

Mobile Routing

CSE 561 Lecture 11, Spring 2002.
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The Mobility Problem: What happens if hosts move?

- Implicit assumption that Internet hosts are fixed
 - IP addresses used to name hosts; cached by higher layers
 - IP routing breaks if addresses change location. **Why?**
- Unfortunately, the buying public likes mobility



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Problems

- How does a mobile host get a local IP address?
- How do you know which IP address to use when sending to a mobile host?
- If a host moves during communication how do you know how to migrate state to the new IP address?
- Backwards compatibility

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Mobile IP: Johnson96

- **Current IETF proposed standard for mobility**
 - Dates back to research in the early 90s
 - IPv4 (RFC 2002), IPv6 version is roughly the same
- **Design constraints**
 - Network layer solution
 - Only requires changes to mobile hosts
 - Stationary hosts oblivious to mobility
 - Incrementally deployable

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Mobile IP Approach

- **Mobile Host (MH) has two addresses**
- **Home address**
 - Never changes, uniquely identifies the host
 - In “home network”
 - Correspondent host (CH) addresses all packets to the home address
- **Care-of address**
 - Will change, perhaps frequently
 - In “foreign network”
 - Related to current location (IP routing gets it to the right place)

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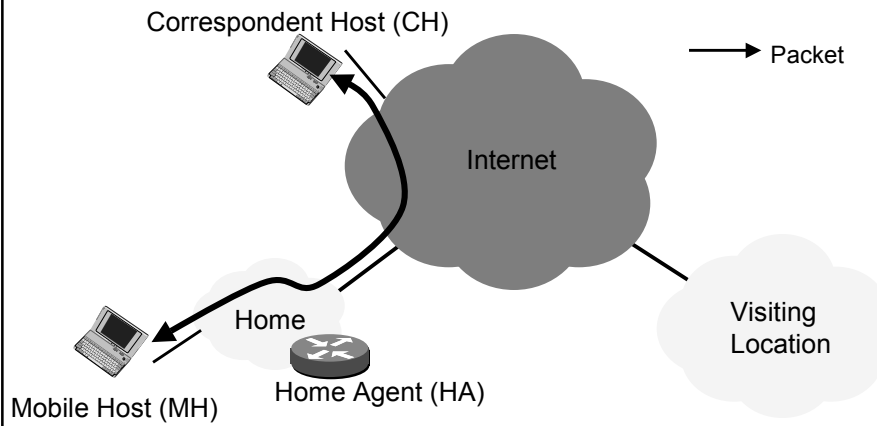
Home and Foreign Agents

- **Home agent (HA)** implements level of indirection between the mobile host and correspondents
 - Accepts traffic sent to home address
 - What about requests *from* home network?
 - Tunnels traffic to the mobile host (using care-of address)
 - And vice versa, correspondent none the wiser
- **Foreign agent (FA)** represents mobile in foreign network
 - Foreign agent can be care-of address
 - Mobile host does not need its own address in foreign network
 - Potential advantage: deal with local mobility locally

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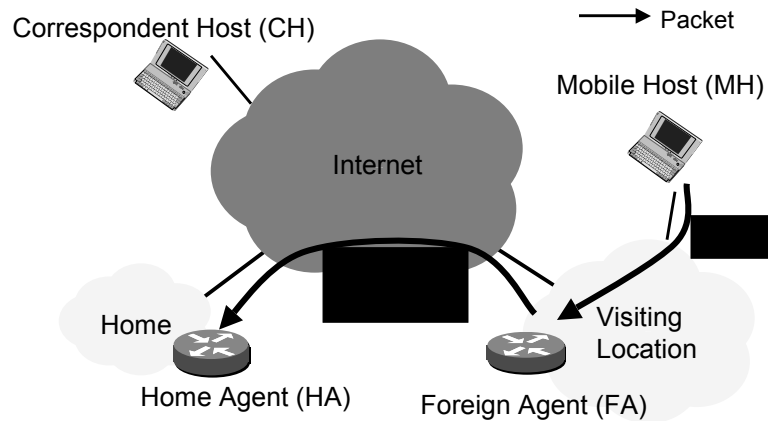
Mobile IP (MH at Home)



