

Robotics Spring 2023

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Courtesy of Maxim Likhachev, CMU, Dieter Fox, UW

Recap: Course Overview



What we have seen so far?



What does full stack robotics involve?



Section 2 of this course

Given an accurate estimate of the state – how do we decide what actions to take?

Motion Planning

Trajectory Optimization

Certification







Task:

 find a feasible (and cost-minimal) path/motion from the current configuration of the robot to its goal configuration (or one of its goal configurations)

Two types of constraints:

- environmental constraints (e.g., obstacles)
- dynamics/kinematics constraints of the robot

Generated motion/path should (objective):

- be any feasible path
- minimize cost such as distance, time, energy, risk, ...

Examples (of what is usually referred to as path planning):













Examples (of what is usually referred to as motion planning):









Examples (of what is usually referred to as motion planning):



Piano Movers' problem

the example above is borrowed from www.cs.cmu.edu/~awm/tutorials

Examples (of what is usually referred to as motion planning):





Planned motion for a 6DOF robot arm





Why is motion planning non-trivial?

- Searching/Optimization through a complex non-convex space
- Combination of discrete/continuous optimization



Scales poorly with dimensionality of space and number of obstacles – PSPACE complete

Uncertainty and Planning

- Uncertainty can be in:
 - prior environment (i.e., door is open or closed)
 - execution (i.e., robot may slip)
 - sensing environment (i.e., seems like an obstacle but not sure)
 - pose
- Planning approaches:
 - deterministic planning:
 - assume some (i.e., most likely) environment, execution, pose
 - plan a single least-cost trajectory under this assumption
 - re-plan as new information arrives
 - planning under uncertainty:
 - associate probabilities with some elements or everything
 - -plan a policy that dictates what to do for each outcome of sensing/action and minimizes expected cost-to-goal
 - re-plan if unaccounted events happen

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re-plan every time sensory data arrives or robot deviates off its path



Uncertainty and Planning

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 - planning under uncertainty:
 - associate probabilities with some elements or everything

-plan a policy that dictates what to do for each outcome of sensing/action and minimizes expected cost-to-goal

- re-plan if unaccounted events happen

computationally MUCH harder

Example



Urban Challenge Race, CMU team, planning with Anytime D*

Lecture Outline



Defining the Motion Planning Problem

• Problem:

- Given start state xS, goal state xG
- Asked for: a sequence of control inputs that leads from start to goal
- Why tricky?
 - Need to avoid obstacles
 - For systems with underactuated dynamics: can't simply move along any coordinate at will
 - E.g., car, helicopter, airplane, but also robot manipulator hitting joint limits

Configuration Space



Configuration Space

Configuration space: space of joint configurations of the robot

Obstacles/constraints do not live in the joint space of the robot but in the world space \rightarrow non-convex when projected into configuration space

Finding collision free paths is a non-trivial search problem.



Motion Planning in Configuration Space

Cannot directly use optimization techniques like gradient descent, must solve a non-convex optimization problem.





Overview

Idea 1: Modeling as discrete search

Idea 2: Sequential convexification of non-convex problems

Planning as Search





Recasting Planning as Search



Can use efficient techniques for **discrete** graph search

Which ones?

Motion Planning via Cell Decomposition

Approximate Cell Decomposition:

- construct a graph and search it for a least-cost path



Planning via Cell Decomposition

- Approximate Cell Decomposition:
 - construct a graph and search it for a least-cost path



Planning via Cell Decomposition

- Approximate Cell Decomposition:
 - construct a graph and search it for a least-cost path
 - VERY popular due to its simplicity and representation of
 - arbitrary obstacles

- Problem: transitions difficult to execute on non-holonomic robots



How can we connect states for non-holonomic robots?

Requires solving a 2 point boundary value problem on kinematics



Differentially Constrained Mobile Robot Motion Planning

in State Lattices

s₀ and s₁ are connected if there exists a u such that $|f(s_0, u) - s_1| < \varepsilon$

On the Reachability of Quantized Control Systems

Antonio Bicchi, Alessia Marigo, Benedetto Piccoli

Mihail Pivtoraiko, Ross A. Knepper, and Alonzo Kelly Robotics Institute Carnegie Mellon University Pittsburgh, Pennsylvania 15213 e-mail: mihail@cs.cmu.edu, rak@ri.cmu.edu, alonzo@ri.cmu.edu

Received 6 August 2008; accepted 4 January 200

Can be extended to dynamics systems too!



