

CSE 544

Data Models

Lecture #4

Announcements

- **Cancelled:**
 - Lecture on Wednesday, April 4
- **Next Lecture:**
 - Monday, April 9
- **Homework 1**
 - Due on Sunday, 11:59pm
- **Projects:**
 - Decide on a topic
 - Sign up to meet with me next week, Tuesday morning or Thursday morning

<http://www.doodle.com/ugtg59qa6b3cnu35>

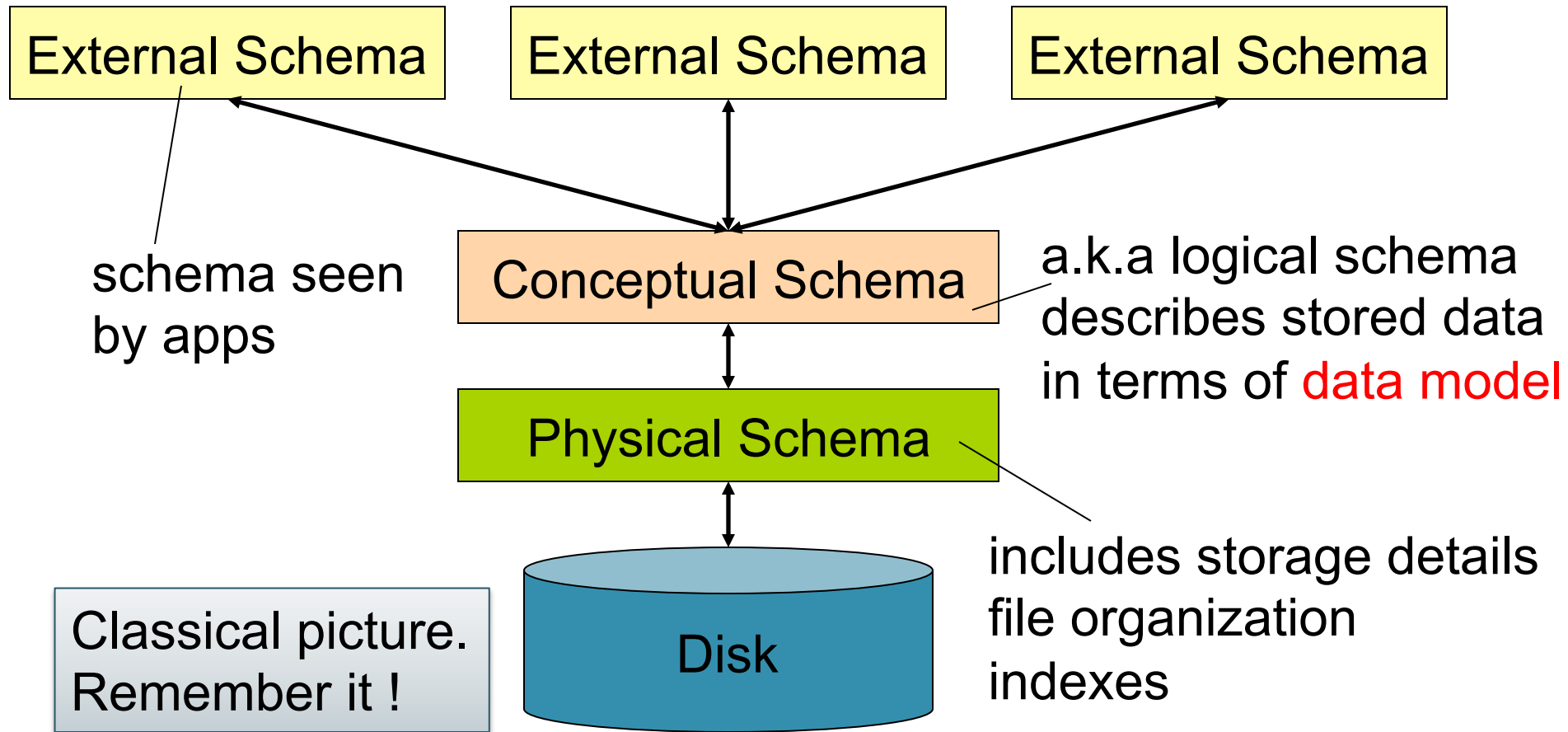
Data Models

- M. Stonebraker and J. Hellerstein. What Goes Around Comes Around. In "Readings in Database Systems" (aka the Red Book). 4th ed.