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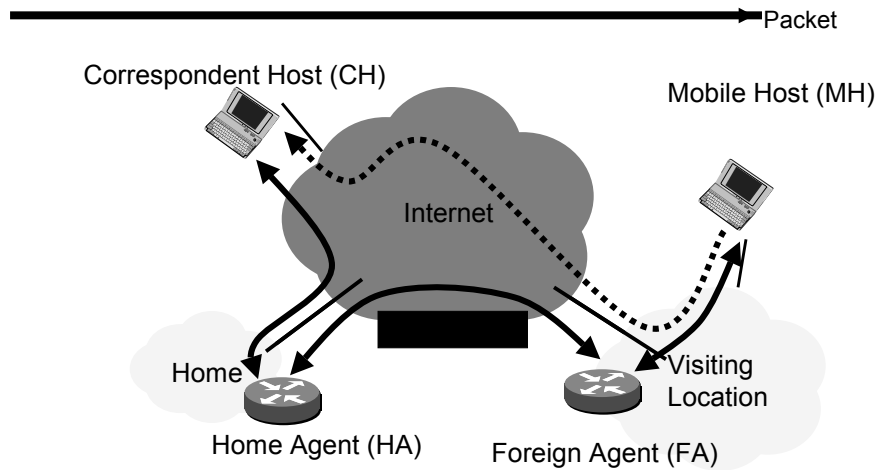
Mobile IP (MH Moving)



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Mobile IP (MH Away)



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Mobile IP Issues

- To make all this happen, a number of issues have to be addressed
 - Discovering agents
 - Registering addresses with agents (establishing bindings)
 - Authentication
 - Tunneling
 - Performance (!)

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Agent Discovery

- **Agent discovery** enables a mobile host
 - To notice when it changes networks
 - To notice when it is home again
 - When home, take down the tunnel
 - To find a foreign agent to register with
- Agents multicast **agent advertisements** locally
 - Beacons that tell the mobile who it can hear
 - Start in network A, move to network B
 - Lack of A's beacons and presence of B's tells mobile it has switched networks
- Mobile can also multicast an **agent solicitation**
- Why does multicast work here?

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Registration

- Mobiles must register care-of addresses with their home agents
 - So that the home agent knows where to tunnel packets
 - Registration needs to be updated when location changes
- **Multiple steps**
 - Registration requests first go to foreign agent, then to home agent, which replies to foreign agent, which forwards back to the mobile
- **Lifetimes**
 - Registrations have TTLs
 - Why?

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Registration Authentication

- Registration requests can be used by attackers to hijack tunnels from home agent
 - Hey, send all the mobile's traffic to me now
- Need to authenticate that a registration
 - Came from mobile host (authenticity)
 - Has not been altered (integrity)
 - Is not a replay attack (freshness)
- Mechanisms
 - Shared keys (mobile and home are from same admin domain)
 - MD5 digests
 - Nonces or timestamps

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Tunneling

- Home agent and mobile communicate using a tunnel
 - IP in IP encapsulation
- Original packet
 - Correspondent address (src) → mobile home address (dest)
 - Gets sent to home agent
- Tunnel packet
 - Encapsulates original packet
 - Home agent (src) → care-of address (dest)
 - Gets sent to foreign agent (or mobile, depending on care-of)
 - Same forward and reverse. Why no short-cut on reverse?
- Asides
 - Bit of overhead (20 byte header for every packet...poor telnet)

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Performance

- The good: No overhead in local operation
 - Home agent out of picture, no longer intercepts packets
 - The common case?
- The bad: Significant overhead in remote operation
 - Triangle routing: Packets between two hosts separated by inches can travel 1000s of miles
 - Wide-area effects determine “local” connection performance
 - The uncommon case? Even so, a steep price to pay
- Hence: Route optimization

Route Optimization

- Route optimization shortcuts the triangle
 - Correspondents can learn and use mobile care-of addresses
 - Tunnel packets directly to care-of address, skip home agent
 - Requires changes to correspondents
 - Or to routers: How likely do you think this is?
- Issues
 - Binding cache updates (consistency)
 - Binding update authentication (more trust)
 - Yet more complexity
 - Necessary for scalability?

