### CSE 544 Parallel Databases

#### Tuesday, February 17th, 2011

### Final Thoughts on Optimization: Parameters !



Chaudhuri "Rethinking the Contract"

Dan Suciu -- 544, Winter 2011

# **Overview of Today's Lecture**

- Parallel databases (Chapter 22.1 22.5)
- Map/reduce
- Pig-Latin
  - Some slides from Alan Gates (Yahoo! Research)

### Parallel v.s. Distributed Databases

- Parallel database system:
  - Improve performance through parallel implementation
  - Will discuss in class
- Distributed database system:
  - Data is stored across several sites, each site managed by a DBMS capable of running independently
  - Will not discuss in class

# Parallel DBMSs

Goal

 Improve performance by executing multiple operations in parallel

- Key benefit
  - Cheaper to scale than relying on a single increasingly more powerful processor
- Key challenge
  - Ensure overhead and contention do not kill performance

### Performance Metrics for Parallel DBMSs

- Speedup
  - More processors  $\rightarrow$  higher speed
  - Individual queries should run faster
  - Should do more transactions per second (TPS)
- Scaleup
  - More processors  $\rightarrow$  can process more data
  - Batch scaleup
    - Same query on larger input data should take the same time
  - Transaction scaleup
    - N-times as many TPS on N-times larger database
    - But each transaction typically remains small



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### Linear v.s. Non-linear Scaleup



### Challenges to Linear Speedup and Scaleup

- Startup cost
  - Cost of starting an operation on many processors
- Interference
  - Contention for resources between processors
- Skew
  - Slowest processor becomes the bottleneck

### Architectures for Parallel Databases

- Shared memory
- Shared disk
- Shared nothing

### **Shared Memory**



### Shared Disk



### **Shared Nothing**



# **Shared Nothing**

- Most scalable architecture
  - Minimizes interference by minimizing resource sharing
  - Can use commodity hardware
- Also most difficult to program and manage
- Processor = server = node
- P = number of nodes

We will focus on shared nothing

### Taxonomy for Parallel Query Evaluation

Inter-query parallelism

- Each query runs on one processor

- Inter-operator parallelism
  - A query runs on multiple processors
  - An operator runs on one processor
- Intra-operator parallelism
  - An operator runs on multiple processors

We study only intra-operator parallelism: most scalable

# Horizontal Data Partitioning

- Relation R split into P chunks R<sub>0</sub>, ..., R<sub>P-1</sub>, stored at the P nodes
- Round robin: tuple t<sub>i</sub> to chunk (i mod P)
- Hash based partitioning on attribute A:
   Tuple t to chunk h(t.A) mod P
- Range based partitioning on attribute A:
   Tuple t to chunk i if v<sub>i-1</sub> < t.A < v<sub>i</sub>

### **Parallel Selection**

Compute  $\sigma_{A=v}(R)$ , or  $\sigma_{v1 < A < v2}(R)$ 

Conventional database:

-Cost = B(R)

Parallel database with P processors:
 – Cost = B(R) / P

### **Parallel Selection**

Different processors do the work:

- Round robin partition: all servers do the work
- Hash partition:

– One server for  $\sigma_{A=v}(R)$ ,

– All servers for  $\sigma_{v1 < A < v2}(R)$ 

Range partition: one server does the work

# Data Partitioning Revisited

What are the pros and cons?

- Round robin
  - Good load balance but always needs to read all the data
- Hash based partitioning
  - Good load balance but works only for equality predicates and full scans
- Range based partitioning
  - Works well for range predicates but can suffer from data skew

# Parallel Group By: $\gamma_{A, sum(B)}(R)$

Step 1: server i partitions chunk R<sub>i</sub> using a hash function h(t.A): R<sub>i0</sub>, R<sub>i1</sub>, ..., R<sub>i,P-1</sub>

Step 2: server i sends partition R<sub>ii</sub> to server j

Step 3: server j computes  $\gamma_{A, sum(B)}$  on  $R_{0j}, R_{1j}, ..., R_{P-1,j}$ 

# Cost of Parallel Group By

Recall conventional cost = 3B(R)

- Step 1: Cost = B(R)/P I/O operations
- Step 2: Cost = (P-1)/P B(R) blocks are sent
   Network costs << I/O costs</li>
- Step 3: Cost = 2 B(R)/P

– When can we reduce it to 0?

