2 Announcements

Piazza

Recitation session:
- Review of linear algebra
  - **Location: Thursday, April 11, from 3:30-5:20 pm in SIG 134 (here)**

Deadlines next Thu, 11:59 PM:
- HW0, HW1

How to find teammates for project?
- Piazza Team Search
- Make sure you have a good dataset **accessible**

Theory of Locality Sensitive Hashing

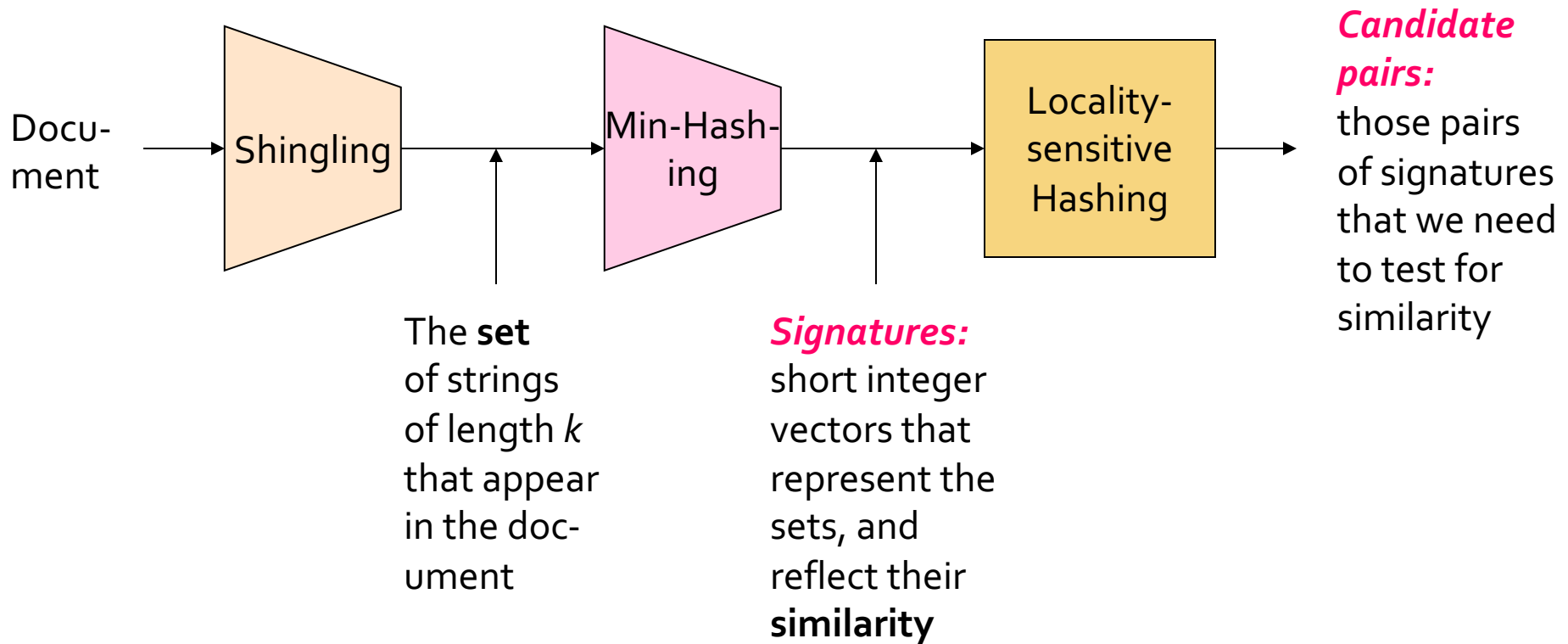**CS547 Machine Learning for Big Data**
Tim Althoff

W PAUL G. ALLEN SCHOOL
OF COMPUTER SCIENCE & ENGINEERING

Recap: Finding similar documents

- **Task:** Given a large number (N in the millions or billions) of documents, find "near duplicates"

- **Problem:**
  - Too many documents to compare all pairs

- **Solution:** Hash documents so that similar documents hash into the same bucket
  - Documents in the same bucket are then **candidate pairs** whose similarity is then evaluated

Recap: The Big Picture

Docu-
ment

$\longrightarrow$

Shingling

$\longrightarrow$

Min-Hash-
ing

$\longrightarrow$

Locality-
sensitive
Hashing

$\longrightarrow$

Candidate pairs:
those pairs of signatures that we need to test for similarity

The **set** of strings of length k that appear in the doc-ument

Signatures:
short integer vectors that represent the sets, and reflect their **similarity**

Recap: Shingles

- A *k*-shingle (or *k*-gram) is a sequence of *k* tokens that appears in the document

  - Example: **k=2**; **D$_1$** = abcab
    Set of 2-shingles: **C$_1$** = S(D$_1$) = {ab, bc, ca}

- Represent a doc by a set of hash values of its *k*-shingles

- A natural **similarity measure** is then the **Jaccard similarity:**

  $$sim(D_1, D_2) = |C_1 \cap C_2| / |C_1 \cup C_2|$$

  - Similarity of two documents is the Jaccard similarity of their shingles

Recap: Minhashing

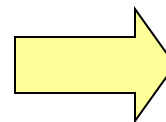- *Min-Hashing*: Convert large sets into short signatures, while preserving similarity: $Pr[h(C_1) = h(C_2)] = sim(D_1, D_2)$

Permutation π

2	4	3
3	2	4
7	1	7
6	3	2
1	6	6
5	7	1
4	5	5

Input matrix (Shingles x Documents)

1	0	1	0
1	0	0	1
0	1	0	1
0	1	0	1
0	1	0	1
1	0	1	0
1	0	1	0

Signature matrix M

2	1	2	1
2	1	4	1
1	2	1	2

Similarities of columns and signatures (approx.) match!

	1-3	2-4	1-2	3-4
Col/Col	0.75	0.75	0	0
Sig/Sig	0.67	1.00	0	0

Recap: LSH

- **Hash columns of the signature matrix *M:* Similar columns likely hash to same bucket**
  - Divide matrix *M* into ***b*** bands of ***r*** rows (M=b·r)
  - *Candidate* column pairs are those that hash to the same bucket for ≥ **1** band
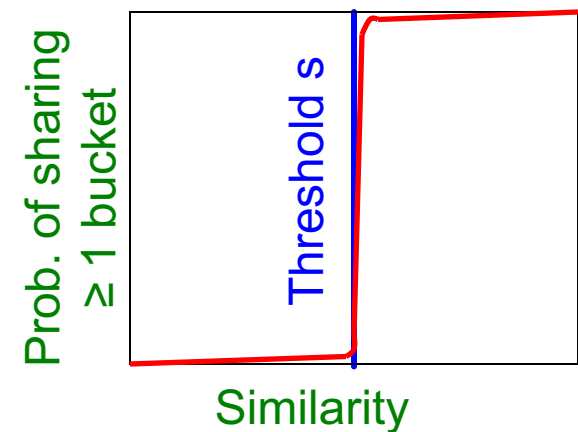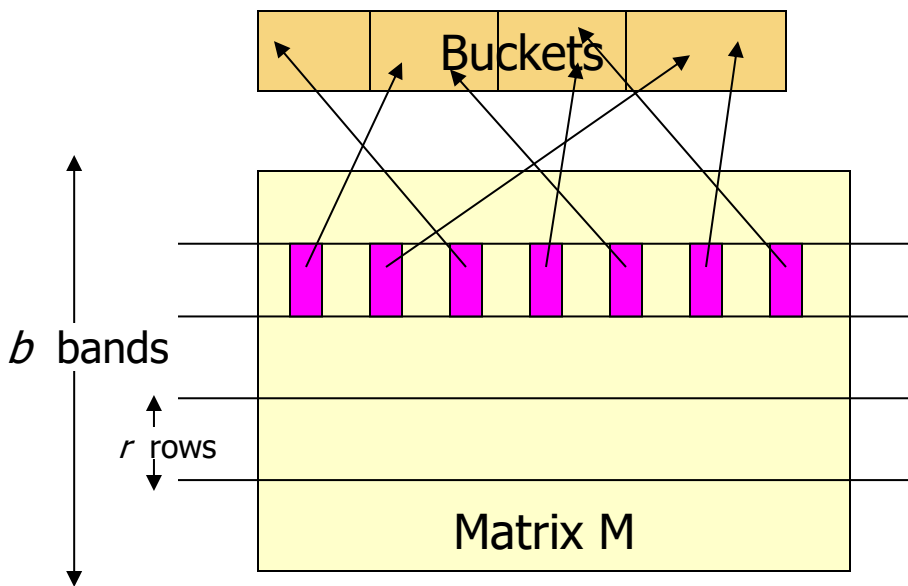
Today: Generalizing Min-hash

Signatures: short integer signatures that reflect point similarity

Candidate pairs: those pairs of signatures that we need to test for similarity

Points → **Hash func.** → Locality-sensitive Hashing →
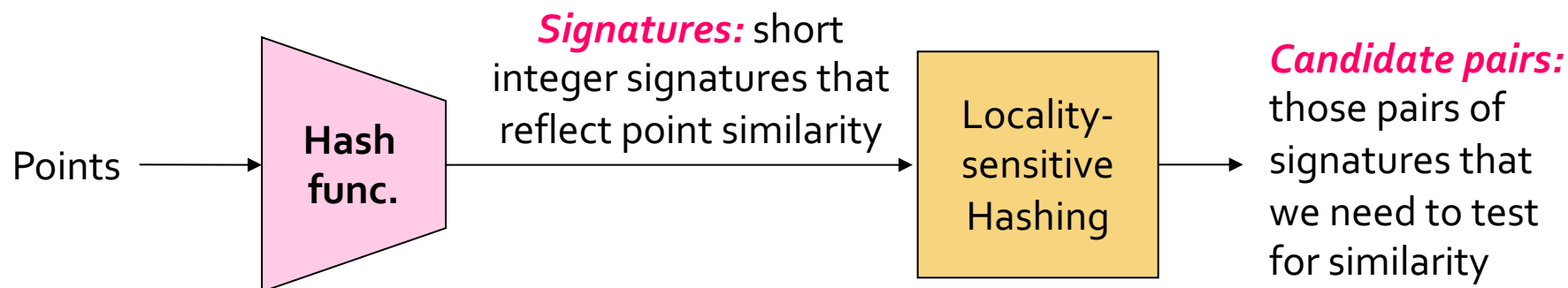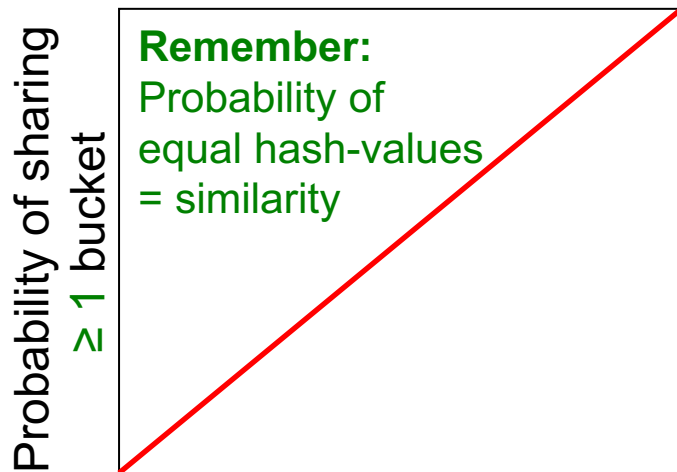
Design a **locality sensitive hash function (for a given distance metric)**