“Data Model”

- Apps need to model real-world data
 - Typically includes entities and relationships between them
 - Entities: e.g. students, courses, products, clients
 - Relationships: e.g. course registrations, product purchases
- **Data model** enables a user to define the data **using high-level constructs** without worrying about many low-level details of how data will be stored on disk

Levels of Abstraction



Outline

- Different types of data
- Early data models
 - IMS
 - CODASYL
- Relational model
- Other data models: E/R Diagrams, XML

Different Types of Data

- **Structured data**
 - What is this ? Examples ?
- **Semistructured data**
 - What is this ?
 - Examples ?
- **Unstructured data**
 - What is this ? Examples ?

Different Types of Data

- **Structured data**
 - All data conforms to a schema. Ex: business data
- **Semistructured data**
 - Some structure in the data but implicit and irregular
 - Ex: resume, ads
- **Unstructured data**
 - No structure in data. Ex: text, sound, video, images
- **Our focus: structured data & relational DBMSs**

Early Proposal 1: IMS

- What is it ?

Early Proposal 1: IMS

- **Hierarchical data model**
- **Record**
 - **Type**: collection of named fields with data types (+)
 - **Instance**: must match type definition (+)
 - Each instance must have a **key** (+)
 - Record types must be arranged in a **tree** (-)
- **IMS database** is collection of instances of record types organized in a tree

IMS Example

- See Figure 2 in paper “What goes around comes around”

Data Manipulation Language: DL/1

- How does a programmer retrieve data in IMS ?

Data Manipulation Language: DL/1

- Each record has a hierarchical sequence key (HSK)
 - Records are totally ordered: depth-first and left-to-right
- HSK defines semantics of commands:
 - `get_next`
 - `get_next_within_parent`
- **DL/1 is a record-at-a-time language**
 - Programmer constructs an algorithm for solving the query
 - Programmer must worry about query optimization

Data storage

- How is the data physically stored in IMS ?

Data storage

- Root records
 - Stored sequentially (sorted on key)
 - Indexed in a B-tree using the key of the record
 - Hashed using the key of the record
- Dependent records
 - Physically sequential
 - Various forms of pointers
- Selected organizations restrict DL/1 commands
 - No updates allowed with sequential organization
 - No “get-next” for hashed organization

Data Independence

- What is it ?

Data Independence

- **Physical data independence**: Applications are insulated from changes in **physical storage details**
- **Logical data independence**: Applications are insulated from changes to **logical structure of the data**
- **Why are these properties important?**
 - Reduce program maintenance as
 - Logical database design changes over time
 - Physical database design tuned for performance

IMS Limitations

- **Tree-structured data model**
 - Redundant data, existence depends on parent, artificial structure
- **Record-at-a-time** user interface
 - User must specify **algorithm** to access data
- **Very limited physical independence**
 - Phys. organization limits possible operations
 - Application programs break if organization changes
- Provides **some logical independence**
 - DL/1 program runs on logical database
 - Difficult to achieve good logical data independence with a tree model

Early Proposal 2: CODASYL

- What is it ?

Early Proposal 2: CODASYL

- **Networked data model**
- Primitives are also **record types** with **keys (+)**
- Network model is **more flexible than hierarchy(+)**
 - Ex: no existence dependence
- Record types are organized into **network (-)**
 - A record can have multiple parents
 - Arcs between records are named
 - At least one entry point to the network
- **Record-at-a-time** data manipulation language (-)

CODASYL Example

- See Figure 5 in paper “What goes around comes around”

CODASYL Limitations

- **No physical data independence**
 - Application programs break if organization changes
- **No logical data independence**
 - Application programs break if organization changes
- Very **complex**
- Programs must “**navigate** the hyperspace”
- Load and recover as **one gigantic object**

Relational Model Overview

- Proposed by Ted Codd in 1970
- Motivation: better logical and physical data independence

Relational Model Overview

- Defines logical schema only
 - No physical schema
- Set-at-a-time query language

Physical Independence

- **Definition:** Applications are insulated from changes in physical storage details
- Early models (IMS and CODASYL): **No**
- Relational model: **Yes**
 - Yes through set-at-a-time language: algebra or calculus
 - No specification of what storage looks like
 - Administrator can optimize physical layout

Logical Independence

- **Definition:** Applications are insulated from changes to logical structure of the data
- Early models
 - IMS: **some** logical independence
 - CODASYL: **no** logical independence
- Relational model
 - **Yes** through views

Great Debate

- Pro relational
 - What were the arguments ?
- Against relational
 - What were the arguments ?
- How was it settled ?

