

Chapel: Features

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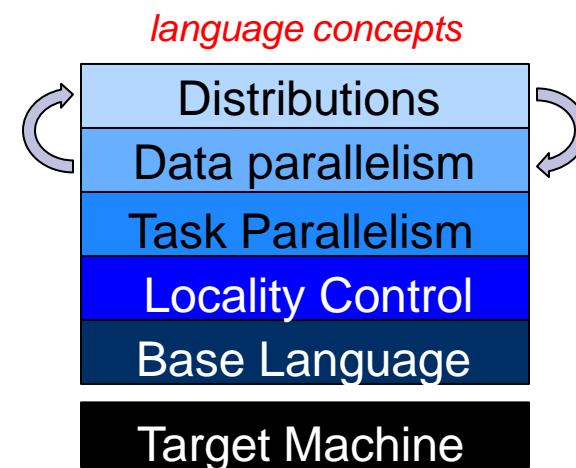
CSEP 524
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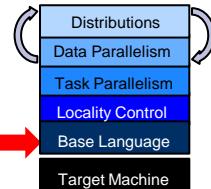


Outline

➤ Language Overview

- Base Language
- ❑ Task Parallelism
- ❑ Data Parallelism
- ❑ Locality
- ❑ Distributions





Base Language: Design

- Block-structured, imperative programming
- Intentionally not an extension to an existing language
- Instead, select attractive features from others:

ZPL, HPF: data parallelism, index sets, distributed arrays
(see also APL, NESL, Fortran90)

Cray MTA C/Fortran: task parallelism, lightweight synchronization

CLU: iterators (see also Ruby, Python, C#)

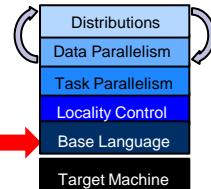
ML: latent types (see also Scala, Matlab, Perl, Python, C#)

Java, C#: OOP, type safety

C++: generic programming/templates (without adopting its syntax)

C, Modula, Ada: syntax

- Follow lead of C family of languages when useful
(C, Java, C#, Perl, ...)



Base Language: My Favorite Features

▪ Rich compile-time language

- parameter values (compile-time constants)
- folded conditionals, unrolled for loops, tuple expansions
- type and parameter functions – evaluated at compile-time

▪ Latent types

- ability to omit type specifications for convenience or code reuse
- type specifications can be omitted from...
 - ...variables (inferred from initializers)
 - ...class members (inferred from constructors)
 - ...function arguments (inferred from callsite)
 - ...function return types (inferred from return statements)

▪ Configuration variables (and parameters)

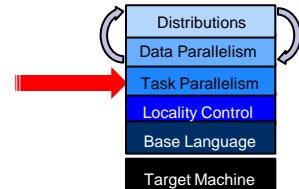
```
config const n = 100; // override with ./a.out --n=100000
```

▪ Tuples

▪ Iterators (in the CLU, Ruby sense, not C++/Java-style)

▪ Declaration Syntax: more like Pascal/Modula/Scala than C

Task Parallelism: Task Creation



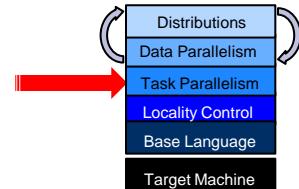
begin: creates a task for future evaluation

```
begin DoThisTask();  
WhileContinuing();  
TheOriginalThread();
```

sync: waits on all begins created within its dynamic scope

```
sync {  
    begin treeSearch(root);  
}  
  
def treeSearch(node) {  
    if node == nil then return;  
    begin treeSearch(node.right);  
    begin treeSearch(node.left);  
}
```

Task Parallelism: Structured Tasks



cobegin: creates a task per component statement:

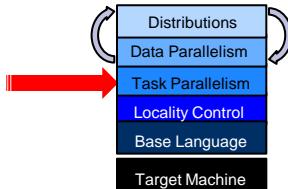
```
computePivot(lo, hi, data);  
cobegin {  
    Quicksort(lo, pivot, data);  
    Quicksort(pivot, hi, data);  
} // implicit join here
```

```
cobegin {  
    computeTaskA(...);  
    computeTaskB(...);  
    computeTaskC(...);  
} // implicit join
```

coforall: creates a task per loop iteration

```
coforall e in Edges {  
    exploreEdge(e);  
} // implicit join here
```

Task Parallelism: Task Coordination



sync variables: store full/empty state along with value

```
var result$: sync real;      // result is initially empty
sync {
    begin ... = result$;        // block until full, leave empty
    begin result$ = ...;       // block until empty, leave full
}
result$.readXX();           // read value, leave state unchanged;
                           // other variations also supported
```

