

Fix  $x_1, \dots, x_n \in \mathbb{R}^d$  and  $\theta_* \in \mathbb{R}^d$ .

At  $i$  observe

$$y_i = x_i^T \theta_* + \varepsilon_i, \quad \{\varepsilon_i\} \text{ are IID mean-zero } 1\text{-sub-Gaussian R.V.}$$

Least squares estimate:

$$\begin{aligned} \hat{\theta} &= \arg \min_{\theta} \sum_{i=1}^n (\langle x_i, \theta \rangle - y_i)^2 \\ &= \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} \sum_{i=1}^n x_i y_i \\ &= \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} \left( \sum_{i=1}^n x_i x_i^T \right) \theta_* + \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} \sum_{i=1}^n x_i \varepsilon_i \\ &= \theta_* + \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} \sum_{i=1}^n x_i \varepsilon_i \end{aligned}$$

Observe

$$\{(x_i, y_i)\}_{i=1}^n$$

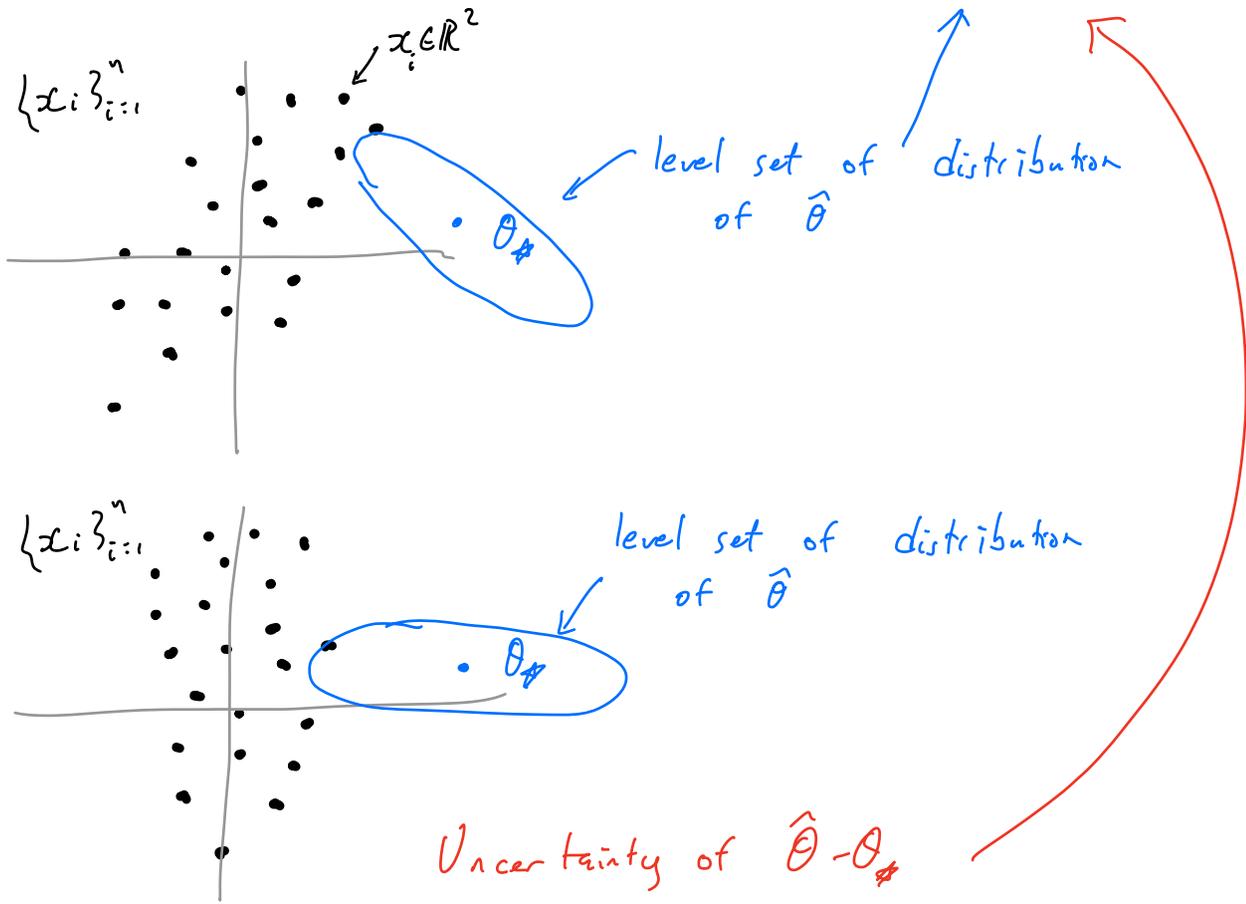
$$E[\hat{\theta}] = \theta_* \quad \text{Cov}(\hat{\theta}) = \left( \sum_{i=1}^n x_i x_i^T \right)^{-1}$$

