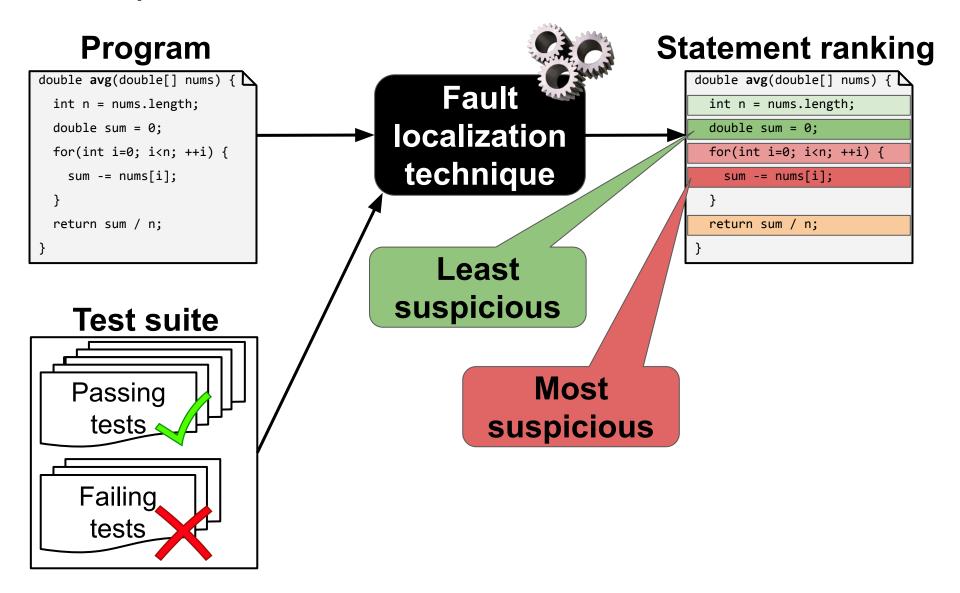
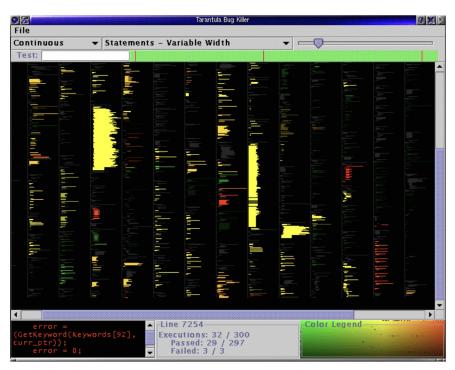
Static analysis

UW CSE P 504

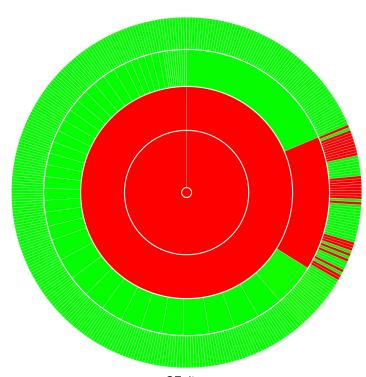
Today

- Recap: statistical fault localization
- Static Analysis
 - Motivation
 - Examples
 - Intro to Abstract Interpretation





Jones et al., Visualization of test information to assist fault localization, ICSE'02



GZoltar

Developer in the loop

- Which granularity is most useful?
 - file level
 - method level
 - statement level
- What context do you need to reason about?
 - a file
 - a method
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Developer in the loop

- Which granularity is most useful?
 - file level
 - method level
 - statement level
- What context do you need to reason about?
 - a file
 - a method
 - a statement
- Processing FL output
 - How useful is color coding (heatmap) vs. ranking?
 - How realistic is "sequential debugging"?

Static Analysis

Dynamic analysis

- Reason about the program based on some program executions.
- Observe concrete behavior at run time.
- Improve confidence in correctness.
- Unsound* but precise.

^{*} Some static analyses are unsound; dynamic analyses can be sound.

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$$y = x++$$

[
$$y$$
:=2, x :=3]

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- Reason about the program without executing it.
- Build an abstraction of run-time states.
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- Prove a property of the program.
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$$[y:=2, x:=2]$$

$$V = X++$$

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$$y = x++$$

[y:=prime, x:=anything]

^{*} Some static analyses are unsound; dynamic analyses can be sound.