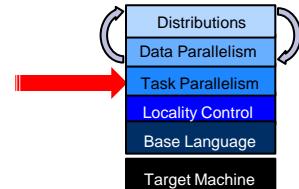
single-assignment variables: writeable once only

```
var result$: single real = begin f(); // result initially empty
...                         // do some other things
total += result$;          // block until f() has completed
```

atomic sections: support transactions against memory

```
atomic {
    newnode.next = insertpt;
    newnode.prev = insertpt.prev;
    insertpt.prev.next = newnode;
    insertpt.prev = newnode;
}
```

Producer/Consumer example



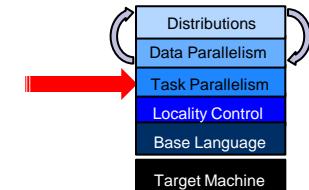
```
var buff$: [0..buffersize-1] sync int;
```

```
cobegin {
    producer();
    consumer();
}
```

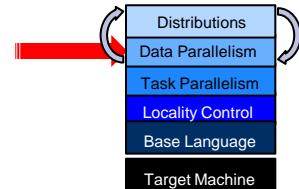
```
def producer() {
    var i = 0;
    for ... {
        i = (i+1) % buffersize;
        buff$(i) = ...;
    }
}
```

```
def consumer() {
    var i = 0;
    while {
        i = (i+1) % buffersize;
        ...buff$(i)...;
    }
}
```

QuickSort in Chapel



```
def quickSort(arr: [],
              thresh: int,
              low: int = arr.domain.low,
              high: int = arr.domain.high) {
    if high - low < 8 {
        bubbleSort(arr, low, high);
    } else {
        const pivotVal = findPivot(arr, low, high);
        const pivotLoc = partition(arr, low, high, pivotVal);
        serial thresh <= 0 do cobegin {
            quickSort(arr, thresh-1, low, pivotLoc-1);
            quickSort(arr, thresh-1, pivotLoc+1, high);
        }
    }
}
```

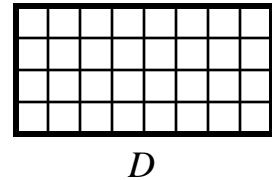


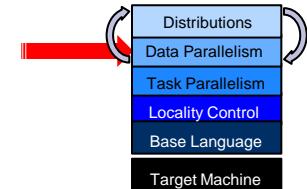
Data Parallelism: Domains

domain: a first-class index set

```
var m = 4, n = 8;
```

```
var D: domain(2) = [1..m, 1..n];
```

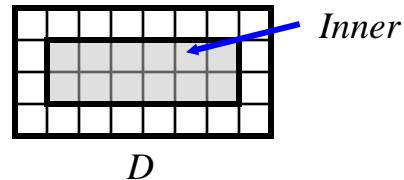




Data Parallelism: Domains

domain: a first-class index set

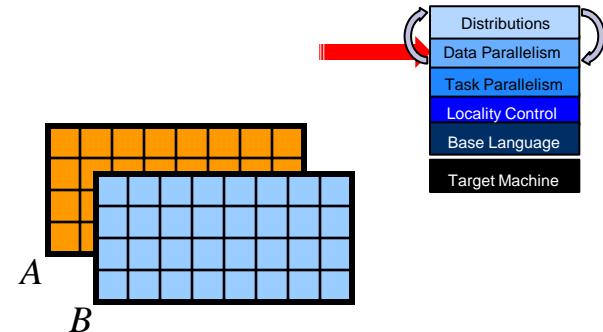
```
var m = 4, n = 8;  
  
var D: domain(2) = [1..m, 1..n];  
var Inner: subdomain(D) = [2..m-1, 2..n-1];
```



Domains: Some Uses

- Declaring arrays:

```
var A, B: [D] real;
```

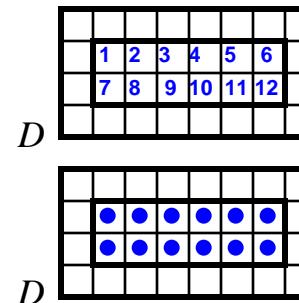


- Iteration (sequential or parallel):

```
for ij in Inner { ... }
```

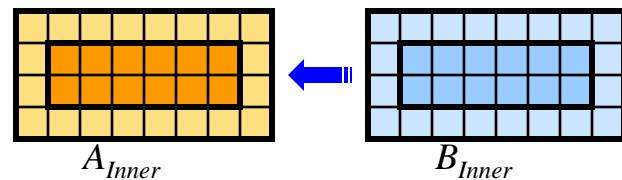
or: **forall** ij **in** Inner { ... }

or: ...



