

# Chapel: Features

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Cray Inc.

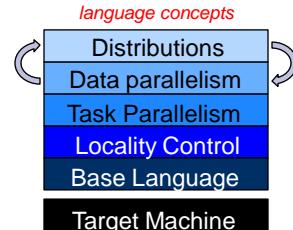
CSEP 524  
May 20, 2010

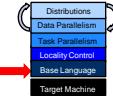


CRAY

## Outline

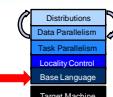
- Language Overview
  - Base Language
  - Task Parallelism
  - Data Parallelism
  - Locality
  - Distributions





## Base Language: Design

- Block-structured, imperative programming
- Intentionally not an extension to an existing language
- Instead, select attractive features from others:
  - ZPL, HPF:** data parallelism, index sets, distributed arrays  
(see also APL, NESL, Fortran90)
  - Cray MTA C/Fortran:** task parallelism, lightweight synchronization
  - CLU:** iterators (see also Ruby, Python, C#)
  - ML:** latent types (see also Scala, Matlab, Perl, Python, C#)
  - Java, C#:** OOP, type safety
  - C++:** generic programming/templates (without adopting its syntax)
  - C, Modula, Ada:** syntax
- Follow lead of C family of languages when useful  
(C, Java, C#, Perl, ...)



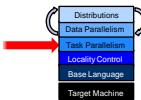
## Base Language: My Favorite Features

- **Rich compile-time language**
  - parameter values (compile-time constants)
  - folded conditionals, unrolled for loops, tuple expansions
  - type and parameter functions – evaluated at compile-time
- **Latent types**
  - ability to omit type specifications for convenience or code reuse
  - type specifications can be omitted from...
    - ...variables (inferred from initializers)
    - ...class members (inferred from constructors)
    - ...function arguments (inferred from callsite)
    - ...function return types (inferred from return statements)
- **Configuration variables** (and parameters)
 

```
config const n = 100; // override with ./a.out --n=100000
```
- **Tuples**
- **Iterators** (in the CLU, Ruby sense, not C++/Java-style)
- **Declaration Syntax:** more like Pascal/Modula/Scala than C



## Task Parallelism: Task Creation



**begin:** creates a task for future evaluation

```
begin DoThisTask();
  WhileContinuing();
  TheOriginalThread();
```

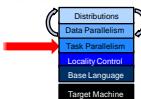
**sync:** waits on all begins created within its dynamic scope

```
sync {
  begin treeSearch(root);
}

def treeSearch(node) {
  if node == nil then return;
  begin treeSearch(node.right);
  begin treeSearch(node.left);
}
```



## Task Parallelism: Structured Tasks



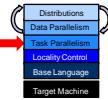
**cobegin:** creates a task per component statement:

<pre>computePivot(lo, hi, data); cobegin {   Quicksort(lo, pivot, data);   Quicksort(pivot, hi, data); } // implicit join here</pre>	<pre>cobegin {   computeTaskA(...);   computeTaskB(...);   computeTaskC(...); } // implicit join</pre>
--	--

**coforall:** creates a task per loop iteration

```
coforall e in Edges {
  exploreEdge(e);
} // implicit join here
```





## Task Parallelism: Task Coordination

*sync variables:* store full/empty state along with value

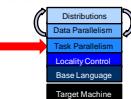
```
var result$: sync real; // result is initially empty
sync {
    begin ... = result$; // block until full, leave empty
    begin result$ = ...; // block until empty, leave full
}
result$.readXX(); // read value, leave state unchanged;
// other variations also supported
```

*single-assignment variables:* writeable once only

```
var result$: single real = begin f(); // result initially empty
... // do some other things
total += result$; // block until f() has completed
```

*atomic sections:* support transactions against memory

```
atomic {
    newnode.next = insertpt;
    newnode.prev = insertpt.prev;
    insertpt.prev.next = newnode;
    insertpt.prev = newnode;
}
```



## Producer/Consumer example

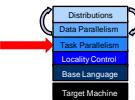
```
var buff$: [0..buffersize-1] sync int;

cobegin {
    producer();
    consumer();
}

def producer() {
    var i = 0;
    for ... {
        i = (i+1) % buffersize;
        buff$(i) = ...;
    }
}

def consumer() {
    var i = 0;
    while ... {
        i = (i+1) % buffersize;
        ...buff$(i)...;
    }
}
```

