Chapel: Heat Transfer (+ X10/Fortress)

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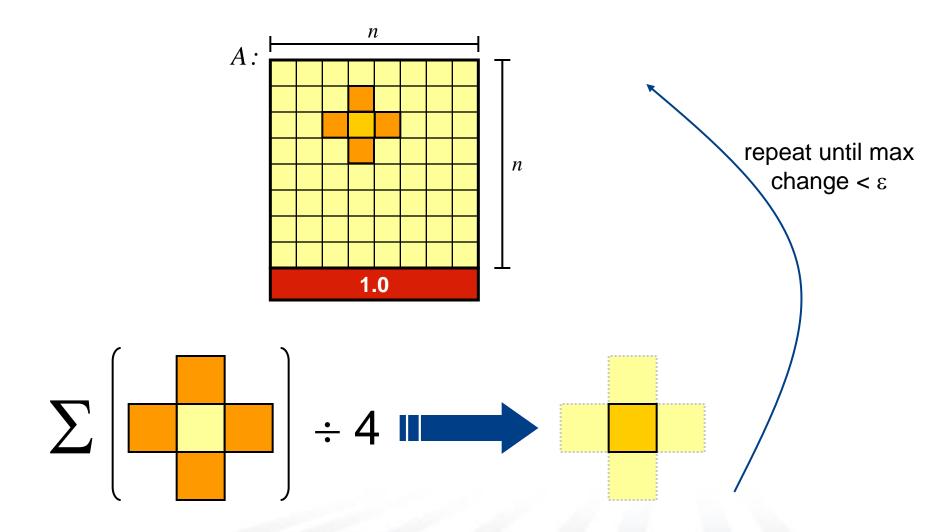








Heat Transfer in Pictures





```
config const n = 6,
              epsilon = 1.0e-5;
const BiqD: domain(2) = [0..n+1, 0..n+1],
         D: subdomain (BiqD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do {
  [(i,j) \text{ in } D] \text{ Temp}(i,j) = (A(i-1,j) + A(i+1,j))
                            + A(i,j-1) + A(i,j+1)) / 4;
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```



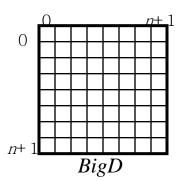
```
config const n = 6,
               epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1],
          D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
var A, Temp : [BiqD] real;
                     Declare program parameters
A[Las
       const ⇒ can't change values after initialization
do {
       config ⇒ can be set on executable command-line
                prompt> jacobi --n=10000 --epsilon=0.0001
  con
      note that no types are given; inferred from initializer
                n \Rightarrow integer (current default, 32 bits)
 whi
                epsilon ⇒ floating-point (current default, 64 bits)
writeIn(A);
```

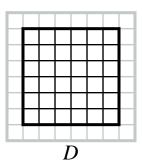


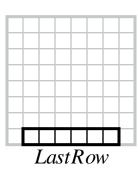
Declare domains (first class index sets)

domain(2) ⇒ 2D arithmetic domain, indices are integer 2-tuples

subdomain(P**)** \Rightarrow a domain of the same type as P whose indices are guaranteed to be a subset of P's







exterior ⇒ one of several built-in domain generators

4;



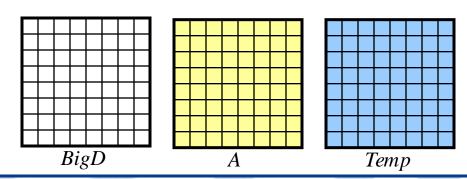
Declare arrays

var ⇒ can be modified throughout its lifetime

: $T \Rightarrow$ declares variable to be of type T

: **[D]** $T \Rightarrow$ array of size D with elements of type T

(no initializer) ⇒ values initialized to default value (0.0 for reals)

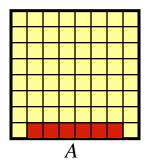


4;



Set Explicit Boundary Condition

indexing by domain ⇒ slicing mechanism array expressions ⇒ parallel evaluation



4



Compute 5-point stencil

 $[(i,j) \text{ in } D] \Rightarrow \text{ parallel forall expression over } D$'s indices, binding them to new variables i and j

Note: since $(i,j) \in D$ and $D \subseteq BigD$ and Temp: [BigD] \Rightarrow no bounds check required for Temp(i,j) with compiler analysis, same can be proven for A's accesses





Compute maximum change

op reduce ⇒ collapse aggregate expression to scalar using *op*

Promotion: abs() and – are scalar operators, automatically promoted to work with array operands