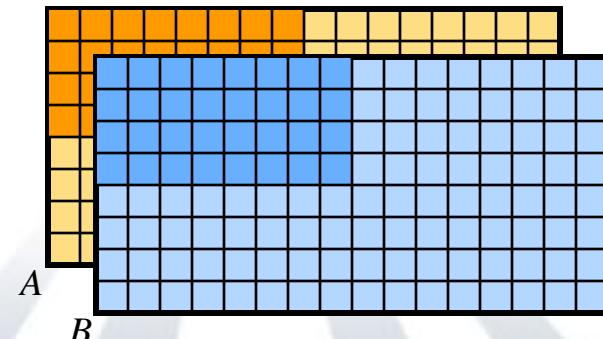
- Array Slicing:

```
A[Inner] = B[Inner];
```



- Array reallocation:

```
D = [1..2*m, 1..2*n];
```



Forall vs. For vs. Coforall

for loops:

- Use the current task to execute the loop serially

coforall loops:

- Execute the loop using a distinct task per iteration
- Can have synchronization between iterations

forall loops:

- Use some number of tasks between these two extremes
- Must be legally executable by a single task
- How many tasks are used in practice?

Data Parallelism Throttles

--dataParTasksPerLocale=#

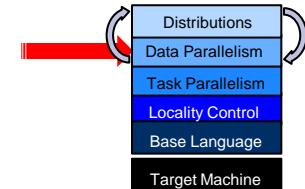
- Specify # of tasks to execute forall loops
- Default: number of cores (*in current implementation*)

--dataParIgnoreRunningTasks=[true|false]

- If false, reduce # of forall tasks by # of running tasks
- Default: true (*in current implementation*)

--dataParMinGranularity=#

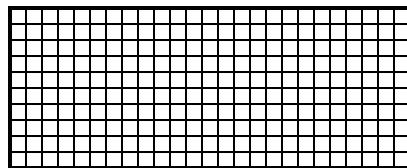
- reduce # of tasks if any task has fewer iterations
- Default: 1 (*in current implementation*)



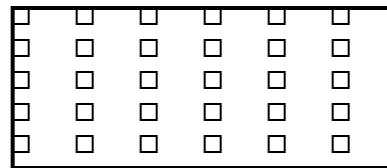
Data Parallelism: Domain Types

Chapel supports several domain types...

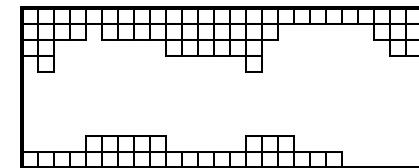
```
var OceanSpace = [0..#lat, 0..#long],  
    AirSpace = OceanSpace by (2,4),  
    IceSpace: sparse subdomain(OceanSpace) = genCaps();
```



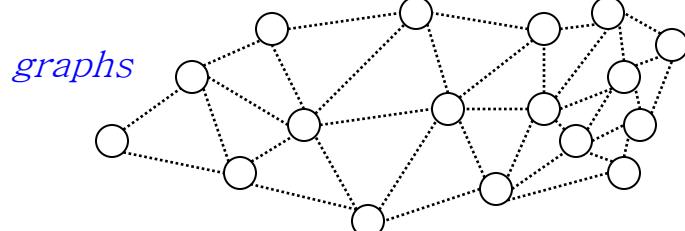
dense



strided

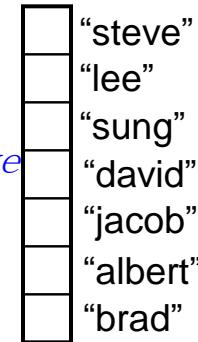


sparse

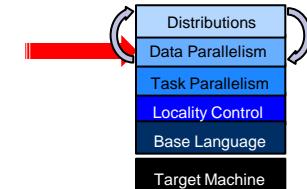


graphs

associative



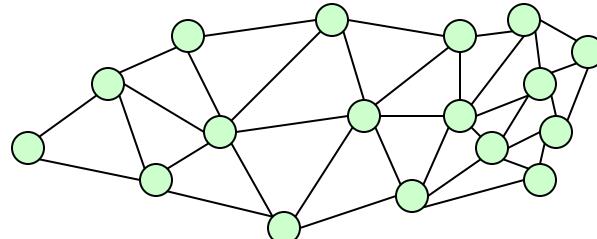
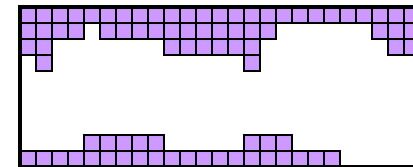
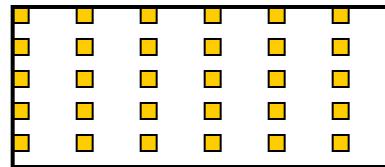
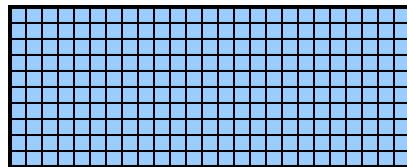
```
var Vertices: domain(opaque) = ..., People: domain(string) = ...;
```



Data Parallelism: Domain Uses

All domain types can be used to declare arrays...