Total = 3B(R) / P + communication costs

# Parallel Join: $R \bowtie_{A=B} S$

Step 1

- For all servers in [0,k], server i partitions chunk R<sub>i</sub> using a hash function h(t.A): R<sub>i0</sub>, R<sub>i1</sub>, ..., R<sub>i,P-1</sub>
- For all servers in [k+1,P], server j partitions chunk S<sub>i</sub> using a hash function h(t.A): S<sub>j0</sub>, S<sub>j1</sub>, ..., R<sub>j,P-1</sub>

Step 2:

- Server i sends partition R<sub>iu</sub> to server u
- Server j sends partition S<sub>ju</sub> to server u

Steps 3: Server u computes the join of R<sub>iu</sub> with S<sub>iu</sub>

### Cost of Parallel Join

- Step 1: Cost = (B(R) + B(S))/P
- Step 2: 0
  - (P-1)/P (B(R) + B(S)) blocks are sent, but we assume network costs to be << disk I/O costs</p>
- Step 3:
  - Cost = 0 if small table fits in memory: B(S)/P <=M</p>
  - Cost = 4(B(R)+B(S))/P otherwise

# Parallel Query Plans

- Same relational operators
- Add special split and merge operators

   Handle data routing, buffering, and flow control
- Example: exchange operator
  - Inserted between consecutive operators in the query plan

# Map Reduce

- Google: paper published 2004
- Free variant: Hadoop
- Map-reduce = high-level programming model and implementation for large-scale parallel data processing

### Data Model

Files !

#### A file = a bag of (key, value) pairs

A map-reduce program:

- Input: a bag of (inputkey, value) pairs
- Output: a bag of (outputkey, value) pairs

# Step 1: the MAP Phase

User provides the MAP-function:

- Input: one (input key, value)
- **Ouput: bag of** (intermediate key, value) **pairs**

System applies the map function in parallel to all (input key, value) pairs in the input file

# Step 2: the REDUCE Phase

User provides the REDUCE function:

- Input: (intermediate key, bag of values)
- Output: bag of output values
   System groups all pairs with the same intermediate key, and passes the bag of values to the REDUCE function

# Example

 Counting the number of occurrences of each word in a large collection of

map(String key, String value): // key: document name // value: document contents for each word w in value: EmitIntermediate(w, "1"):

reduce(String key, Iterator values):
 // key: a word
 // values: a list of counts
 int result = 0;
 for each v in values:
 result += ParseInt(v);
 Emit(AsString(result));



### Map = GROUP BY, Reduce = Aggregate

R(documentKey, word)

SELECT word, sum(1) FROM R GROUP BY word

# Implementation

- There is one master node
- Master partitions input file into *M* splits, by key
- Master assigns *workers* (=servers) to the *M* map tasks, keeps track of their progress
- Workers write their output to local disk, partition into *R regions*
- Master assigns workers to the *R reduce tasks*
- Reduce workers read regions from the map workers' local disks

### **MR** Phases



# Interesting Implementation Details

- Worker failure:
  - Master pings workers periodically,
  - If down then reassigns its splits to all other workers → good load balance
- Choice of M and R:
  - Larger is better for load balancing
  - Limitation: master needs O(M×R) memory

### Interesting Implementation Details

Backup tasks:

- Straggler = a machine that takes unusually long time to complete one of the last tasks.
   Eg:
  - Bad disk forces frequent correctable errors (30MB/s → 1MB/s)
  - The cluster scheduler has scheduled other tasks on that machine
- Stragglers are a main reason for slowdown
- Solution: pre-emptive backup execution of the last few remaining in-progress tasks

# Map-Reduce Summary

- Hides scheduling and parallelization details
- However, very limited queries
  - Difficult to write more complex tasks
  - Need multiple map-reduce operations
- Solution:


# Following Slides courtesy of: Alan Gates, Yahoo!Research

# What is Pig?

- An engine for executing programs on top of Hadoop
- It provides a language, Pig Latin, to specify these programs
- An Apache open source project
   <u>http://hadoop.apache.org/pig/</u>





# **Map-Reduce**

- Computation is moved to the data
- A simple yet powerful programming model
  - Map: every record handled individually
  - Shuffle: records collected by key
  - Reduce: key and iterator of all associated values
- User provides:
  - input and output (usually files)
  - map Java function
  - key to aggregate on
  - reduce Java function
- Opportunities for more control: partitioning, sorting, partial aggregations, etc.