Apply the "Bands" technique

The S-Curve

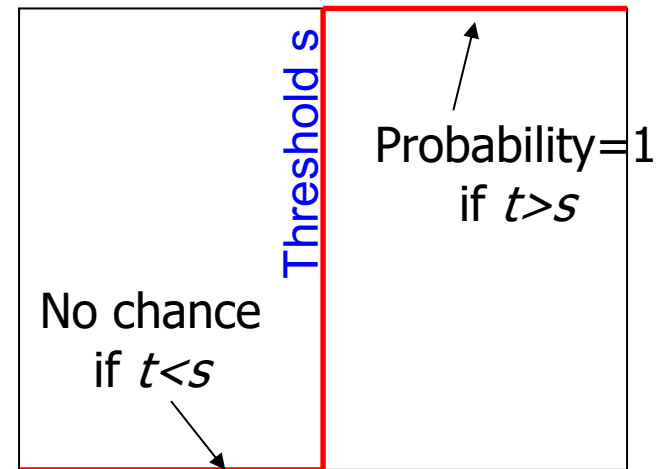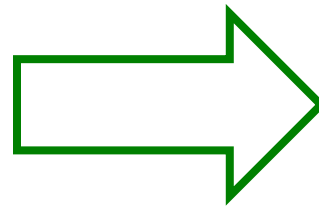- **The S-curve is where the "magic" happens**

Probability of sharing ≥ 1 bucket

Remember: Probability of equal hash-values = similarity

Similarity t of two sets

This is what 1 hash-code gives you

$$\Pr[h_\pi(C_1) = h_\pi(C_2)] = sim(D_1, D_2)$$

Threshold s

Probability=1 if $t>s$

No chance if $t<s$

Similarity t of two sets

This is what we want!
How to get a step-function?
**By choosing r and b!**

How Do We Make the S-curve?

- **Remember:** b bands, r rows/band
- Let $sim(C_1, C_2) = s$

What's the prob. that at least 1 band is equal?

- Pick some band (r rows)
  - Prob. that elements in a single row of columns C_1 and C_2 are equal $= s$
  - Prob. that all rows in a band are equal $= s^r$
  - Prob. that some row in a band is not equal $= 1 - s^r$
- Prob. that all bands are not equal $= (1 - s^r)^b$
- Prob. that at least 1 band is equal $= 1 - (1 - s^r)^b$

$$P(C_1, C_2 \text{ is a candidate pair}) = 1 - (1 - s^r)^b$$
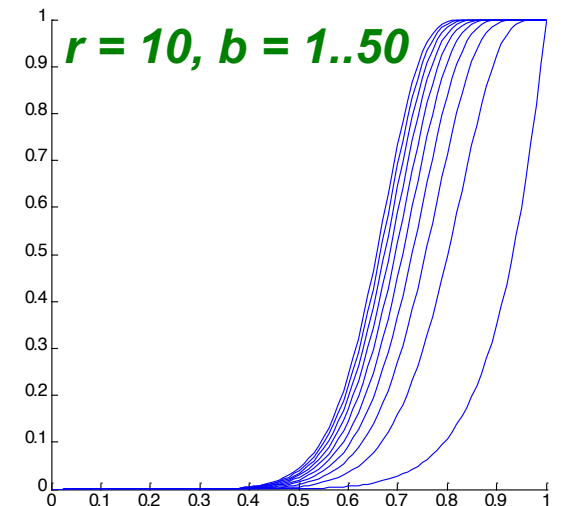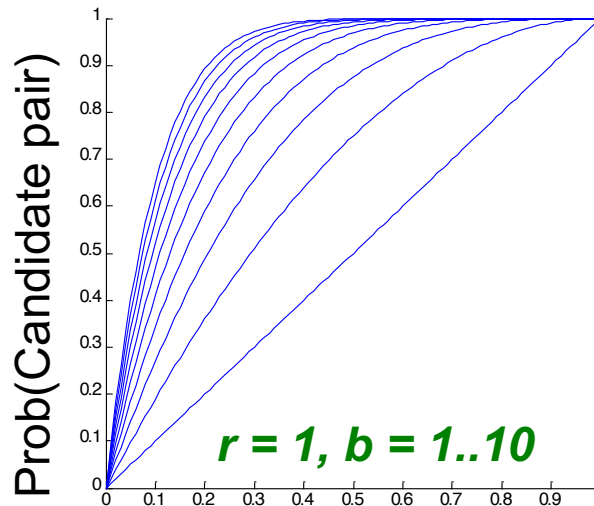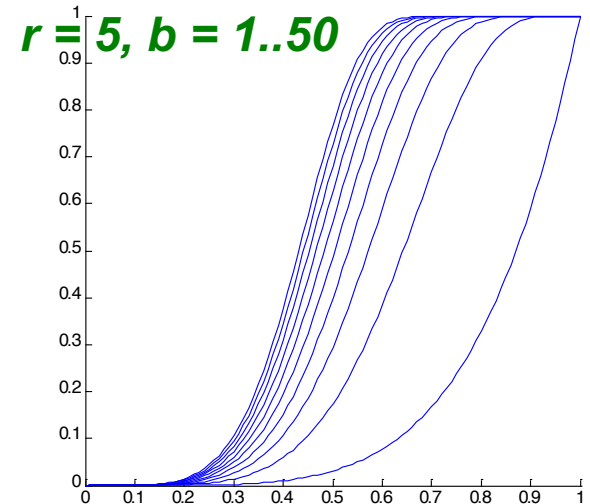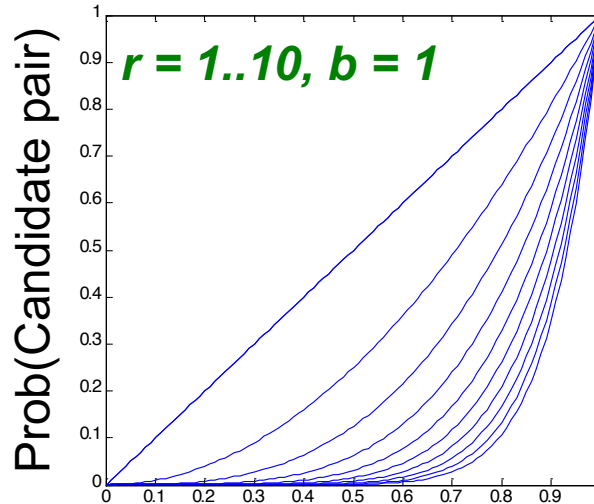
Picking *r* and *b*: The S-curve

- **Picking *r* and *b* to get the best S-curve**
  - 50 hash-functions (r=5, b=10)
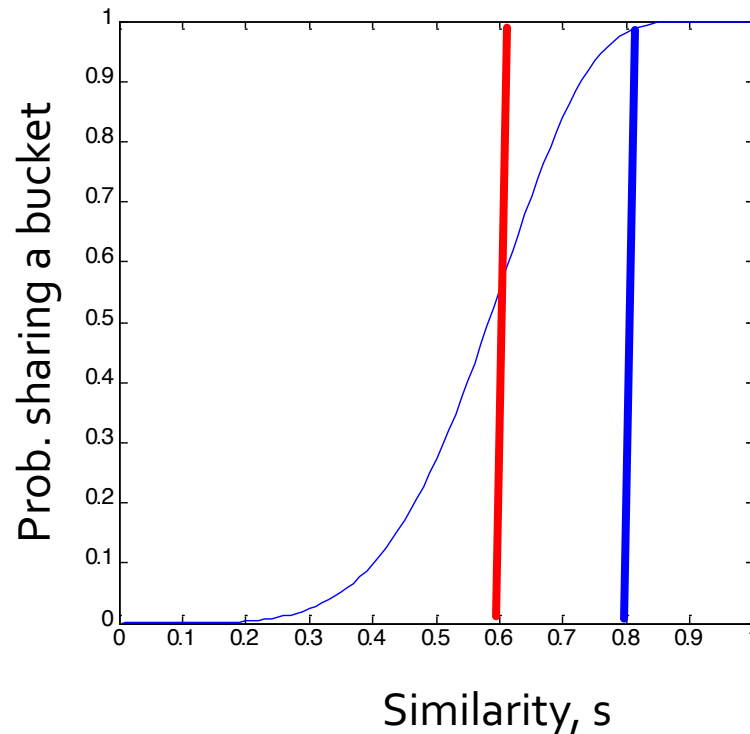
S-curves as a func. of *b* and *r*

Given a fixed threshold *s*.

We want choose *r* and *b* such that the *P(Candidate pair)* has a "step" right around *s*.

$$prob = 1 - (1 - t^{\,r})^b$$

Min-Hash-ing

Locality-sensitive Hashing

Signatures: short vectors that represent the sets, and reflect their similarity

Candidate pairs: those pairs of signatures that we need to test for similarity
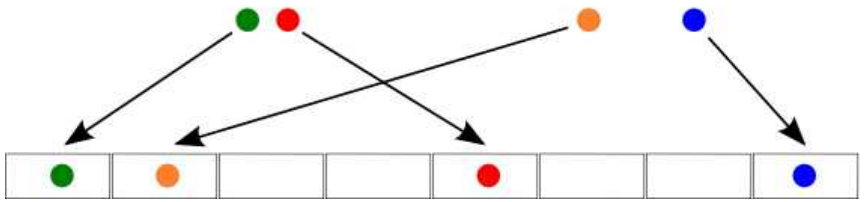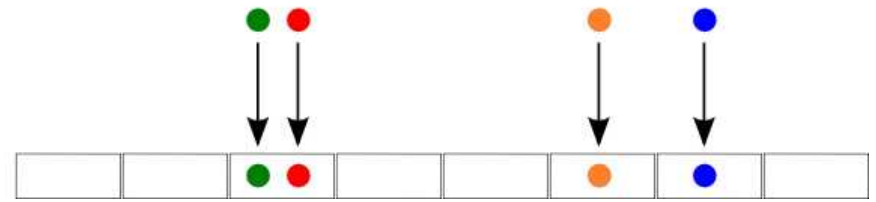
Theory of LSH

general hashing

locality-sensitive hashing

Theory of LSH

- **We have used LSH to find similar documents**
  - More generally, we found similar columns in large sparse matrices with high Jaccard similarity

- **Can we use LSH for other distance measures?**
  - e.g., Euclidean distances, Cosine distance
  - **Let's generalize what we've learned!**