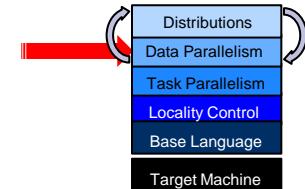
```
var Ocean: [OceanSpace] real,  
Air: [AirSpace] real,  
IceCaps[IceSpace] real;
```



	"steve"
	"lee"
	"sung"
	"david"
	"jacob"
	"albert"
	"brad"

```
var Weight: [Vertices] real,
```

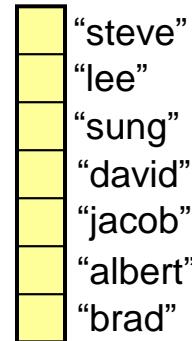
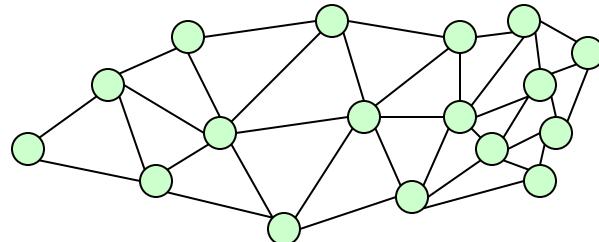
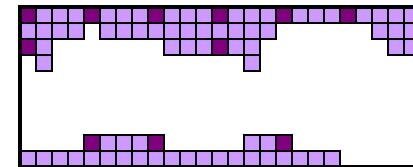
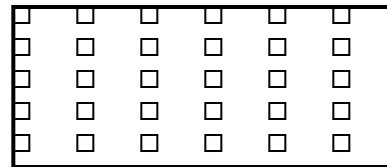
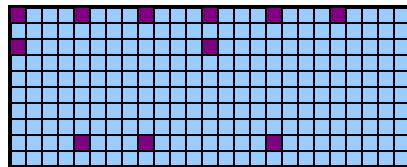
```
Age: [People] int;
```



Data Parallelism: Domain Uses

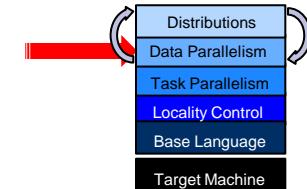
...to iterate over index sets...

```
forall ij in AirSpace do
    Ocean(ij) += IceCaps(ij);
```



```
forall v in Vertices do
    Weight(v) = numEdges(v);
```

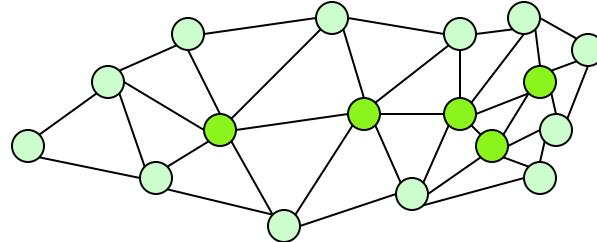
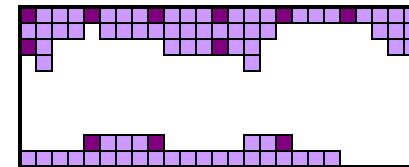
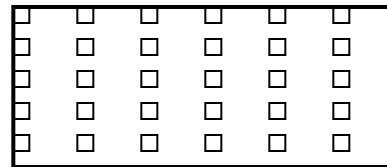
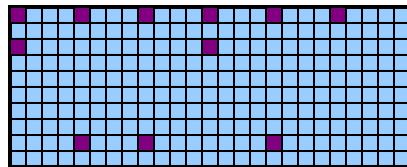
```
forall p in People do
    Age(p) += 1;
```



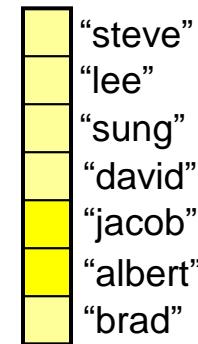
Data Parallelism: Domain Uses

...to slice arrays...

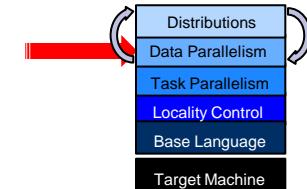
```
Ocean[AirSpace] += IceCaps[AirSpace];
```



...Vertices[Interior]...



