CSE P 501 – Compilers

Code Shape I – Basic Constructs
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Agenda

- Mapping source code to x86
  - Mapping for other common architectures follows same basic pattern
- Now: basic statements and expressions
- Next: Object representation, method calls, and dynamic dispatch

Review: Variables

- For us, all data will be in either:
  - A stack frame for method local variables
  - An object for instance variables
- Local variables accessed via ebp

```
mov  eax,[ebp+12]
```

- Instance variables accessed via an object address in a register
  - Details later

Conventions for Examples

- Examples show code snippets in isolation
- Real code generator needs to worry about things like
  - Which registers are busy at which point in the program
  - Which registers to spill into memory when a new register is needed and no free ones are available
    - (x86: temporaries are often pushed on the stack, but can also be stored in a stack frame)
- Register eax used below as a generic example
  - Rename as needed for more complex code involving multiple registers

Peephole Optimizations

- A class of optimizations involving small numbers of instructions
- We'll point out a few of these along the way

Constants

- Source

```
17
```

- x86

```
mov   eax,17
```

  - Idea: realize constant value in a register
- Optimization: if constant is 0

```
xor   eax,eax
```

Assignment Statement

- Source
  - var = exp;
- x86
  - <code to evaluate exp into, say, eax>
  - mov [ebp+offset$_{var}$],eax

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Unary Minus

- Source
  - -exp
- x86
  - <code evaluating exp into eax>
  - neg eax
- Optimization
  - Collapse -(-exp) to exp
- Unary plus is a no-op

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Binary +

- Source
  - exp1 + exp2
- x86
  - <code evaluating exp1 into eax>
  - <code evaluating exp2 into edx>
  - add eax,edx

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Binary +

- Optimizations
  - If exp2 is a simple variable or constant
    - add eax,exp2
  - Change exp1 + -exp2 into exp1-exp2
  - If exp2 is 1
    - inc eax

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Binary -, *

- Same as +
  - Use sub for –
  - Use imul for *
- Optimizations
  - Use left shift to multiply by powers of 2
  - Use x+x instead of 2*x, etc. (faster)
  - Use dec for x-1

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Integer Division

- Ghastly on x86
  - Only works on 64 bit int divided by 32-bit int
  - Requires use of specific registers
- Source
  - exp1 / exp2
- x86
  - <code evaluating exp1 into eax **ONLY**>
  - <code evaluating exp2 into ebx>
  - cdq ; extend to edx:eax, clobbers edx
  - idiv ebx ; quotient in eax; remainder in edx

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Control Flow

- Basic idea: decompose higher level operation into conditional and unconditional gotos
- In the following, j_{false} is used to mean jump when a condition is false
  - No such instruction on x86
  - Will have to realize with appropriate sequence of instructions to set condition codes followed by conditional jumps
  - Normally won't actually generate the value "true" or "false" in a register

While

- Source

  while (cond) stmt
- x86

```
test:   <code evaluating cond>
        jfalse done
        <code for stmt>
        jmp  test
done:
```


Labels

- In x86 assembly language we'll need to produce unique labels for each if, while, etc.
- Some assemblers allow for "local" labels that can be reused
- Ignore for now – concentrate on code shape

Optimization for While

- Put the test at the end

```
        jmp  test
loop:   <code for stmt>
test:   <code evaluating cond>
        jtrue  loop
```

- Why bother?
  - Pulls one instruction (jmp) out of the loop
  - Avoids a pipeline stall on jmp on each iteration
    - Although modern processors can often predict control flow and avoid the stall
- Easy to do from IR; not so easy if generating code on the fly (e.g., recursive descent 1-pass compiler)

Do-While

- Source

  do stmt while(cond);
- x86

```
loop:   <code for stmt>
        <code evaluating cond>
        jtrue  loop
```


If

- Source

  if (cond) stmt
- x86

```
        <code evaluating cond>
        jfalse skip
        <code for stmt>
skip:
```

If-Else

- Source

  if (cond) stmt1 else stmt2

- x86

```
        <code evaluating cond>
        j_false else
        <code for stmt1>
        jmp  done
else:   <code for stmt2>
done:
```

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Jump Chaining

- Observation: naïve implementation can produce jumps to jumps
- Optimization: if a jump has as its target an unconditional jump, change the target of the first jump to the target of the second
  - Repeat until no further changes

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Boolean Expressions

- What do we do with this?

  x > y

- It is an expression that evaluates to true or false
  - Could generate the value (0/1 or whatever the local convention is)
  - But normally we don't want/need the value; we're only trying to decide whether to jump

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Code for exp1 > exp2

- Basic idea: designate jump target, and whether to jump if the condition is true or if it is false
- Example: exp1 > exp2, target L123, jump on false

```
<evaluate exp1 to eax>
<evaluate exp2 to edx>
cmp  eax,edx
jng   L123
```

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Boolean Operators: !