Dynamic analysis

- Concrete domain
- Precise but unsound
- Slow if exhaustive

Static analysis

- Abstract domain
- Sound but imprecise
- Slow if precise

Dynamic analysis

- Concrete domain
- Precise but unsound
- Slow if exhaustive

Concrete domain

```
int getValue(int a) {
  return (a % 3) * 2;
}
int x = getValue(7);
```

Static analysis

- Abstract domain
- Sound but imprecise
- Slow if precise

Abstract domain

What possible value(s) does getValue() return?

Dynamic analysis

- Concrete domain
- Precise but unsound
- Slow if exhaustive

Concrete domain

```
0, 2, 4, 6, 8, 10, \dots
```

```
int getValue(int a) {
  return (a % 3) * 2;
}
int x = getValue(7);
```

Static analysis

- Abstract domain
- Sound but imprecise
- Slow if precise

Abstract domain

even, odd, anything

What possible value(s) does getValue() return?

Recall the following terms:

- Precision vs. Recall (and FP/FN/TP/TN)
- 2. Soundness vs. Completeness
- 3. Accuracy vs. Precision

Analysis result Pos Neg Neg Neg

Concrete domain vs. Abstract domain

0, 2, 4, 6, 8, 10, ... even, odd, anything

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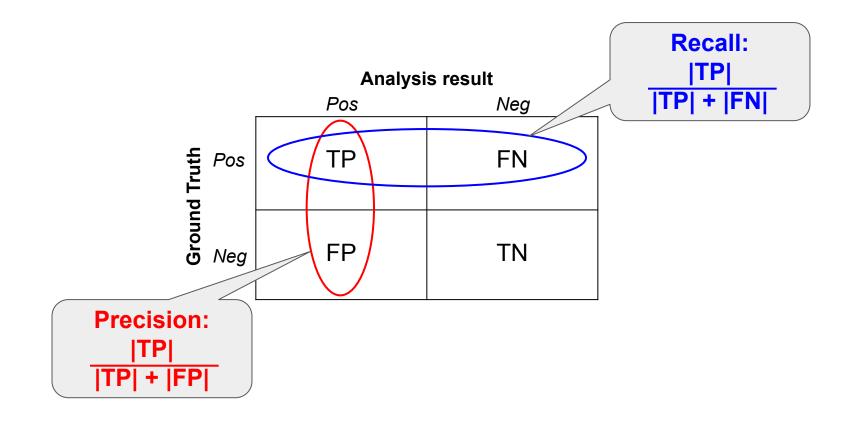
1. Precision vs. Recall (and FP/FN/TP/TN)

		Analysis result		
		Pos	Neg	
d Truth	Pos Neg			
Groun	Neg			

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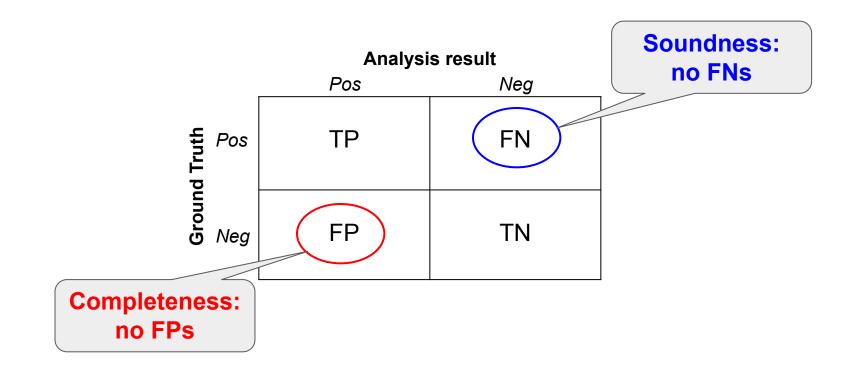
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- 2. Soundness vs. Completeness

		Analysis result	
		Pos	Neg
Ground Truth	Pos	TP	FN
	Neg	FP	TN

- 1. Precision vs. Recall (and FP/FN/TP/TN)
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Concrete domain

0, 2, 4, 6, 8, 10, ...