Fix any  $z \in \mathbb{R}^d$ . With probability at least  $1 - \delta$

$$\langle z, \hat{\theta} - \theta_* \rangle \leq \|z\|_{\left( \sum_{i=1}^n x_i x_i^T \right)^{-1}} \sqrt{2 \log(1/\delta)}. \quad (\|x\|_A^2 = x^T A x)$$

Proof: Show  $\langle z, \hat{\theta} - \theta_* \rangle$  is  $\|z\|_{\left( \sum_{i=1}^n x_i x_i^T \right)^{-1}}$ -sub-Gaussian. Apply Chernoff bound.  $\square$

Intuition: If  $\epsilon_i \sim \mathcal{N}(0, 1)$  then  $\tilde{\theta} \sim \mathcal{N}(\theta_*, (\sum_{i=1}^n x_i x_i^T)^{-1})$



Uncertainty of  $\hat{\theta} - \theta_*$   
 depends only on  $\{x_i\}_{i=1}^n$  and not  $\theta_*$  or  $\{y_i\}$

Experimental Design: Exploits this observation  
 to choose  $\{x_i\}$  in order to obtain a  
 desired covariance shape.

Given choice of  $\{x_i\}_{i=1}^n$ , and then observe  $y_i = \langle x_i, \theta_* \rangle + \epsilon_i$   
 results in  $\hat{\theta}$  w/  $\mathbb{E}[\hat{\theta}] = \theta_*$   $\text{Cov}(\hat{\theta}) = (\sum_i x_i x_i^T)^{-1}$

Given some pool of points  $\mathcal{X} \subset \mathbb{R}^d$

choose  $\{x_i\}_{i=1}^n$

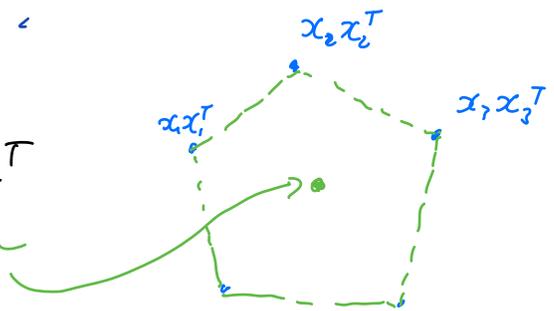
For every choice of points  $\{x_i\}_{i=1}^n \subset \mathcal{X}$ ,  $\exists \lambda \in \Delta_{\mathcal{X}}$

$$\Delta_{\mathcal{X}} = \left\{ p \in \mathbb{R}^{|\mathcal{X}|} : \sum_{x \in \mathcal{X}} p_x = 1, p_x \geq 0 \forall x \in \mathcal{X} \right\}$$

s.t. 
$$\sum_{i=1}^n x_i x_i^T = n \sum_{x \in \mathcal{X}} \lambda_x x x^T$$

why? 
$$\lambda_x = \frac{1}{n} \sum_{i=1}^n \mathbb{1}\{x_i = x\}$$

Define 
$$A(\lambda) = \sum_{x \in \mathcal{X}} \lambda_x x x^T$$



Exp. Design Objectives

- A-optimality 
$$f_A(\lambda) = \text{Tr} \left( A(\lambda)^{-1} \right)$$

$$\mathbb{E} \left[ \|\hat{\theta} - \theta_*\|_2^2 \right] = \text{Tr} \left( \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} \right)$$

