

- sub-Gaussian R.V.s

A mean-zero R.V. Z is σ^2 -sub-Gaussian if $\mathbb{E}[\exp(\lambda Z)] \leq \exp(\sigma^2 \lambda^2 / 2) \quad \forall \lambda \in \mathbb{R}$

- Chernoff bound w/ sub-Gaussian R.V.s:

If X_1, \dots, X_n are IID R.V. w/ $\mathbb{E}[X_i] = \mu$, and $X_i - \mu$ is σ^2 -sub-Gaussian, then

$$\mathbb{P}\left(\frac{1}{n} \sum_{i=1}^n X_i - \mu > \varepsilon\right) \leq \exp\left(-\frac{n\varepsilon^2}{2\sigma^2}\right) = \delta$$

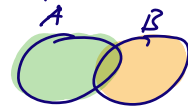
- Hoeffding's lemma: (If $X \in [a, b]$, then take $a=0, b=1 \Rightarrow X - \mathbb{E}[X]$ is $\frac{1}{4}$ -sub-Gaussian)

If $X \in [a, b]$ then $X - \mathbb{E}[X]$ is $\frac{(b-a)^2}{4}$ -sub-Gaussian

- Union Bound (or, sub-additivity of probability)

Let A_1, \dots, A_n be events in a well-defined probability space,

then $\mathbb{P}\left(\bigcup_{i=1}^n \{A_i\}\right) \leq \sum_{i=1}^n \mathbb{P}(A_i)$



Given n treatments., Input horizon $T \in \mathbb{N}$
 for $t=1, 2, \dots, T$
 Patient t arrives w/ illness
 Agent prescribes treatment $I_t \in [n] = \{1, \dots, n\}$
 Observe $X_{I_t, t} = \mathbb{1}\{\text{patient } t \text{ recovers given } I_t\}$

Let $X_{i,t} = \mathbb{1}\{\text{patient } t \text{ recovers given } i\}$.

Assume for each $i \in [n]$, $\{X_{i,t}\}_{t=1}^T$ are IID w/ $\mathbb{E}[X_{i,1}] = \theta_i$

perhaps $X_{i,t} = \mathbb{1}\{e_{i,t} > 1\}$

Questions:

- How long does it take to identify the best treatment $\arg \max_{i \in [n]} \theta_i$ with probability at least $1 - \delta$?

- If $R_T = \max_{i \in [n]} \mathbb{E}\left[\sum_{t=1}^T X_{i,t} - X_{I_t, t}\right]$, how can we minimize R_T ?

- Expected Regret

$$R_T = \max_{i \in [n]} \mathbb{E} \left[\sum_{t=1}^T X_{i,t} - X_{I_t,t} \right] = \sum_{i=1}^n \Delta_i \mathbb{E}[T_i]$$

where $\Delta_i = \max_{j \in [n]} \theta_j - \theta_i$, $T_i = \sum_{t=1}^T \mathbb{1}\{I_t = i\}$

Warm-up: $n=2$. Consider a strategy that plays action $i \in [n]$ round-robin until $T_i = 3$, then plays $\operatorname{argmax}_{i \in [n]} \hat{\theta}_{i,3}$ for the remainder $T - n3$ times.

↖ empirical mean of first 3 plays.

$i \in [n]$

Define $\mathcal{E}_i = \left\{ |\hat{\theta}_{i,3} - \theta_i| \leq \sqrt{\frac{2\sigma^2 \log(2n/\delta)}{3}} \right\}$, $\sigma^2 = \frac{1}{4}$

Lemma $\mathbb{P}(\bigcap_{i=1}^n \mathcal{E}_i) = 1 - \mathbb{P}(\bigcup_{i=1}^n \mathcal{E}_i^c) \geq 1 - \delta$.