# Making Parallelism Simple

- Sequential reads = good read speeds
- In large cluster failures are guaranteed; Map Reduce handles retries
- Good fit for batch processing applications that need to touch all your data:
  - data mining
  - model tuning
- Bad fit for applications that need to find one particular record
- Bad fit for applications that need to communicate between processes; oriented around independent units of work



# Why use Pig?

Suppose you have user data in one file, website data in another, and you need to find the top 5 most visited sites by users aged 18 -25.





#### In Map-Reduce

import java.io.IOException; import java.util.ArrayList; import java.util.Iterator; import java.util.List; import org.apache.hadoop.fs.Path; import org.apache.hadoop.io.LongWritable; import org.apache.hadoop.io.Text; import org.apache.hadoop.io.Writable; import org.apache.hadoop.io.WritableComparable; import org.apache.hadoop.mapred.FileInputFormat; import org.apache.hadoop.mapred.FileOutputFormat; import org.apache.hadoop.mapred.JobConf; import org.apache.hadoop.mapred.jobCont; import org.apache.hadoop.mapred.sobCont; import org.apache.hadoop.mapred.KeyValueTextInputFormat; import org.apache.hadoop.mapred.MapRedCollast; import org.apache.hadoop.mapred.MapRedCollast; import org.apache.hadoop.mapred.RedCollast; import org.apache.hadoop.mapred.SequencefileOutputFormat; import org.apache.hadoop.mapred.isCollast; import org.apache.hadop public class MRExample {
 public static class LoadPages extends MapReduceBase implements Mapper<LongWritable, Text, Text, Text> { Reporter reporter) throws IOException { Reporter reporter) throws IOException {
 // Pull the key out
 String line = val.tosting();
 String key = line.aubstring() firstComma);
 String value = line.substring() firstComma);
 String value = line.substring();
 // Prepend an index to the value so we know which file
 // it came from.
 Text outVal = new Text("1 = value);
 oc.collect(outExey, outVal); }
public static class LoadAndFilterUsers extends MapReduceBase
 implements Mapper<LongWritable, Text, Text, Text> { Reporter reporter) throws IOException {
 // Pull the key out
 String line = val.toString();
 int firstComma = line.indexof(');
 String = Integer.parteInt(value);
 if (age < 18 || age > 25) return;
 String key = line.substring(0, firstComma);
 Text outkey = new Text(key);
 // Prepend an index to the value so we know which file
 // it came from.
 Text('2' + value);
 gc.collect(outKey, outVal); }
public static class Join extends MapReduceBase
implements Reducer<Text, Text, Text, Text, Text> { public void reduce(Text key, Iterator<Text> iter, OutputCollector<Text, Text> oc, Reporter reporter) throws IOException { // For each value, figure out which file it's from and store it // accordingly. List<String> first = new ArrayList<String>(); List<String> second = new ArrayList<String>(); 

```
reporter.setStatus("OK");
                          3
                          }
                  }
         public void map(
                         lic void map(
    Text k,
    Text val,
    Outputcollector<Text, LongWritable> oc,
    Reporter reporter) throws IOException {
    // Find the url
    // Find the url
    int scondocomma = line.indexOf(',');
    int firstComma = line.indexOf(',');
    int scondocomma = line.indexOf(',');
    int scondocomma = line.indexOf(',');
    firstComma, secondComma);
    // drop the rest of the record, I don't need it anymore,
    // just pass a 1 for the combiner/reducer to sum instead.
    Text outRey = new Text(key);
    oc.collect(outKey, new LongWritable(lL));

                  }
          public static class ReduceUrls extends MapReduceBase
implements Reducer<Text, LongWritable, WritableComparable,
Writable> {
                  public void reduce(
                          iic void reduce(
    Text Key,
    Tetrator:LongWritable> iter,
    OutputCollector:WritableComparable, Writable> oc,
    Reporter reporter) throws IOException {
    // Add up all the values we see
                           long sum = 0;
while (iter.hasNext()) {
    sum += iter.next().get();
    reporter.setStatus("OK");
                          oc.collect(key, new LongWritable(sum));
                  }
         }
public static class LoadClicks extends MapReduceBase
    implements Mapper<WritableComparable, Writable, LongWritable,</pre>
Text> {
                 public void amp(
    writableComparable key,
    writable val,
    OutputCollector<LongWritable, Text> oc,
    Reporter reporter) throws IOException {
    oc.collect((LongWritable)val, (Text)key);
                  }
         public static class LimitClicks extends MapReduceBase
    implements Reducer<LongWritable, Text, LongWritable, Text> {
                   int count = 0;
public void reduce(
                           LongWritable key,
Iterator<Text> iter,
OutputCollector<LongWritable, Text> oc,
                           Reporter reporter) throws IOException {
                           count++;
                          }
                  }
         jubics static void main(String[] args) throws IOException {
    JobConf lp = new JobConf(MEKExample.class);
    lp.setJobName("Load Pages");
    lp.setJobName("Load Pages");
                                                                                                                                                                    3
```