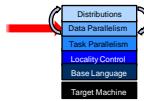
## QuickSort in Chapel



```
def quickSort(arr: [],
              thresh: int,
              low: int = arr.domain.low,
              high: int = arr.domain.high) {
    if high - low < 8 {
        bubbleSort(arr, low, high);
    } else {
        const pivotVal = findPivot(arr, low, high);
        const pivotLoc = partition(arr, low, high, pivotVal);
        serial thresh <= 0 do cobegin {
            quickSort(arr, thresh-1, low, pivotLoc-1);
            quickSort(arr, thresh-1, pivotLoc+1, high);
        }
    }
}
```

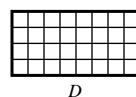


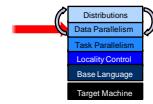
## Data Parallelism: Domains



*domain*: a first-class index set

```
var m = 4, n = 8;
var D: domain(2) = [1..m, 1..n];
```

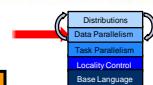
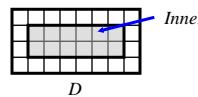




## Data Parallelism: Domains

*domain*: a first-class index set

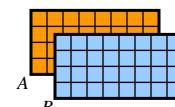
```
var m = 4, n = 8;
var D: domain(2) = [1..m, 1..n];
var Inner: subdomain(D) = [2..m-1, 2..n-1];
```



## Domains: Some Uses

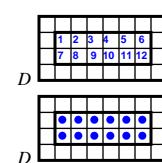
- Declaring arrays:

```
var A, B: [D] real;
```



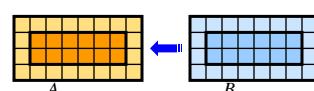
- Iteration (sequential or parallel):

```
for ij in Inner { ... }
or: forall ij in Inner { ... }
or: ...
```



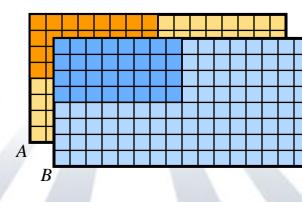
- Array Slicing:

```
A[Inner] = B[Inner];
```



- Array reallocation:

```
D = [1..2*m, 1..2*n];
```



## Forall vs. For vs. Coforall

### for loops:

- Use the current task to execute the loop serially

### coforall loops:

- Execute the loop using a distinct task per iteration
- Can have synchronization between iterations

### forall loops:

- Use some number of tasks between these two extremes
- Must be legally executable by a single task
- How many tasks are used in practice?



## Data Parallelism Throttles

### --dataParTasksPerLocale=#

- Specify # of tasks to execute forall loops
- Default: number of cores (*in current implementation*)

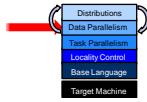
### --dataParIgnoreRunningTasks=[true|false]

- If false, reduce # of forall tasks by # of running tasks
- Default: true (*in current implementation*)

### --dataParMinGranularity=#

- reduce # of tasks if any task has fewer iterations
- Default: 1 (*in current implementation*)

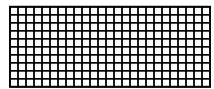




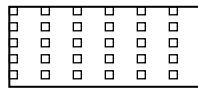
## Data Parallelism: Domain Types

Chapel supports several domain types...

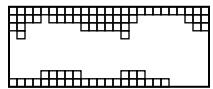
```
var OceanSpace = [0..#lat, 0..#long],  
    AirSpace = OceanSpace by (2,4),  
    IceSpace: sparse subdomain(OceanSpace) = genCaps();
```



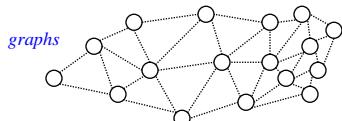
dense



strided



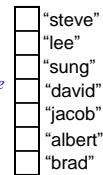
sparse



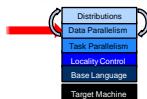
graphs



associative



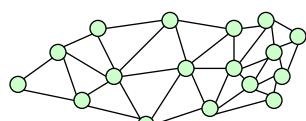
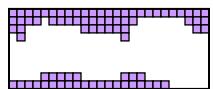
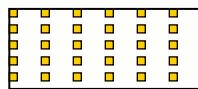
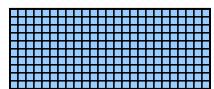
```
var Vertices: domain(opaque) = ..., People: domain(string) = ...;
```



## Data Parallelism: Domain Uses

All domain types can be used to declare arrays...

```
var Ocean: [OceanSpace] real,  
    Air: [AirSpace] real,  
    IceCaps[IceSpace] real;
```



```
var Weight: [Vertices] real,
```