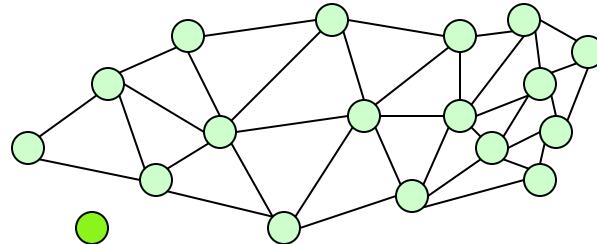
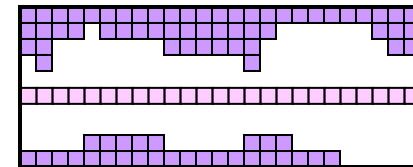
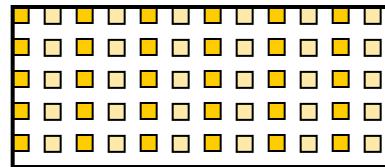
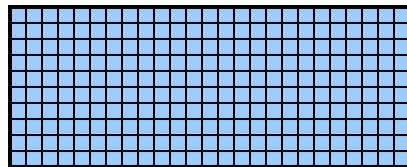
...People[Interns]...



Data Parallelism: Domain Uses

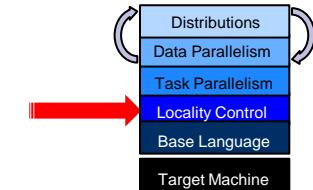
...and to reallocate arrays

```
AirSpace = OceanSpace by (2,2);  
IceSpace += genEquator();
```



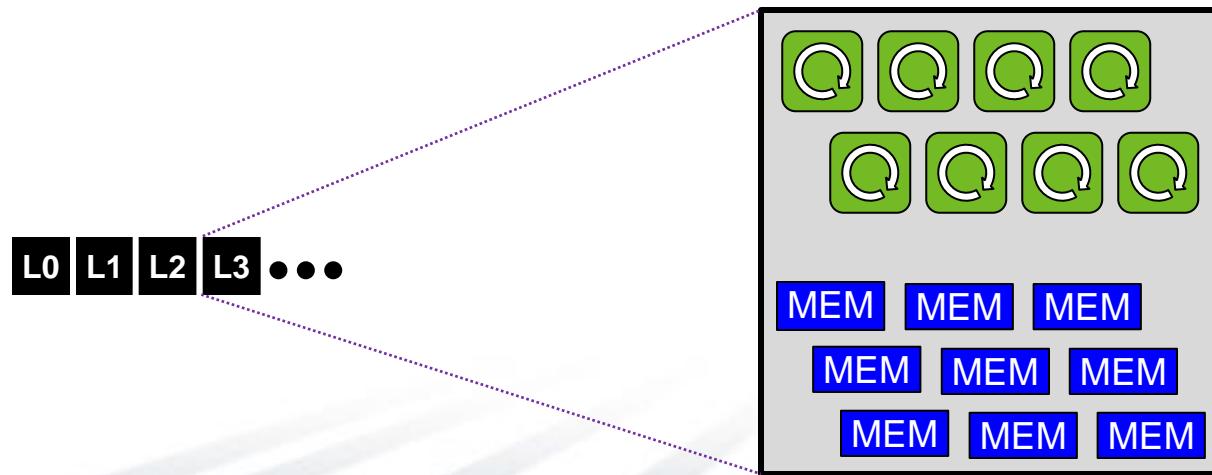
```
newnode = Vertices.create(); People += "srini";
```

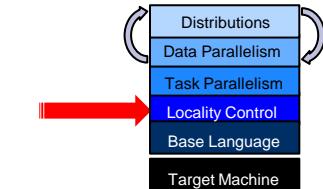
Locality: Locales



locale: An abstract unit of the target architecture

- supports reasoning about locality
- has capacity for processing and storage
- two threads in a given locale have similar access to a given address
 - addresses in that locale are ~uniformly accessible
 - addresses in other locales are also accessible, but at a price
- locales are defined for a given architecture by a Chapel compiler
 - e.g., a multicore processor or SMP node could be a locale





Locales and Program Startup

- Chapel users specify # locales on executable command-line

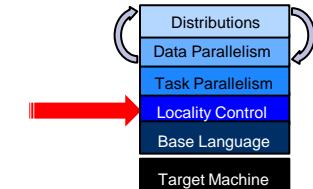
```
prompt> myChapelProg -n1=8      # run using 8 locales
```



- Chapel launcher bootstraps program execution:

- obtains necessary machine resources
 - e.g., requests 8 nodes from the job scheduler
- loads a copy of the executable onto the machine resources
- starts running the program. *Conceptually...*
 - ...locale #0 starts running program's entry point (`main()`)
 - ...other locales wait for work to arrive

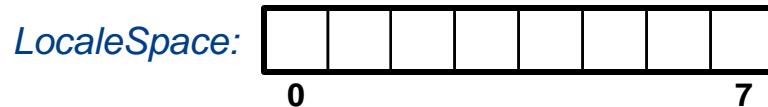
Locale Variables



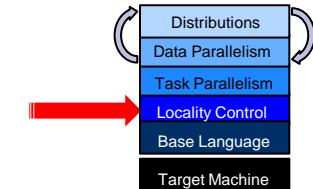
Built-in variables represent a program's locale set:

```
config const numLocales: int;           // number of locales
const LocaleSpace = [0..numLocales-1],   // locale indices
    Locales: [LocaleSpace] locale;        // locale values
```

numLocales: 8



Locale Views



Using standard array operations, users can create their own locale views:

```
var TaskALocs = Locales[..numTaskALocs];
```

