

Lecture 4: Graph Coloring Using Path Coupling

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October 13

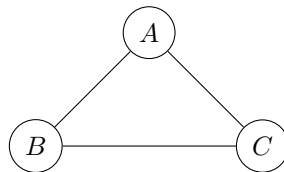
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In this lecture we use coupling to design a Markov chain to generate a uniform random proper coloring of a graph with q colors. Recall that a coloring of vertices of G with q colors is *proper* if any two adjacent vertices have distinct colors.

Let G be a graph with maximum degree Δ . Obviously, any such graph has a coloring with $q = \Delta + 1$ colors: Greedily color every vertex with a color which is not used by any of its neighbors. There is a well-known Brooks theorem which says that any graph has a coloring with Δ colors unless it has a $\Delta + 1$ clique or $\Delta = 2$ and it has an odd cycle. For $q < \Delta$ even the decision problem is NP-hard.

So, let us focus on the case where $q \geq \Delta + 1$. We want to see if the Metropolis chain mixes rapidly. Note that the chain is very simple to describe: At each time we choose a uniformly random vertex and a uniformly random color and we color that vertex with that color if possible.

It turns out that if $q = \Delta + 1$, then the above chain is not irreducible as shown below: Even though the



graph has 6 proper colorings we cannot move out of the above state. In other words, the Metropolis chain is not irreducible in this case.

First, we argue that the Metropolis chain is irreducible if $q \geq \Delta + 2$. It is enough to show that we can go from any proper coloring to any fixed proper coloring, because the chain is reversible. We show that if $q \geq \Delta + 2$ this is possible. Consider a vertex v and suppose it has a color c in the target coloring. If we can re-color v with c we just do so and we get one step closer to the target; otherwise some neighbors of v are colored c . We re-color each neighbor of v colored c with a color different from c . That is always possible because every neighbor vertex u has $\leq \Delta$ neighbors; so there is a color different from all neighbors of u and c (since $q \geq \Delta + 2$). After re-coloring all neighbors of v with a color different from c , we color v with c . Once v is colored, we can practically ignore v and repeat this until we get to the destination.

The following is a major open problem in the field of counting/sampling:

Conjecture 6.1. *The Metropolis chain mixes in time $O(n \log n)$ if $q \geq \Delta + 2$.*

It turns out that even though there have been a huge amount of efforts we are still very far from proving the above conjecture. In this lecture and the next we will discuss several of the techniques to bound mixing time for different ranges of q .

6.1 Bounding Mixing Time by Coupling

In this section we prove the following theorem.

Theorem 6.2. *For any $q \geq 4\Delta + 1$ the mixing time is $O(n \log n)$.*

Let us construct a coupling between two copies of the chain. Let X_t, Y_t be two colorings of graph G . Consider the natural metric where $d(X_t, Y_t)$ is equal to the number of vertices with different colors. We construct a coupling of X_{t+1}, Y_{t+1} such that

$$\mathbb{E}[d(X_{t+1}, Y_{t+1})] \leq \mathbb{E}[d(X_t, Y_t)] \left(1 - \frac{q - 4\Delta}{qn}\right). \quad (6.1)$$

Note that once we have the above inequality by a repeated application we can show that

$$\mathbb{E}[d(X_t, Y_t)] \leq d(X_0, Y_0) \left(1 - \frac{q - 4\Delta}{qn}\right)^t \leq e^{-\frac{(q-4\Delta)t}{qn}}.$$

Now, note that $d(X_0, Y_0) \leq n$ because in the worst case all vertices have different colors. Therefore, for $t = 8\Delta n \log n$, we have $\mathbb{E}[d(X_t, Y_t)] \leq 1/n$ which implies that $\mathbb{P}[X_t \neq Y_t] \leq 1/n$.

So, it remains to prove (6.1). We use the natural coloring, i.e., we use the same coin for both chains. Both X_t, Y_t will choose the same vertex v and try to color it with c if possible. Now, let us consider under what circumstances the distance of X_t, Y_t will decrease or increase.

