# Online Convex Optimization



## **Convex surrogate loss functions**

Previous section for the adversarial case suggested using multiplicative weights over the |H| hypotheses, which is completely intractable in practice.

And in the stochastic case we used  $h_t \in \arg\min_{h \in \mathcal{H}} \sum_{s=1}^{t-1} \mathbf{1}\{h(x_s) \neq y_s\}$  which is also intractable to compute!

So it seems we have no practical algorithm! Solution: relax the objective.

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Instead of 
$$\max_{h \in \mathcal{H}} \sum_{t=1}^T \mathbf{1}\{h_t(x_t) \neq y_t\} - \mathbf{1}\{h(x_t) \neq y_t\}$$
  
We use  $\max_{h \in \mathcal{H}} \sum_{t=1}^T \ell(h_t, (x_t, y_t)) - \ell(h, (x_t, y_t))$  with  $\mathcal{H}$  convex

**Example:** Linear classification takes  $\mathcal{H} \subset \mathbb{R}^d$  and  $\ell(h, (x_t, y_t)) = \log(1 + \exp(-y_t h^\top x_t))$ 

# **Convex surrogate loss functions**

Goal: 
$$\max_{h \in \mathcal{H}} \sum_{t=1}^T \ell(h_t, (x_t, y_t)) - \ell(h, (x_t, y_t))$$
 with  $\mathcal{H}$  convex

#### Online gradient descent

Input:  $\mathcal{H} \subset \mathbb{R}^d$ , convex loss function  $\ell$ , step size  $\eta > 0$ 

Initialize: Choose any  $h_1 \in \mathcal{H}$ 

for  $t = 1, 2, \ldots$ 

Player plays  $h_t \in \mathcal{H}$ 

Adversary simultaneously reveals  $(x_t, y_t)$ 

Player pays loss  $\ell_t(h_t) := \ell(h_t, (x_t, y_t))$ 

Player updates  $w_{t+1} = \Pi_{\mathcal{H}}(w_t - \eta \nabla_h \ell_t(h_t))$ 

**Theorem** Online gradient descent satisfies for any  $h_* \in \mathcal{H}$ 

$$\sum_{t=1}^{T} \ell(h_t, (x_t, y_t)) - \ell(h_*, (x_t, y_t)) \le \frac{\|h_*\|_2^2}{2\eta} + \frac{\eta}{2} \sum_{t=1}^{T} \|\nabla_h \ell_t(h_t)\|_2^2$$

#### **Proof**

**Theorem** Online gradient descent satisfies for any  $h_* \in \mathcal{H}$ 

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$$||h_{t+1} - h_*||_2^2 = ||\Pi_{\mathcal{H}}(h_{t+1}) - \Pi_{\mathcal{H}}(h_*)||_2^2$$

$$= ||\Pi_{\mathcal{H}}(h_t - \eta \nabla \ell_t(h_t)) - \Pi_{\mathcal{H}}(h_*)||_2^2$$

$$\leq ||h_t - \eta \nabla \ell_t(h_t) - h_*||_2^2$$

$$= ||h_t - h_*||_2^2 - 2\eta \nabla \ell_t(h_t)^{\top} (h_t - h_*) + \eta^2 ||\nabla \ell_t(h_t)||_2^2$$

$$\leq ||h_t - h_*||_2^2 - 2\eta (\ell_t(h_t) - \ell_t(h_*)) + \eta^2 ||\nabla \ell_t(h_t)||_2^2$$

$$\sum_{t=1}^{T} \left( \ell_{t}(h_{t}) - \ell_{t}(h_{*}) \right) \leq \sum_{t=1}^{T} \frac{\|h_{t} - h_{*}\|_{2}^{2} - \|h_{t+1} - h_{*}\|_{2}^{2}}{2\eta} + \sum_{t=1}^{T} \frac{\eta}{2} \|\nabla \ell_{t}(h_{t})\|_{2}^{2}$$

$$\leq \frac{\|h_{1} - h_{*}\|_{2}^{2}}{2\eta} + \sum_{t=1}^{T} \frac{\eta}{2} \|\nabla \ell_{t}(h_{t})\|_{2}^{2}$$



Given a collection of stocks, let the *i*th stock have price  $S_t(i)$  over time *t*.

You start with  $v_1$  dollars and fractionally invest it into d stocks according to  $p_1 \in \triangle_d$  .