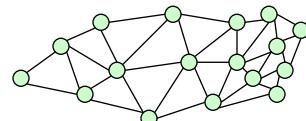
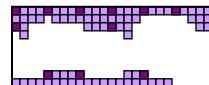
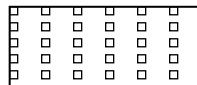
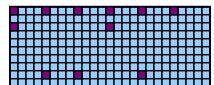
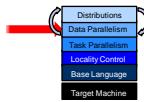
```
Age: [People] int;
```



## Data Parallelism: Domain Uses

...to iterate over index sets...

```
forall ij in AirSpace do
    Ocean(ij) += IceCaps(ij);
```



- “steve”
- “lee”
- “sung”
- “david”
- “jacob”
- “albert”
- “brad”

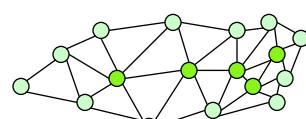
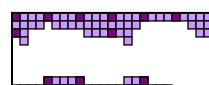
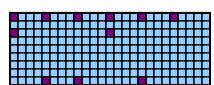
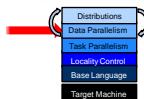
```
forall v in Vertices do
    Weight(v) = numEdges(v);
```

```
forall p in People do
    Age(p) += 1;
```

## Data Parallelism: Domain Uses

...to slice arrays...

```
Ocean[AirSpace] += IceCaps[AirSpace];
```

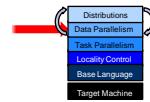


- “steve”
- “lee”
- “sung”
- “david”
- “jacob”
- “albert”
- “brad”

...Vertices[Interior]...

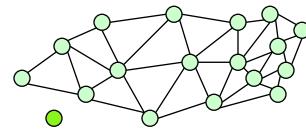
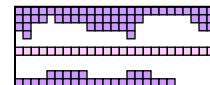
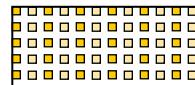
...People[Interns]...

## Data Parallelism: Domain Uses



...and to reallocate arrays

```
AirSpace = OceanSpace by (2,2);
IceSpace += genEquator();
```

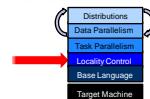


- “steve”
- “lee”
- “sung”
- “david”
- “jacob”
- “albert”
- “brad”
- “srini”

```
newnode = Vertices.create();    People += “srini”;
```

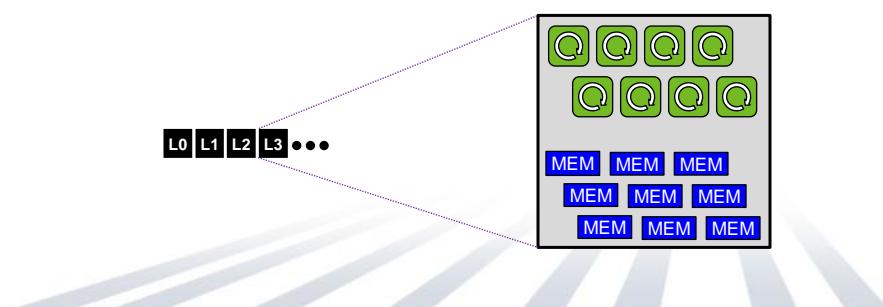


## Locality: Locales

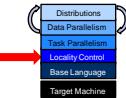


*locale*: An abstract unit of the target architecture

- supports reasoning about locality
- has capacity for processing and storage
- two threads in a given locale have similar access to a given address
  - addresses in that locale are ~uniformly accessible
  - addresses in other locales are also accessible, but at a price
- locales are defined for a given architecture by a Chapel compiler
  - e.g., a multicore processor or SMP node could be a locale



## Locales and Program Startup



- Chapel users specify # locales on executable command-line

prompt> `myChapelProg -n1=8` # run using 8 locales

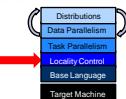
`L0 | L1 | L2 | L3 | L4 | L5 | L6 | L7`

- Chapel launcher bootstraps program execution:

- obtains necessary machine resources
  - e.g., requests 8 nodes from the job scheduler
- loads a copy of the executable onto the machine resources
- starts running the program. *Conceptually...*
  - ...locale #0 starts running program's entry point (`main()`)
  - ...other locales wait for work to arrive



## Locale Variables



Built-in variables represent a program's locale set:

```

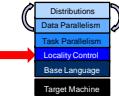
config const numLocales: int;           // number of locales
const LocaleSpace = [0..numLocales-1], // locale indices
    Locales: [LocaleSpace] locale; // locale values
  