```
config const n = 6,
              epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1],
          D: subdomain (BiqD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
         Copy data back & Repeat until done
var
A [La uses slicing and whole array assignment
     standard do...while loop construct
do
  [(i,j) \text{ in } D] \text{ Temp}(i,j) = (A(i-1,j) + A(i+1,j))
                             + A(i,j-1) + A(i,j+1)
  const delta = max reduce abs(A[D]         Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```



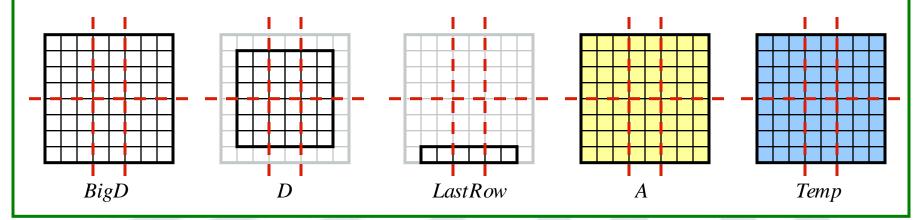
```
config const n = 6,
              epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1],
         D: subdomain (BiqD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do
               Write array to console
     If written to a file, parallel I/O would be used
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```



With this change, same code runs in a distributed manner

Domain distribution maps indices to *locales*

⇒ decomposition of arrays & default location of iterations over locales Subdomains inherit parent domain's distribution





```
config const n = 6,
              epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1] dmapped Block,
         D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do {
  [(i,j) \text{ in } D] \text{ Temp}(i,j) = (A(i-1,j) + A(i+1,j))
                            + A(i,j-1) + A(i,j+1)) / 4;
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```

Heat Transfer in Chapel (Variations)









Heat Transfer in Chapel (double buffered version)

```
config const n = 6,
             epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1] dmapped Block,
         D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
var A : [1...2] [BigD] real;
A[..][LastRow] = 1.0;
var src = 1, dst = 2;
do {
  [(i,j) \text{ in } D] A(dst)(i,j) = (A(src)(i-1,j) + A(src)(i+1,j)
                             + A(src)(i,j-1) + A(src)(i,j+1)) / 4;
  const delta = max reduce abs(A[src] - A[dst]);
  src <=> dst;
} while (delta > epsilon);
writeln(A);
```



Heat Transfer in Chapel (ZPL style)

```
config const n = 6,
             epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1] dmapped Block,
         D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
const north = (-1,0), south = (1,0), east = (0,1), west = (0,-1);
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do {
  [ind in D] Temp(ind) = (A(ind + north) + A(ind + south)
                        + A(ind + east) + A(ind + west)) / 4;
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```



Heat Transfer in Chapel (array of offsets version)

```
config const n = 6,
             epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1] dmapped Block,
         D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
param offset : [1..4] (int, int) = ((-1,0), (1,0), (0,1), (0,-1));
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do {
  [ind in D] Temp(ind) = (+ reduce [off in offset] A(ind + off))
                        / offset.numElements;
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```