- **Good Moves:** This is a move in which the distance between X_t, Y_t decreases. Suppose we choose a vertex of disagreement, say v , and we choose a color that belongs to none of the neighbors of v in X_t, Y_t . In such a case the distance decreases by 1. There are $d(X_t, Y_t)$ such vertices. Note that v has at most Δ neighbors so at most 2Δ colors are used on the neighbors of v . So, if the color c is different from these 2Δ , v will have the same color in X_{t+1}, Y_{t+1} , so $d(X_{t+1}, Y_{t+1}) = d(X_t, Y_t) - 1$. The probability of this event is $\frac{d(X_t, Y_t)(q - 2\Delta)}{qn}$.
- **Bad Moves:** This is a move in which the distance between X_t, Y_t increases. Now, suppose we choose a vertex of agreement, say v and a color c . A bad event is if c is an invalid color for v in exactly one of the chains. In such a case the distance increases by 1. Observe that such an event happens if v is a neighbor of a disagreement vertex, say u , and c is the color of u in one of the two chains. The disagreement vertices have at most $\Delta d_t(X_t, Y_t)$ neighbors, and for any such neighbors there are at 2 bad colors. Therefore,

$$\mathbb{P}[d(X_{t+1}, Y_{t+1}) = d(X_t, Y_t) + 1] \leq \frac{2\Delta d_t(X_t, Y_t)}{qn}$$

- **Neutral Moves:** In any other move the distance remains invariant.

Putting these together we have,

$$\mathbb{E}[d(X_{t+1}, Y_{t+1}) | X_t, Y_t] \leq d(X_t, Y_t) - \frac{d(X_t, Y_t)(q - 2\Delta)}{qn} + \frac{d(X_t, Y_t)2\Delta}{qn}$$

Taking expectation from both sides proves (6.1).

6.2 Path Coupling

Path coupling is a technique that simplifies the construction of couplings, when the underlying metric is a shortest path on a graph. The main consequence is that while in a normal coupling we have couple all pairs of states $x, y \in \Omega$, the path-coupling technique shows that it is enough to construct a coupling only between *neighboring states*. As we will see this simplifies the task of coupling and in many cases allows us to prove stronger mixing bounds.

Definition 6.3 (pre-metric). *A pre-metric on Ω is a weighted connected undirected graph such that for any edge (x, y) , the edge is a shortest path between x and y . In other words, the length of every other path from x to y is at least the weight of the edge (x, y) .*

Note that we can naturally extend a pre-metric to a metric by considering the shortest path metric. The path-coupling technique says that when defining the coupling it is enough to define it only for adjacent states in the pre-metric.

Theorem 6.4 ([BD97]). *Suppose there exists a coupling $(X, Y) \rightarrow (X', Y')$ defined for all adjacent pair of states in the pre-metric such that for all adjacent X, Y*

$$\mathbb{E}[d(X', Y') | X, Y] \leq (1 - \alpha)d(X, Y)$$

where d is the shortest path metric. Then, this coupling can be extended to a coupling between all pair of states that also satisfy the above inequality.

Proof. Fix a pair of states (X, Y) and let us define a coupling between X, Y . Consider the shortest path from X to Y ,

$$X = Z_0, Z_1, \dots, Z_\ell = Y.$$

We construct a coupling by composing couplings of all pairs (Z_i, Z_{i+1}) along this path.

- First we map (Z_0, Z_1) to (Z'_0, Z'_1) according to the promised coupling.
- Iteratively, for each $i \geq 1$, we map (Z_i, Z_{i+1}) to (Z'_i, Z'_{i+1}) according to the promised coupling conditioned on the Z'_i chosen in the previous step of the iteration.

Observe that the above is a valid coupling. In particular, let us just look at the Z_1, Z_2 pair. First, observe that Z'_1 is chosen in the very first step, so Z'_1 is obtained by running the Markov chain from Z_1 for one step. Conditioned on the coin tosses that sends $Z_1 \rightarrow Z'_1$, we sample Z'_2 . Note that it could be that conditioned on those coin tosses Z'_2 is deterministically defined given Z_2 . This is a valid coupling because in principal one can run any coupling this way; first run an honest step of the first chain and then by looking at the result of the coin tosses (and perhaps some additional coin tosses) move the second chain.