```

`numLocales: 8`

`LocaleSpace:` 

`Locales:` `L0 | L1 | L2 | L3 | L4 | L5 | L6 | L7`





## Locale Views

Using standard array operations, users can create their own locale views:

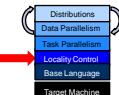
```
var TaskALocs = Locales[..numTaskALocs];
var TaskBLocs = Locales[numTaskALocs+1..];
```

L0	L1
L2	L3

L4	L5	L6	L7
----	----	----	----

```
var CompGrid = Locales.reshape([1..gridRows,
                                1..gridCols]);
```

L0	L1	L2	L3
L4	L5	L6	L7



## Locale Methods

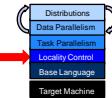
- The locale type supports built-in methods:

```
def locale.id: int;           // index in LocaleSpace
def locale.name: string;       // similar to uname -n
def locale.numCores: int;      // # of processor cores
def locale.physicalMemory(...): ...; // amount of memory
...
```

- Locale queries can also be made:

```
...myvar.locale... // query the locale where myvar is stored
...here...         // query where the current task is running
```





## Locality: Task Placement

*on clauses:* indicate where statements should execute:

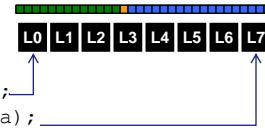
Either by naming locales explicitly...

```
cobegin {
    on TaskALocs do computeTaskA(...);
    on TaskBLocs do computeTaskB(...);
    on Locales(0) do computeTaskC(...);
}
```



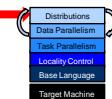
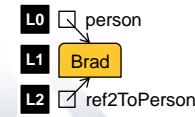
...or in a data-driven manner:

```
const pivot = computePivot(lo, hi, data);
cobegin {
    on data[lo] do Quicksort(lo, pivot, data);
    on data[hi] do Quicksort(pivot+1, hi, data);
}
```



They can also control where data is allocated:

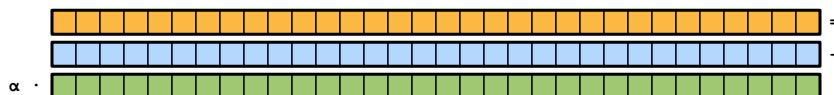
```
var person: Employee;
on Locales(1) do person = new Employee("Brad");
on Locales(2) do var ref2ToPerson = person;
```



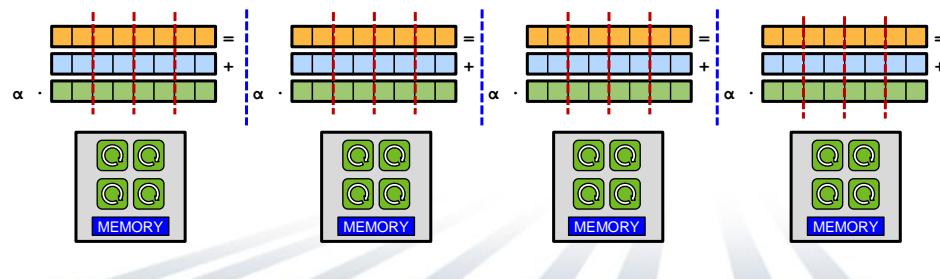
## Chapel Distributions

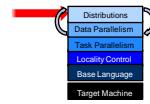
*Distributions:* “Recipes for parallel, distributed arrays”

- help the compiler map from the computation’s global view...



...down to the *fragmented, per-processor implementation*

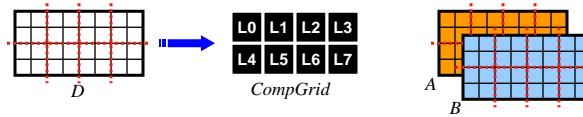




## Domain Distribution

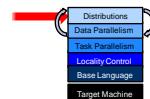
Domains may be distributed across locales

```
var D: domain(2) dmapped Block(CompGrid, ...) = ...;
```



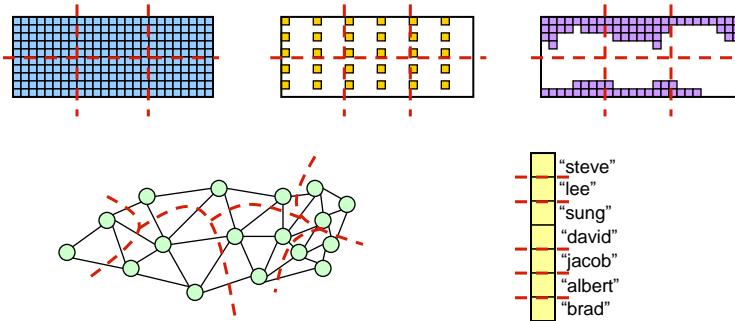
A distribution defines...

