

Lower Bounds

$$(\Omega, \mathcal{F}, P) \quad (\Omega, \mathcal{F}, Q)$$

$$H_0: X_t \sim P = \mathcal{N}(0, 1) \\ H_1: X_t \sim Q = \mathcal{N}(\Delta, 1) \quad t = 1, 2, \dots, T$$

How big does T need to be to decide whether H_0 , or H_1 is true?

If we define $\hat{\mu}_T = \frac{1}{T} \sum_{t=1}^T X_t$ then w.p. $\geq 1 - \delta$

$$|\hat{\mu}_T - \mathbb{E}[X_t]| \leq \sqrt{\frac{2 \log(2/\delta)}{T}} = \Delta/2$$

Choose $T = \underline{8\Delta^{-2} \log(2/\delta)}$, compute $\hat{\mu}_T$

and decide H_1 if $\hat{\mu}_T > \Delta/2$, H_0 if $\hat{\mu}_T < \Delta/2$.

This is correct w.p. $\geq 1 - \delta$.

Def The total Variation between P, Q is

$$TV(P, Q) = \sup_{A \in \mathcal{F}} |P(A) - Q(A)| = \int |p(x) - q(x)| dx$$

Def The Kullback-Liebler Divergence from P, Q is

$$KL(P|Q) = \int \log\left(\frac{dP}{dQ}(\omega)\right) dP(\omega) = \int \log\left(\frac{p(x)}{q(x)}\right) p(x) dx$$

Lemma (Pinsker's inequality) $TV(P, Q) \leq \sqrt{KL(P||Q)/2}$.

$$\leq \sqrt{1 - e^{-KL(P||Q)/2}}$$

Aside: Likelihood of $X_1, \dots, X_T$ under $\mathcal{N}(\mu, 1)$

is $\prod_{t=1}^T \frac{1}{\sqrt{2\pi}} \exp\left(-\frac{(X_t - \mu)^2}{2}\right)$ so likelihood

ratio of H_1 to H_0 is $\prod_{t=1}^T \frac{\exp\left(-\frac{(X_t - \Delta)^2}{2}\right)}{\exp\left(-\frac{X_t^2}{2}\right)} = \frac{q(\{X_t\}_t^T)}{q(\{X_t\}_{t=1}^T)}$

The likelihood ratio test (LRT) says

pick $q(H_1)$ if $\prod_{t=1}^T \frac{q(X_t)}{p(X_t)} > \frac{1}{\delta}$ and $p(H_0)$ otherwise.

Make an error under P (i.e. $X_t \sim P \forall t$)

w/ prob

$$P_{X_t \sim P} \left(\prod_{t=1}^T \frac{q(X_t)}{p(X_t)} > \frac{1}{\delta} \right) \leq \delta \cdot \mathbb{E}_P \left[\prod_{t=1}^T \frac{q(X_t)}{p(X_t)} \right] = \delta.$$

$$X_1 \sim P \quad \mathbb{E}_{X_1 \sim P} \left[\frac{q(X_1)}{p(X_1)} \right] = \int p(x) \cdot \frac{q(x)}{p(x)} dx = 1$$

$$P_{X_t \sim q} \left(\prod_{t=1}^T \frac{q(X_t)}{p(X_t)} < \delta \right) \leq \delta.$$

$$P\left(\prod_{t=1}^T \frac{q(x_t)}{p(x_t)} > \frac{1}{\delta}\right) = P_{x_t \sim p}\left(\sum_{t=1}^T \log\left(\frac{q(x_t)}{p(x_t)}\right) > \log(1/\delta)\right)$$

$$\mathbb{E}_{x_t \sim q}\left[\sum_{t=1}^T \log\left(\frac{q(x_t)}{p(x_t)}\right)\right] = T \text{KL}(q|p)$$