Distance Measures

- **d()** is a **distance measure** if it is a function from pairs of points **x,y** to real numbers such that:
  - $d(x, y) \geq 0$
  - $d(x, y) = 0 \;\; iff \;\; x = y$
  - $d(x, y) = d(y, x)$
  - $d(x, y) \leq d(x, z) + d(z, y)$ (triangle inequality)

- Jaccard distance for sets = 1 - Jaccard similarity
- Cosine distance for vectors = angle between the vectors
- Euclidean distances:
  - L_2 *norm*: d(x,y) = square root of the sum of the squares of the differences between x and y in each dimension
    - The most common notion of "distance"
  - L_1 *norm*: sum of absolute value of the differences in each dimension
    - *Manhattan distance* = distance if you travel along coordinates only

Families of Hash Functions

- **For Min-Hashing signatures, we got a Min-Hash function for each permutation of rows**
- A "hash function" is any function that allows us to say whether two elements are **"equal"**
  - Shorthand: *h(x) = h(y)* means "*h says x and y are equal*"

- A *family* of hash functions is any set of hash functions from which we can *pick one at random efficiently*

  - **Example:** The set of Min-Hash functions generated from permutations of rows

Locality-Sensitive (LS) Families

- **Suppose we have a space *S* of points with a <u>distance</u> measure *d(x,y)***
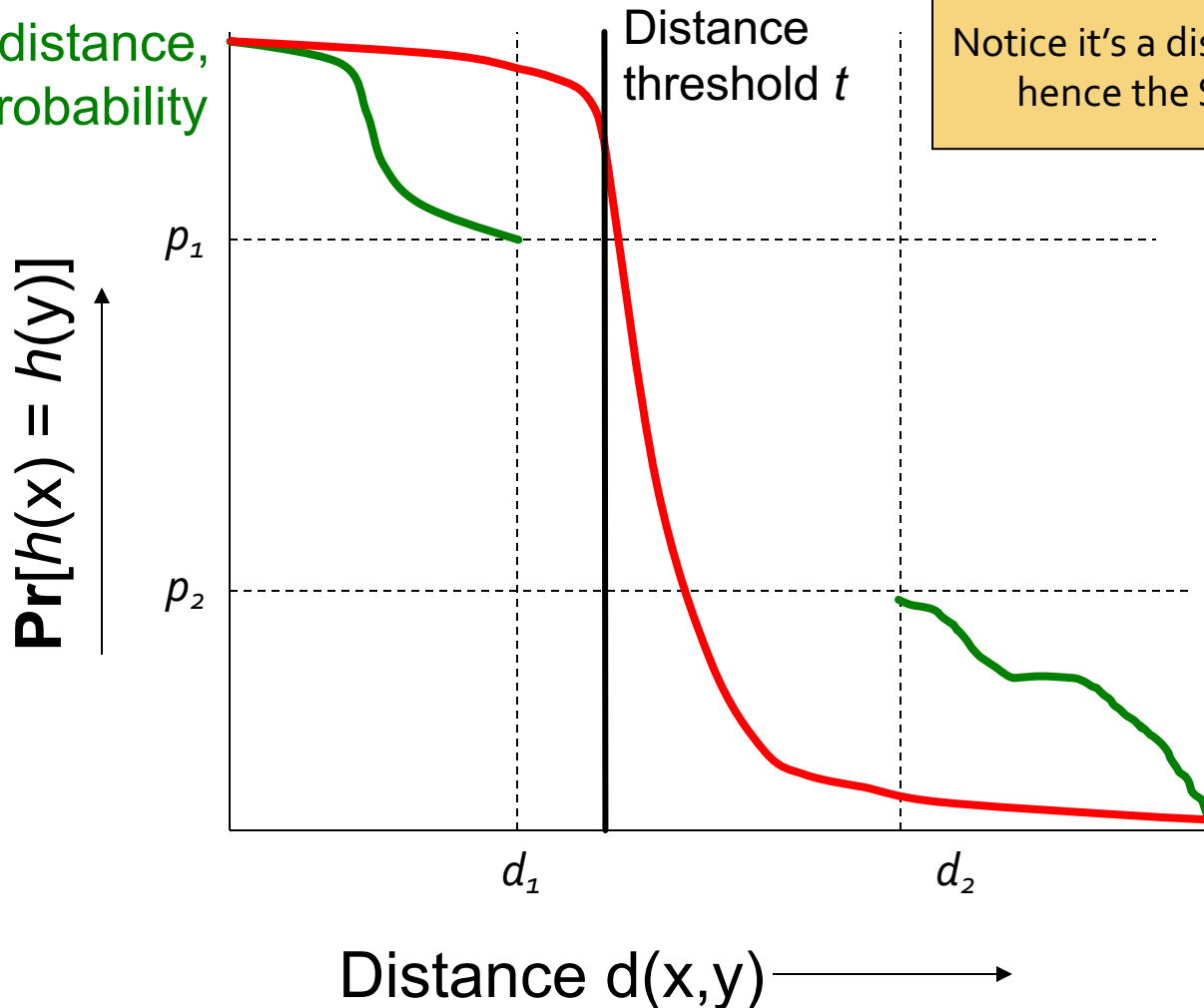
  Critical assumption

- A family *H* of hash functions is said to be (d_1, d_2, p_1, p_2)-*sensitive* if for any *x* and *y* in *S*:

  1. If $d(x, y) \leq d_1$, then the probability over all $h \in H$, that $h(x) = h(y)$ is at least p_1

  2. If $d(x, y) \geq d_2$, then the probability over all $h \in H$, that $h(x) = h(y)$ is at most p_2

With a LS Family we can do LSH!

A (d_1, d_2, p_1, p_2)-sensitive function

Small distance, high probability

Distance threshold t

Notice it's a distance, not similarity, hence the S-curve is flipped!

p_1

$\mathbf{Pr}[h(x) = h(y)]$

p_2

Large distance, low probability of hashing to the same value

d_1

d_2

Distance d(x,y)

Example of LS Family: Min-Hash

- **Let:**
  - *S* = space of all sets,
  - *d* = Jaccard distance,
  - *H* is family of Min-Hash functions for all permutations of rows
- Then for any hash function $h \in H$:

$$Pr[h(x) = h(y)] = 1 - d(x, y)$$

  - Simply restates theorem about Min-Hashing in terms of distances rather than similarities

Example: LS Family – (2)

- **Claim: Min-hash H is a (1/3, 2/3, 2/3, 1/3)-sensitive family for S and d.**

If distance $\leq$ 1/3 (so similarity $\geq$ 2/3)

Then probability that Min-Hash values agree is $\geq$ 2/3

- **For Jaccard similarity, Min-Hashing gives a $(d_1, d_2, (1-d_1), (1-d_2))$-sensitive family for any $d_1 < d_2$**

Amplifying a LS-Family

Prob. of sharing a bucket

Similarity t

- **Can we reproduce the "S-curve" effect we saw before for any LS family?**

- The "**bands**" technique we learned for signature matrices carries over to this more general setting
- **Can do LSH with any (d_1, d_2, p_1, p_2)-sensitive family!**

- Two constructions:
  - **AND** construction like "rows in a band"
  - **OR** construction like "many bands"
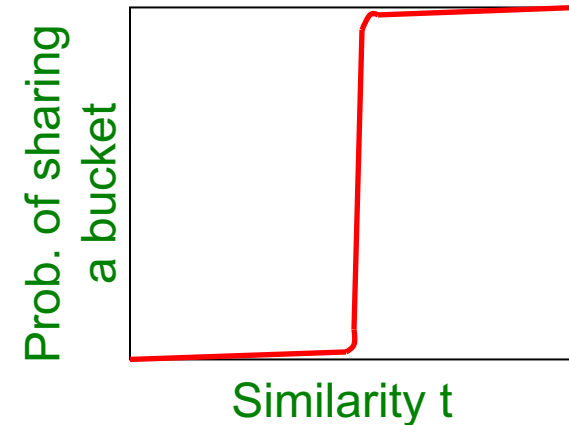
Amplifying Hash Functions: AND and OR

AND of Hash Functions

- Given family **H**, construct family **H'** consisting of **r** functions from **H**
- For $h = [h_1,\ldots,h_r]$ in **H'**, we say $h(x) = h(y)$ if and only if $h_i(x) = h_i(y)$ for **all** i $\quad 1 \le i \le r$
  - Note this corresponds to creating a band of size **r**

- **Theorem:** If **H** is (d_1, d_2, p_1, p_2)-sensitive, then **H'** is $(d_1, d_2, (p_1)^r, (p_2)^r)$-sensitive
- **Proof:** Use the fact that h_i's are **independent**

Also lowers probability for small distances (Bad)

Lowers probability for large distances (Good)

Subtlety Regarding Independence

- **Independence of hash functions (HFs) really means that the prob. of two HFs saying "yes" is the product of each saying "yes"**
  - **But** two particular hash functions could be highly correlated
    - For example, in Min-Hash if their permutations agree in the first one million entries
  - **However**, the probabilities in definition of a LSH-family are over all possible members of *H*, *H'* (i.e., average case and not the worst case)
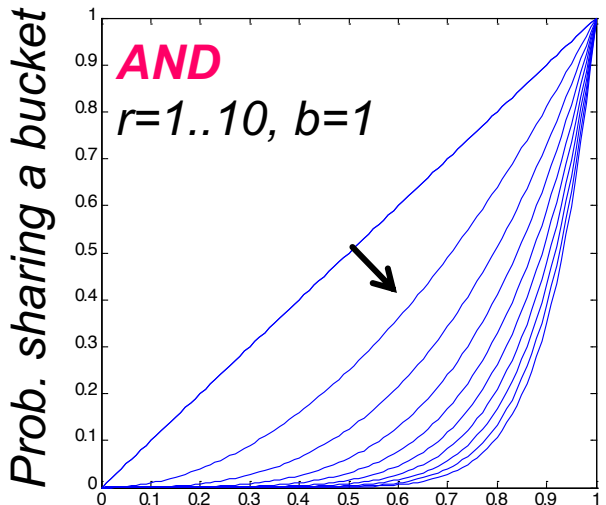
OR of Hash Functions

- Given family *H*, construct family *H'* consisting of *b* functions from *H*

- For $h = [h_1, ..., h_b]$ in *H'*,
  $h(x) = h(y)$ if and only if $h_i(x) = h_i(y)$ for **at least 1** *i*

- **Theorem:** If *H* is (d_1, d_2, p_1, p_2)-sensitive,
  then *H'* is $(d_1, d_2, 1-(1-p_1)^b, 1-(1-p_2)^b)$-sensitive

- **Proof:** Use the fact that h_i's are **independent**

Raises probability for small distances (Good)

Raises probability for large distances (Bad)

Effect of AND and OR Constructions

- **AND** makes all probs. **shrink**, but by choosing *r* correctly, we can make the lower prob. approach 0 while the higher does not

- **OR** makes all probs. **grow**, but by choosing *b* correctly, we can make the higher prob. approach 1 while the lower does not