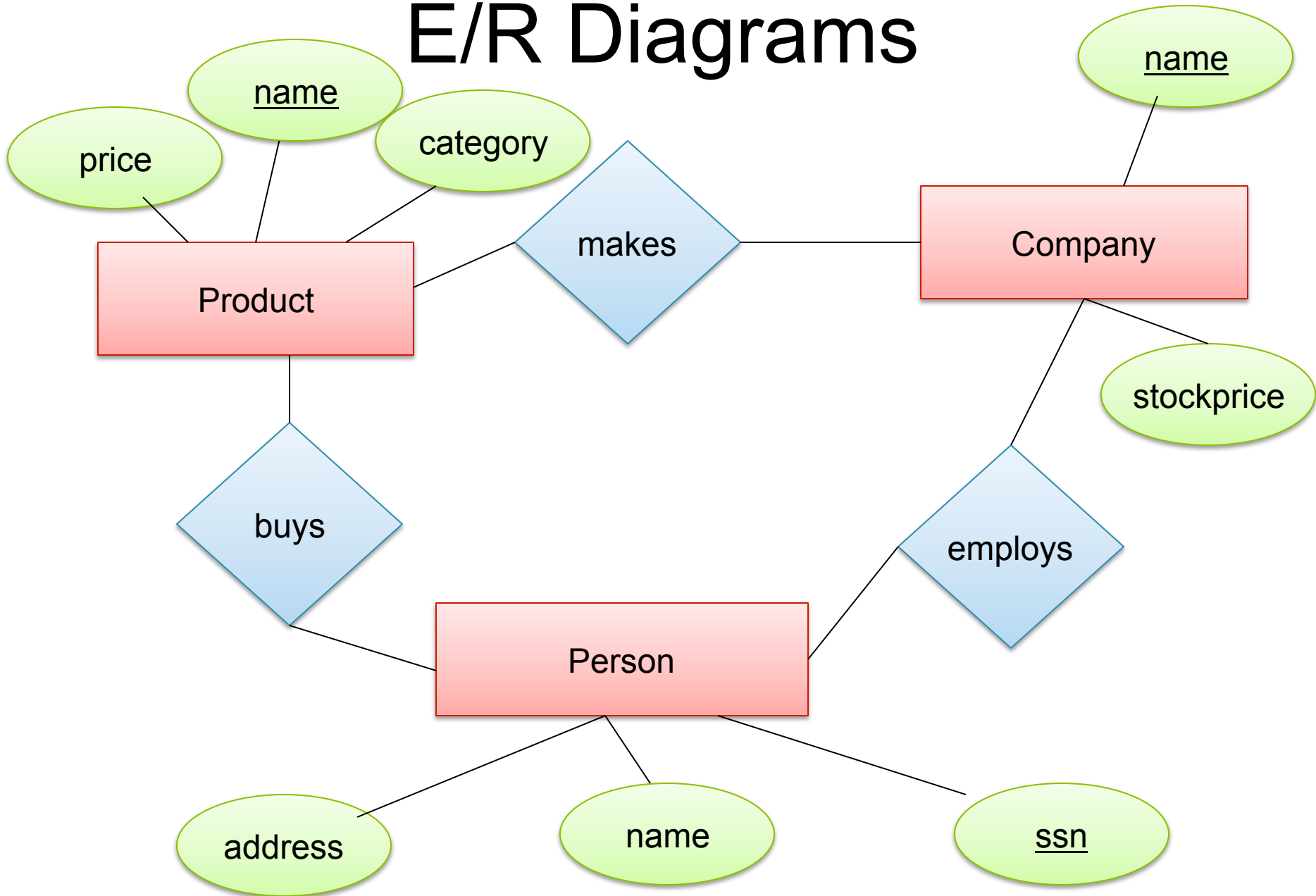
Great Debate

- Pro relational
 - CODASYL is too complex
 - CODASYL does not provide sufficient data independence
 - Record-at-a-time languages are too hard to optimize
 - Trees/networks not flexible enough to represent common cases
- Against relational
 - COBOL programmers cannot understand relational languages
 - Impossible to represent the relational model efficiently
 - CODASYL can represent tables
- Ultimately settled by the market place