Your portfolio at time 2 is worth 
$$v_2 := \sum_{i=1}^d v_1 p_1(i) r_1(i) = v_1 \langle p_1, r_1 \rangle$$
 dollars

where 
$$r_t(i) = \frac{S_{t+1}(i)}{S_t(i)} = \frac{\text{price of GOOG at time t+1}}{\text{price of GOOG at time t}}$$
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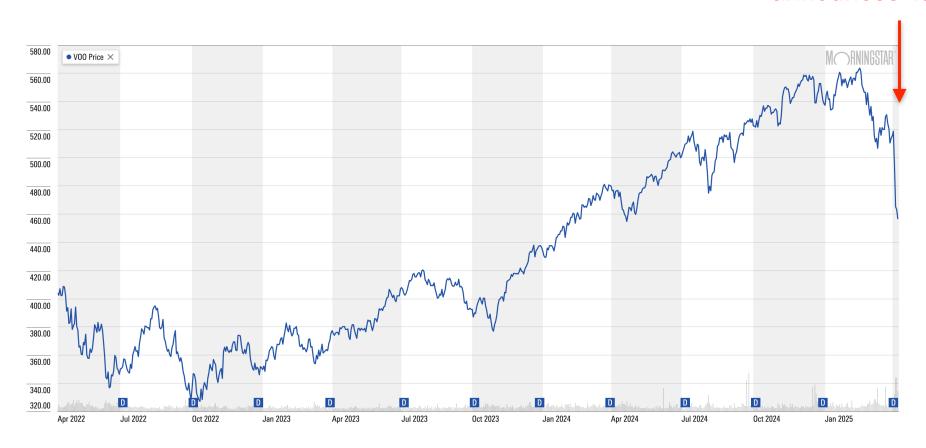
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.

Classical Portfolio Theory (Markowitz 1952): Assume returns  $r_t \in \mathbb{R}^n_+$  are IID with mean  $\mu = \mathbb{E}[r_t]$  and covariance  $\Sigma = \mathbb{E}[(r_t - \mu)(r_t - \mu)^\top]$ . The for a return target  $\bar{r} \geq 0$  solve

$$\min_{p \in \triangle_d} p^{\top} \Sigma p \quad \text{ subject to } \quad p^{\top} \mu \geq \bar{r}$$

In practice, estimate  $\mu$ ,  $\Sigma$  from data. What could possibly go wrong?

# Trump administration announces Tariffs



#### Returns are not an IID stochastic random walk!

Can we model the stock market as an online learning problem and develop an algorithm that is robust to even adversarial returns?

You start with  $v_1$  dollars and fractionally invest it into d stocks according to  $p_1 \in \triangle_d$  .

Your portfolio at time 2 is worth 
$$v_2:=\sum_{i=1}^d v_1p_1(i)r_1(i)=v_1\langle p_1,r_1\rangle$$
 dollars

where 
$$r_t(i) = \frac{\text{price of GOOG at time t+1}}{\text{price of GOOG at time t}}$$
.

After 
$$T$$
 times your portfolio is worth  $v_T = v_1 \prod_{t=1}^{T-1} \langle p_t, r_t \rangle$ .

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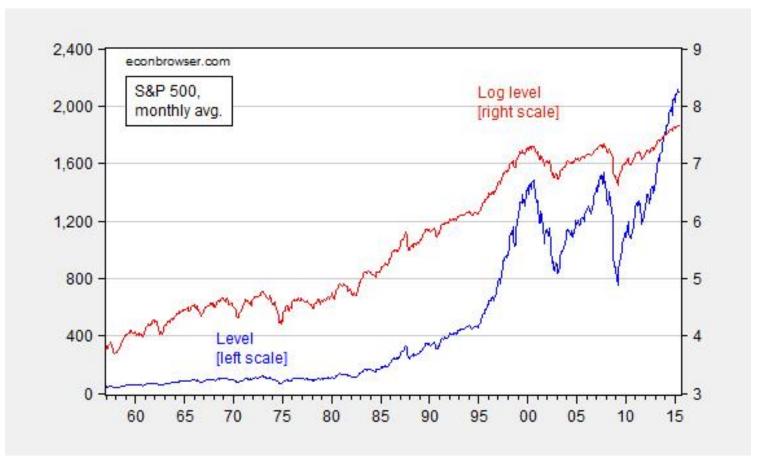
After 
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 times your portfolio is worth  $v_T = v_1 \prod_{t=1}^{T-1} \langle p_t, r_t \rangle$ .

