CSE 544 Principles of Database Management Systems

Magdalena Balazinska Winter 2009 Lecture 15 - Data Warehousing: CStore

References

- Column-Stores vs Row-Stores: How Different Are They Really? D. Abadi, S. Madden, and N. Hachem. SIGMOD'08
- Sections on the final: 1, 4, 5.1-5.3, and 6

From Row-Store to Column-Store



Rows stored contiguously on disk (+ tuples headers) Columns stored contiguously on disk (no headers needed) 3

More Detailed Example



Column-Store Optimizations

- Late tuple materialization (3X improvement)
 - Process individual columns as long as possible
 - Merge columns into complete tuples as late as possible
- Block iteration (1.5X)
 - Pass blocks of values between ops instead of individual tuples
- Compression: e.g., run-length encoding of columns (10X)
- Invisible joins (1.5X)

Late Tuple Materialization

Ex:SELECT R.b from R where R.a=X and R.d=Y

Early materialization

Late materialization





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Compression Example



Simulating a Column-Store DBMS in a Row-Store DBMS

- Vertical partitioning
 - Two-column tables: (key, attribute)
- Index-only plans
 - Create a B+ tree index on each attribute
 - Answer queries using indexes only, without reading actual data
- Materialized views
 - Each view contains a subset of columns

Performance Comparison

• See Section 6 in the paper

Conclusion

- Column-store DBMS outperforms row-store DBMS
 - Measured on a data warehousing benchmark (SSBM)
- Late materialization and compression are key factors
- Difficult to simulate a column-store in a row-store
 - Tuple overheads cause data blow-up
 - Column joins are expensive
 - Hard to get the DBMS to "do the right thing" (e.g., index plans)
- Not the end of the story, however, ... see CIDR'09 paper

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