Other Data Models

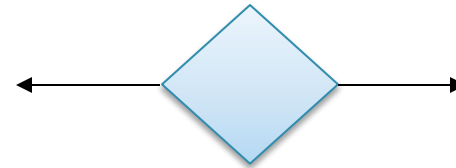
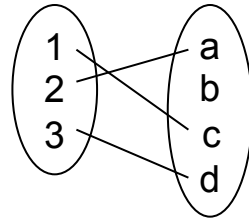
- Entity-Relationship: 1970's
 - Successful in logical database design (you'll use it in hw1)
- Extended Relational: 1980's
- Semantic: late 1970's and 1980's
- Object-oriented: late 1980's and early 1990's
 - Address impedance mismatch: relational dbs \leftrightarrow OO languages
 - Interesting but ultimately failed (several reasons, see paper)
- Object-relational: late 1980's and early 1990's
 - User-defined types, ops, functions, and access methods
- Semi-structured: late 1990's to the present

E/R Diagrams

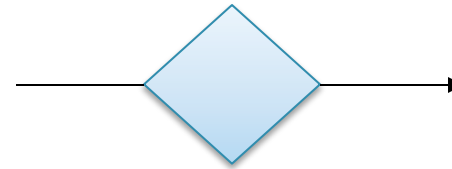
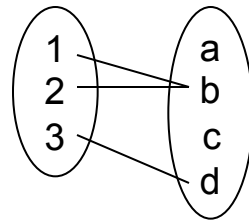


Multiplicity of E/R Relations

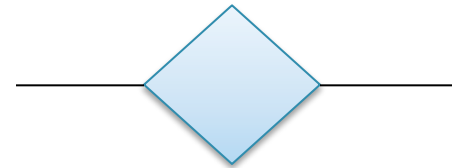
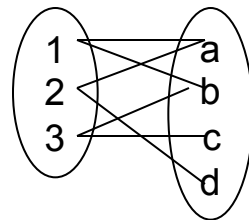
- one-one:

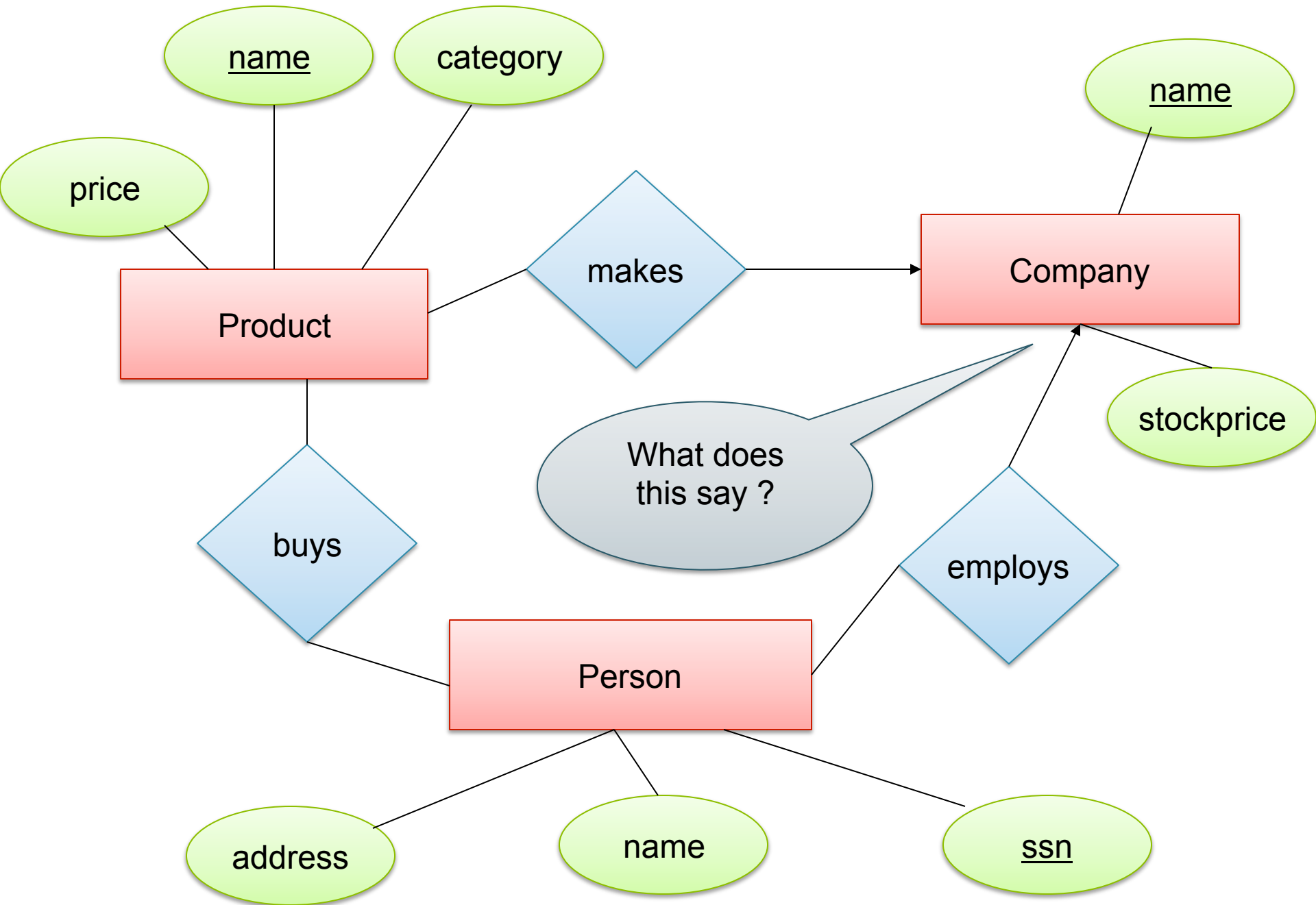


- many-one

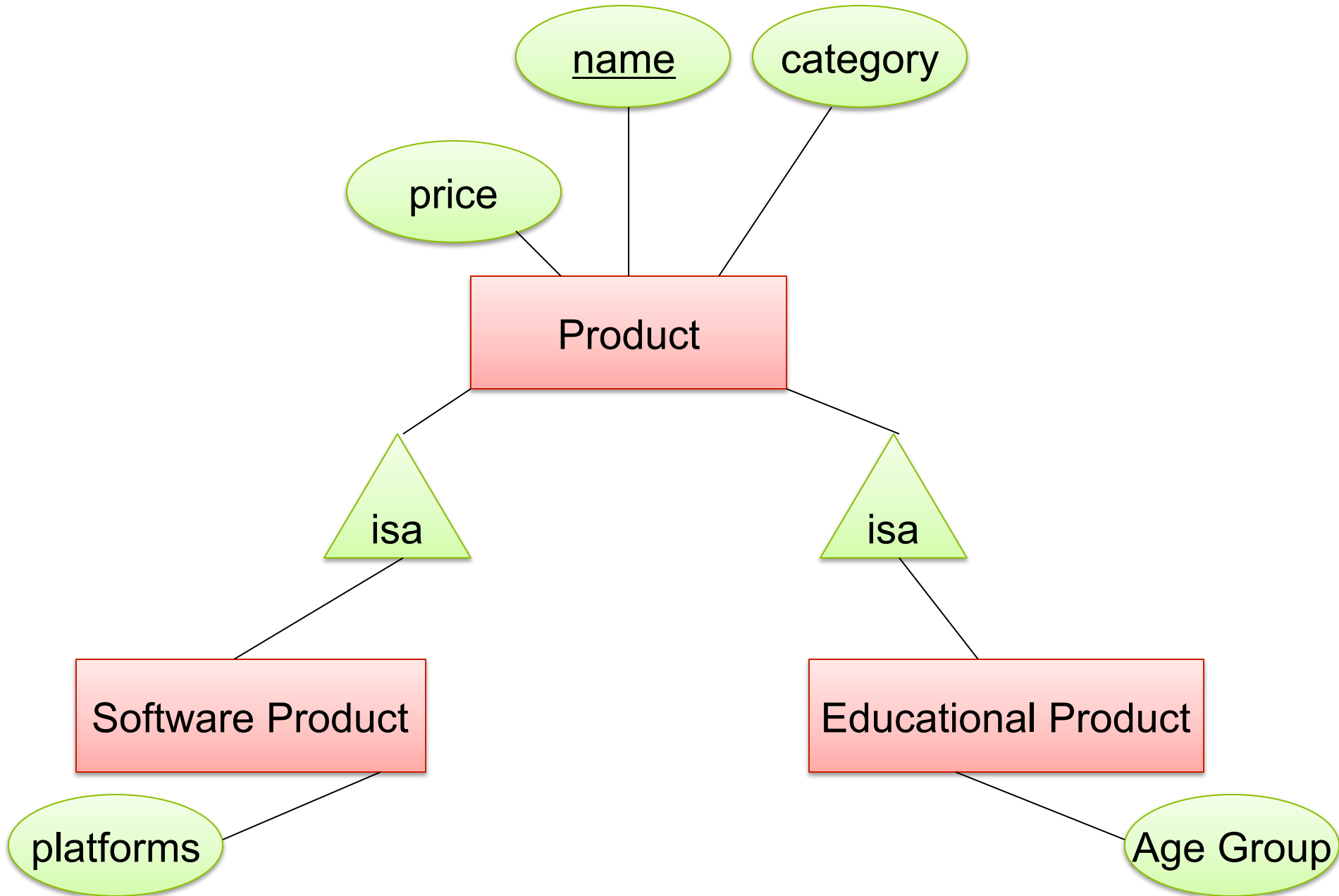


- many-many





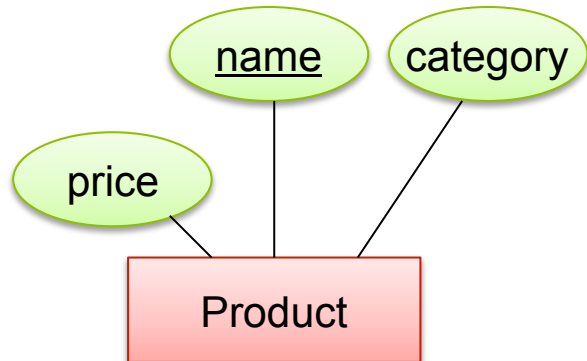
Subclasses



Subclasses to Relations

Product

<u>Name</u>	Price	Category
Gizmo	99	gadget
Camera	49	photo
Toy	39	gadget



Sw.Product

<u>Name</u>	platforms
Gizmo	unix



Ed.Product

<u>Name</u>	Age Group
Gizmo	todler
Toy	retired

Other ways to convert are possible

XML Syntax

```
<bibliography>  
  <book>  <title> Foundations... </title>  
          <author> Abiteboul </author>  
          <author> Hull </author>  
          <author> Vianu </author>  
          <publisher> Addison Wesley </publisher>  
          <year> 1995 </year>  
  
  </book>  
  ...  
</bibliography>
```

XML Terminology

- Tags: `book`, `title`, `author`, ...
- Start tag: `<book>`, end tag: `</book>`
- Elements: `<book>...</book>`, `<author>...</author>`
- Elements are nested
- Empty element: `<red></red>` abbrev. `<red/>`
- An XML document: single *root element*

Well formed XML document

- Has matching tags
- A short header
- And a root element

Well-Formed XML

```
<? xml version="1.0" encoding="utf-8" standalone="yes" ?>  
<SomeTag>  
    ...  
</SomeTag>
```

Parsing and processing XML Documents:

- DOM = Document Object Model = main memory
- SAX = Simple API for XML = event driven = we use it in HW1

More XML: Attributes

```
<book price = "55" currency = "USD">  
  <title> Foundations of Databases </title>  
  <author> Abiteboul </author>  
  ...  
  <year> 1995 </year>  
</book>
```

Attributes v.s. Elements

```
<book price = "55" currency = "USD">  
  <title> Foundations of DBs </title>  
  <author> Abiteboul </author>  
  ...  
  <year> 1995 </year>  
</book>
```

```
<book>  
  <title> Foundations of DBs </title>  
  <author> Abiteboul </author>  
  ...  
  <year> 1995 </year>  
  <price> 55 </price>  
  <currency> USD </currency>  
</book>
```

