#### **Multicast Transport**

CSE 561 Lecture 14, Spring 2002. David Wetherall

#### **Overview**

- Last time: multicast routing
  - How to get packets from a sender to a set of receivers
- This time: multicast transport
  - What's the equivalent of "TCP" for multicast?
- Case study of two multimedia apps and protocols
  - *vic*, video conferencing; key issue: congestion control
  - *wb*, a shared whiteboard; key issue: reliability

### **Multicast Congestion Control**

- What are the bandwidth needs of multicast applications?
  - E.g., software distribution versus conferencing
  - They must still be matched to the network even if not elastic.
- Key Issue: Heterogeneity
  - Different receiver bandwidths mean no single answer is sufficient
  - So how do we match receivers w/o separate unicasting?
- Key Approach: Layered Coding
  - Send at several rates and let receivers select the best
  - Rates can carry separate or layered information

### **RLM (McCanne 95)**

- How do receivers select "the best" layers?
  - Want to avoid overwhelming the network
- One solution:
  - Imagine if routers implemented priority drop (and FQ) ...
  - Source could just send and "best layers" would fall out
  - But routers are best effort drop-tail with one class of service!
- RLM approach:
  - Have receivers learn (by join experiments) what layers suit them
  - Implement using one IP multicast group per layer

# **Binary Exponential Backoff (BEB)**

- Ethernet collisions are the classic example
  - Double interval over which retransmission timer is chosen
  - Reset interval once successful
- The technique is generally useful for adapting to an environment (e.g., network conditions)
  - TCP timeouts
  - RLM
  - SRM
  - Damping flapping links (BGP, AutoNet skeptics)?

# **Application Level Framing (ALF)**

- Clark and Tennenhouse (SIGCOMM'90)
- A design principle that calls for applications to send data in terms of units meaningful to them and lower layers to preserve these boundaries.
- Why?
  - Consider lost/reordered data
  - Consider manipulation inside the network  $\ensuremath{\textcircled{\odot}}$

### **RLM Discussion**

- Scalability
  - Shared join experiments
  - One receiver can learn when a layer will fail, but not succeed.
- Security
  - What are the interactions?

### **Multicast Reliability**

- Scaling problems: why is multicast reliability hard, different from TCP?
  - Straightforward use of ACKs doesn't scale
  - Nor do NACKs due to *implosion*
  - Centralized retransmissions become a scaling bottleneck
  - Receiver orientation if IP multicast semantics
- Approaches to distribute work and hence scale
  - Use all group members for error recovery
  - Randomization (to avoid implosion)
  - FEC/parity coding (one retransmission for different losses)

# SRM (Floyd et. al. 96)

- Approach is to *distribute* retransmissions over group
  Challenge is to minimize repair requests/responses
- Consider different topologies:
  - Chain use network distance to suppress duplicates
  - Star use randomization to suppress duplicates
  - Trees a mixture
- Adaptive learning
  - Tune timer parameters to network conditions

### **SRM Discussion**

- Scalability
  - What are the problems?
  - How well does local recovery work?
  - Do we need network support for local recovery?
- Security
  - Cooperation is an underlying assumption
  - What about incentives?