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“minimize average error over all directions”

- E-optimality  $f_E(\lambda) = \sup_{u: \|u\|_2 \leq 1} u^T A(\lambda)^{-1} u$

$$\sup_{u: \|u\|_2 \leq 1} \mathbb{E} \left[ \langle u, \hat{\theta} - \theta_* \rangle^2 \right] = \sup_u u^T \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} u$$

"minimize worst-case error direction"

- G-optimality  $f_G(\lambda) = \max_{x \in \mathcal{X}} x^T A(\lambda)^{-1} x$

$$\begin{aligned} \max_{x \in \mathcal{X}} \mathbb{E} \left[ \langle x, \hat{\theta} - \theta_* \rangle^2 \right] &= \max_{x \in \mathcal{X}} x^T \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} x \\ &= \max_{x \in \mathcal{X}} \|x\|_{\left( \sum_{i=1}^n x_i x_i^T \right)^{-1}}^2 \end{aligned}$$

- D-optimality  $f_D(\lambda) = -\log(|A(\lambda)|)$

$$= \log(|A(\lambda)^{-1}|)$$

Recall: if  $Z \sim \mathcal{N}(0, \Sigma)$

then entropy of  $Z = \frac{1}{2} \log(2\pi e |\Sigma|)$

"minimize entropy of  $\hat{\theta}$ "

Lemma (Kiefer-Wolfowitz 1960)

For any  $X \subset \mathbb{R}^d : \dim(\text{span}(X)) = d$ , there exists a  $\lambda^* \in \Delta_X$  such that

-  $\min_{\lambda} f_0(\lambda) = f_0(\lambda^*)$

-  $\min_{\lambda} f_G(\lambda) = f_G(\lambda^*)$

-  $f_G(\lambda^*) = d$

-  $\text{support}(\lambda^*) \leq \frac{(d+1)d}{2}$

Caratheodory  
Theorem.

$$\min_{\lambda} f_G(\lambda) = \min_{\lambda} \max_{x \in X} x^T A(\lambda)^{-1} x$$

$$\geq \min_{\lambda} \sum_{x \in X} \lambda_x x^T A(\lambda)^{-1} x$$

$$= \min_{\lambda} \text{Tr} \left( \underbrace{\sum_{x \in X} \lambda_x x x^T}_{= I} A(\lambda)^{-1} \right)$$

$x \in \mathbb{R}^d$  is  
 $s$ -sparse if

$|\{i \in [d] : x_i \neq 0\}| = s$ .

$$= d$$

Proposition Fix  $\mathcal{X} \subset \mathbb{R}^d$ . Assume  $Y = \langle X, \theta_* \rangle + \varepsilon$  w/

$\varepsilon$  1-sub-Gaussian. If  $\lambda^* = \min_{\lambda \in \Delta_{\mathcal{X}}} f_G(\lambda)$  is a  $\frac{(d+1)d}{2}$ -sparse solution, and we pull arm  $x \in \mathcal{X}$  exactly

$\lceil 3\lambda_x^* \rceil$  times, then w.p.  $\geq 1 - \delta$

$$|\langle x, \hat{\theta} - \theta_* \rangle| \leq \sqrt{\frac{2d \log(2/\delta)}{3}} \quad \forall x \in \mathcal{X}$$

and total # pulls  $\leq 3 + \frac{(d+1)d}{2}$ .  
 $= n$

From above, for any  $x \in \mathcal{X}$  w.p.  $\geq 1 - \delta'$

$$\langle x, \hat{\theta} - \theta_* \rangle \leq \|x\| \left( \sum_{i=1}^n x_i x_i^T \right)^{-1} \sqrt{2 \log(1/\delta')}$$

$$\sum_{i=1}^n x_i x_i^T$$

$$= \sum_{x \in \mathcal{X}} \lceil 3\lambda_x^* \rceil x x^T$$

$$\geq \sum_{x \in \mathcal{X}} 3\lambda_x^* x x^T$$

$$\leq \|x\| \left( 3 \sum_{x \in \mathcal{X}} \lambda_x^* x x^T \right)^{-1} \sqrt{2 \log(1/\delta')}$$

$$\leq \max_{x' \in \mathcal{X}} \|x\| \underbrace{\left( \sum_{x \in \mathcal{X}} \lambda_x^* x x^T \right)^{-1}}_{= \sqrt{d}} \sqrt{\frac{2 \log(1/\delta')}{3}}$$