- Source

  ! exp

- Context: evaluate exp and jump to L123 if false (or true)
- To compile !, reverse the sense of the test: evaluate exp and jump to L123 if true (or false)

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Boolean Operators: && and ||

- In C/C++/Java/C#, these are *short-circuit* operators
  - Right operand is evaluated only if needed
- Basically, generate the if statements that jump appropriately and only evaluate operands when needed

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Example: Code for &&

- Source

  if (exp1 && exp2) stmt

- x86

  <code for exp1>
  j_{false} skip
  <code for exp2>
  j_{false} skip
  <code for stmt>
  skip:

Example: Code for ||

- Source

  if (exp1 || exp2) stmt

- x86

  <code for exp1>
  j_{true} doit
  <code for exp2>
  j_{false} skip
  doit: <code for stmt>
  skip:

Realizing Boolean Values

- If a boolean value needs to be stored in a variable or method call parameter, generate code needed to actually produce it
- Typical representations: 0 for false, +1 or -1 for true
  - C uses 0 and 1; we'll use that
  - Best choice can depend on machine architecture; normally some convention is established during the primeval history of the architecture

Boolean Values: Example

- Source

  var = bexp ;

- x86

  <code for bexp>
  j_{false} genFalse
  mov eax,1
  jmp storeIt
  genFalse:
  mov eax,0
  storeIt: mov [ebp+$offset_{var}$],eax ; generated by asg stmt

Faster, If Enough Registers

- Source

  var = bexp ;

- x86

  xor eax,eax
  <code for bexp>
  j_{false} storeIt
  inc eax
  storeIt: mov [ebp+$offset_{var}$],eax ; generated by asg stmt

  - Or use conditional move (movecc) instruction if available

Other Control Flow: switch

- Naïve: generate a chain of nested if-else if statements
- Better: switch is designed to allow an O(1) selection, provided the set of switch values is reasonably compact
- Idea: create a 1-D array of jumps or labels and use the switch expression to select the right one
  - Need to generate the equivalent of an if statement to ensure that expression value is within bounds

Switch

- Source

```
switch (exp) {
  case 0: stmts0;
  case 1: stmts1;
  case 2: stmts2;
}
```

- X86

```
        <put exp in eax>
        "if (eax < 0 || eax > 2)
            jmp defaultLabel"
        mov eax,swtab[eax*4]
        jmp eax
                .data
swtab           dd L0
                dd L1
                dd L2
                .code
L0:             <stmts0>
L1:             <stmts1>
L2:             <stmts2>
```


x86 Addressing Modes

- A memory address in x86 can be
  - register
  - +register optionally scaled by *2, *4, or *8
  - +constant offset
- Assemblers have many syntax variations involving labels, register values in brackets, etc.

Arrays

- Several variations
- C/C++/Java
  - 0-origin; an array with n elements contains variables a[0]...a[n-1]
  - 1 or more dimensions; row major order
- Key step is to evaluate a subscript expression and calculate the location of the corresponding element

0-Origin 1-D Integer Arrays

- Source

```
exp1[exp2]
```

- x86

```
<evaluate exp1 (array address) in eax>
<evaluate exp2 in edx>
address is [eax+4*edx]   ; 4 bytes per element
```


Fortran Arrays

- Subscripts start with 1 (default)
- Column-major order
  - E.g., an array with 3 rows and 2 columns is stored in this sequence: a(1,1), a(2,1), a(3,1), a(1,2), a(2,2), a(3,2)

a(i,j) in Fortran

- To find a(i,j), we need to know
  - Values of i and j
  - How many *rows* the array has
- Location of a(i,j) is
  - Location of a + (j-1)*(#of rows) + (i-1)
- Factor to pull out load-time constant part and evaluate that at load time – no recalculating at runtime
  - [Loc. of a – (#rows) – 1] + [j*(#rows) + i]

Coming Attractions

- Code Generation for Objects
  - Representation
  - Method calls
  - Inheritance and overriding
- Strategies for implementing code generators
- Code improvement - optimization

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