## CSE 599s: Modern Spectral Graph Theory

Winter 2022

## Problem Set 1

Deadline: Feb 9 (at 11:59 PM) in gradescope

P1) a) Show that the cycle  $C_n$  with vertices  $1, \ldots, n$  does not satisfy local-to-global properties. In particular, construct a function  $f: [n] \to \mathbb{R}$  such that f is locally correlated,

$$\mathbb{E}_{\{i,j\} \sim \pi_1} (f(i) - f(j))^2 \le 1$$

but

$$\mathbb{E}_{i,j \sim \pi_0} (f(i) - f(j))^2 \ge \Omega(n^2).$$

- b) Given a graph G = (V, E) suppose that there are k disjoint sets  $S_1, \ldots, S_k \subseteq V$  such that  $\phi(S_i) \leq \epsilon$  for all  $1 \leq i \leq k$ . Prove that  $\lambda_k(P) \geq 1 \Omega(\epsilon)$ , where  $\lambda_k(.)$  is the k-th largest eigenvalue of P.
- c) Prove that  $\lambda_k(C_n) \geq 1 \Omega(k/n)$ .

Furthermore, assume  $P \succeq_{\pi_0} 0$ . Show that

- P2) Given a graph G = (V, E) with a random walk operator P, in class we argued that  $\phi(S) = \mathbb{P}[X_1 \notin S | X_0 \sim_{\pi_0} S]$  is the probability that a random walk started at a vertex v chosen according to  $\pi_0$  from S, leaves S in 1-step.
  - a) Let  $A \in \mathbb{R}^{n \times n}$  be a positive semi-definite matrix. Show that (for any inner-product  $\langle ., .rangle \rangle$  any unit vector  $x \in \mathbb{R}^n$  and any  $k \geq 1$ ,  $\langle Ax, x \rangle^k < \langle A^k x, x \rangle$ .
  - b) Suppose that P is the random walk operator on a graph G with respective distributions  $\pi_0, \pi_1$ .

$$\mathbb{P}[X_k \in S | X_0 \sim_{\pi_0} S] \ge (1 - \phi(S))^k.$$

c) Recall that the L2 mixing time of the (lazy) walk P on G is the smallest time t such that for any starting vertex u,

$$\left\| \frac{P^t(u, \cdot)}{\pi_0(\cdot)} - 1 \right\|_{\pi_0} = \sum_{v} \pi_0(v) \left( \frac{P^t(u, v)}{\pi_0(v)} - 1 \right)^2 \le 1/4.$$

Suppose that G has a set S with  $\pi_0(S) = 1/\sqrt{n}$  and  $\phi(S) \ll 1$ . Show that the L2 mixing time of the (lazy) random walk on G is at least  $\Omega(\frac{\log n}{\phi(S)})$ . Note that such a set is not necessarily a a barrier to L1 mixing and L1 mixing of G could be  $O(\log n)$ .

P3) Given two symmetric matrices  $A, B \in \mathbb{R}^{n \times n}$ ; recall that  $A \bullet B = \operatorname{trace}(AB)$ . Prove that

$$A \bullet B \le \sum_{i=1}^{n} \lambda_i(A) \cdot \lambda_i(B),$$

where  $\lambda_i(A)/\lambda_i(B)$  are the *i*-th smallest eigenvalues of A/B respectively.