Planning via Cell Decomposition

- Graph construction:
 - lattice graph
 - pros: sparse graph, feasible paths
 - cons: possible incompleteness



Lecture Outline



Techniques for Search



Goal is to avoid obstacles and reach a particular goal with:

- 1. As few node expansions as possible
- 2. Lowest cost path

Techniques for Search



Breadth First Search

Uniform Cost Search

A* Search







Search Attempt 1: Breadth First Search

Breadth First Search

Expand the search uniformly in all directions from start

frontier = Queue()
frontier.put(start)
came_from = dict()
came_from[start] = None

```
while not frontier.empty():
    current = frontier.get()
```

```
if current == goal: X
    break
```

```
for next in graph.neighbors(current):
    if next not in came_from:
        frontier.put(next)
        came_from[next] = current
```



Search Attempt 1: Breadth First Search

Expand the search uniformly in all directions from start



Pro: Guaranteed to find shortest paths

Cons:

- 1. Doesn't take costs into account
- 2. May expand way more nodes than necessary

Search Attempt 2: Uniform Cost Search

Expand the search according to lowest cost from the start

```
frontier = PriorityQueue()
frontier.put(start, 0)
came_from = dict()
cost_so_far = dict()
came_from[start] = None
cost_so_far[start] = 0
```

```
while not frontier.empty():
    current = frontier.get()
```

```
if current == goal:
    break
```

```
for next in graph.neighbors(current):
    new_cost = cost_so_far[current] + graph.cost(current, next)
    if next not in cost_so_far or new_cost < cost_so_far[next]:
        cost_so_far[next] = new_cost
        priority = new_cost
        frontier.put(next, priority)
        came_from[next] = current</pre>
```



Uniform Cost Search
Search Attempt 2: Uniform Cost Search

Expand the search according to lowest cost from the start



Pro: Guaranteed to find lowest cost paths

Cons: 1. May expand way more nodes than necessary

Informed Search

What if we knew some (approximate) information about how far a node is from the goal? \rightarrow Heuristics



Example: for shortest path goal reaching around obstacles, reasonable heuristics are:

1. Euclidean distance

2. Manhattan distance

Incorporate domain knowledge while always <u>underestimating</u> cost

Informed Search Attempt 1: Best-First Search

Choose the next node to expand as the one that has the lowest heuristic – "greedy best first"

```
frontier = PriorityQueue()
frontier.put(start, 0)
came from = dict()
came from[start] = None
while not frontier.empty():
   current = frontier.get()
   if current == goal:
      break
   for next in graph.neighbors(current):
      if next not in came from:
         priority = heuristic(goal, next)
         frontier.put(next, priority)
         came from[next] = current
```

Informed Search Attempt 1: Best-First Search

Greedy Best-First Search

Dijkstra's Algorithm





Con: Can return suboptimal paths with obstacles

Informed Search Attempt 2: A* Search

Choose the next node to expand as the one that has the lowest heuristic + cost so far Greedy best first Uniform cost search frontier = PriorityQueue() frontier.put(start, 0) came from = dict() cost so far = dict() came from[start] = None cost so far[start] = 0 while not frontier.empty(): current = frontier.get() if current == goal: break for next in graph.neighbors(current): new cost = cost so far[current] + graph.cost(current, next) if next not in cost_so_far or new_cost < cost_so_far[next]:</pre> cost so far[next] = new cost priority = new cost + heuristic(goal, next) frontier.put(next, priority) came from[next] = current

A* Search: Setup

Computes optimal g-values for relevant states at any point of time



g-value: shortest path so far from the start to a particular state

A* Search: Setup

Computes optimal g-values for relevant states at any point of time



one popular heuristic function – Euclidean distance

Why A* Search?

Combines the best of both greedy best first search and uniform cost search

Dijkstra's Algorithm 16 17 18 19 20 21 22 23 13 14 15 12 13 14 15 16 17 18 19 20 21 🗙 11 12 13 14 10 11 12 13 14 15 16 17 18 19 20 12 13 14 15 16 17 18 11 19 11 12 13 14 15 16 17 18 22 10 11 12 13 14 15 16 17 21 22 10 11 12 13 14 15 16 20 21 9 10 11 12 13 14 15 19 20 8 9 10 11 12 13 14 18 19 7 8 9 10 11 12 13 17 18 3 4 5 6 7 8 9 10 11 12 2 16 17 15 16 5 6 7 8 9 10 11 12 13 14 15 5 6 7 8 9 10 11 12 13 14 15 16





Small number of node expansions

Guaranteed lowest cost path (assuming positive costs)

Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s)+h(s)] from *OPEN*; insert *s* into *CLOSED*;

for every successor *s* ' of *s* such that *s* ' not in *CLOSED*

if
$$g(s') > g(s) + c(s,s')$$

 $g(s') = g(s) + c(s,s');$
insert *s*' into *OPEN*;

 $CLOSED = \{\}$ $OPEN = \{s_{start}\}$ next state to expand: s_{start}



Computes optimal g-values for relevant states

ComputePath function

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 $g(s') = g(s) + c(s,s');$
insert s' into OPEN;