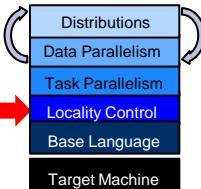


```
var TaskBLocs = Locales[numTaskALocs+1..];
```



```
var CompGrid = Locales.reshape([1..gridRows,  
                                1..gridCols]);
```





Locale Methods

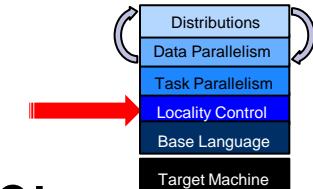
- The locale type supports built-in methods:

```
def locale.id: int;           // index in LocaleSpace  
def locale.name: string;      // similar to uname -n  
def locale.numCores: int;      // # of processor cores  
def locale.physicalMemory(...): ...; // amount of memory  
...
```

- Locale queries can also be made:

```
...myvar.locale... // query the locale where myvar is stored  
...here...        // query where the current task is running
```

Locality: Task Placement



on clauses: indicate where statements should execute:

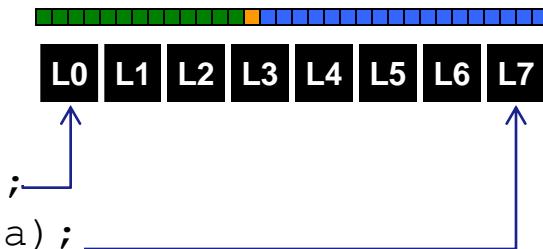
Either by naming locales explicitly...

```
cobegin {
    on TaskALocs do computeTaskA(...);
    on TaskBLocs do computeTaskB(...);
    on Locales(0) do computeTaskC(...);
}
```



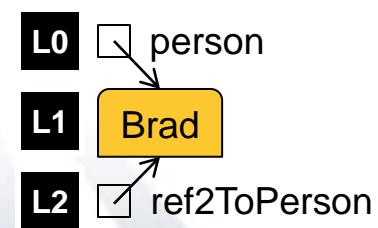
...or in a data-driven manner:

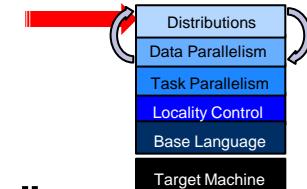
```
const pivot = computePivot(lo, hi, data);
cobegin {
    on data[lo] do Quicksort(lo, pivot, data);
    on data[hi] do Quicksort(pivot+1, hi, data);
}
```



They can also control where data is allocated:

```
var person: Employee;
on Locales(1) do person = new Employee("Brad");
on Locales(2) do var ref2ToPerson = person;
```

