#### **Case Study 2: Document Retrieval**

# Task Description: Finding Similar Documents

Machine Learning for Big Data CSE547/STAT548, University of Washington Sham Kakade April 11, 2017

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1

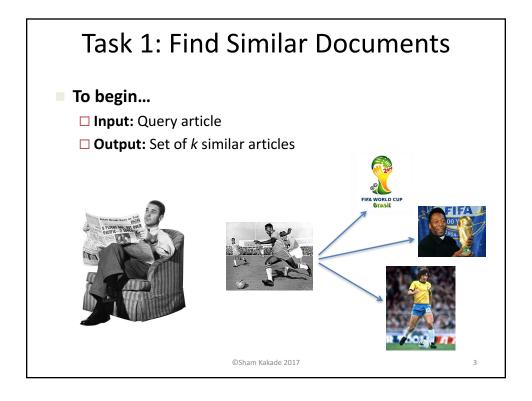
#### **Document Retrieval**

- Goal: Retrieve documents of interest
- Challenges:
  - ☐ Tons of articles out there
  - ☐ How should we measure similarity?





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# **Document Representation**

Bag of words model



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# 1-Nearest Neighbor

- Articles
- Query:
- 1-NN
  - ☐ Goal:
  - □ Formulation:

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5

# k-Nearest Neighbor

- $\qquad \text{Articles} \quad X = \{x^1, \dots, x^N\}, \quad x^i \in \mathbb{R}^d$
- lacksquare Query:  $x \in \mathbb{R}^d$
- k-NN
  - ☐ Goal:
  - □ Formulation:

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#### Distance Metrics – Euclidean

$$d(u, v) = \sqrt{\sum_{i=1}^{d} (u_i - v_i)^2}$$

Or, more generally, 
$$d(u,v) = \sqrt{\sum_{i=1}^d \sigma_i^2 (u_i - v_i)^2}$$

Equivalently,

$$d(u,v) = \sqrt{(u-v)'\Sigma(u-v)}$$
 where 
$$\Sigma = \begin{bmatrix} \sigma_1^2 & 0 & \cdots & 0 \\ 0 & \sigma_2^2 & \cdots & 0 \\ \vdots & \vdots & \cdots & \vdots \\ 0 & 0 & \cdots & \sigma_d^2 \end{bmatrix}$$

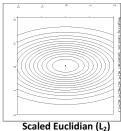
Other Metrics...

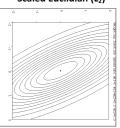
Mahalanobis, Rank-based, Correlation-based, cosine similarity...

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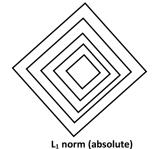
#### **Notable Distance Metrics** (and their level sets)

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Mahalanobis ( $\Sigma$  is general sym pos def matrix, on previous slide = diagonal)





#### **Euclidean Distance + Document Retrieval**

Recall distance metric

$$d(u,v) = \sqrt{\sum_{i=1}^{d} (u_i - v_i)^2}$$

- lacksquare What if each document were lpha times longer?
  - □ Scale word count vectors
  - □ What happens to measure of similarity?
- Good to normalize vectors

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9

### Issues with Document Representation

Words counts are bad for standard similarity metrics





- Term Frequency Inverse Document Frequency (tf-idf)
  - □ Increase importance of rare words

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#### TF-IDF

Term frequency:

$$tf(t,d) =$$

- $\square$  Could also use  $\{0,1\}, 1 + \log f(t,d), \dots$
- Inverse document frequency:

$$idf(t, D) =$$

tf-idf:

$$tfidf(t, d, D) =$$

□ High for document d with high frequency of term t (high "term frequency") and few documents containing term t in the corpus (high "inverse doc frequency")

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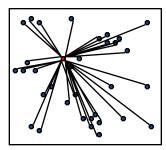
11

## Issues with Search Techniques

Naïve approach:

#### **Brute force search**

- $\square$  Given a query point  ${\mathcal X}$
- $\hfill\Box$  Scan through each point  $x^i$
- □ O(N) distance computations per 1-NN query!
- □ O(*N*log*k*) per *k*-NN query!



33 Distance Computations

What if N is huge??? (and many queries)

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# Think about Web Search/Image Search

- How big is N?
- How fast do we desire to do recall?

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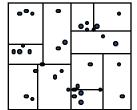
# Intuition (?): NN in 1D and Sorting

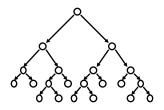
- How do we do 1-NN searches in 1 dim?
- Pre-processing time:
- Query time:

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#### **KD-Trees**

- Smarter approach: kd-trees
  - ☐ Structured organization of documents
    - Recursively partitions points into axis aligned boxes.
  - ☐ Enables more efficient pruning of search space
    - Examine nearby points first.
    - Ignore any points that are further than the nearest point found so far.
- kd-trees work "well" in "lowmedium" dimensions
  - □ We'll get back to this...