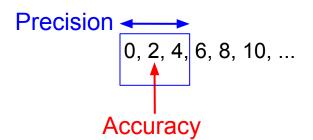
Abstract domain

even, odd, anything

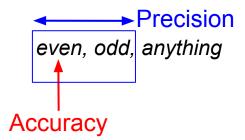
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Concrete domain



Abstract domain



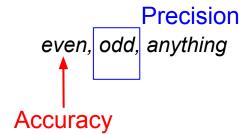
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Concrete domain

Precision 0, 2, 4, 6, 8, 10, ... Accuracy

Abstract domain



An analysis/measure can be precise and inaccurate at the same time!

Static analysis: applications

Compiler checks and optimizations

- Liveness analysis (register reallocation)
- Reachability analysis (dead code elimination)
- Code motion (while(cond) {x = comp(); ...})

Static analysis: code examples

Liveness

```
public class Liveness {
  public void liveness() {
    int a;
    if (alwaysTrue()) {
        a = 1;
    }
    System.out.println(a);
  }
}
```

Reachability

```
public void deadCode() {
   return;
   System.out.println("Here!");
}
```

Common static analyses

Live examples

- Definite assignment
- Dead code
- Linter warnings

Challenges to adopting static analysis

- Not integrated into the developer's workflow.
- Reported issues are not actionable.
- Developers do not trust the results (FPs).
- Fixing an issue is too expensive or risky.
- Developers do not understand the reported issues.
- Issues theoretically possible but don't manifest in practice.

"Produce less than 10% effective false positives. Developers should feel the check is pointing out an actual issue at least 90% of the time."

Effective false positive

- We consider an issue to be an "effective false positive" if developers did not take positive action after seeing the issue.
- If an analysis incorrectly reports an issue, but developers make the fix anyway to improve code readability or maintainability, that is not an effective false positive.
- If an analysis reports an actual fault, but the developer did not understand the fault and therefore took no action, that is an effective false positive.

Effective false positive: example (mutation testing)

Petrovic et al., ICSTW'18

Effective false positive: discussion

- We consider an issue to be an "effective false positive" if developers did not take positive action after seeing the issue.
- If an analysis incorrectly reports an issue, but developers make the fix anyway to improve code readability or maintainability, that is not an effective false positive.
- If an analysis reports an actual fault, but the developer did not understand the fault and therefore took no action, that is an effective false positive.

Do you agree with this characterization? Is effective false positive rate an adequate measure?

Abstract Interpretation

Properties of an ideal program analysis

- Soundness
- Completeness
- Termination

```
int x = 0;
while (!isDone()) {
    x = x + 1;
}
```

Properties of an ideal program analysis

- Soundness
- Completeness
- Termination

```
int x = 0;
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Abstract interpretation sacrifices completeness (precision)

A first example

Program

```
x = 0;
y = read_even();
x = y + 1;
y = 2 * x;
x = y - 2;
y = x / 2;
```

Are all statements necessary?

A first example: SSA form

Program

```
x = 0;
y = read_even();
x = y + 1;
y = 2 * x;
x = y - 2;
y = x / 2;
```

SSA form

 X_1 is never read.

A first example: one concrete execution

Program

```
x = 0;
y = read_even();
x = y + 1;
y = 2 * x;
x = y - 2;
y = x / 2;
```

Concrete execution

```
{x=0; y=undef}
{x=0; y=8}
{x=9; y=8}
{x=9; y=18}
{x=16; y=18}
{x=16; y=8}
```

A first example: symbolic reasoning



Program