$$\text{Under } Q, \quad T \approx \frac{\log(1/\delta)}{\text{KL}(Q|P)}$$

$$P = \mathcal{N}(0, 1)$$

$$Q = \mathcal{N}(\Delta, 1)$$

$$\text{KL}(Q, P) = \Delta^2/2$$

Sequential Probability Ratio Test (SPRT)

$$\mathcal{T} = \min\left\{T: \sum_{t=1}^T \log\left(\frac{q(x_t)}{p(x_t)}\right) > \log(1/\delta)\right\}$$

$$\log(1/\delta) \approx \underbrace{\mathbb{E}_q\left[\sum_{t=1}^{\mathcal{T}} \log\left(\frac{q(x_t)}{p(x_t)}\right)\right]}_{\text{Wald's Identity.}} = \mathbb{E}_q[\mathcal{T}] \text{KL}(q|p)$$

$$\mathbb{E}_q[\mathcal{T}] \leq \frac{\log(1/\delta)}{\text{KL}(q|p)}.$$

Wald's Identity.

$$\mathbb{E}_p[\mathcal{T}] \leq \frac{\log(1/\delta)}{\text{KL}(p|q)}$$

Theorem For every $T \geq n$ $\exists$ bandit instance $\mathcal{V}_\theta$ where
 $\mathcal{V}_{\theta,i} = \mathcal{N}(\theta_i, 1)$ such that for a policy π :

$$\inf_{\pi} \max_{\theta} \max_i \mathbb{E}_{\theta} \left[\sum_{t=1}^T x_{t,i} - X_{t,I_t} \right] = \inf_{\pi} \max_{\theta} \max_i \mathbb{E}_{\theta} \left[\sum_{t=1}^T \theta_i - \theta_{I_t} \right] \geq \sqrt{\frac{(n-1)T}{256}}.$$

Proof Let $\theta = (\Delta, 0, \dots, 0)$. For any π

$$\exists \hat{i} \in [n] \setminus \{1\} \text{ s.t. } \mathbb{E}_{\theta} [T_{\hat{i}}] \leq \frac{T}{n-1}.$$

Define $\theta' = (\Delta, 0, \dots, 2\Delta, 0, \dots, 0)$

$$P(X > \alpha) \leq \frac{\mathbb{E}[X]}{\alpha}$$

$\uparrow \hat{i}$ location.

$$\theta = \begin{matrix} \dots & \dots & \dots & \dots & \dots \\ & & & & \downarrow \\ \theta' = & \dots & \dots & & \dots \end{matrix}$$

$$\begin{aligned} \mathbb{E}_{\theta'} [R_T] &= \sum_{i=1}^n (\theta'_i - \theta_i) \mathbb{E}_{\theta'} [T_i] \geq (\theta'_{\hat{i}} - \theta_{\hat{i}}) \mathbb{E}_{\theta'} [T_{\hat{i}}] \\ &\geq (\theta'_{\hat{i}} - \theta_{\hat{i}}) P_{\theta'}(T_{\hat{i}} > T/2) T/2 \\ &= \Delta P_{\theta'}(T_{\hat{i}} > T/2) T/2 \end{aligned}$$

$$\mathbb{E}_{\theta} [R_T] = \sum_{i=1}^n (\theta_1 - \theta_i) \mathbb{E}_{\theta} [T_i] = \sum_{i=2}^n \Delta \mathbb{E}_{\theta} [T_i] = \Delta \mathbb{E}_{\theta} \left[\sum_{i=1}^n T_i \right]$$