AND
r=1..10, b=1

Prob. sharing a bucket

Similarity of a pair of items

OR
r=1, b=1..10

Prob. sharing a bucket
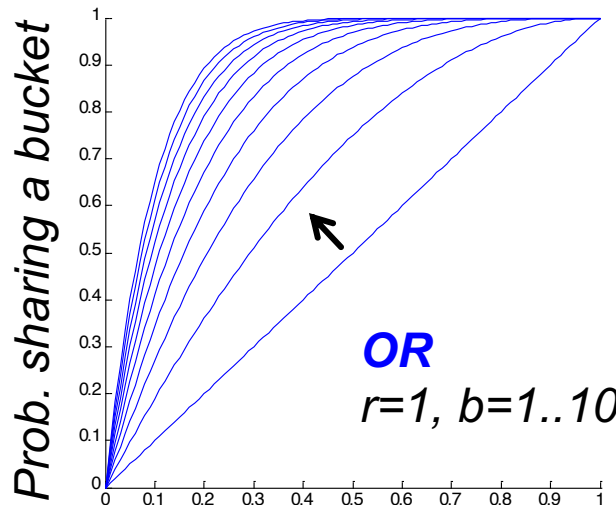
*Similarity of a pair of items*

Combine AND and OR Constructions

- By choosing **b** and **r** correctly, we can make the lower probability approach 0 while the higher approaches 1

- As for the signature matrix, we can use the AND construction followed by the OR construction
  - Or vice-versa
  - Or any sequence of AND's and OR's alternating

Composing Constructions

- *r*-way **AND** followed by *b*-way **OR** construction
  - **Exactly what we did with Min-Hashing**
    - **AND:** If bands match in **all *r*** values hash to same bucket
    - **OR:** Cols that have $\geq$ 1 common bucket $\rightarrow$ **Candidate**

- Take points *x* and *y* s.t. *Pr[h(x) = h(y)] = s*
  - *H* will make **(x,y)** a candidate pair with prob. **s**
- Construction makes **(x,y)** a candidate pair with probability *$1-(1-s^r)^b$*   **The S-Curve!**
  - **Example:** Take **H** and construct **H'** by the **AND** construction with *r* = 4. Then, from **H'**, construct **H''** by the **OR** construction with *b* = 4

Table for Function $1-(1-s^4)^4$

s	$p=1-(1-s^4)^4$
.2	**.0064**
.3	.0320
.4	.0985
.5	.2275
.6	.4260
.7	.6666
.8	**.8785**
.9	.9860
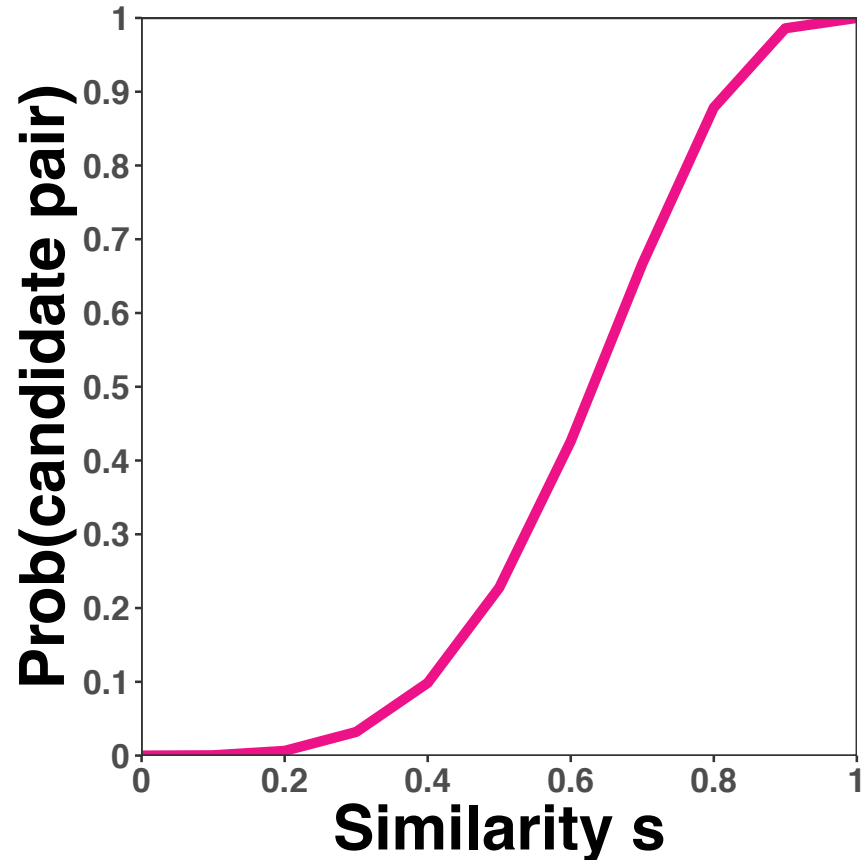
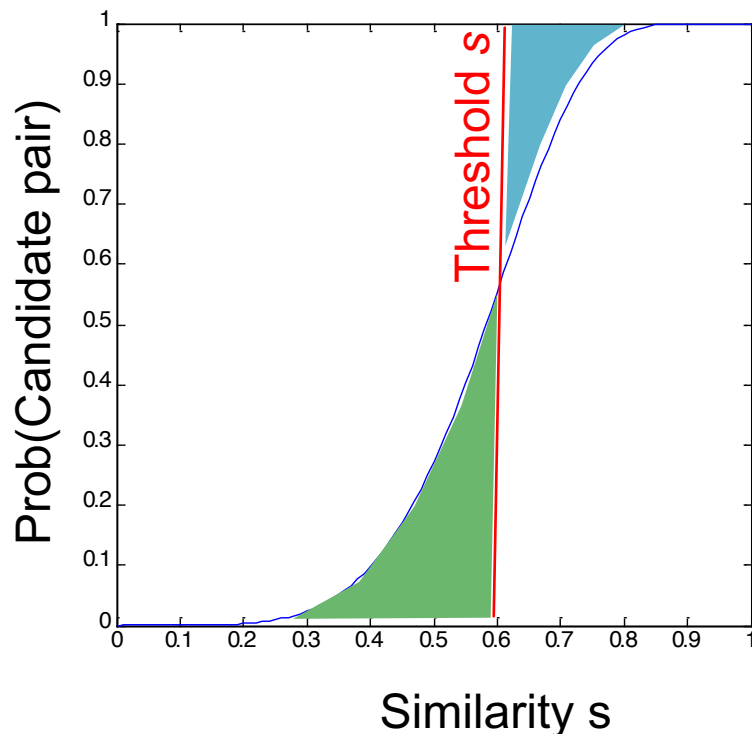*r = 4, b = 4* transforms a (.2,.8,.8,.2)-sensitive family into a (.2,.8,.8785,.0064)-sensitive family.

How to choose *r* and *b*

Picking *r* and *b*: The S-curve

- ## Picking *r* and *b* to get desired performance
  - ### 50 hash-functions (*r = 5, b = 10*)

Blue area *X*: False Negative rate
These are pairs with ***sim > s*** but the ***X***
fraction won't share a band and then
will **never become candidates.** This
means we will never consider these
pairs for (slow/exact) similarity
calculation!

Green area Y: False Positive rate
These are pairs with ***sim < s*** but
we will consider them as candidates.
This is not too bad, we will consider
them for (slow/exact) similarity
computation and discard them.

Picking *r* and *b*: The S-curve

- **Picking *r* and *b* to get desired performance**
  - **50 hash-functions (*r* * *b* = 50)**

r=2, b=25
r=5, b=10
r=10, b=5
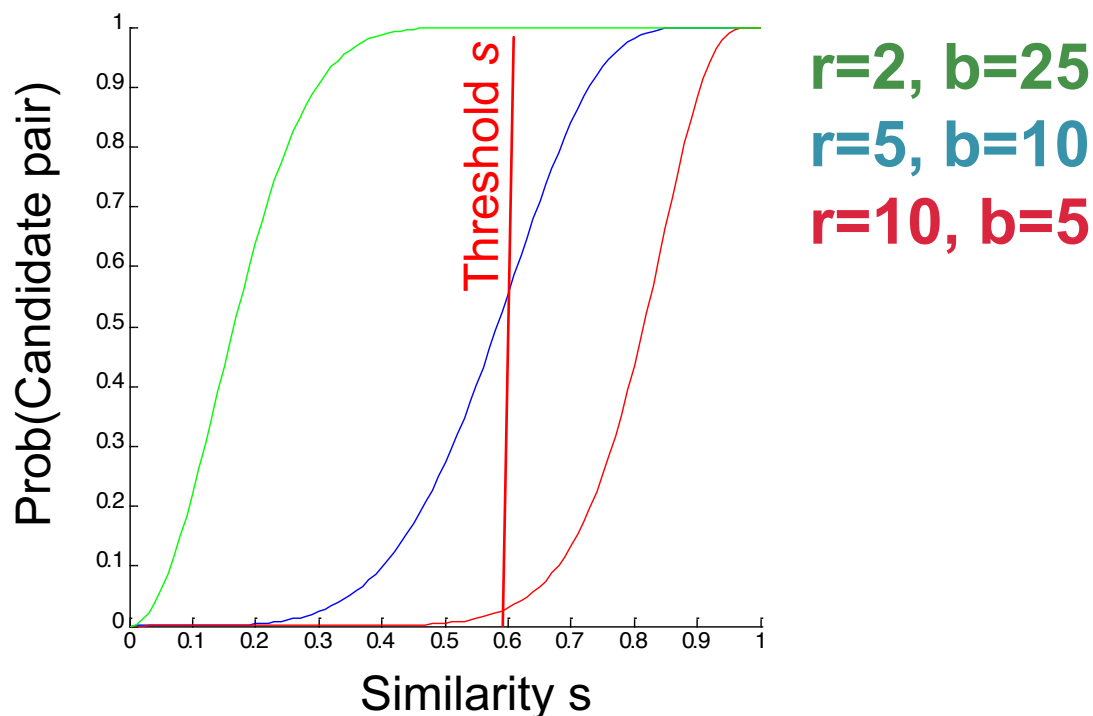
OR-AND Composition
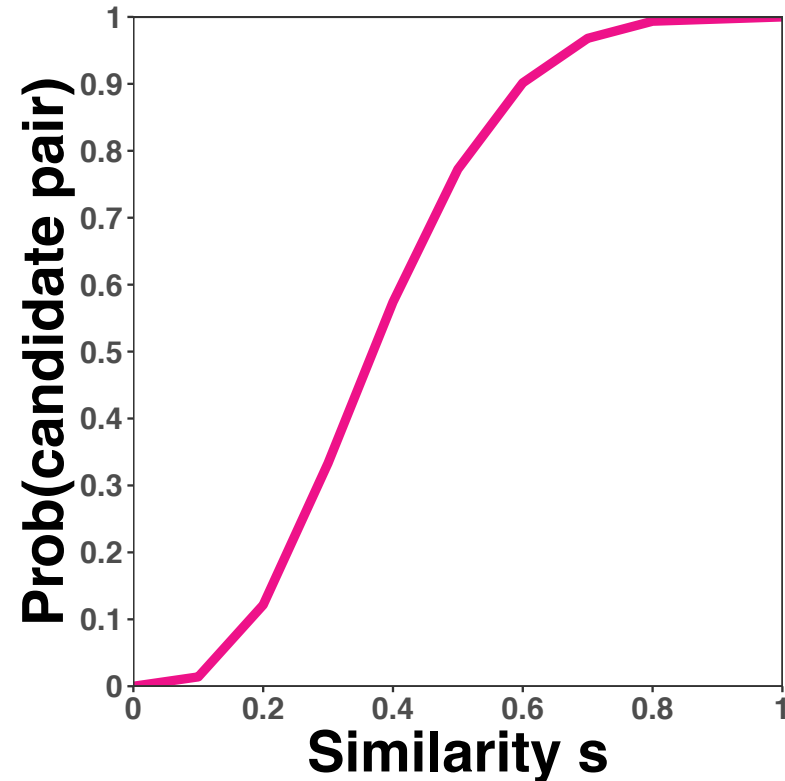
- Apply a **b**-way **OR** construction followed by an **r**-way **AND** construction
- Transforms similarity **s** (probability p) into **(1-(1-s)^b)^r**
  - **The same S-curve, mirrored horizontally and vertically**