```
x = 0;
y = read_even();
x = y + 1;
y = 2 * x;
x = y - 2;
y = x / 2;
```

SSA form

What facts can you deduce about y and x after execution?

A first example: symbolic reasoning

Program

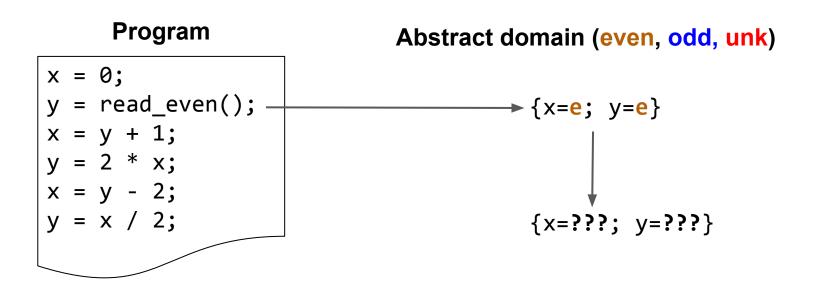
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x = 0;
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x = y + 1;
y = 2 * x;
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```

SSA form

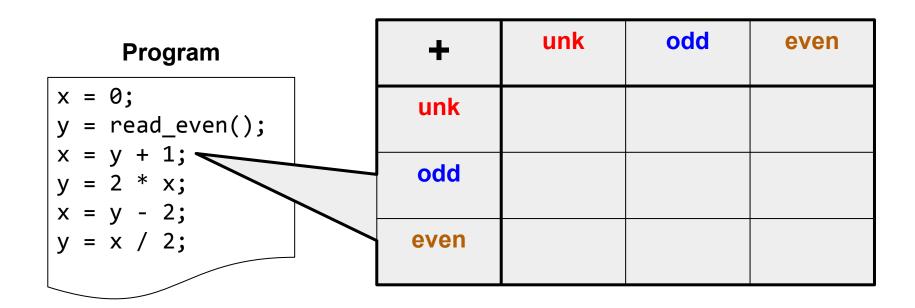
$$y_3 = x_3 / 2$$

 $y_3 = (y_2 - 2) / 2$
 $y_3 = (2 * x_2 - 2) / 2$
 $y_3 = (2 * (y_1 + 1) - 2) / 2$
 $y_3 = (2 * y_1 + 2 - 2) / 2$
 $y_3 = y_1$
 $x_3 = y_1 * 2$

Symbolic reasoning shows simplification potential.



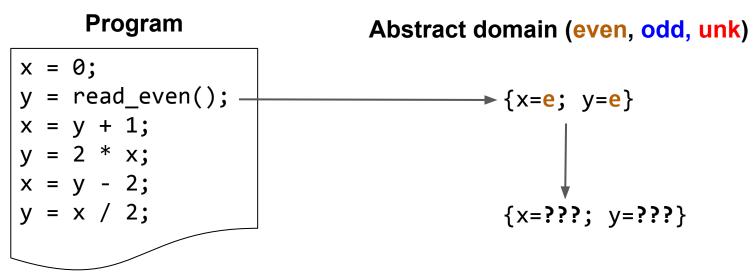
A first example: "abstract execution"



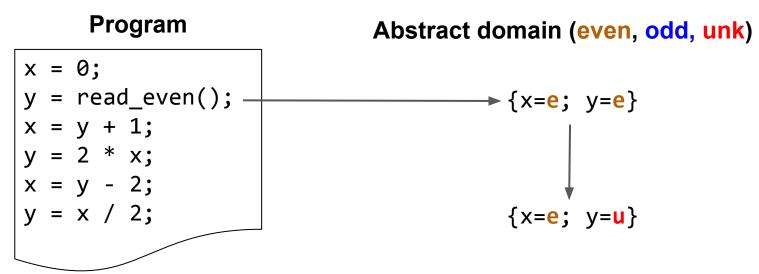
A first example: "abstract execution"

Program	+	unk	odd	even
<pre>x = 0; y = read_even();</pre>	unk	unk	unk	unk
x = y + 1; y = 2 * x; x = y - 2;	odd	unk	even	odd
y = x / 2;	even	unk	odd	even

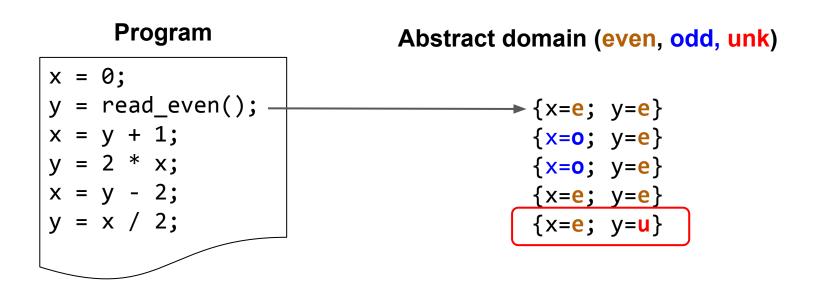








Convince yourself that this is true.



This result is accurate but imprecise.



```
Program

Abstract domain (even, odd, unk)

x = 0;
y = read_even();
x = y + 1;
y = 2 * x;
x = y - 2;
y = x / 2;

Abstract domain (even, odd, unk)

{x=e; y=e}
{x=o; y=e}
{x=e; y=e}
{x=e; y=u}
```

What abstract domain would allow us to conclude that y is even?