Attributes are alternative ways to represent data

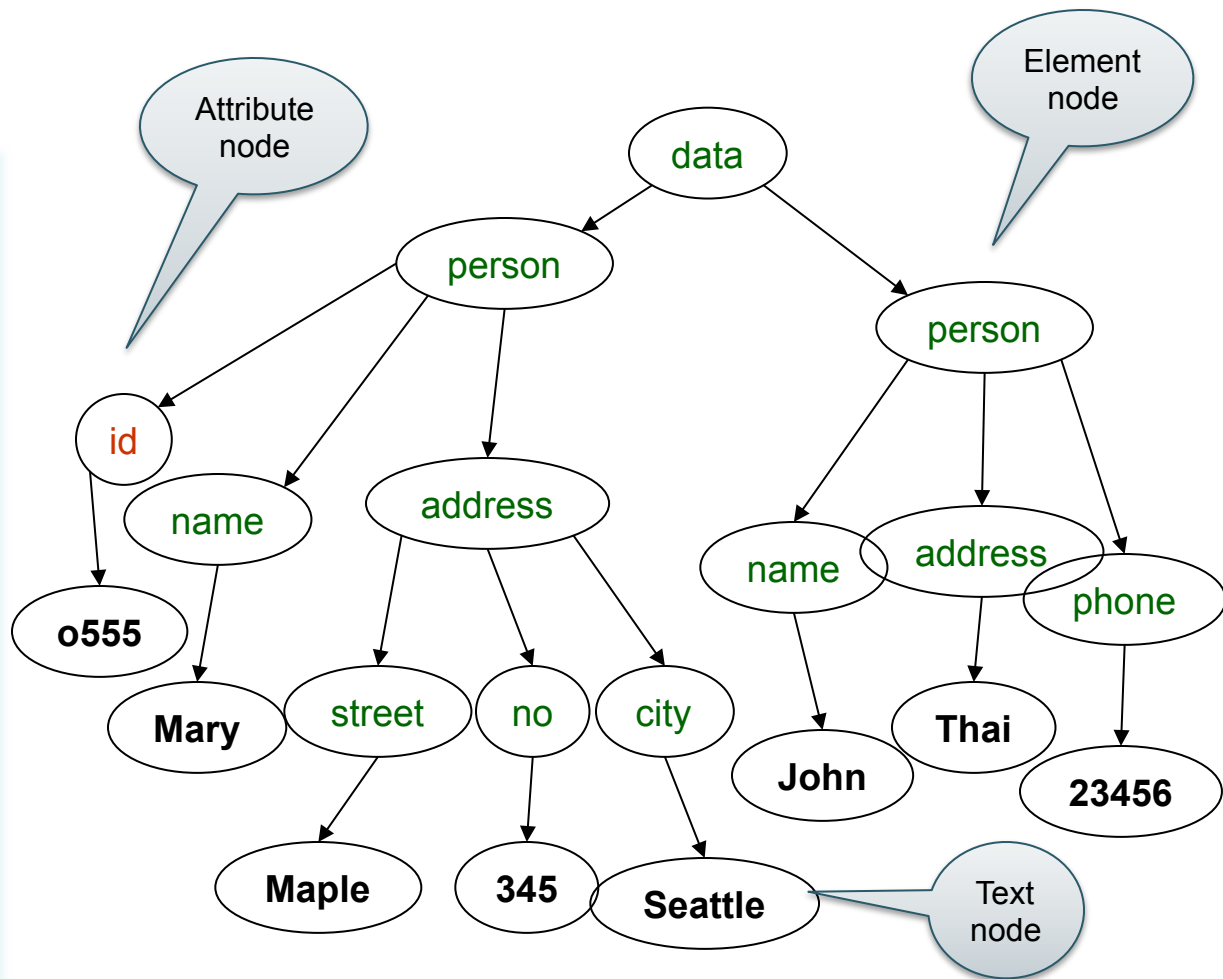
Comparison

Elements	Attributes
Ordered	Unordered
May be repeated	Must be unique
May be nested	Must be atomic

XML Semantics: a Tree !

DOM = Document Object Model

```
<data>
  <person id="o555" >
    <name> Mary </name>
    <address>
      <street>Maple</street>
      <no> 345 </no>
      <city> Seattle </city>
    </address>
  </person>
  <person>
    <name> John </name>
    <address>Thailand
    </address>
    <phone>23456</phone>
  </person>
</data>
```



Order matters !!!