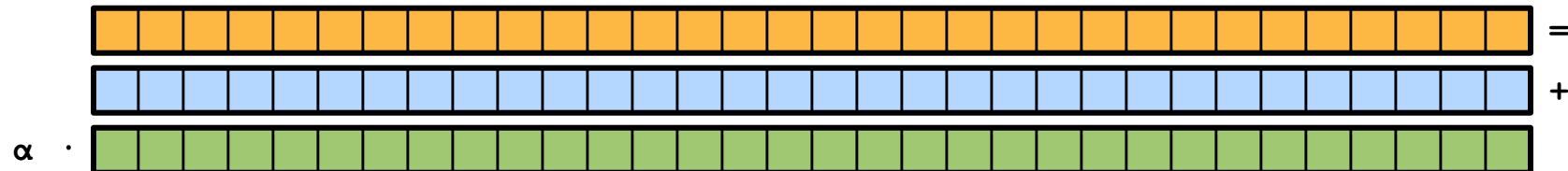




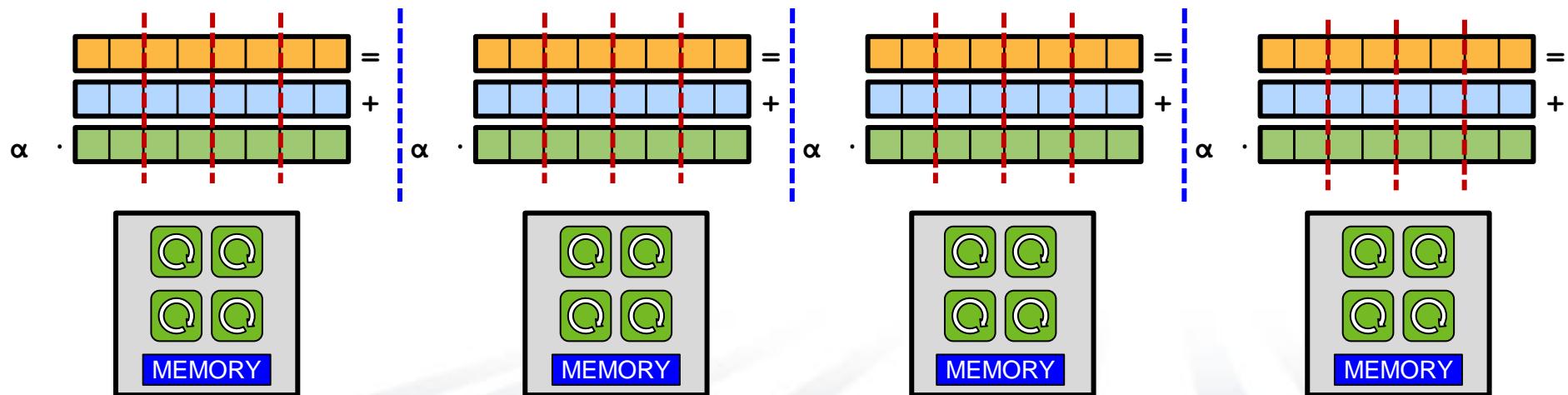
Chapel Distributions

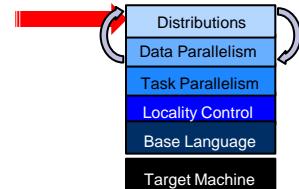
Distributions: “Recipes for parallel, distributed arrays”

- help the compiler map from the computation’s global view...



...down to the *fragmented*, per-processor implementation

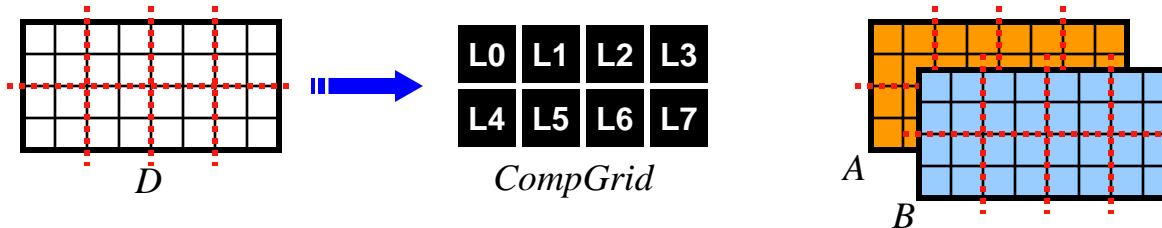




Domain Distribution

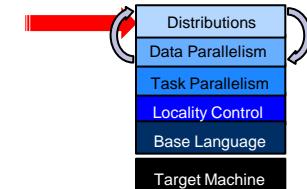
Domains may be distributed across locales

```
var D: domain(2) dmapped Block(CompGrid, ...) = ...;
```



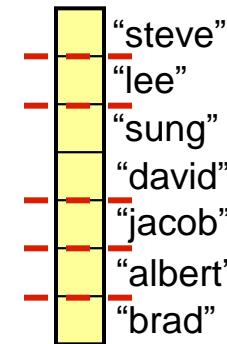
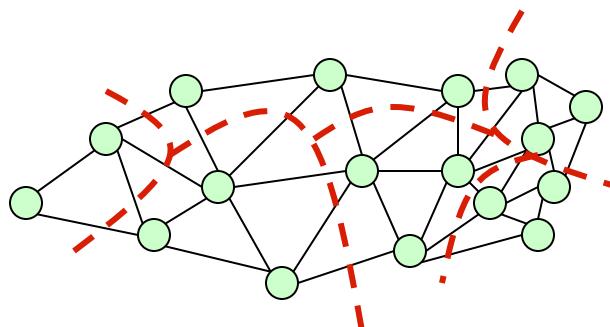
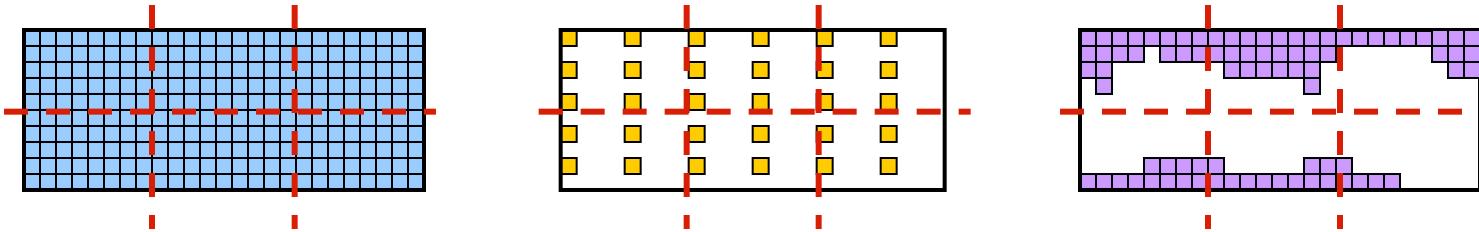
A distribution defines...