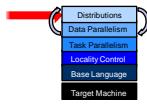
- ...ownership of the domain's indices (and its arrays' elements)
- ...default work ownership for operations on the domains/arrays
  - e.g., forall loops or promoted operations
- ...memory layout/representation of array elements/domain indices
- ...implementation of operations on its domains and arrays
  - e.g., accessors, iterators, communication patterns, ...



## Domain Distributions

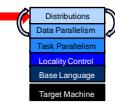
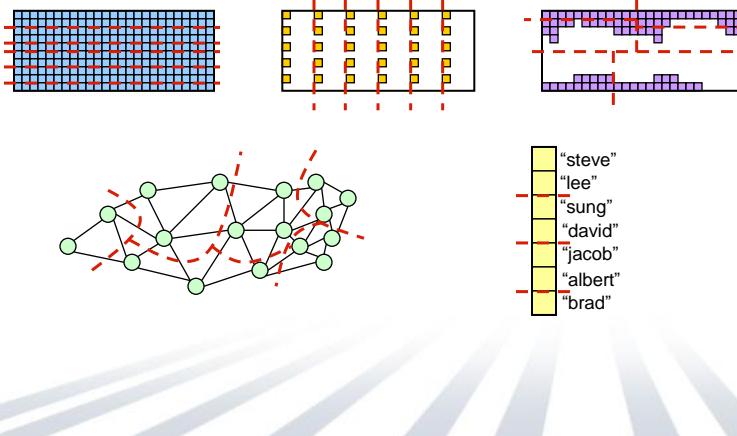
- Any domain type may be distributed
- Distributions do not affect program semantics
  - only implementation details and therefore performance





## Domain Distributions

- Any domain type may be distributed
- Distributions do not affect program semantics
  - only implementation details and therefore performance



## Distributions: Goals & Research

- Advanced users can write their own distributions
  - specified in Chapel using lower-level language features
- Chapel will provide a standard library of distributions
  - written using the same user-defined distribution mechanism

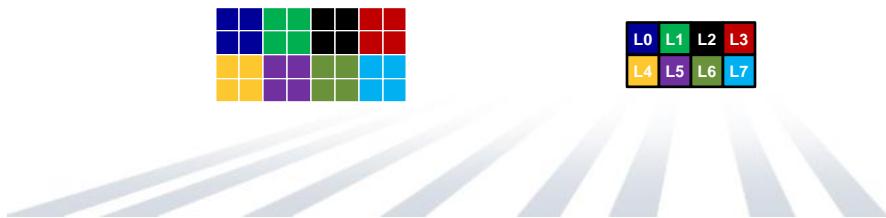
*(Draft paper describing user-defined distribution strategy available by request)*

## The Block Distribution

The Block Distribution maps the indices of a domain in a dense fashion across the target Locales according to the `boundingBox` argument

```
const Dist = new dmap(new Block(boundingBox=[1..4, 1..8]));

var Dom: domain(2) dmapped Dist = [1..4, 1..8];
```

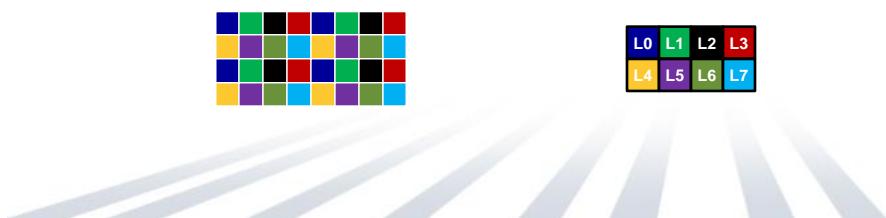


## The Cyclic Distribution

The Cyclic Distribution maps the indices of a domain in a round-robin fashion across the target Locales according to the `startIdx` argument

```
const Dist = new dmap(new Cyclic(startIdx=(1,1)));

var Dom: domain(2) dmapped Dist = [1..4, 1..8];
```



## Other Features

- zippered and `tensor` flavors of iteration and promotion
- `subdomains` and `index types` to help reason about indices
- reductions and scans (standard or user-defined operators)