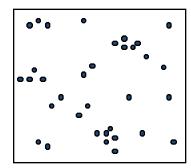




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15

#### **KD-Tree Construction**

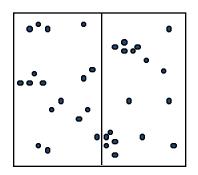


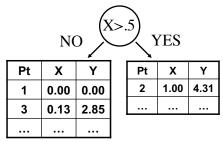
Pt	х	Υ
1	0.00	0.00
2	1.00	4.31
3	0.13	2.85

Start with a list of *d*-dimensional points.

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#### **KD-Tree Construction**



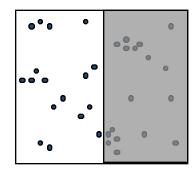


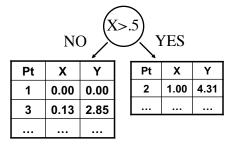
- Split the points into 2 groups by:
  - $\Box$  Choosing dimension  $d_i$  and value V (methods to be discussed...)
  - $\hfill\Box$  Separating the points into  $x_{dj}^i \!\!\!> \!\!\!\! \text{V}$  and  $x_{dj}^i \!\!\!< \!\!\!\!= \!\!\!\!\!\text{V}.$

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17

#### **KD-Tree Construction**

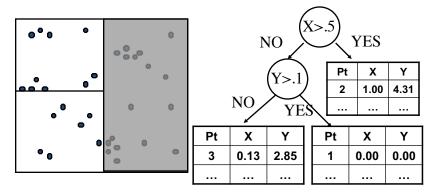




- Consider each group separately and possibly split again (along same/different dimension).
  - □ Stopping criterion to be discussed...

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#### **KD-Tree Construction**

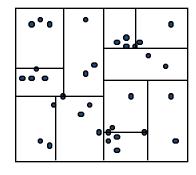


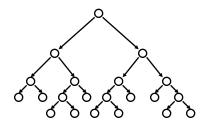
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19

#### **KD-Tree Construction**

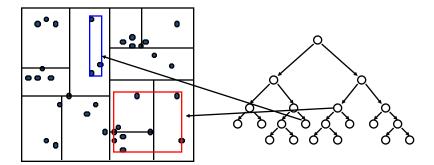




- Continue splitting points in each set
  - □ creates a binary tree structure
- Each leaf node contains a list of points

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#### **KD-Tree Construction**



- Keep one additional piece of information at each node:
  - ☐ The (tight) bounds of the points at or below this node.

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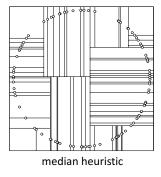
21

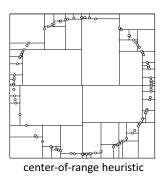
#### **KD-Tree Construction**

- Use heuristics to make splitting decisions:
- Which dimension do we split along?
- Which value do we split at?
- When do we stop?

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# Many heuristics...

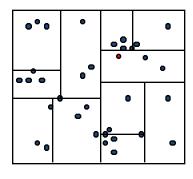


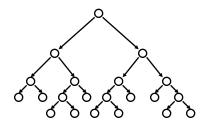


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23

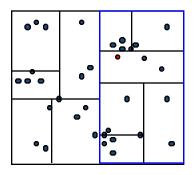
# Nearest Neighbor with KD Trees

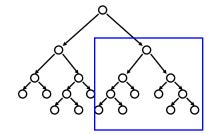




 Traverse the tree looking for the nearest neighbor of the query point.

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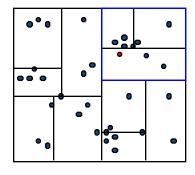


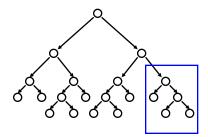
- Examine nearby points first:
  - □ Explore branch of tree closest to the query point first.

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25

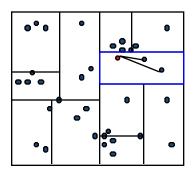
# Nearest Neighbor with KD Trees

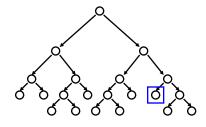




- Examine nearby points first:
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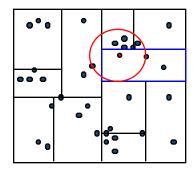


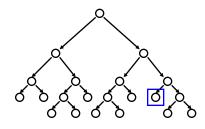
- When we reach a leaf node:
  - □ Compute the distance to each point in the node.

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27

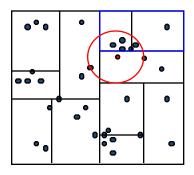
# Nearest Neighbor with KD Trees

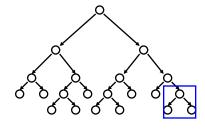




- When we reach a leaf node:
  - □ Compute the distance to each point in the node.