- **Example:** Take **H** and construct **H'** by the **OR** construction with **b** = 4. Then, from **H'**, construct **H''** by the **AND** construction with **r** = 4

Table for Function $(1-(1-s)^4)^4$

s	$p=(1-(1-s)^4)^4$
.1	.0140
.2	**.1215**
.3	.3334
.4	.5740
.5	.7725
.6	.9015
.7	.9680
.8	**.9936**

The example transforms a (.2,.8,.8,.2)-sensitive family into a (.2,.8,.9936,.1215)-sensitive family

Cascading Constructions

- **Example:** Apply the **(4,4) OR-AND** construction followed by the **(4,4) AND-OR** construction

- **Transforms a (.2, .8, .8, .2)-sensitive family into a (.2, .8, .9999996, .0008715)-sensitive family**

  - **Note this family uses 256 (=4*4*4*4) of the original hash functions**

General Use of S-Curves

- Fixpoint: For each AND-OR S-curve $1-(1-s^r)^b$, there is a *threshold* t, for which $1-(1-t^r)^b = t$
- Above t, high probabilities are increased; below t, low probabilities are decreased
- You improve the sensitivity as long as the low probability is less than t, and the high probability is greater than t
  - Iterate as you like
- Similar observation for the OR-AND type of S-curve: $(1-(1-s)^b)^r$

Visualization of Threshold

Prob(Candidate pair)

Probability
Is raised

Threshold
t

Probability
Is lowered
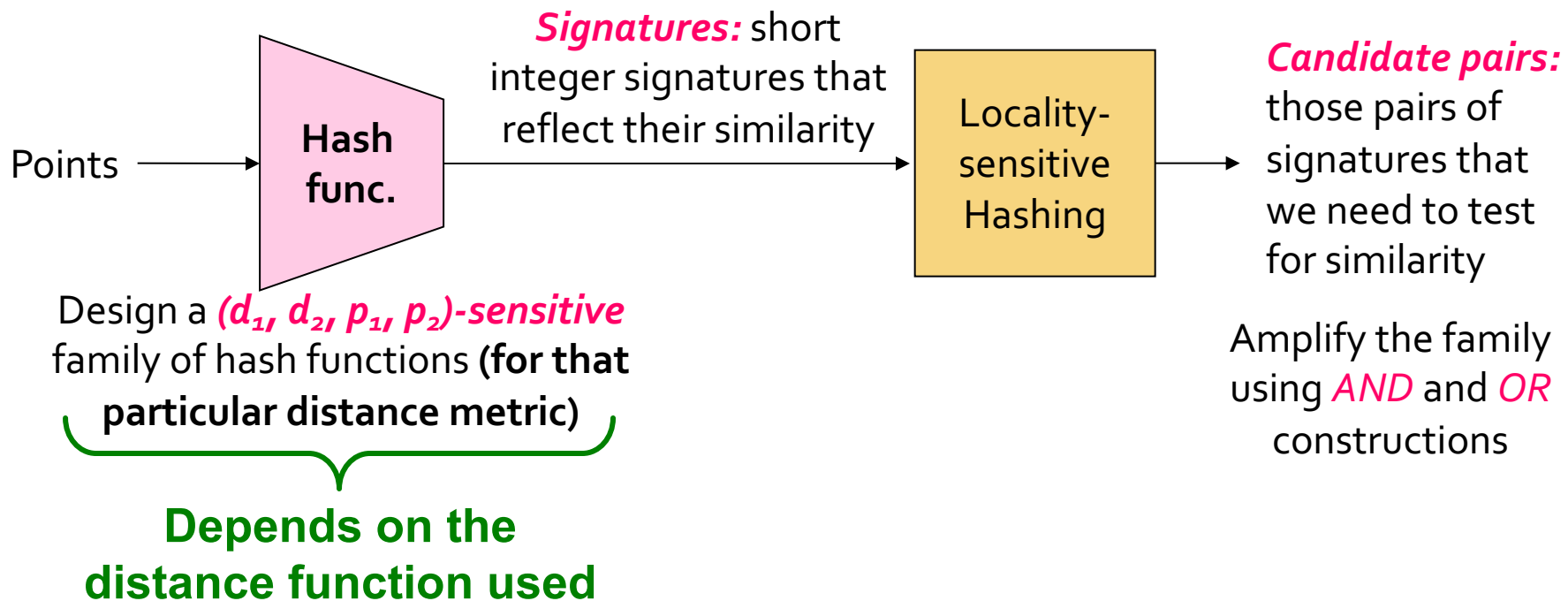
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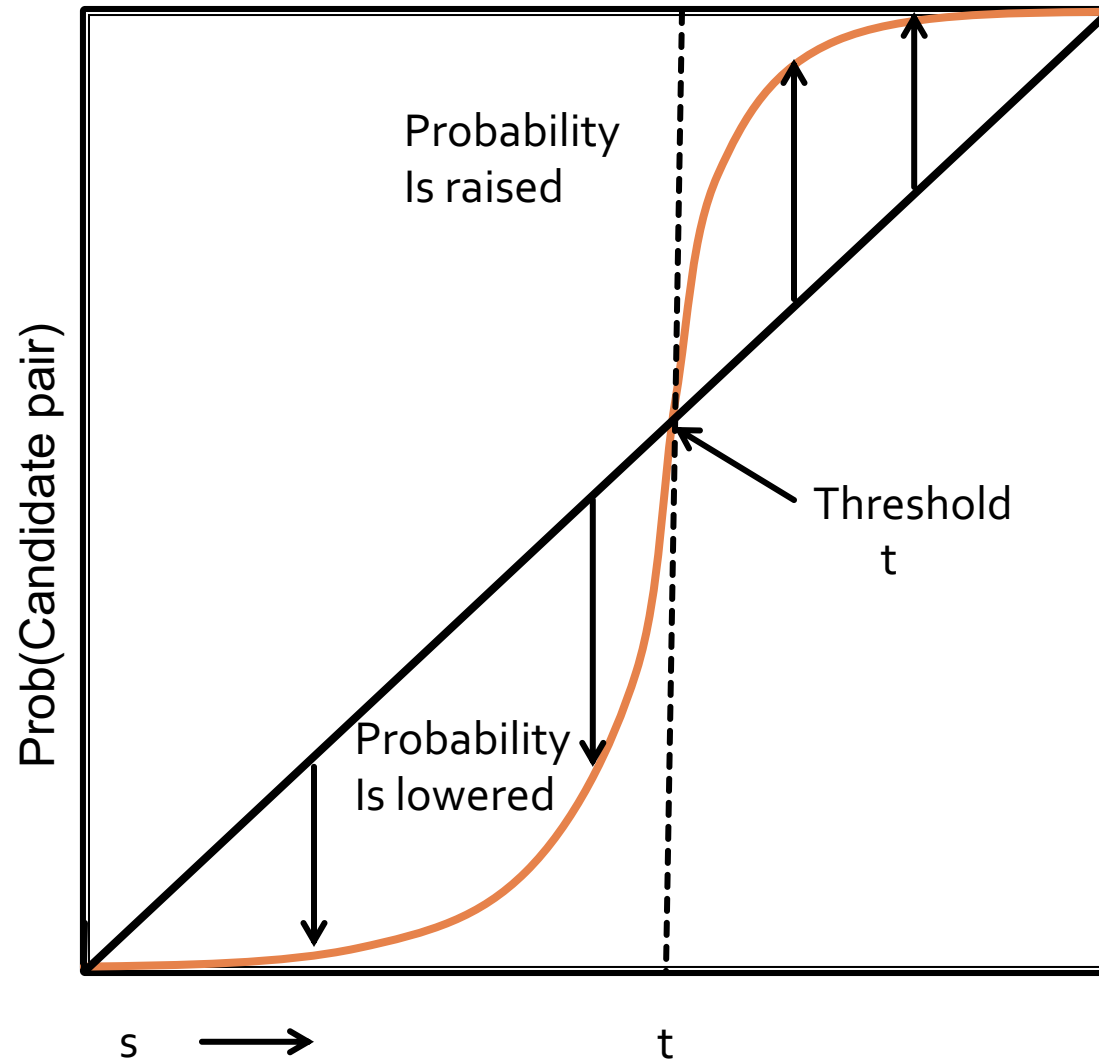
Summary

- Pick any two distances $d_1 < d_2$

- Start with a $(d_1, d_2, (1- d_1), (1- d_2))$-sensitive family

- Apply constructions to **amplify** (d_1, d_2, p_1, p_2)-sensitive family, where p_1 is almost 1 and p_2 is almost 0