$A \succeq B$  if

$A - B \succeq 0$  (PSD)

$$= \sqrt{\frac{2d \log(1/\delta')}{3}}$$

$$\begin{aligned} \|x\|_{\left(\sum_{x' \in X} \lambda_{x'} x' x'^T\right)^{-1}} &= \sqrt{x^T \left(\sum_{x' \in X} \lambda_{x'} x' x'^T\right)^{-1} x} \\ &= \frac{1}{\sqrt{3}} \cdot \|x\|_{A(G^*)^{-1}} \end{aligned}$$

**Input:** Finite set  $\mathcal{X} \subset \mathbb{R}^d$ , confidence level  $\delta \in (0, 1)$ .

Let  $\mathcal{X}_1 \leftarrow \mathcal{X}, \ell \leftarrow 1$

**while**  $|\mathcal{X}_\ell| > 1$  **do**

Let  $\hat{\lambda}_\ell \in \Delta_{\mathcal{X}_\ell}$  be a  $\frac{d(d+1)}{2}$ -sparse minimizer of  $f(\lambda) = \max_{x \in \mathcal{X}_\ell} \|x\|_{(\sum_{x \in \mathcal{X}_\ell} \lambda_x x x^\top)^{-1}}^2$

$\epsilon_\ell = 2^{-\ell}, \tau_\ell = 2d\epsilon_\ell^{-2} \log(4\ell^2|\mathcal{X}|/\delta)$

Pull arm  $x \in \mathcal{X}$  exactly  $\lceil \hat{\lambda}_{\ell,x} \tau_\ell \rceil$  times and construct the least squares estimator  $\hat{\theta}_\ell$  using only the observations of this round

$\mathcal{X}_{\ell+1} \leftarrow \mathcal{X}_\ell \setminus \{x \in \mathcal{X}_\ell : \max_{x' \in \mathcal{X}_\ell} \langle x' - x, \hat{\theta}_\ell \rangle > 2\epsilon_\ell\}$

$\ell \leftarrow \ell + 1$

**Output:**  $\mathcal{X}_\ell$

$$\begin{aligned} \text{Regret } R_T &= \max_{x \in \mathcal{X}} \mathbb{E} \left[ \sum_{t=1}^T \langle x, \theta_* \rangle - \langle x_t, \theta_* \rangle \right] \\ &= \sum_{x \in \mathcal{X}} \Delta_x \mathbb{E}[T_x] \end{aligned}$$

$\uparrow$   
 arm algorithm played  
 at time  $t$

$$\Delta_x = \max_{x' \in \mathcal{X}} \langle x' - x, \theta_* \rangle$$

$$T_x = \sum_{t=1}^T \mathbb{1}\{x = x_t\}$$

Assume  $\max_{x \in \mathcal{X}} |\langle x, \theta_* \rangle| \leq 1$ . Let  $x^* = \arg \max_{x \in \mathcal{X}} \langle x, \theta_* \rangle$

Lemma With prob at least  $1 - \delta$ , we have  $x^* \in \mathcal{X}_\ell$

and  $\max_{x \in \mathcal{X}_\ell} \Delta_x \leq 8\epsilon_\ell$  for all  $\ell \in \mathbb{N}$ .

What does linear structure get us?

- If we ignored structure our elim. alg of last time has  $R_T \leq \sqrt{|\mathcal{X}| T \log(T)}$ .

- This algorithm that exploits linear structure,

we will show  $R_T \leq \sqrt{d T \log(|X|T)}$ .

I can "cover"  $X$  with just  $O(2^d)$  landmarks

$$\Rightarrow R_T \leq \sqrt{d^2 T \log(T)}$$