 $CLOSED = \{\}$ $OPEN = \{s_{start}\}$ next state to expand: s_{start}

Computes optimal g-values for relevant states

ComputePath function

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$$g(s') > g(s) + c(s,s')$$

 $g(s') = g(s) + c(s,s');$
insert *s*' into *OPEN*;

 $CLOSED = \{s_{start}\}$ $OPEN = \{s_2\}$ next state to expand: s_2



Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s)+h(s)] from *OPEN*; insert *s* into *CLOSED*;

for every successor *s* ' of *s* such that *s* ' not in *CLOSED*

if
$$g(s') > g(s) + c(s,s')$$

 $g(s') = g(s) + c(s,s');$
insert *s*' into *OPEN*;

 $CLOSED = \{s_{start}, s_2\}$ $OPEN = \{s_1, s_4\}$ next state to expand: s_1



Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s)+h(s)] from *OPEN*; insert *s* into *CLOSED*;

for every successor *s* ' of *s* such that *s* ' not in *CLOSED*

if
$$g(s') > g(s) + c(s,s')$$

 $g(s') = g(s) + c(s,s');$
insert *s*' into *OPEN*;

 $CLOSED = \{s_{start}, s_2, s_1\}$ $OPEN = \{s_4, s_{goal}\}$ $next state to expand: s_4$



Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s)+h(s)] from *OPEN*; insert *s* into *CLOSED*;

for every successor *s* ' of *s* such that *s* ' not in *CLOSED*

 $CLOSED = \{s_{start}, s_2, s_1, s_4\}$ $OPEN = \{s_3, s_{goal}\}$ $next state to expand: s_{goal}$



Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s)+h(s)] from *OPEN*; insert *s* into *CLOSED*;

for every successor *s* ' of *s* such that *s* ' not in *CLOSED*

if
$$g(s') > g(s) + c(s,s')$$

 $g(s') = g(s) + c(s,s');$
insert s' into OPEN;

$$CLOSED = \{s_{start}, s_{2}, s_{1}, s_{4}, s_{goal}\} \begin{cases} g=0 \\ h=3 \\ S_{start} \end{cases} \begin{cases} g=0 \\ s_{2} \\ S_{2} \\ S_{2} \\ S_{2} \\ S_{3} \\ S_{3} \\ g=2 \\ h=2 \\ h=1 \\ S_{3} \\ S_{3} \\ g=5 \\ h=1 \\ S_{3} \\ S_{3}$$

Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s) + h(s)] from *OPEN*; insert *s* into *CLOSED*;

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if
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insert *s*' into *OPEN*;





Computes optimal g-values for relevant states

ComputePath function

while(s_{goal} is not expanded) remove *s* with the smallest [f(s) = g(s) + h(s)] from *OPEN*; insert *s* into *CLOSED*;

for every successor *s* ' of *s* such that *s* ' not in *CLOSED*

for every expanded state g(s) is optimal for every other state g(s) is an upper bound we can now compute a least-cost path



- Is guaranteed to return an optimal path (in fact, for every expanded state) – optimal in terms of the solution
- Performs provably minimal number of state expansions required to guarantee optimality – optimal in terms of the computations g=1 g=3h=1



 Is guaranteed to return an optimal path (in fact, for every expanded state) – optimal in terms of the solution

helps with robot deviating off its path if we search with A* backwards (from goal to start)

Performs provably minimal number of state expansions required to guarantee optimality – optimal in terms of the computations g=1g=3h=2



Connecting A* Search back to Motion Planning



A* Search: expands states in the order of f = g + h values



A* Search: expands states in the order of f = g + h values

for large problems this results in A* quickly running out of memory (memory: O(n))



Weighted A* Search: expands states in the order of *f* = *g*+*ɛh* values, *ɛ* > *1* = bias towards states that are closer to goal





Weighted A* Search: expands states in the order of $f = g + \epsilon h$ values, $\epsilon > 1$ = bias towards states that are closer to goal

20DOF simulated robotic arm state-space size: over 10²⁶ states



planning with ARA* (anytime version of weighted A*)

- planning in 8D (<x,y> for each foothold)
- heuristic is Euclidean distance from the center of the body to the goal location
- cost of edges based on kinematic stability of the robot and quality of footholds



planning with R* (randomized version of weighted A*)

joint work with Subhrajit Bhattacharya, Jon Bohren, Sachin Chitta, Daniel D. Lee, Aleksandr Kushleyev, Paul Vernaza

Is A* always optimal for all heuristics?