$$\geq \Delta P_{\theta} \left(\sum_{i=1}^n T_i > T/2 \right) T/2$$

$$= \Delta P_{\theta} (T_1 < T/2) T/2$$

$$= \Delta (1 - P_{\theta}(T_1 > T/2)) T/2$$

$$\begin{aligned}
\max_{\mu \in \{\theta, \theta'\}} \mathbb{E}_{\mu} [R_T] &\geq \frac{1}{2} (\mathbb{E}_{\theta'} [R_T] + \mathbb{E}_{\theta} [R_T]) \\
&\geq \frac{\Delta T}{4} (\mathbb{P}_{\theta'}(T_1 > T/2) + 1 - \mathbb{P}_{\theta}(T_1 > T/2)) \\
&\geq \frac{\Delta T}{4} (1 - |\mathbb{P}_{\theta}(T_1 > T/2) - \mathbb{P}_{\theta'}(T_1 > T/2)|) \\
&\geq \frac{\Delta T}{4} (1 - TV(\mathbb{P}_{\theta}, \mathbb{P}_{\theta'})) \\
&\geq \frac{\Delta T}{4} (1 - \sqrt{KL(\mathbb{P}_{\theta} | \mathbb{P}_{\theta'}) / 2}) \\
&\geq \frac{\Delta T}{4} (1 - \sqrt{\Delta^2 \mathbb{E}_{\theta}[T_1^2]}) \\
&\geq \frac{\Delta T}{4} (1 - \sqrt{\Delta^2 \frac{T}{n-1}}) \quad \Delta = \sqrt{\frac{n-1}{4T}} \\
&\geq \sqrt{\frac{(n-1)T}{256}}
\end{aligned}$$

Lemma Let $\nu = (\nu_1, \dots, \nu_n)$ and ν' be two bandit instances, Fix policy π and let $\mathbb{P}_{\nu}$ be measure under ν , and $\mathbb{P}_{\nu'}$ under ν' . Then

$$KL(\mathbb{P}_{\nu} | \mathbb{P}_{\nu'}) = \sum_{i=1}^n \mathbb{E}_{\nu}[T_i(\pi)] KL(\nu_i | \nu'_i)$$

Furthermore, if $\mathcal{S}$ is a stopping time, the above holds w/ $T = \mathcal{S}$.

Proof $(I_1, X_1, I_2, X_2, I_3, X_3, \dots, I_T, X_T) = \mathcal{Z}$

$$dP_v(z) = \prod_{t=1}^T \nu_{I_t}(x_t) \pi(I_t | \{I_s, x_s\}_{s < t})$$

$$KL(P_v | P_{v'}) = \mathbb{E}_v \left[\log \left(\frac{dP_v(z)}{dP_{v'}(z)} \right) \right]$$

$$= \mathbb{E}_v \left[\log \left(\frac{\prod_t \nu_{I_t}(x_t)}{\prod_t \nu'_{I_t}(x_t)} \right) \right]$$

$$= \mathbb{E}_v \left[\sum_{t=1}^T \log \left(\frac{\nu_{I_t}(x_t)}{\nu'_{I_t}(x_t)} \right) \right]$$

$$\mathbb{E}_v \left[\log \left(\frac{\nu_{I_t}(x_t)}{\nu'_{I_t}(x_t)} \right) \right] = \mathbb{E}_v \left[\mathbb{E} \left[\log \left(\frac{\nu_i(x_t)}{\nu'_i(x_t)} \right) \mid I_t = i \right] \right]$$

$$= \mathbb{E}_v \left[KL(\nu_{I_t} | \nu'_{I_t}) \right]$$

$$= \mathbb{E}_v \left[\sum_{i=1}^n \sum_{t=1}^T \mathbb{1}\{I_t = i\} KL(\nu_i | \nu'_i) \right]$$

$$= \sum_{i=1}^n KL(\nu_i | \nu'_i) \mathbb{E}_v \left[\sum_t \mathbb{1}\{I_t = i\} \right]$$

$$= \sum_i \mathbb{E}_v [T_i] KL(\nu_i | \nu'_i),$$

Best arm identification

for $t = 1, 2, \dots, T$

Player chooses $I_t \in [n]$

Nature reveals $X_{t, I_t} \sim \mathcal{N}(\theta_{I_t}, 1)$

If player stops:

Output $\hat{i} \in [n]$

Def We say an algorithm is δ -correct if

$$\mathbb{P}_{\theta_*}(\hat{i} = \arg\max_i \theta_{*,i}) \geq 1 - \delta \quad \text{for all } \theta_* \in \mathbb{R}^d.$$

Recall Action Elim. alg identifies $i_* = \arg\max_i \theta_i$ w.p. $\geq 1 - \delta$

after just $T = c \sum_{i \neq i_*} \Delta_i^{-2} \log\left(\frac{n \log(n)}{\delta}\right)$ pulls $\Delta_i = \theta_{i_*} - \theta_i$.