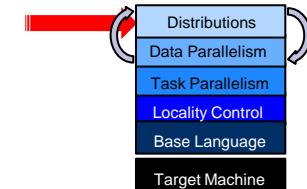
- ...ownership of the domain's indices (and its arrays' elements)
- ...default work ownership for operations on the domains/arrays
 - e.g., forall loops or promoted operations
- ...memory layout/representation of array elements/domain indices
- ...implementation of operations on its domains and arrays
 - e.g., accessors, iterators, communication patterns, ...



Domain Distributions

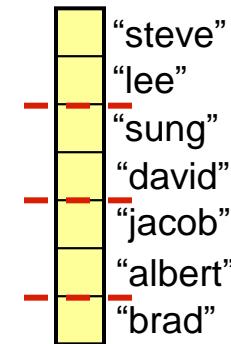
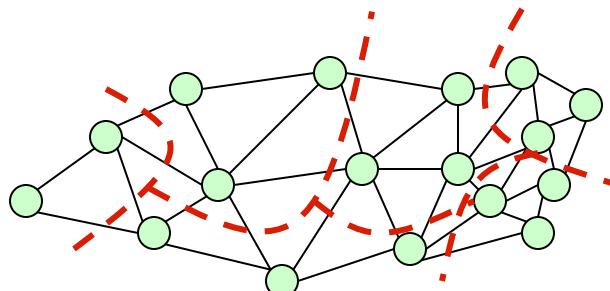
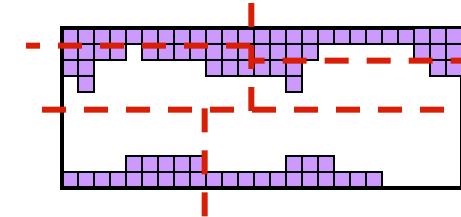
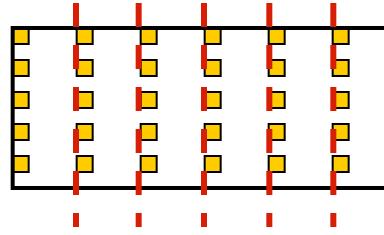
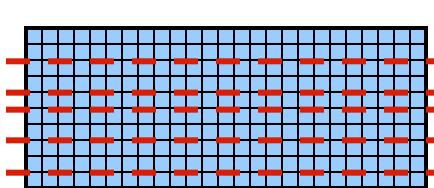
- Any domain type may be distributed
- Distributions do not affect program semantics
 - only implementation details and therefore performance

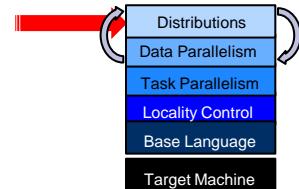




Domain Distributions

- Any domain type may be distributed
- Distributions do not affect program semantics
 - only implementation details and therefore performance





Distributions: Goals & Research

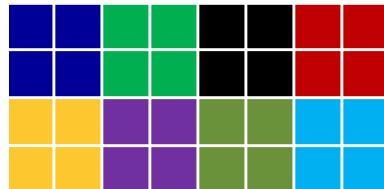
- Advanced users can write their own distributions
 - specified in Chapel using lower-level language features
- Chapel will provide a standard library of distributions
 - written using the same user-defined distribution mechanism

(Draft paper describing user-defined distribution strategy available by request)

The Block Distribution

The Block Distribution maps the indices of a domain in a dense fashion across the target Locales according to the `boundingBox` argument

```
const Dist = new dmap(new Block(boundingBox=[1..4, 1..8]));  
  
var Dom: domain(2) dmapped Dist = [1..4, 1..8];
```



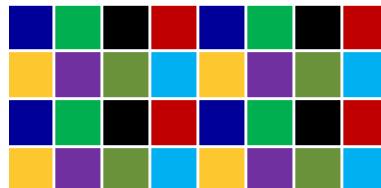
L0	L1	L2	L3
L4	L5	L6	L7

The Cyclic Distribution

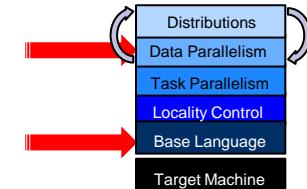
The Cyclic Distribution maps the indices of a domain in a round-robin fashion across the target Locales according to the `startIdx` argument

```
const Dist = new dmap(new Cyclic(startIdx=(1,1)));
```

```
var Dom: domain(2) dmapped Dist = [1..4, 1..8];
```



Other Features



- zippered and **tensor** flavors of iteration and promotion
- *subdomains* and *index types* to help reason about indices
- reductions and scans (standard or user-defined operators)