- **The closer to 0 and 1 we get, the more hash functions must be used!**
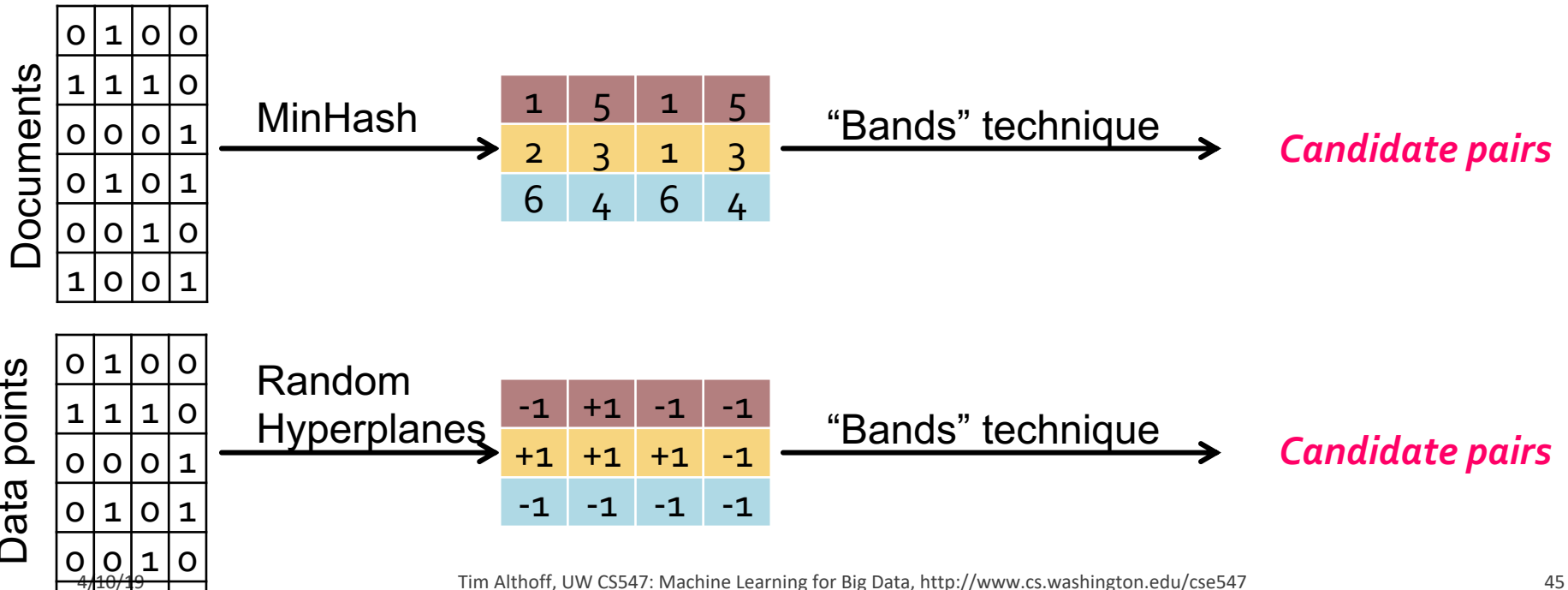
LSH for other distance metrics

LSH for other Distance Metrics

- **LSH methods for other distance metrics:**
  - **Cosine distance: Random hyperplanes**
  - **Euclidean distance: Project on lines**

Points →

Hash func.

Signatures: short integer signatures that reflect their similarity →

Locality-sensitive Hashing →

Candidate pairs: those pairs of signatures that we need to test for similarity

Design a (d_1, d_2, p_1, p_2)-*sensitive* family of hash functions **(for that particular distance metric)**

Amplify the family using *AND* and *OR* constructions

Depends on the distance function used

Summary of what we will learn

Signatures: short integer signatures that reflect their similarity

Data → **Hash func.** → Locality-sensitive Hashing → **Candidate pairs:** those pairs of signatures that we need to test for similarity

Documents:

0	1	0	0
1	1	1	0
0	0	0	1
0	1	0	1
0	0	1	0
1	0	0	1

MinHash →

1	5	1	5
2	3	1	3
6	4	6	4

"Bands" technique → *Candidate pairs*

Data points:

0	1	0	0
1	1	1	0
0	0	0	1
0	1	0	1
0	0	1	0

Random Hyperplanes →

-1	+1	-1	-1
+1	+1	+1	-1
-1	-1	-1	-1

"Bands" technique → *Candidate pairs*
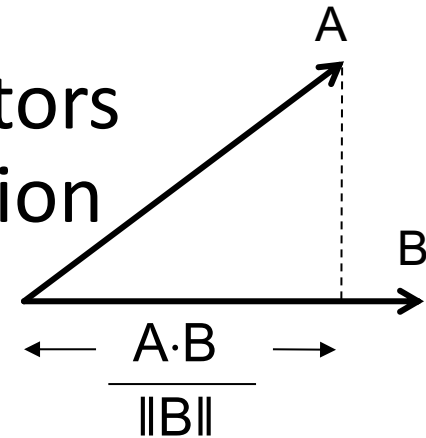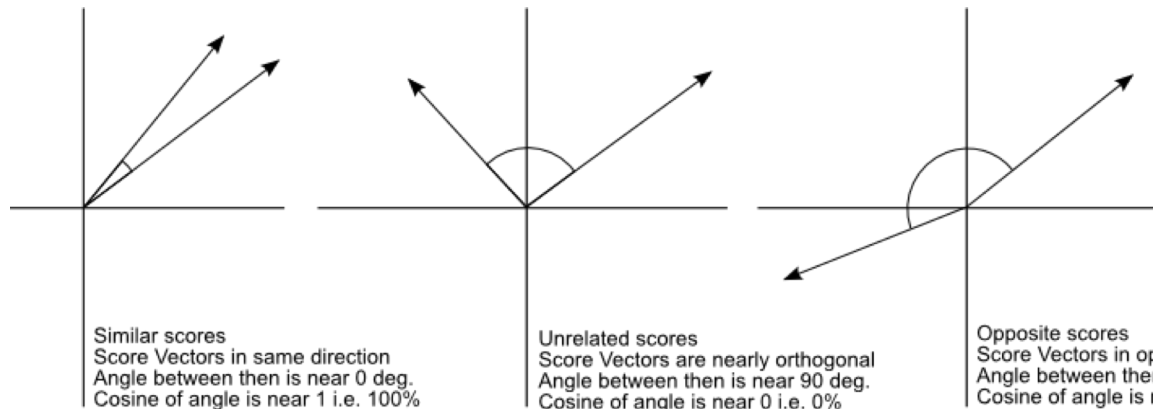
Cosine Distance

■ ***Cosine distance*** = angle between vectors from the origin to the points in question
d(A, B) = θ = arccos(A·B / ‖A‖·‖B‖)

■ Has range $[0, \pi]$ (equivalently [0,180°])

■ Can divide θ by π to have distance in range [0,1]

■ **Cosine similarity = 1-d(A,B)**

■ But often defined as **cosine sim:** $\cos(\theta) = \dfrac{A \cdot B}{\|A\| \|B\|}$

- Has range -1…1 for general vectors
- Range 0..1 for non-negative vectors (angles up to 90°)

Similar scores
Score Vectors in same direction
Angle between then is near 0 deg.
Cosine of angle is near 1 i.e. 100%

Unrelated scores
Score Vectors are nearly orthogonal
Angle between then is near 90 deg.
Cosine of angle is near 0 i.e. 0%

Opposite scores
Score Vectors in opposite direction
Angle between then is near 180 deg.
Cosine of angle is near -1 i.e. -100%
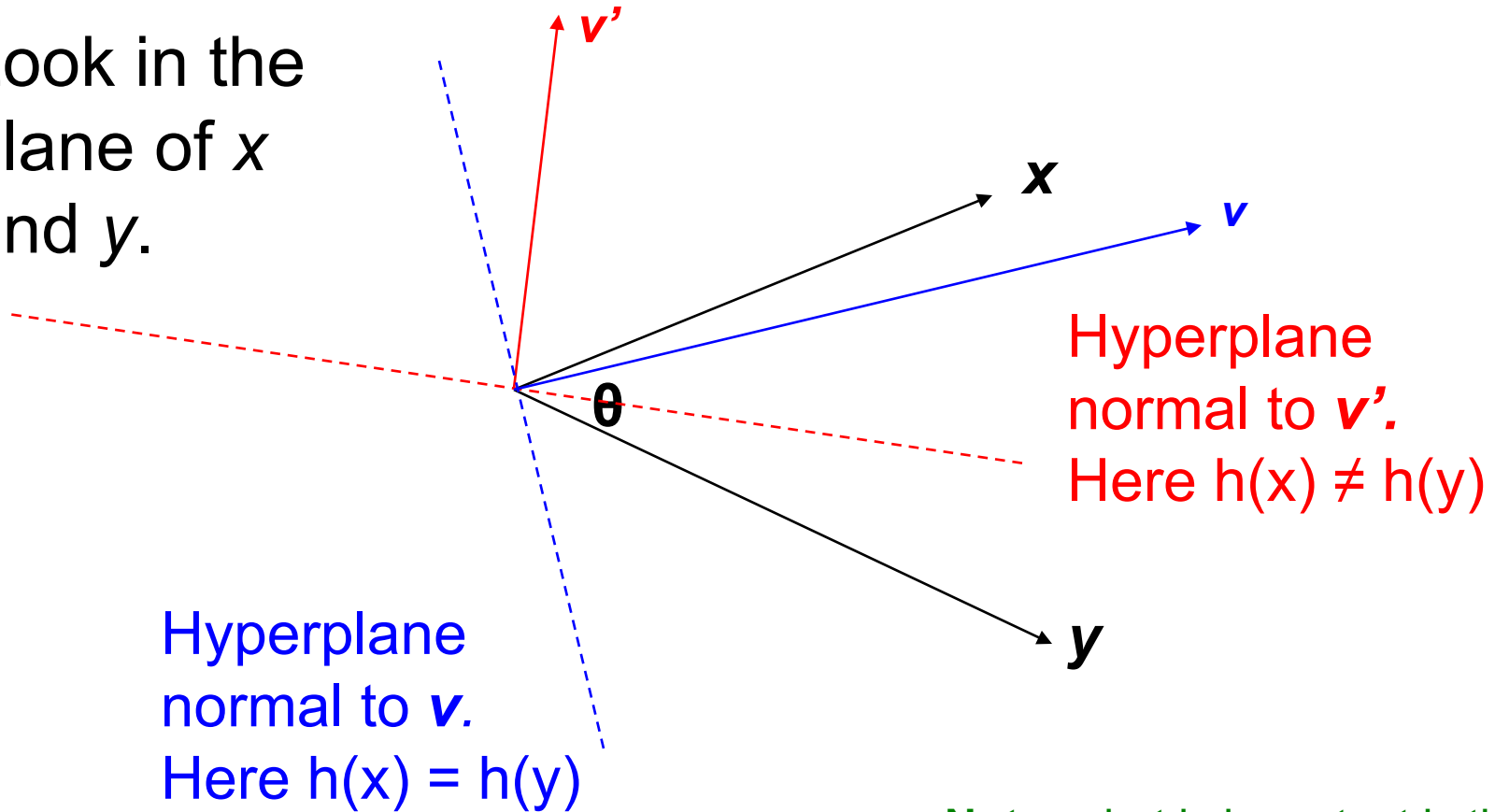
LSH for Cosine Distance

- For **cosine distance**, there is a technique called **Random Hyperplanes**

  - Technique similar to Min-Hashing

- **Random Hyperplanes** method is a $(d_1, d_2, (1-d_1/\pi), (1-d_2/\pi))$-*sensitive* family for any d_1 and d_2

- **Reminder:** (d_1, d_2, p_1, p_2)-*sensitive*
  1. If $d(x,y) \leq d_1$, then prob. that $h(x) = h(y)$ is at least p_1
  2. If $d(x,y) \geq d_2$, then prob. that $h(x) = h(y)$ is at most p_2

Random Hyperplanes

- Each vector v determines a hash function h_v with two buckets

- $h_v(x) = +1$ if $v \cdot x \geq 0$; $= -1$ if $v \cdot x < 0$

- LS-family H = set of all functions derived from any vector

- **Claim:** For points x and y,
$$Pr[h(x) = h(y)] = 1 - d(x,y) / \pi$$

Proof of Claim

Look in the plane of *x* and *y*.

v'

x

v

θ

Hyperplane normal to ***v'*.**
Here h(x) ≠ h(y)

y

Hyperplane normal to ***v*.**
Here h(x) = h(y)

Note: what is important is that hyperplane is outside the angle, not that the vector is inside.

