# Topics in Probabilistic and Statistical Databases

# Lecture 10: Sampling and Review

Dan Suciu University of Washington

## References

- Towards Estimation Error Guarantees for Distinct Values, Charikar, Chaudhuri, Motwani, Narasayya, PODS 2000
- Sampling-Based Estimation of the Number of Distinct Values of an Attribute, Haas, Naughton, Seshadri, Stokes, VLDB 1995

## **Distinct** Values

- Problem definition:
- Data set with n tuples
- Column of interest has values {1,...,D}
- Let  $n_i =$  number of times value i occurs
- $n = \sum_{i=1,D} n_i$
- Goal: estimate D, denote the estimate  $\check{D}$
- Error is  $\check{D}/D$ , or  $D/\check{D}$ , whichever is > 1

# Negative Result

**Theorem** [Charikar'00] Consider any (possibly adaptive and randomized) estimator Ď for the number of distinct values D that examines at most r rows in a table with n rows. Then, for any  $\gamma > \exp(-r)$ , there exist a choice of the input data such that with probability at least  $\gamma$ :

 $\operatorname{error}(\check{D}) \ge \operatorname{sqrt}((n-r)/2r * \ln(1/\gamma))$ 

Proof in class

## Estimators

- Goodman's unbiased estimator
- Many specialized estimators from the statistics literature (won't discuss; see [Haas'95])
- GEE [Charikar'95]; will discuss because it matches the lower bound

### Notations

- Select random sample of size r
- d=number of distinct values in the sample
- f<sub>i</sub>=number of distinct values that occur exactly i times
- Thus:  $d = \sum_{i=1,r} f_i$   $r = \sum_{i=1,r} i^* f_i$

## Goodman's Unbiased Estimator

Goodman proved in 1949 that:

• If r ≥ max(n<sub>1</sub>, ..., n<sub>D</sub>) then there exists only one unbiased estimator:

$$\widehat{D}_{\text{Good}} = d + \sum_{i=1}^{n} (-1)^{i+1} \frac{(N-r+i-1)! (r-i)!}{(N-r-1)! r!} f_i$$

- If r < max(n<sub>1</sub>, ..., n<sub>D</sub>) then there exists no unbiased estimator
- Very unstable, with errors of 20,000%

#### The GEE Estimator

**Definition** The GEE is: 
$$\check{D} = \operatorname{sqrt}(n/r) f_1 + \sum_{i=2,r} f_i$$

#### **Theorem**. Expected ratio error is O(sqrt(n/r))

## Review of this Course

Three areas in Probabilistic and Statistical Databases

- Explicit probabilities
- Implicit probabilities
- Statistics

- "Classical" probabilistic databases
- Each tuple has a probability value
  - "maybe-tuple"
  - "x-tuple"
- Possible worlds semantics

- What are some key applications ?
- What is lineage and why is it important?

- Rule of thumb 1:
  - ProbDB = IncompleteDB + Probabilities
- Rule of thumb 2:
  - ProbDB = Disjoint/IndependentDB + Joins
- Rule of thumb 3:
  - GM Factorization = DB-normalization + probidentities

Query Evaluation is #P hard in general:

- General methods: Monte Carlo, OBDDs, ...
- Safe queries and safe plans
- Top k query answering

• Major Open Research Problems [IN CLASS]

- All tuples have the same probability
- What are the major differences from explicit probabilistic data ?

- Dense random graphs  $- Pr(t) = \frac{1}{2}$
- Fagin's 0/1 law for FO

- For every sentence  $\varphi$ , lim Pr( $\varphi$ ) = 0 or =1

- "Theory of almost certain sentences" = ?
- "THE random graph" = ?

- Material random graphs:  $- Pr(t) = \beta / n^{arity(R)}$
- Every conjunctive query has an explicit asymptotic formula:

 $- Pr(q) = C(q) / n^{exp(q)} + O(n^{exp(q)+1})$ 

- General Random Graphs: G(n,p) [WHAT IS THAT ?]
- Erdos and Renyi's theorem
- Random graphs  $G(n, \beta/n^{\alpha})$ :
  - Threshold values for α (no 0/1 laws):
    2, 1+1/2, 1+1/3, ..., 1+1/k, ... 1, [rationals], 0
  - Everywhere else: 0/1 Law for FO

• The major applications today:

• ... but great theory !

-?

• Research topics: [IN CLASS]

#### Data Statistics

• What is their main usage in database systems ?

### Data Statistics

• Histograms

- Eqwidth, eqdepth, V-optimal

- Sampling
  - Sequential sampling techniques
  - Join synopses

#### Data Statistics

• Limitations of how data statistics are used today: [IN CLASS]

• Major research topics in data statistics: [IN CLASS]

# Final Thoughts

- Computer Science in the past:
   Driven by better algorithms
- Computer Science today:
  - Driven by massive amounts of data
  - Processed with approximate methods
  - Data itself is often imprecise
- Computer Science tomorrow:
  - Probabilistic databases 😳