Motivating a simple solution ...

- **Just travel with your laptop and boot it up elsewhere**
 - Essentially need a way to access the network: DHCP
- **Dynamic Host Configuration Protocol (DHCP)**
 - Request IP address dynamically (special broadcast address)
 - How do you get contacted at new IP address?
 - One solution: dynamic DNS
 - Authentication issues (who can use the local 802.11?)
 - Pro: Great over longer time scales...
 - Con: What happens during a session?

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E2E Host Mobility: Snoeren00

- **E2E Host Mobility: Same goals**
 - Do not disrupt connections when network address changes
- **Different approach**
 - Combination of DNS and connection migration
 - Naming + transport (vs. network-layer w/ Mobile IP)
 - Based upon observation of how connections are made from mobile
- **Three components**
 - Addressing
 - Locating mobile hosts
 - Connection migration

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Addressing

- Mobiles obtain an network-local IP address
 - No home agent, no home address
 - No foreign agent
 - No tunneling
 - Communication between correspondent and mobile uses addresses directly
- Problem: How does the correspondent learn the mobile's address?
 - Client case: if the mobile initiates the connection, the mobile tells the correspondent its address with the SYN packet
 - What about mobile servers?

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Locating Mobiles

- Observation: Whenever connections are established, a DNS lookup is performed
- Idea: Use the DNS lookup to return latest mobile address to correspondent
 - In Mobile IP, home address is used to unique identify mobile
 - In E2E, DNS name is used for this purpose
 - Force lookups by setting DNS response TTL to 0
 - What are the performance implications?
 - When mobile moves and obtains a new IP address, it updates its DNS entry (using secure DNS)
 - Opportunity for a race condition
 - Proposed solution: Application-level retries

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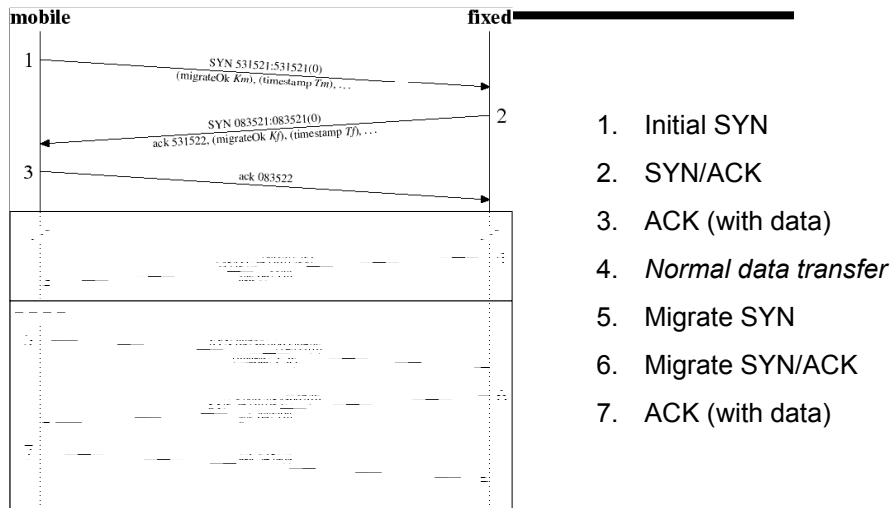
Connection Migration

- Problem: What about existing open connections?
- Solution: TCP Connection Migration
 - New IP Option: Migrate
 - Negotiated with Migrate-Permitted option in SYN
 - Requires modification to TCP stacks at both ends
 - A solution for TCP – what about other transports?