**Goal:** Maximize your return 
$$\frac{v_T}{v_1}$$
, equivalent to  $\log(\frac{v_T}{v_1}) = \sum_{t=1}^{T-1} \log\langle p_t, r_t \rangle$ 

$$\text{Regret} = \max_{p \in \triangle_d} \sum_{t=1}^{T-1} \log \langle p, r_t \rangle - \sum_{t=1}^{T-1} \log \langle p_t, r_t \rangle$$

Regret = 
$$\max_{p \in \triangle_d} \sum_{t=1}^{T-1} \log \langle p, r_t \rangle - \sum_{t=1}^{T-1} \log \langle p_t, r_t \rangle$$

The SP500 (VOO) is an index that weights 500 stocks by their market capitalization. An alternative index (RSP) weights these 500 stocks uniformly  $p=(\frac{1}{500},...,\frac{1}{500})$ .



Adaptive Regnet Bounds

$$\sum_{t} \mathcal{L}(P_{t}, Z_{t}) - \sum_{t} \mathcal{L}(u_{t}, Z_{t}) \leq \sqrt{\left(\sum_{t} ||u_{t} - u_{t-1}||_{1} + 1\right)T}$$

$$\sum_{t} \mathcal{L}(\rho_{t}, z_{t}) - \sum_{t} \mathcal{L}(\rho, z_{t}) \leq \sqrt{\sum_{t} \mathcal{L}(\rho, z_{t}) + i}$$

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109(2)

for t = 1, 2, ...

Player picks  $p_t \in \triangle_d$ 

Adversary simultaneously reveals  $r_t \in \mathbb{R}^d_+$ 

Player pays loss  $\ell_t(p_t) = -\log\langle p_t, r_t \rangle$ 

#### **Exponential weights algorithm**

Initialize:  $w_1 = (1, \ldots, 1) \in \mathbb{R}^d$ 

for 
$$t = 1, 2, ...$$

Player plays 
$$p_t(i) = w_t(i) / \sum_{j=1}^d w_t(j)$$

Adversary simultaneously reveals convex loss  $\ell_t(\cdot)$ Player pays loss  $\ell_t(p_t)$ 

Player updates weights  $w_{t+1}(i) = w_t(i) \exp(-\eta \ell_t(\mathbf{e}_i))$ 

# עחiversal Portfolio Optimization | אין אַרָּיָּי בּיִּי (פּיִּ, רָּיִּ) | Universal Portfolio Optimization |

Regret = 
$$\max_{p \in \triangle_d} \sum_{t=1}^{T-1} \log \langle p, r_t \rangle - \sum_{t=1}^{T-1} \log \langle p_t, r_t \rangle$$

Competes with the single best stock in hindsight!

Theorem: With 
$$\eta = 1$$
 and  $l_t(p) = -\log\langle p, r_t \rangle$ ,  $\max_{i \in [d]} \sum_{t=1}^{t-1} \log\langle \mathbf{e}_i, r_t \rangle - \log\langle p_t, r_t \rangle \leq \log(d)$ 

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$$\log \frac{W_{T+1}}{W_1} = \sum_{t=1}^{T} \log \frac{W_{t+1}}{W_t} \qquad \log \frac{W_{T+1}}{W_1} \ge \log \frac{w_{T+1}(i)}{W_1}$$

$$= \sum_{t=1}^{T} \log \left( \sum_{i=1}^{d} \frac{w_{t+1}(i)}{W_t} \right) \qquad = -\log(d) + \log \left( \prod_{t=1}^{T} \exp(-\eta \ell_t(\mathbf{e}_i)) \right)$$

$$= \sum_{t=1}^{T} \log \left( \sum_{i=1}^{d} \frac{w_t(i) \exp(-\eta \ell_t(\mathbf{e}_i))}{W_t} \right) \qquad = -\log(d) - \sum_{t=1}^{T} \eta \ell_t(\mathbf{e}_i)$$

$$= \sum_{t=1}^{T} \log \left( \sum_{i=1}^{d} p_t(i) \exp(-\eta \ell_t(\mathbf{e}_i)) \right) \qquad = -\log(d) + \sum_{t=1}^{T} \log\langle \mathbf{e}_i, r_t \rangle$$

$$= \sum_{t=1}^{T} \log \left( \sum_{i=1}^{d} p_t(i) \exp(\log\langle \mathbf{e}_i, r_t \rangle) \right)$$

$$= \sum_{t=1}^{T} \log \langle p_t, r_t \rangle \qquad \Longrightarrow \quad \max_{i \in [d]} \sum_{t=1}^{T} \log\langle \mathbf{e}_i, r_t \rangle - \log\langle p_t, r_t \rangle \le \log(d)$$