Fix any $\theta \in \mathbb{R}^n$, assume WLOG $i_* = 1$

$$v_i = \mathcal{N}(\theta_i, 1) \quad \forall i$$

$$v_i^{(j)} = \begin{cases} v_i & \text{if } i \neq j \\ \mathcal{N}(\theta_i + \Delta_i + \varepsilon, 1) & \text{if } i = j \end{cases}$$

Lemma (Data Processing Inequality) Given

a. probability space $(\Omega, \mathcal{F})$ and dist. P, Q

R.V. Z w/ push-forward measure P_Z, Q_Z

$$KL(P_Z | Q_Z) \leq KL(P | Q)$$

$$KL(\underbrace{P_v(\hat{i}=1)}_{\geq 1-\delta} | \underbrace{P_{v^{(i)}}(\hat{i}=1)}_{< \delta}) \leq KL(P_v | P_{v^{(i)}})$$

$$P_v(\hat{i}=1) = p$$

$$P_{v^{(i)}}(\hat{i}=1) = q$$

$$= p \log\left(\frac{p}{q}\right) + (1-p) \log\left(\frac{1-p}{1-q}\right)$$

$$\geq (1-\delta) \log\left(\frac{1-\delta}{\delta}\right) + \delta \log\left(\frac{\delta}{1-\delta}\right)$$

$$\geq \log\left(\frac{1}{2.4\delta}\right)$$

$$\Rightarrow \log\left(\frac{1}{2.4\delta}\right) \leq KL(P_v | P_{v^{(i)}})$$

$$= \mathbb{E}_v[T_j | \mathcal{I}] \underbrace{KL(v_j | v_j^{(i)})}_{\uparrow}$$

$$(\Delta_j + \varepsilon)^2 / 2 =$$

$$\mathcal{N}(\theta_j + \Delta_j + \varepsilon, 1)$$

$$\Rightarrow \mathbb{E}_\nu [T_j(\delta)] \geq \frac{2}{(\Delta_j + \varepsilon)^2} \log\left(\frac{1}{2.4\delta}\right)$$

$$\mathbb{E}_\nu[\mathfrak{J}] \geq \sum_{j=2}^n \mathbb{E}_\nu [T_j(\delta)] \geq \sum_{j=2}^n \frac{2 \log(1/2.4\delta)}{(\Delta_j + \varepsilon)^2}$$

$\varepsilon > 0$ was arbitrary, take $\varepsilon \rightarrow 0$.

Thm Any δ -correct algorithm requires

$$\mathbb{E}_\nu[\mathfrak{J}] \geq \sum_{i \neq i_0} \frac{2 \log(1/2.4\delta)}{\Delta_i^2}.$$

Lemma Law of the Iterated Logarithm (LIL)

Let $Z_t \sim \mathcal{N}(0, 1)$. and $S_t = \sum_{s=1}^t Z_s$

Then $\limsup_{t \rightarrow \infty} \frac{S_t}{\sqrt{2t \log \log(t)}} = 1$ a.s.



$$\mathbb{E}[|S_t|] \asymp \sqrt{t}$$

Using similar techniques one can show:

Thm] Fix any alg s.t. $\max_{\theta} \mathbb{E}_{\theta}[R_T] \leq T^{1-d}$ for any $d > 0$

$$\text{then } \lim_{T \rightarrow \infty} \frac{\mathbb{E}_{\theta}[R_T]}{\log(T)} \geq \sum_{i \in I_p} \Delta_i^{-1}$$