lp.setOutputKeyClass(Text.class); lp.setApperClass(LoadPages.class); PileInputFormat.addInputPath(lp, new user/gates/pages")); FileOutputFormat.setOutputPath(lp, new Path("/user/gates/tmp/indexed\_pages")); lp.setNumReduceTasks(0); Job loadPages = new Job(lp); Path("/ JobConf lfu = new JobConf(MRExample.class); lfu.setJobName("Load and Filter Users"); lfu.setInputFormat(TextInputFormat.class); JobConf join = new JobConf(MRExample.class); join.setJobName("Join Users and Pages"); join.setInputFormat(KeyValuefextInputFormat.class); join.setOutputKeyClass(Tet.class); join.setMapperClass(IdentityMapper.class); join.setMapperClass(Join.class); FileInputFormat.addInputPath(join, new Path("Juser/gates/tmp/indexed\_pages")); FileInputFormat.addInputFath(join, new Path(\*\userVates/userVates/vate JobConf group = new JobConf(MRExample.class); group.setJobName("Group URLs"); group.setUobName("Group URLs"); group.setOutputKey(Lass[Text.clts]); group.setOutputValueClass(LongWritable.class); group.setOutputValueClass(LongWritable.class); group.setMapperClass(LongJoined.class); group.setKapperClass(LongJoined.class); group.setCombinerClass(ReduceDris.class); group.setLeducerClass(ReduceDris.class); Path("/user/gates/tmp/joined")); FileOutputFormat.setOutputFath(group, new Path("/user/gates/tmp/grouped")); group.setNumReduceTasks(50); Job groupJob = new Job(group); groupJob.addDependingJob(joinJob); globpol.adu/geniningOo(); JobConf topl00 = new JobConf(MKExample.class); topl00.setJobName("Top 100 sites"); topl00.setJoutputValue(SequenceFileInputPormat.class); topl00.setOutputValueClass(Text.class); topl00.setVoutputValueClass(Text.class); topl00.setKapperClass(LoadClicks.class); topl00.setKapperClass(LoadClicks.class); topl00.setKeducerClass(LimitClicks.class); topl00.setKeducerClass(LimitClicks.class); Path("/user/gates/tmp/grouped")); Path("/user/gates/tmp/grouped")); Joblimit = new JobtCoruserIs(LimitClicks.class); path("/user/gates/tmp/grouped")); Joblimit = new JobtCoruserIs(LimitClicks.class); Joblimit = new JobtCoruserIs(LimitClicks); JobtCoruseLimitClicks); JobtCoruseLimitClicks; JobtCoruseLi Job limit = new Job(top100); limit.addDependingJob(groupJob); 