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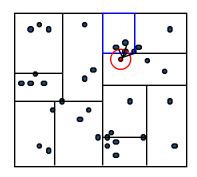


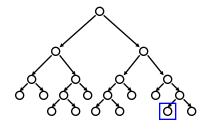
Then backtrack and try the other branch at each node visited

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29

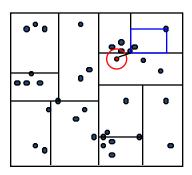
# Nearest Neighbor with KD Trees

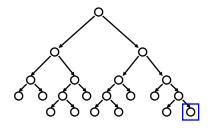




Each time a new closest node is found, update the distance bound

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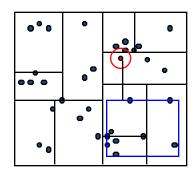


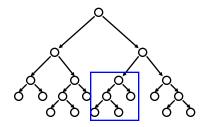
- Using the distance bound and bounding box of each node:
  - □ Prune parts of the tree that could NOT include the nearest neighbor

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31

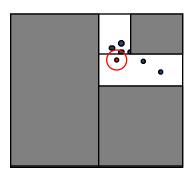
## Nearest Neighbor with KD Trees

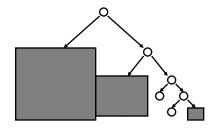




- Using the distance bound and bounding box of each node:
  - ☐ Prune parts of the tree that could NOT include the nearest neighbor

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- Using the distance bound and bounding box of each node:
  - ☐ Prune parts of the tree that could NOT include the nearest neighbor

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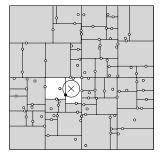
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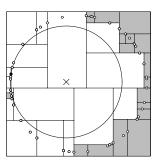
### Complexity

- For (nearly) balanced, binary trees...
- Construction
  - ☐ Size:
  - □ Depth:
  - ☐ Median + send points left right:
  - ☐ Construction time:
- 1-NN query
  - ☐ Traverse down tree to starting point:
  - ☐ Maximum backtrack and traverse:
  - □ Complexity range:
- Under some assumptions on distribution of points, we get O(logN) but exponential in d (see citations in reading)

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# Complexity





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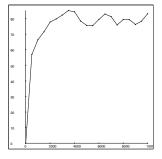
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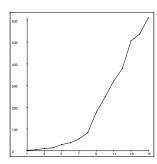
# Complexity for N Queries

- Ask for nearest neighbor to each document
- Brute force 1-NN:
- kd-trees:

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# Inspections vs. N and d

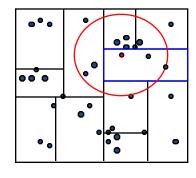


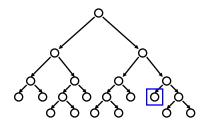


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37

#### K-NN with KD Trees

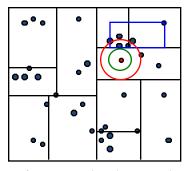


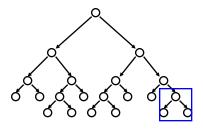


- Exactly the same algorithm, but maintain distance as distance to furthest of current *k* nearest neighbors
- Complexity is:

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# Approximate K-NN with KD Trees





- Before: Prune when distance to bounding box >
- Now: Prune when distance to bounding box >
- Will prune more than allowed, but can guarantee that if we return a neighbor at distance r, then there is no neighbor closer than  $r/\alpha$ .
- In practice this bound is loose...Can be closer to optimal.
- Saves lots of search time at little cost in quality of nearest neighbor.

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39

## Cover trees (+ ball trees)

- What about exact NNs searches in high dimensions?
- Idea: utilize triangle inequality of metric (so allow for arbitrary metric)
- cover-tree guarantees:

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# Cover trees: what does the triangle inequality imply?

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## Cover trees: data structure

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### Wrapping Up – Important Points

#### kd-trees

- Tons of variants
  - □ On construction of trees (heuristics for splitting, stopping, representing branches...)
  - Other representational data structures for fast NN search (e.g.,cover trees, ball trees,...)

#### **Nearest Neighbor Search**

Distance metric and data representation are crucial to answer returned

#### For both...

- High dimensional spaces are hard!
  - □ Number of kd-tree searches can be exponential in dimension
    - Rule of thumb...  $N >> 2^d$ ... Typically useless.
  - □ Distances are sensitive to irrelevant features
    - Most dimensions are just noise → Everything equidistant (i.e., everything is far away)
    - Need technique to learn what features are important for your task

#### What you need to know

- Document retrieval task
  - □ Document representation (bag of words)
  - □ tf-idf
- Nearest neighbor search
  - □ Formulation
  - ☐ Different distance metrics and sensitivity to choice
  - ☐ Challenges with large N
- kd-trees for nearest neighbor search
  - □ Construction of tree
  - □ NN search algorithm using tree
  - □ Complexity of construction and query
  - ☐ Challenges with large *d*

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# Acknowledgment

- This lecture contains some material from Andrew Moore's excellent collection of ML tutorials:
  - □ http://www.cs.cmu.edu/~awm/tutorials
- In particular, see:
  - □ <a href="http://grist.caltech.edu/sc4devo/.../files/sc4devo scalable">http://grist.caltech.edu/sc4devo/.../files/sc4devo scalable</a>
    \_datamining.ppt

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