Proof of Claim
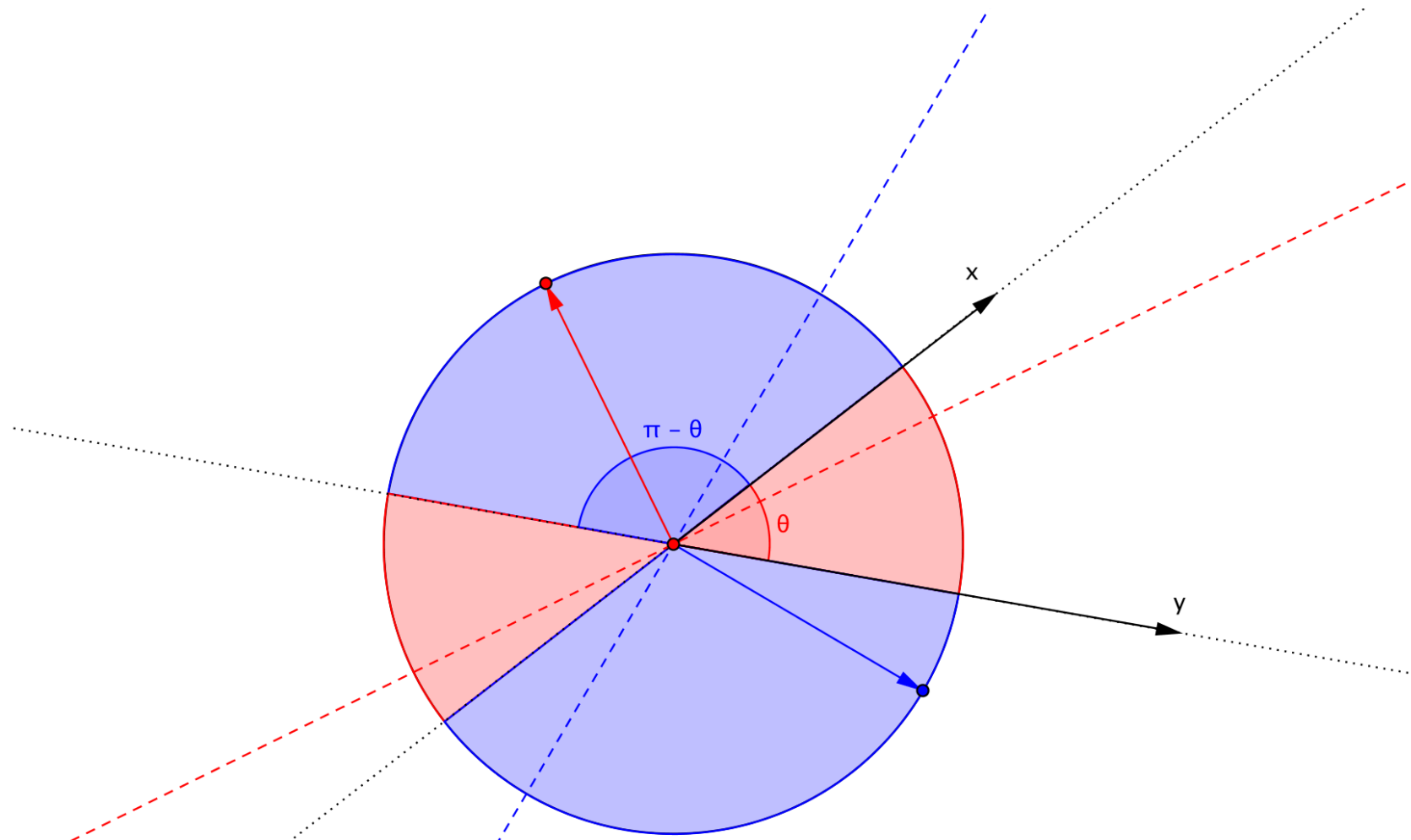
So: Prob[Red case] = θ / π
So: *P[h(x)=h(y)] = 1- θ/π = 1-d(x,y)/π*

Signatures for Cosine Distance

- Pick some number of random vectors, and hash your data for each vector

- The result is a signature (*sketch*) of **+1**'s and **−1**'s for each data point

- Can be used for LSH like we used the Min-Hash signatures for Jaccard distance

- Amplify using **AND**/**OR** constructions

# How to pick random vectors?

- Expensive to pick a random vector in M dimensions for large M

  - Would have to generate M random numbers

- **A more efficient approach**

  - It suffices to consider only vectors v consisting of +1 and −1 components

    - **Why?** Assuming data is random, then vectors of +/-1 cover the entire space evenly (and does not bias in any way)

LSH for Euclidean Distance

- **Idea:** **Hash functions correspond to lines**

- Partition the line into buckets of size **a**

- **Hash each point to the bucket containing its projection onto the line**
  - An element of the "Signature" is a bucket id for that given projection line

- **Nearby points are always close**; distant points are rarely in same bucket
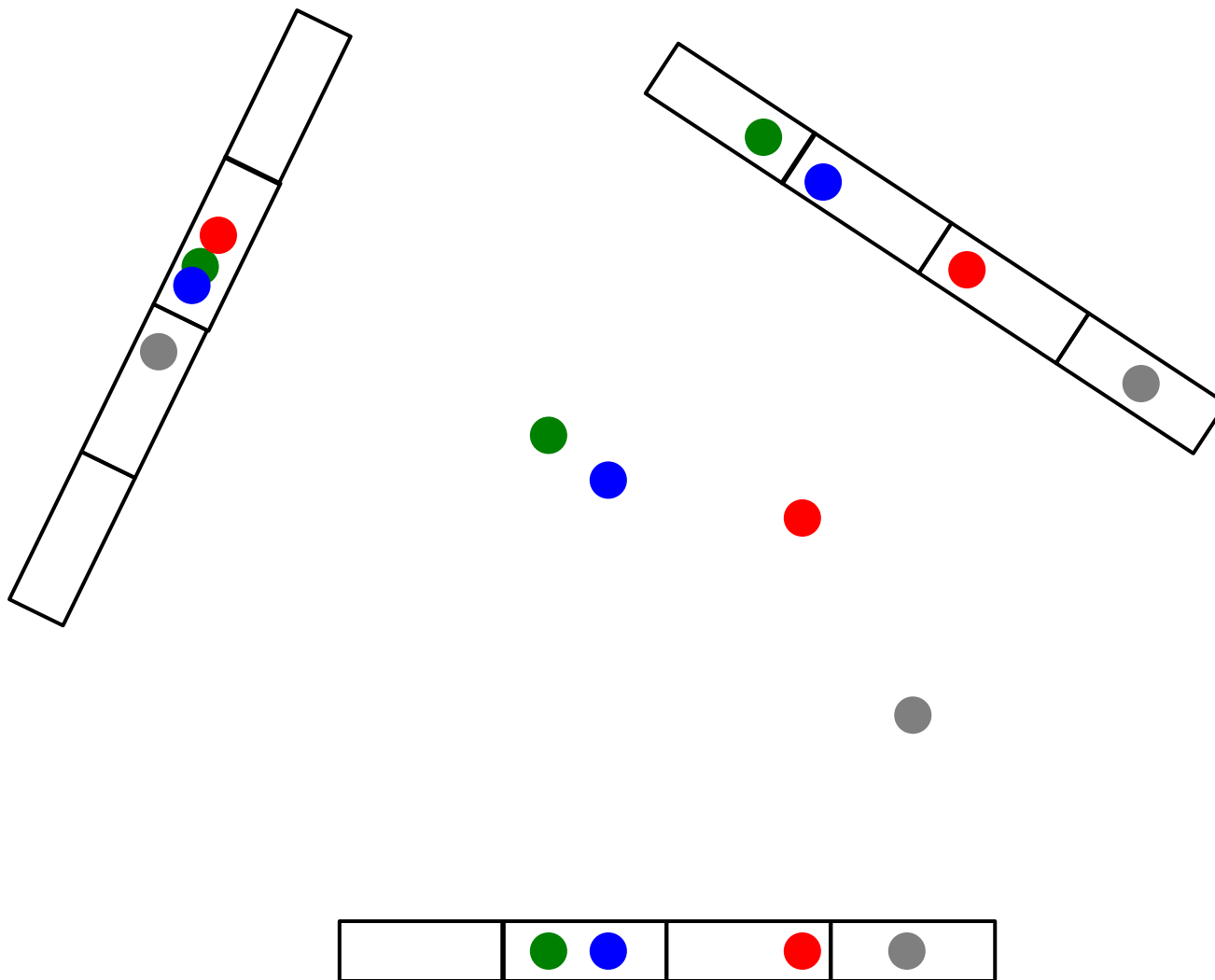
Projection of Points

Line

Buckets of size *a*
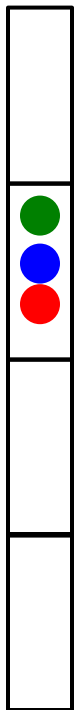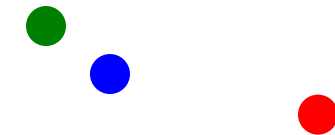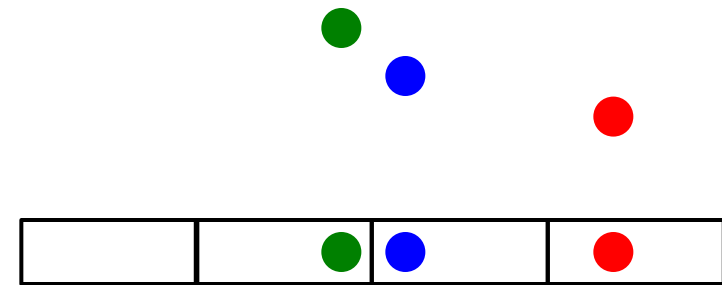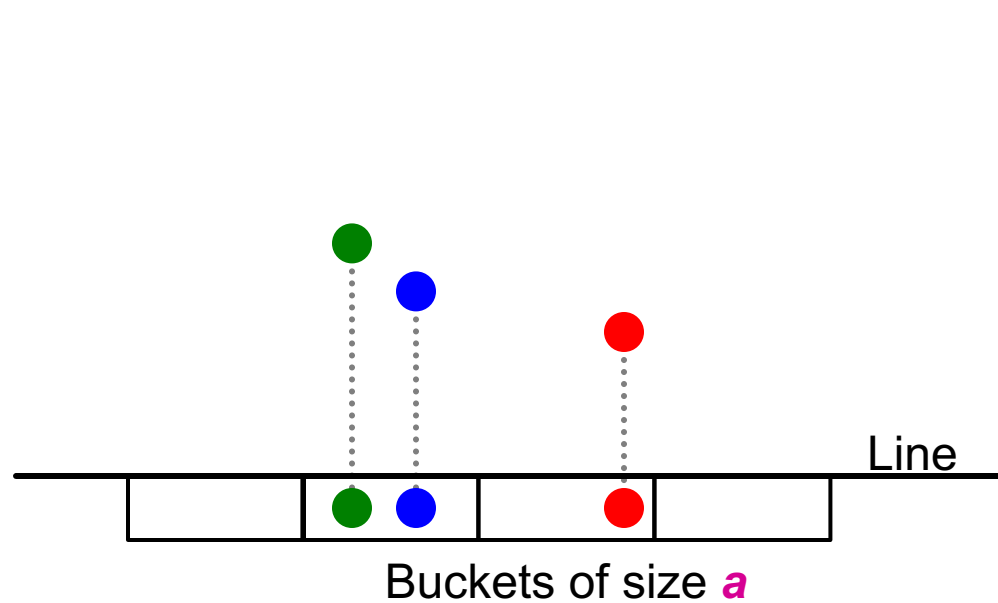
- **"Lucky" case:**
  - Points that are close hash in the same bucket
  - Distant points end up in different buckets