#### 170 lines of code, 4 hours to write



# In Pig Latin

```
Users = load 'users' as (name, age);
Fltrd = filter Users by
        age >= 18 and age <= 25;
Pages = load 'pages' as (user, url);
Jnd = join Fltrd by name, Pages by user;
Grpd = group Jnd by url;
Smmd = foreach Grpd generate group,
       COUNT (Jnd) as clicks;
Srtd = order Smmd by clicks desc;
Top5 = limit Srtd 5;
store Top5 into `top5sites';
```

# 9 lines of code, 15 minutes to write



## But can it fly?





## **Essence of Pig**

- Map-Reduce is too low a level to program, SQL too high
- Pig Latin, a language intended to sit between the two:
  - Imperative
  - Provides standard relational transforms (join, sort, etc.)
  - Schemas are optional, used when available, can be defined at runtime
  - User Defined Functions are first class citizens
  - Opportunities for advanced optimizer but optimizations by programmer also possible



# How It Works





# **Cool Things We've Added In the Last Year**

- Multiquery Ability to combine multiple group bys into a single MR job (0.3)
- Merge join If data is already sorted on join key, do join via merge in map phase (0.4)
- Skew join Hash join for data with skew in join key. Allows splitting of key across multiple reducers to handle skew. (0.4)
- Zebra Contrib project that provides columnar storage of data (0.4)
- Rework of Load and Store functions to make them much easier to write (0.7, branched but not released)
- Owl, a metadata service for the grid (committed, will be released in 0.8).



# **Fragment Replicate Join**

Aka "Broakdcast Join"







Pages = load 'pages' as (user, url); Jnd = join Pages by user, Users by name using "replicated";







Pages = load 'pages' as (user, url); Jnd = join Pages by user, Users by name using "replicated";













```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Users by name, Pages by user;
```





```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Users by name, Pages by user;
```



















```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Pages by user, Users by name using "skewed";
```





```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Pages by user, Users by name using "skewed";
```























```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Pages by user, Users by name using "merge";
```





```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Pages by user, Users by name using "merge";
```





```
Users = load 'users' as (name, age);
Pages = load 'pages' as (user, url);
Jnd = join Pages by user, Users by name using "merge";
```



# **Multi-store script**





## Multi-Store Map-Reduce Plan





# What are people doing with Pig

- At Yahoo ~70% of Hadoop jobs are Pig jobs
- Being used at Twitter, LinkedIn, and other companies
- Available as part of Amazon EMR web service and Cloudera Hadoop distribution
- What users use Pig for:
  - Search infrastructure
  - Ad relevance
  - Model training
  - User intent analysis
  - Web log processing
  - Image processing
  - Incremental processing of large data sets



# What We're Working on this Year

- Optimizer rewrite
- Integrating Pig with metadata
- Usability our current error messages might as well be written in actual Latin
- Automated usage info collection
- UDFs in python



# **Research Opportunities**

- Cost based optimization how does current RDBMS technology carry over to MR world?
- Memory Usage given that data processing is very memory intensive and Java offers poor control of memory usage, how can Pig be written to use memory well?
- Automated Hadoop Tuning Can Pig figure out how to configure Hadoop to best run a particular script?
- Indices, materialized views, etc. How do these traditional RDBMS tools fit into the MR world?
- Human time queries Analysts want access to the petabytes of data available via Hadoop, but they don't want to wait hours for their jobs to finish; can Pig find a way to answer analysts question in under 60 seconds?
- Map-Reduce-Reduce Can MR be made more efficient for multiple MR jobs?
- How should Pig integrate with workflow systems?
- See more: <a href="http://wiki.apache.org/pig/PigJournal">http://wiki.apache.org/pig/PigJournal</a>



## Learn More

- Visit our website: <u>http://hadoop.apache.org/pig/</u>
- On line tutorials
  - From Yahoo, <a href="http://developer.yahoo.com/hadoop/tutorial/">http://developer.yahoo.com/hadoop/tutorial/</a>
  - From Cloudera, <u>http://www.cloudera.com/hadoop-training</u>
- A couple of Hadoop books are available that include chapters on Pig, search at your favorite bookstore
- Join the mailing lists:
  - pig-user@hadoop.apache.org for user questions
  - <u>pig-dev@hadoop.apache.com</u> for developer issues
- Contribute your work, over 50 people have so far