XML Data

- XML is **self-describing**
- Schema elements become part of the data
 - Relational schema: `person(name,phone)`
 - In XML `<person>`, `<name>`, `<phone>` are part of the data, and are repeated many times
- Consequence: XML is much more flexible
- XML = **semistructured** data

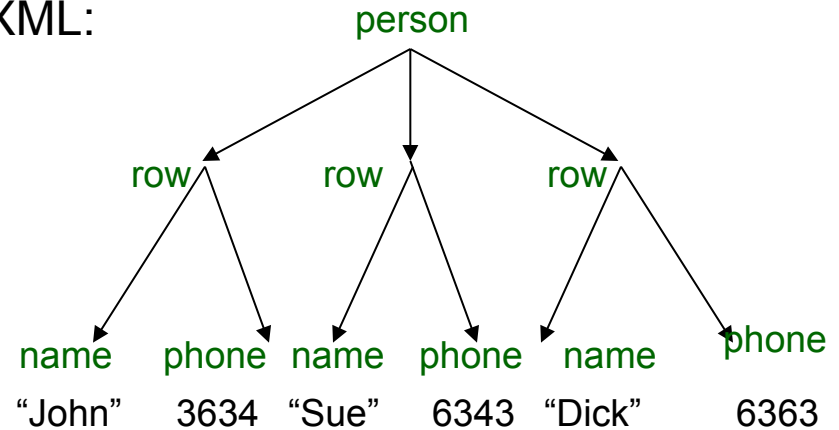
Mapping Relational Data to XML Data

The canonical mapping:

Person

Name	Phone
John	3634
Sue	6343
Dick	6363

XML:



```
<person>
  <row> <name>John</name>
    <phone> 3634</phone></row>
  <row> <name>Sue</name>
    <phone> 6343</phone></row>
  <row> <name>Dick</name>
    <phone> 6363</phone></row>
</person>
```

Mapping Relational Data to XML Data

Application specific mapping

Person

Name	Phone
John	3634
Sue	6343

Orders

PersonName	Date	Product
John	2002	Gizmo
John	2004	Gadget
Sue	2002	Gadget

XML

```
<people>
  <person>
    <name> John </name>
    <phone> 3634 </phone>
    <order> <date> 2002 </date>
      <product> Gizmo </product>
    </order>
    <order> <date> 2004 </date>
      <product> Gadget </product>
    </order>
  </person>
  <person>
    <name> Sue </name>
    <phone> 6343 </phone>
    <order> <date> 2004 </date>
      <product> Gadget </product>
    </order>
  </person>
</people>
```

XML=Semi-structured Data (1/3)

- Missing attributes:

```
<person> <name> John</name>  
          <phone>1234</phone>  
</person>  
  
<person> <name>Joe</name>  
</person>
```

no phone !

- Could represent in a table with nulls

name	phone
John	1234
Joe	-

XML=Semi-structured Data (2/3)

- Repeated attributes

```
<person> <name> Mary</name>  
          <phone>2345</phone>  
          <phone>3456</phone>  
</person>
```

Two phones !

- Impossible in tables:

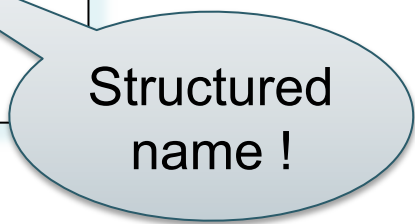
name	phone	
Mary	2345	3456

???

XML=Semi-structured Data (3/3)

- Attributes with different types in different objects

```
<person> <name> <first> John </first>  
          <last> Smith </last>  
        </name>  
        <phone>1234</phone>  
</person>
```



Structured name !

- Nested collections
- Heterogeneous collections:
 - <db> contains both <book>s and <publisher>s

Summary

- **Data independence is desirable**
 - Both physical and logical
 - Early data models provided very limited data independence
 - Relational model facilitates data independence
 - Set-at-a-time languages facilitate phys. indep. [more next lecture]
 - Simple data models facilitate logical indep. [more next lecture]
- **Flat models are also simpler, more flexible**
- **User should specify what they want not how to get it**
 - Query optimizer does better job than human
- **New data model proposals must**
 - Solve a “major pain” or provide significant performance gains