- **Two "unlucky" cases:**
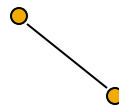 - **Top:** unlucky quantization
  - **Bottom:** unlucky projection
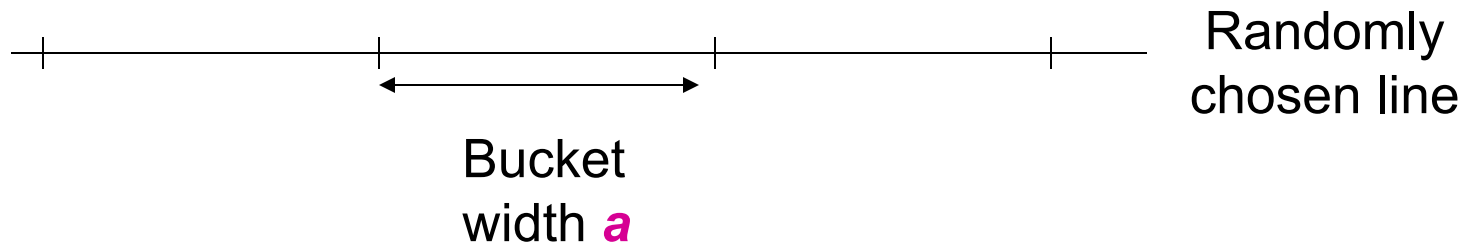
Multiple Projections

Projection of Points

Points at distance **d**

If **d** << **a**, then the chance the points are in the same bucket is at least **1 – d/a**.

Bucket width **a**

Randomly chosen line

Projection of Points

If **d** >> **a**, θ must be close to 90° for there to be any chance points go to the same bucket.

Points at distance **d**

θ

d cos θ

Randomly chosen line

Bucket width **a**

Tim Althoff, UW CS547: Machine Learning for Big Data, http://www.cs.washington.edu/cse547
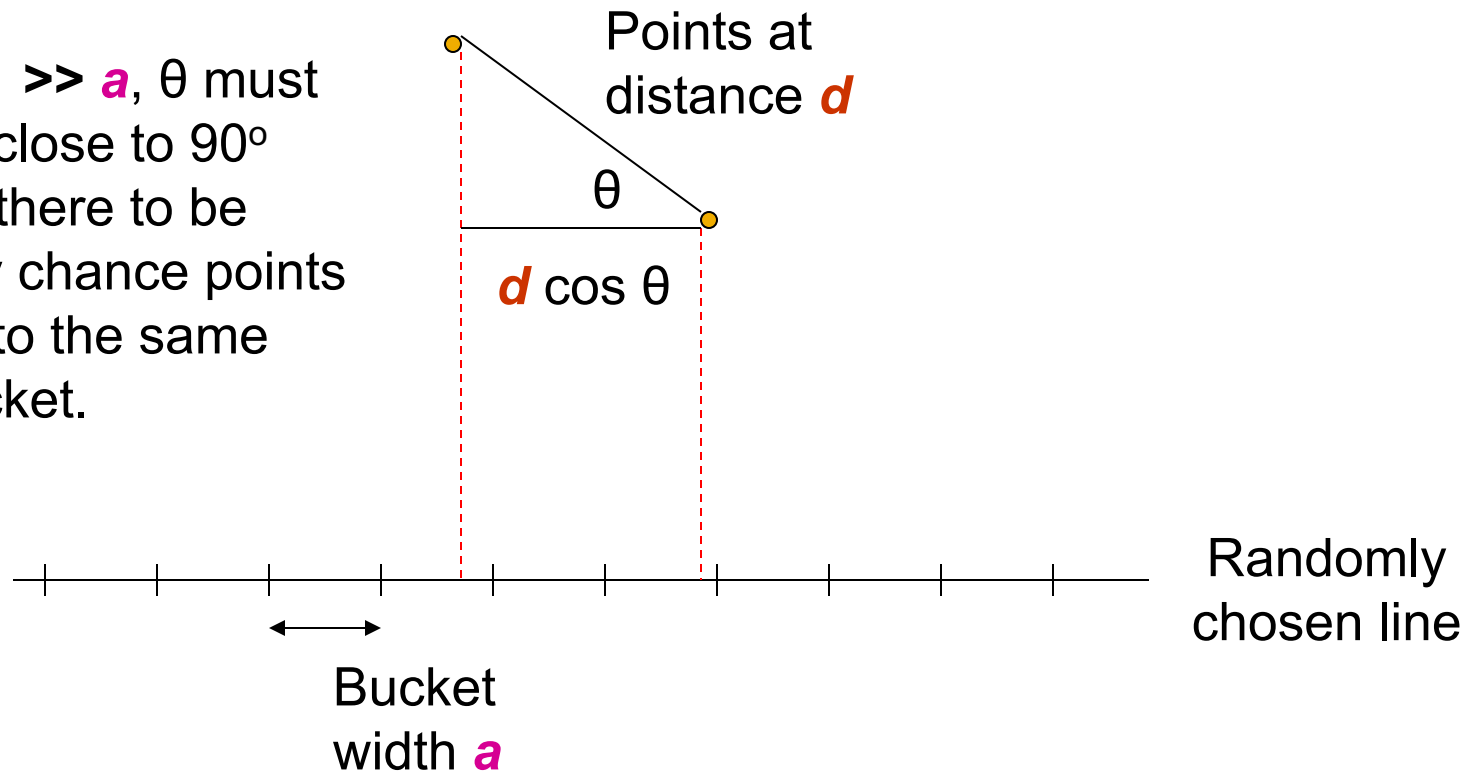
A LS-Family for Euclidean Distance

- If points are distance $d \leq a/2$, prob. they are in same bucket $\geq 1- d/a = ½$
- If points are distance $d \geq 2a$ apart, then they can be in the same bucket only if $d \cos θ \leq a$
  - $\cos θ \leq ½$
  - $60 \leq θ \leq 90$, i.e., at most 1/3 probability

- Yields a *(a/2, 2a, 1/2, 1/3)-sensitive* family of hash functions for any a
- **Amplify using AND-OR cascades**

Summary
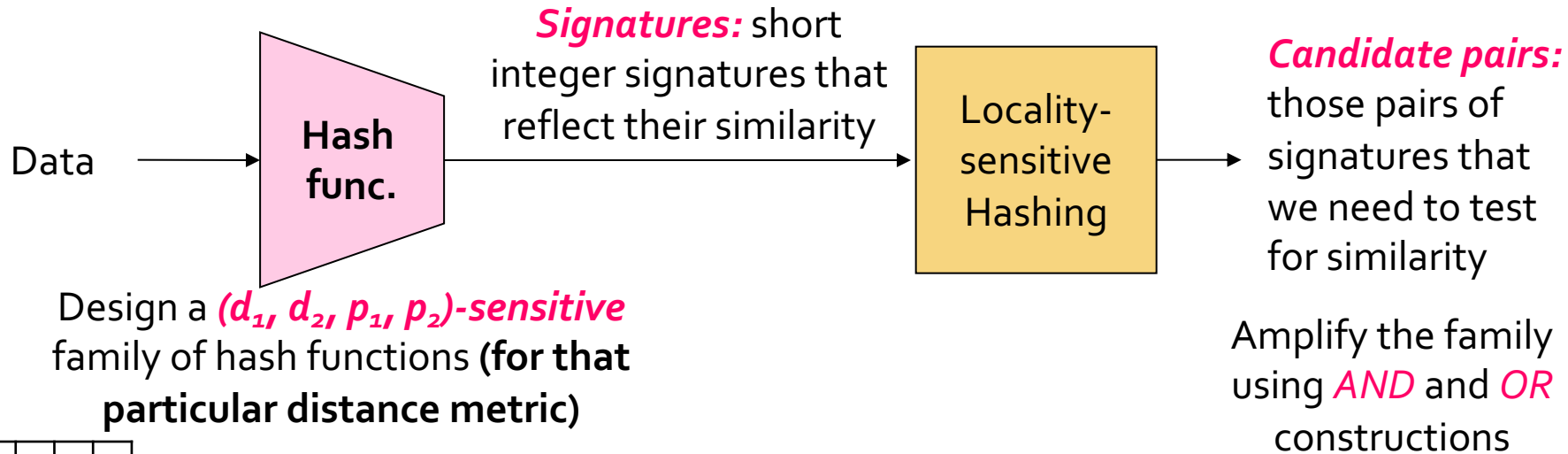
Data → **Hash func.** → *Signatures:* short integer signatures that reflect their similarity → Locality-sensitive Hashing → *Candidate pairs:* those pairs of signatures that we need to test for similarity

Design a (d_1, d_2, p_1, p_2)-*sensitive* family of hash functions **(for that particular distance metric)**

Amplify the family using *AND* and *OR* constructions

Documents

0	1	0	0
1	1	1	0
0	0	0	1
0	1	0	1
0	0	1	0
1	0	0	1

MinHash →

1	5	1	5
2	3	1	3
6	4	6	4

"Bands" technique → *Candidate pairs*

Data points

0	1	0	0
1	1	1	0
0	0	0	1
0	1	0	1
0	0	1	0

Random Hyperplanes →

-1	+1	-1	-1
+1	+1	+1	-1
-1	-1	-1	-1

"Bands" technique → *Candidate pairs*

Two Important Points

- **Property $P(h(C_1)=h(C_2))=sim(C_1,C_2)$ of hash function h is the essential part of LSH, without which we can't do anything**

- **LS-hash functions transform data to signatures so that the bands technique (AND, OR constructions) can then be applied**