Regret = 
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Is competing against single best stock a good benchmark? Consider just 2 stocks:

$$r_t(1) = (2, \frac{1}{2}, 2, \frac{1}{2}, 2, \frac{1}{2}, \dots)$$
  
 $r_t(2) = (\frac{1}{2}, 2, \frac{1}{2}, 2, \frac{1}{2}, 2 \dots)$ 

$$\prod_{t=1}^{T} \langle \mathbf{e}_i, r_t \rangle = 1$$

$$\prod_{t=1}^{T} \left\langle \begin{bmatrix} 1/2 \\ 1/2 \end{bmatrix}, r_t \right\rangle = \left( \left( \frac{1}{2} \right)^2 + 1 \right)^{T/2}$$

How do we compete with any  $p \in \triangle_d$ ?

# Online Gradient Descent

For conver less l, domain A w/  $\sup_{a \in A} ||a||_2 \le R$   $\sum_{t=1}^{T} l(a_t, z_t) - l(a, z_t) \le RGJT$   $\sum_{t=1}^{Sup} ||\nabla l(a, z)||_2 \le G$   $l(a, z) = -\log((a, z))$   $\nabla_a l(a, z) = -\frac{1}{(a, z)} \cdot 2$ 

$$\sum_{\ell=1}^{T} \log(\langle P_{\ell}, r_{\ell} \rangle) - \log(\langle P_{\ell}, r_{\ell} \rangle) \leq \sqrt{dT}$$

 $\sum_{\substack{a \in A \\ e=1}}^{m \cdot v} \frac{1}{2!} \left( a_{\ell}, z_{\ell} \right) - \left( a^{e}, z_{\ell} \right) \leq \sum_{\substack{a \in A \\ e=1}}^{m \cdot v} \frac{1}{2!} \left( a_{\ell}, z_{\ell} \right) \left( a_{\ell}, z$ 

 $= \max_{i \in [J]} \sum_{t=1}^{+} \langle \alpha_{\epsilon}, \widetilde{z}_{t} \rangle - \langle e_{i}, \widetilde{z}_{\epsilon} \rangle$ 

 $a^{\alpha} = \underset{a \in \Delta_{d}}{\operatorname{argman}} \sum_{\ell \in I} \nabla l(a_{\ell}, z_{\ell}) \alpha$   $= \underset{\ell \in I}{\operatorname{ac}} \Delta_{d} \underset{\ell \in I}{\operatorname{Expon}} \underset{\text{losse } I}{\operatorname{Weights}}$   $= \sum_{\ell \in I} log(\langle P_{\ell}, \Gamma_{\ell} \rangle) - log(\langle P_{\ell}, \Gamma_{\ell} \rangle) \leq \sqrt{log(d)} T \qquad \text{who gradies of losse } I$ 



Fix a convex set  $\mathscr{A}$  and a convex loss function  $l(\cdot,z):\mathscr{A}\to\mathbb{R}$  for each  $z\in\mathscr{Z}$ .

#### Theorem:

For any 
$$\eta > 0$$
 and  $l(\cdot, \cdot) \in [0,1]$  we have  $\max_{a \in \mathcal{A}} \sum_{t=1}^{I-1} l(a_t, z_t) - l(a, z_t) \le \frac{d \log(T)}{\eta} + \frac{\eta T}{8} + 1$ 

#### Continuous Exponential weights algorithm

Initialize:  $w_1(a) = 1$  for all  $a \in \mathcal{A}$ 

for 
$$t = 1, 2, ...$$

Player plays  $a_t = \mathbb{E}_{A \sim p_t}[A]$  where  $p_t(a) = \frac{w_t(a)}{\int_{a \in \mathcal{A}} w_t(a) da}$ 

