CSE P 501 – Compilers

Inlining and Devirtualization
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References

• *Adaptive Online Context-Sensitive Inlining*
  Hazelwood and Grove, ICG 2003

• *A Study of Devirtualization Techniques for a Java JIT Compiler*
  Ishizaki, et al, OOPSLA 2000

• Earlier versions of this lecture by Vijay Menon,
  CSE 501, Sp09 & Jim Hogg, CSE P 501 Sp14
Inlining

```c
long res;

void foo(long x) {
    res = 2 * x;
}

void bar() {
    res = foo(5);
}
```

```
long res;

void foo(long x) {
    res = 2 * x;
}

void bar() {
    res = 2 * 5;
}
```

```
long res;

void foo(long x) {
    res = 2 * x;
}

void bar() {
    res = 10;
}
```
Benefits

• Removes overhead of function call
  – No marshalling / unmarshalling parameters and return values
  – Better instruction cache locality

• Bonus: expands optimization opportunities
  – CSE, constant propagation, unreachable code elimination, ...
  – Poor person’s interprocedural optimization
Costs

• Code size
  – Typically expands overall program size
  – Can hurt instruction cache

• Compilation time
  – Larger methods can lead to more expensive compilation, more complex control flow
Language / runtime aspects

• What is the cost of a function call?
  – C: cheap, Java: moderate (virtual dispatch), Python: expensive

• Are targets resolved at compile time or run time?
  – C: compile time; Java, Python: run time

• Is the whole program available for analysis?
  – “separate compilation”

• Is profile information available?
  – If “m” is rarely called, don’t inline it
When to inline?

Jikes RVM (with Hazelwood/Grove adaptations):

- Call Instruction Sequence (CIS) = # of instructions to make call
  - Tiny (function size < 2x call size): Always inline
  - Small (2-5x): Inline subject to space constraints
  - Medium (5-25x): Inline if hot (subject to space constraints)
  - Large: Never inline
Gathering profile info

• Counter-based: Instrument edges in CFG
  – Entry + loop back edges
  – Enough edges (enough to get good results without excessive overhead)
  – Expensive - typically removed in optimized code
  – Depends critically on the “training sets”

• Call stack sampling
  – Periodically walk stack
  – Interrupt-based or instrumentation-based
  – May gather info on what calls what (callsite info)
Object-oriented languages

• OO encourages lots of small methods
  – getters, setters, ...
  – Inlining is a requirement for performance
    • High call overhead wrt total execution
    • Limited scope for compiler optimizations without it
  – For Java, C#, if you’re going to anything, do this!
  – But ... virtual methods are a challenge
Virtual methods

In general, we cannot determine the target until runtime

Some languages (e.g., Java) allow *dynamic class loading*: all subclasses of A may not be visible until runtime

```java
class A {
    int foo() { return 0; }
    int bar() { return 1; }
}

class B extends A {
    int foo() { return 2; }
}

void baz(A x) {
    y = x.foo();
    z = x.bar();
}
```
Virtual tables

• Object layout in a JVM:
Virtual method dispatch

Source:
\[
\begin{align*}
  y &= x.\text{foo}() ; \\
  z &= x.\text{bar}() ; \\
\end{align*}
\]

t1 = ldvtable x
t2 = ldvirtfunaddr t1, A::foo
t3 = call [t2] (x)
t4 = ldvtable x
t5 = ldvirtfunaddr t4, A::bar
t6 = call [t5] (x)

- x is the *receiver* object
- For a receiver object with a runtime type of B, t2 will refer to B::foo.
Devirtualization

• Goal: change virtual calls to static calls in compiler
• Benefits: enables inlining, lowers call overhead, better l-cache performance, better indirect-branch prediction
• Often optimistic:
  – Make guess at compile time
  – Test guess at run time
  – Fall back to virtual call if necessary
Guarded devirtualization

- Guess receiver type is B (based on profile or other information)
- Call to B::foo is statically known - can be inlined
- But guard inhibits optimization

\[
\begin{align*}
t1 &= \text{ldvtable } x \\
t7 &= \text{getvtable } B \\
\text{if } t1 &= t7 \\
    t3 &= \text{call } B::\text{foo}(x) \\
\text{else} \\
    t2 &= \text{ldvirtfunaddr } t1, A::\text{foo} \\
    t3 &= \text{call } [t2] (x) \\
\end{align*}
\]
Guarded by method test

- Guess that method is B::foo outside guard
- More robust, but more overhead
- Harder to optimize redundant guards

```plaintext
t1 = ldvtable x
t2 = ldvirtfunaddr t1
t7 = getfunaddr B::foo
if t2 == t7
    t3 = call B::foo(x)
else
    t2 = ldvirtfunaddr t1, A::foo
    t3 = call [t2] (x)
...
```
How to guess receiver?

• Profile information
  – Record call site targets and / or frequently executed methods at run time
  – “monomorphic” vs. “polymorphic”

• Class hierarchy analysis
  – Walk class hierarchy at compile time

• Type analysis
  – Intra / interprocedural data flow analysis
Class hierarchy analysis

• Walk class hierarchy at compilation time
  – If only one implementation of a method (i.e., in the base class), devirtualize to that target
• Not guaranteed in the presence of class loading
  – Still need runtime test / fallback
Flow sensitive type analysis

• Perform a forward dataflow analysis propagating type information.

• At each use site, compute the possible set of types.

• At call sites, use type information of receiver to narrow targets.

```
A a1 = new B();
a1.foo();
if (a2 instanceof C)
a2.bar();
```
Alternatives to guarding

• Guarding impose overheads
  – run-time test on every call, merge points impede optimization

• Often “know” only one target is invoked
  – call site is monomorphic

• Alternative: compile without guards
  – recover as assumption is violated (e.g., class load)
  – cheaper runtime test vs more costly recovery
Recompilation approach

• Optimistically assume current class hierarchy will never change wrt a call
• Devirtualize and/or inline call sites without guard
• On violating class load, recompile caller method
  – Recompiled code installed before new class
  – New invocations will call de-optimized code
  – What about current invocations?
• Nice match with JIT compiling
Preexistence analysis

• Idea: if the receiver object pre-existed the caller method invocation, then the call site is only affected by a class load in future invocations.

• If new class C is loaded during execution of baz, x cannot have type C:

```java
void baz(A x) {
    ...
    // C loaded here
    x.bar();
}
```
Code-patching

• Pre-generate fallback virtual call out of line
• On invalidating class load, overwrite direct call / inlined code with a jump to the fallback code
  – Must be thread-safe!
  – On x86, single write within a cache line is atomic
• No recompilation necessary
Patching - before

t3 = 2 // B::foo (inlined)
next:
  ...

fallback:
t2 = ldvirtfunaddr t1, A::foo
t3 = call [t2] (x)
goto next
Patching - after

t3 = 2 // B::foo (inlined)  ← goto fallback
next:

... 

fallback:
t2 = ldvirtfunaddr t1, A::foo
t3 = call [t2] (x)
goto next