Admissible \rightarrow underestimate

Consistent \rightarrow monotone

 $h(s) < h^{*}(s)$

h(s) < c(s, s') + h(s')



A* search returns optimal paths on graphs only when the heuristic is admissible and consistent

Common Heuristics in Robotics

Art more than a science – commonly used heuristics are Euclidean/Manhattan distance or distance through coarse/convexified obstacles



Visualization of Search



Motion Planning via Search



Lecture Outline



Incremental version of A* (LPA*)

- Robot needs to re-plan whenever
 - new information arrives (partially-known environments or/and dynamic environments)
 - robot deviates off its path

ATRV navigating initially-unknown environment



planning map and path



Incremental version of A* (LPA*/D*/D* Lite)

- Robot needs to re-plan whenever
 - new information arrives (partially-known environments or/and dynamic environments) incremental planning (re-planning):
 - robot deviates off its path reuse of previous planning efforts planning in dynamic environments



Tartanracing, CMU

Reuse state values from previous searches cost of least-cost paths to s_{goal} initially



cost of least-cost paths to s_{goal} after the door turns out to be closed

14	13	12	11	10	9	8	7	6	6	6	6	6	6	6	6	6	6
14	13	12	11	10	9	8	7	6	5	5	5	5	5	5	5	5	5
14	13	12	11	10	9	8	7	6	5	4	4	4	4	4	4	4	4
14	13	12	11	10	9	8	7	6	5	4	3	3	3	3	3	3	3
14	13	12	11	10	9	8	7	6	5	4	3	2	2	2	2	2	3
14	13	12	11	10	9	8	7	6	5	4	3	2	1	1	1	2	3
14	13	12	11		9		7	6	5	4	3	2	1	Senal	1	2	3
					10				5	4	3	2	1	-1	1	2	3
15	14	13	12	11	11		7	6	5	4	3	2	2	2	2	2	3
15	14	13	12	12	Sstart				5	4	3	3	3	3	3	3	3
15	14	13	13	13	13		7	6	5	4	4	4	4	4	4	4	4
15	14	14	14	14	14		7	6	5	5	5	5	5	5	5	5	5
15	15	15	15	15	15		7	6	6	6	6	6	6	6	6	6	6
					16		7	7	7	7	7	7	7	7	7	7	7
21	20	19	18	17	17		8	8	8	8	8	8	8	8	8	8	8

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14	13	12	11	10	9	8	7	6	6	6	6	6	6	6	6	6	6
14	13	12	11	10	9	8	7	6	5	5	5	5	5	5	5	5	5
14	13	12	11	10	9	8	7	6	5	4	4	4	4	4	4	4	4
14	13	12	11	10	9	8	7	6	5	4	3	3	3	3	3	3	3
14	13	12	11	10	9	8	7	6	5	4	3	2	2	2	2	2	3
14	13	12	11	10	9	8	7	-6	5	4	3	2	1	1	1	2	3
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15	14	13	12	11	11		7	6	5	4	3	2	2	2	2	2	3
15	14	13	12	12	Sstart				5	4	3	3	3	3	3	3	3
15	14	13	13	13	13		7	6	5	4	4	4	4	4	4	4	4
15	14	14	14	14	14		7	6	5	5	5	5	5	5	5	5	5
15	15	15	15	15	15		7	6	6	6	6	6	6	6	6	6	6
					16		7	7	7	7	7	7	7	7	7	7	7
21	20	19	18	17	17		8	8	8	8	8	8	8	8	8	8	8

Reuse state values from previous searches cost of least-cost paths to s_{goal} initially



17	15	14	11	10	2	0	/	0	0	0	0	0	0	0	0	0	0
14	13	12	11	10	9	8	7	6	5	5	5	5	5	5	5	5	5
14	13	12	11	10	9	8	7	6	5	4	4	4	4	4	4	4	4
14	13	12	11	10	9	8	7	6	5	4	3	3	3	3	3	3	3
14	13	12	11	10	9	8	7	6	5	4	3	2	2	2	2	2	3
14	13	12	11	10	9	8	7	6	5	4	3	2	1	1	1	2	3
14	13	12	11		9		7	6	5	4	3	2	1	Secal	1	2	3
					10				5	4	3	2	1	ĩ	1	2	3
15	14	13	12	11	11		7	6	5	4	3	2	2	2	2	2	3
15	14	13	12	12	Sstart				5	4	3	3	3	3	3	3	3
15	14	13	13	13	13		7	6	5	4	4	4	4	4	4	4	4
15	14	14	14	14	14		7	6	5	5	5	5	5	5	5	5	5
15	15	15	15	15	15		7	6	6	6	6	6	6	6	6	6	6
					16		7	7	7	7	7	7	7	7	7	7	7
21	20	19	18	17	17		8	8	8	8	8	8	8	8	8	8	8

Reuse state values from previous searches cost of least-cost paths to s_{goal} initially


Motivation for Incremental Version of A*

Reuse state values from previous searches cost of least-cost paths to s_{goal} initially



Incremental Version of A*

Reuse state values from previous searches



