

- Queries with universal quantifiers or with negation
- Queries that use aggregates in certain ways
 - sum(..) and count(*) are NOT monotone,
 because they do not satisfy set containment
 - select count(*) from R is not monotone!