It follows that

$$\begin{aligned} \mathbb{E}[d(X', Y')] &\leq \sum_{i=0}^{\ell-1} \mathbb{E}[d(Z'_i, Z'_{i+1})] \\ &\leq \sum_{i=0}^{\ell-1} (1 - \alpha)d(Z_i, Z_{i+1}) \\ &= (1 - \alpha)d(X, Y). \end{aligned}$$

The first inequality is because d is a shortest path distance, the second inequality follows by the promise of the theorem and the last equality follows by the definition of a pre-metric. \square

Note that once we have the above theorem, similar to the coloring example, we can argue that the mixing time is $O(\frac{1}{\alpha} \log D)$ where D is the diameter of the shortest path metric.

6.3 Coloring using Path Coupling

In this section we prove the following theorem due to Jerrum [Jer95].

Theorem 6.5. *If $q \geq 2\Delta + 1$, then the Metropolis rule mixes in time $O(n \log n)$.*

Let us start by defining a pre-metric. Ideally, similar to the previous proof we would like to let $d(X, Y)$ be the number of vertices that are colored differently. But, we have to make sure that this is a valid shortest path metric. Note that in general we may not be able to find a path from X to Y by correcting the color of every disagreement vertex. This is because along the way we may find non-proper colorings of G .

The idea is to work with an enlarged state space. Let the state space of the Markov chain be all possible q^n colorings of G . Observe that given any state (with possibly a non-proper coloring) the Metropolis chain only assigns valid colorings to vertices. Therefore, after a number of steps we will reach a state with a proper coloring. This means that all states with non-proper coloring are *transient* states of this chain, meaning that they have probability 0 in the stationary distribution. So, the stationary distribution does not change and a coupling proof for the chain with extended state space gives an upper bound on the mixing time.

Now, let X, Y be two (not necessarily proper) coloring of G that differ in exactly one vertex v . Let $c_X(v), c_Y(v)$ be the color of v in X, Y respectively. Consider the following coupling: Pick the same vertex u in both chains, and if

- i) $u = v$: both chains pick the same color c and color u with c (if possible).
- ii) $u \in N(v)$: The X chain picks a color c . If $c = c_X(v)$, then Y chain picks $c_Y(v)$. If $c = c_Y(v)$, then Y chain picks $c_X(v)$, and otherwise the Y chain picks the color c .
- iii) $u \notin N(v)$: In this case both chains pick the same color c .

Observe that this is clearly a valid coupling.

Note that in case (iii) the distance between X, Y remains invariant. This is because u and all neighbors of u have the same color in both chains; so this is the neutral move.

In the case (i) the distance between X, Y decreases if c is different from the color of all neighbors of v . Note that since v is the only point of disagreement, there are $q - \Delta$ good colors. So, with probability $\frac{q-\Delta}{qn}$ the distance decreases by 1.

Now, let us analyze case (ii): if $c = c_X(v)$, then both options are trying to color u with invalid colors; so both actions are rejected and the distance remains invariant. If $c = c_Y(v)$ in the X chain, X tries to color u with $c_Y(v)$ and Y tries to color u with $c_X(v)$. This can indeed happen and it implies that u will have different colors. In all other cases, i.e., if $c \neq c_X(v), c_Y(v)$ both chains will do exactly the same: Either both accept or both reject c , and the distance remains invariant. This implies that

$$\mathbb{P}[d(X', Y') = 2 | X, Y] \leq \frac{\Delta}{qn}.$$

It follows that

$$\mathbb{E}[d(X', Y') | X, Y] = 1 - \frac{q - \Delta}{qn} + \frac{\Delta}{qn} = 1 - \frac{q - 2\Delta}{qn} \leq 1 - \frac{1}{qn}$$

where the last inequality follows by assuming $q \geq 2\Delta + 1$. Therefore, the chain mixes in time $O(qn \log n)$.

References

- [BD97] R. Bubley and M. Dyer. Path coupling: A technique for proving rapid mixing in markov chains. In *FOCS*, pages 223–231, 1997. 6-3
- [Jer95] M.R. Jerrum. A very simple algorithm for estimating the number of k-colourings of a low- degree graph. *Random Structures and Algorithms*, 7:157–165, 1995. 6-4