Adversary simultaneously reveals  $z_t$  and convex loss  $\ell(\cdot, z_t)$ 

Player pays loss  $\ell(a_t, z_t)$ 

Player updates weights  $w_{t+1}(a) = w_t(a) \exp(-\eta \ell(a, z_t))$ 

Let 
$$\gamma = \frac{1}{T}$$
,  $a^* \in \arg\min_{a \in \mathcal{A}} \sum_{t=1}^{T} \ell(a, z_t)$ ,  $\mathcal{N}_{\gamma} \{ (1 - \gamma)a^* + \gamma a, a \in \mathcal{A} \}$ 

$$\begin{split} \log \frac{W_{T+1}}{W_1} &= \log \left( \frac{\int_{a \in \mathcal{A}} w_{T+1}(a) da}{\int_{a \in \mathcal{A}} 1 da} \right) & \log \frac{W_{t+1}}{W_t} &= \log \left( \int_{\mathcal{A}} \frac{w_t(a)}{W_t} \exp(-\eta \ell(a, z_t)) da \right) \\ &\geq \log \left( \frac{\int_{a \in \mathcal{N}_{\gamma}} w_{T+1}(a) da}{\int_{a \in \mathcal{A}} 1 da} \right) &= \log \left( \mathbb{E} \exp(-\eta \ell(A, z_t)) \right) \text{ where } \mathbb{P}(A = a) = \frac{w_t(a)}{W_t} \\ &= \log \left( \frac{\int_{a \in \mathcal{N}_{\gamma}} \exp\left(-\eta \sum_{t=1}^T \ell(a, z_t)\right) da}{\int_{a \in \mathcal{A}} 1 da} \right) &\leq -\eta \mathbb{E}\ell(A, z_t) + \frac{\eta^2}{8} \text{ (Hoeffding's lemma)} \\ &= \log \left( \frac{\int_{a \in \mathcal{N}_{\gamma}} \exp\left(-\eta \sum_{t=1}^T \ell((1 - \gamma)a^* + a, z_t)\right) da}{\int_{a \in \mathcal{A}} 1 da} \right) &= -\eta \ell(a_t, z_t) + \frac{\eta^2}{8}. \end{split}$$

$$&\geq \log \left( \frac{\int_{a \in \mathcal{A}} \exp\left(-\eta \sum_{t=1}^T \ell((1 - \gamma)a^* + \gamma a, z_t)\right) \gamma^d da}{\int_{a \in \mathcal{A}} 1 da} \right) \\ &\geq \log \left( \frac{\int_{a \in \mathcal{A}} \exp\left(-\eta \sum_{t=1}^T \ell((a^*, z_t) + \gamma \ell(a, z_t))\right) \gamma^d da}{\int_{a \in \mathcal{A}} 1 da} \right) \\ &\geq \log \left( \frac{\int_{a \in \mathcal{A}} \exp\left(-\eta \sum_{t=1}^T \ell(a^*, z_t) + \gamma \ell(a, z_t)\right) \gamma^d da}{\int_{a \in \mathcal{A}} 1 da} \right) \\ &= d \log \gamma - \eta \sum_{t=1}^T \ell(a^*, z_t) - \eta \gamma T \end{split}$$

Proof due to Sebastien Bubeck's Introduction to Online Optimization

Fix a convex set  $\mathscr{A}$  and a convex loss function  $l(\cdot,z):\mathscr{A}\to\mathbb{R}$  for each  $z\in\mathscr{Z}$ .

#### Theorem:

With 
$$\eta = 1$$
 and  $l(a, z) = -\log\langle a, z \rangle$ ,  $\max_{p \in \triangle_d} \sum_{t=1}^{t-1} \log\langle p, r_t \rangle - \log\langle p_t, r_t \rangle \le 1 + d\log(T)$ 

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Let 
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$$\log \frac{W_{T+1}}{W_1} = \log \left( \frac{\int_{a \in \mathcal{A}} w_{T+1}(a) da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

$$= \log \left( \frac{\int_{a \in \mathcal{A}} \exp(-\eta \sum_{t=1}^{T} \ell(a, z_t)) da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

$$= \log \left( \frac{\int_{a \in \mathcal{A}} \prod_{t=1}^{T} \langle a, z_t \rangle da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