Heat Transfer in Chapel (sparse offsets version)

```
config const n = 6,
             epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1] dmapped Block,
         D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
param stencilSpace: domain(2) = [-1..1, -1..1],
      offSet: sparse subdomain(stencilSpace)
             = ((-1,0), (1,0), (0,1), (0,-1));
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do {
  [ind in D] Temp(ind) = (+ reduce [off in offSet] A(ind + off))
                        / offSet.numIndices;
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```

The Other HPCS Languages









X10 in a Nutshell

- Heavily influenced by Java, Scala
 - emphasis on type safety, OOP design, small core language
 - also ZPL: support for global-view domains and arrays
- Similar concepts to what you've heard about today in Chapel
 - yet a fairly different syntax and design aesthetic
- Main differences from Chapel
- For more information:
 - http://x10-lang.org/
 - http://sf.net/projects/x10
 - http://dist.codehaus.org/
 - http://dist.codehaus.org/x10/documentation/presentations/UWMay2010.pdf



X10: Similarities to Chapel

- PGAS memory model
 - plus, language concepts for referring to realms of locality
- more dynamic ("post-SPMD") execution model
 - one logical task executes main()
 - any task can create additional tasks--local or remote
- global-view data structures
 - ability to declare and access distributed arrays holistically rather than piecemeal
- many similar concepts, often with different names/semantics
 - tasks vs. tasks
 - places vs. locales
 - 'at' vs. 'on'
 - 'ateach' vs' 'coforall' + 'on'
 - 'async' vs. 'begin'
 - 'finish' vs. 'sync'



X10: Differences from Chapel

- **X10**:
 - takes a purer object-oriented approach
 - for example, arrays have reference rather than value semantics

```
A = B; // alias or copy if A and B are arrays?
```

- based on Java/Scala rather than ab initio
 - reflects IBM's customer base relative to Cray's
- a bit more minimalist and purer
 - e.g., less likely to add abstractions to the language if expressible using objects
- semantics distinguish between local and remote more strongly
 - e.g., communication is more visible in the code
 - e.g., some operations are not legal on remote objects
 - reflect differing choices on orthogonality vs. performance/safety
- has a stronger story for exceptions



Heat Transfer in X10

```
class HeatTransfer v2 {
  const BiqD = Dist.makeBlock([0..n+1, 0..n+1], 0);
  const D = BigD \mid ([1..n, 1..n] \text{ as Region});
  const LR = [0..0, 1..n] as Region;
  const A = DistArray.make[double] (BigD, (p:Point) => { LR.contains(p) ? 1 : 0 });
  const Temp = DistArray.make[double] (BiqD);
  static def stencil 1((x,y):Point(2)) {
    return ((at(A.dist(x-1,y)) A(x-1,y)) +
            (at(A.dist(x+1,y))) A(x+1,y)) +
             A(x,y-1) + A(x,y+1) / 4;
 def run() {
    val D Base = Dist.makeUnique(D.places());
    var delta:double = 1.0;
    do {
      finish ateach (z in D Base)
      for (p:Point(2) in D | here)
      Temp(p) = stencil 1(p);
      delta = A.lift(Temp, D.region, (x:double, y:double)
            =>Math.abs(x-y)).reduce(Math.max.(Double, Double), 0);
      finish ateach (p in D) A(p) = Temp(p);
    } while (delta > epsilon);
```



```
config const n = 6,
              epsilon = 1.0e-5;
const BigD: domain(2) = [0..n+1, 0..n+1] dmapped Block,
         D: subdomain (BigD) = [1..n, 1..n],
   LastRow: subdomain(BigD) = D.exterior(1,0);
var A, Temp : [BiqD] real;
A[LastRow] = 1.0;
do {
  [(i,j) \text{ in } D] \text{ Temp}(i,j) = (A(i-1,j) + A(i+1,j))
                            + A(i,j-1) + A(i,j+1)) / 4;
  const delta = max reduce abs(A[D] - Temp[D]);
  A[D] = Temp[D];
} while (delta > epsilon);
writeln(A);
```



Fortress in a Nutshell

- The most blue-sky, clean-slate of the HPCS languages
- Goal: define language semantics in libraries, not compiler:
 - data structures and types (including scalars types?)
 - operators, typecasts
 - operator precedence
 - in short, as much as possible to support future changes, languages

Other themes:

- implicitly parallel -- most things are parallel by default
- supports mathematical notation, symbols, operators
- functional semantics
- hierarchical representation of target architecture's structure
- units of measurement in the type system (meters, seconds, miles, ...)

For more information:

- http://research.sun.com/projects/plrg/
- http://projectfortress.sun.com/Projects/Community/