Basic Idea

- We have an open connection between correspondent (src) and mobile (dest)
 - Doesn't matter who initiated the connection
 - Connection represented by
 - <src IP, src port, dest IP, dest port>
 - Mobile moves
 - Now has new <dest IP*, dest port*>
 - Want to change connection to
 - <src IP, src port, dest IP*, dest port*>
 - Mobile creates a new connection to the correspondent, forces correspondent to migrate old connection to new one
 - Uses token to show that connections are connected

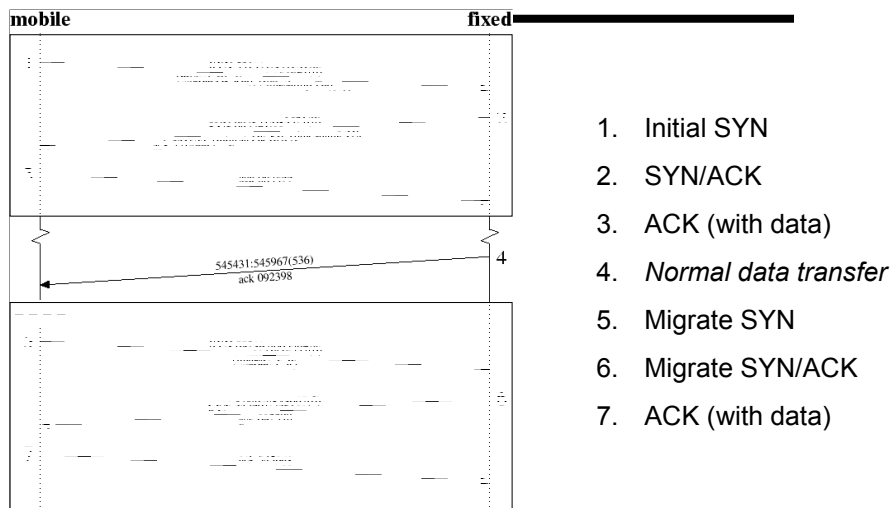
TCP Migration example



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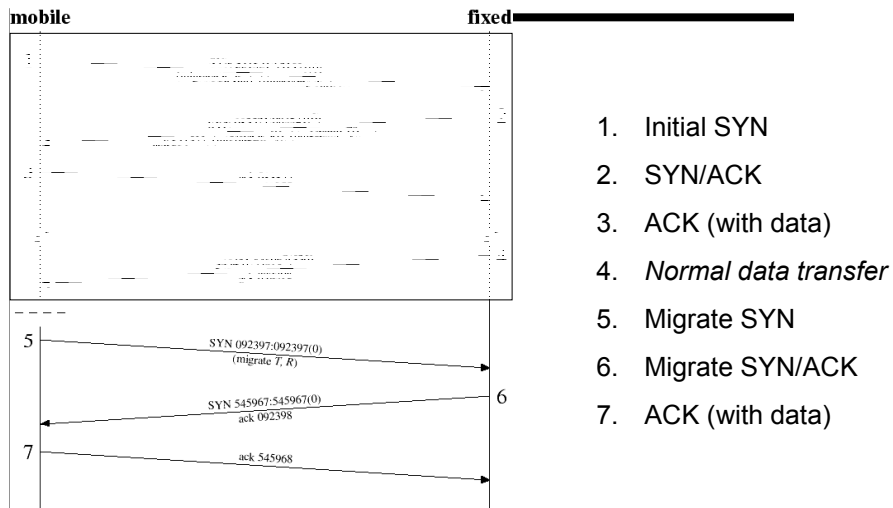
TCP Migration example



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TCP Migration example



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Issues

- **Connection Migration: Devil in the details**
 - Migrate-Permitted option: Negotiate migratability
 - Migrate option: Initiate migration
 - MIGRATE_WAIT state: race between reassignment&migration
 - Security: DoS, hijacking (cf. Mobile IP), keys, IPsec
 - Deployment: Have to change all hosts!
 - Other transport protocols: Can easily generalize?
 - No simultaneous migration (i.e. no ad hoc networks)
 - What happens if disconnected for long period (e.g. 5 mins)? Why?
- **Pluses**
 - No change to routing infrastructure
 - No triangle routes

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GPSR (Karp00)

- **Ad hoc networks vs. hybrids with fixed infrastructure**
 - How realistic? Domain is still being refined.
- **Approaches to make routing scale:**
 - Hierarchy – but not well suited to mobility
 - On-demand routes – relate to forwarding needs
 - Geography – this paper is about it work
- **GPSR: greedy forwarding isn't enough, add RHR.**
 - Simple idea, planarity complications
 - The location service is a large missing component