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$$\geq \log \left( \frac{\int_{a \in \mathcal{N}_{\gamma}} \prod_{t=1}^{T} \langle a, z_t \rangle da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

$$\geq \log \left( \frac{\int_{a \in \mathcal{N}_{\gamma}} \prod_{t=1}^{T} \langle (1 - \gamma) a^* + a, z_t \rangle da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

$$= \log \left( \frac{\int_{a \in \mathcal{A}} \gamma^d \prod_{t=1}^{T} \langle (1 - \gamma) a^* + \gamma a, z_t \rangle da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

$$\geq \log \left( \frac{\int_{a \in \mathcal{A}} \gamma^d \prod_{t=1}^{T} \langle (1 - \gamma) \prod_{t=1}^{T} \langle a^*, z_t \rangle + \gamma \prod_{t=1}^{T} \langle a, z_t \rangle \right) da}{\int_{a \in \mathcal{A}} 1 da} \right)$$

 $\geq -d \log(1/\gamma) + \log(1-\gamma) + \sum \langle a^*, z_t \rangle$ 

$$\log \frac{W_{T+1}}{W_1} = \sum_{t=1}^{T} \log \frac{W_{t+1}}{W_t}$$

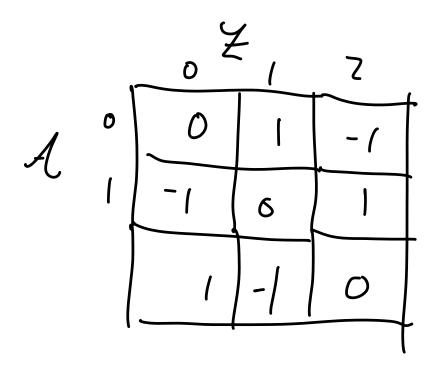
$$= \sum_{t=1}^{T} \log \left( \int_{\mathcal{A}} \frac{w_t(a)}{W_t} \exp(-\eta \ell(a, z_t)) da \right)$$

$$= \sum_{t=1}^{T} \log \left( \mathbb{E}_{A \sim p_t} [\exp(-\eta \ell(A, z_t))] \right)$$

$$= \sum_{t=1}^{T} \log \langle a_t, z_t \rangle$$



Consider a two-player zero-sum game, like matching pennies:



The first player wants to minimize, the second player wants to maximize.

They achieve 
$$\min_{a \in \mathcal{A}} \max_{z \in \mathcal{Z}} L(a, z)$$

But this fails to capture the simultaneous play. Rewrite as

$$\min_{q \in \triangle_{\mathscr{A}}} \max_{z \in \mathscr{Z}} \mathbb{E}_{a \sim q}[L(a, z)]$$

The first player wants to minimize, the second player wants to maximize.

$$\min_{q_1,q_2 \in \triangle_{\mathscr{A}}} \max_{z_1,z_2 \in \mathscr{Z}} \mathbb{E}_{a_1 \sim q(\mathbf{z}_1)} \mathbb{E}_{a_2 \sim q(z_1,\mathbf{z}_2)} [L(\underline{a_1},z_1,a_2,z_2)]$$

We can think of online learning as a two-player game

$$L(a_1, z_1, a_2, z_2) = \sum_{t=1}^{2} \ell(a_t, z_t) - \inf_{a^* \in \mathcal{A}} \sum_{t=1}^{2} \ell(a^*, z_t)$$

The first player wants to minimize, the second player wants to maximize.

$$\mathcal{V}^{obliv}(\mathcal{A},\mathcal{Z}) := \min_{\substack{q_1,\ldots,q_T\\z_1,\ldots,z_T \in \mathcal{Z}}} \max_{\substack{z_1,\ldots,z_T \in \mathcal{Z}}} \mathbb{E}_{a_1 \sim q_1,\ldots,a_T \sim q_T} \left[ \sum_{t=1}^T \ell(a_t,z_t) - \inf_{a^*} \sum_{t=1}^T \ell(a^*,z_t) \right] \leq 7 \text{RM}$$

$$\Re(\mathscr{A},\mathscr{Z}) = \sup_{z_1,\dots,z_T \in \mathscr{Z}} \mathbb{E}_{\epsilon} \left[ \sup_{a \in \mathscr{A}} \sum_{t=1}^T \mathscr{C}(a,z_t) \epsilon_t \right] \qquad \qquad \epsilon_t \in \{-1,1\} \quad \text{with equal probability}$$

$$R(A,Z) \leq V \leq 2R(A,Z)$$