Proof $\hat{\theta}_{i,3} = \frac{1}{3} \sum_{s=1}^3 X_{i,s}$, $X_{i,s} \in \{0,1\}$ are iid $\mathbb{E}[X_{i,1}] = \theta_i$

$\Rightarrow (X_{i,s} - \theta_i)$ is σ^2 -sub-Gaussian with $\sigma^2 = 1/4$.

$$\begin{aligned} \mathbb{P}(\mathcal{E}_i^c) &= \mathbb{P}\left(|\hat{\theta}_{i,3} - \theta_i| > \sqrt{\frac{2\sigma^2 \log(2n/\delta)}{3}}\right) \\ &\leq 2 \exp\left(-\frac{3}{2\sigma^2} \left(\sqrt{\frac{2\sigma^2 \log(2n/\delta)}{3}}\right)^2\right) \\ &= 2 \exp(-\log(2n/\delta)) \\ &= \frac{\delta}{n} \end{aligned}$$

$$P\left(\bigcup_{i=1}^n \{E_i^c\}\right) \leq \sum_{i=1}^n P(E_i^c) \leq \sum_{i=1}^n \frac{\delta}{n} = \delta.$$

Because $P\left(\bigcap_{i=1}^n \{E_i\}\right) \geq 1 - \delta$, assume $\bigcap_{i=1}^n E_i$ occurs.

Assume WLOG $\theta_1 \geq \theta_2 \geq \dots \geq \theta_n$

If $\underline{\gamma} = \lceil 8\sigma^2 \Delta_2^2 \log(2n/\delta) \rceil$ then on events $E_1 \cap E_2$

$$\hat{\theta}_{1,\underline{\gamma}} > \theta_1 - \sqrt{\frac{2\sigma^2 \log(2n/\delta)}{\underline{\gamma}}} \quad (E_1)$$

$$> \theta_1 - \frac{\Delta_2}{2}$$

$$= \theta_2 + \frac{\Delta_2}{2}$$

$$\geq \hat{\theta}_{2,\underline{\gamma}} - \sqrt{\frac{2\sigma^2 \log(2n/\delta)}{\underline{\gamma}}} + \frac{\Delta_2}{2} \quad (E_2)$$

$$\geq \hat{\theta}_{2,\underline{\gamma}}. \quad (\star)$$

\Rightarrow With probability at least $1 - \delta$, strategy identifies $\operatorname{argmax}_{i \in [n]} \theta_i$.

$$\theta_1 \geq \theta_2 \geq \dots \geq \theta_n \quad (n=2 \quad \theta_1 > \theta_2)$$

$$R_T = \sum_{i=1}^n \Delta_i \mathbb{E}[T_i] = \Delta_2 \mathbb{E}[T_2] \quad \text{since } n=2 \text{ and } \Delta_1=0.$$

$$= \Delta_2 \mathbb{E}[T_2 (\mathbb{1}\{E_1 \cap E_2\} + \mathbb{1}\{E_1^c \cup E_2^c\})]$$

$$= \Delta_2 \mathbb{E}[T_2 \mathbb{1}\{E_1 \cap E_2\}] + \Delta_2 \mathbb{E}[T_2 \mathbb{1}\{E_1^c \cup E_2^c\}]$$

$$\leq \Delta_2 \underline{\gamma} \mathbb{E}[\mathbb{1}\{E_1 \cap E_2\}] + \Delta_2 (T - \underline{\gamma}) \mathbb{E}[\mathbb{1}\{E_1^c \cup E_2^c\}]$$

Note: on $\mathcal{E}_1 \cap \mathcal{E}_2$ we proved above $\textcircled{\star}$ $\hat{\theta}_{1,\beta} > \hat{\theta}_{2,\beta}$

If $\hat{\theta}_{1,\beta} > \hat{\theta}_{2,\beta}$ then $T_2 = \beta, T_1 = T - \beta$

$$\mathcal{E}_1 \cap \mathcal{E}_2 \Rightarrow \{T_2 = \beta, T_1 = T - \beta\}$$

$$\begin{aligned} \mathbb{E}[T_2 \mathbb{1}\{\mathcal{E}_1 \cap \mathcal{E}_2\}] &= \mathbb{E}[\beta \mathbb{1}\{\mathcal{E}_1 \cap \mathcal{E}_2\}] \\ &= \beta \cdot \mathbb{E}[\mathbb{1}\{\mathcal{E}_1 \cap \mathcal{E}_2\}] \\ &\leq \beta \end{aligned}$$

$$T_2 \mathbb{1}\{\mathcal{E}_1 \cap \mathcal{E}_2\} = \beta \mathbb{1}\{\mathcal{E}_1 \cap \mathcal{E}_2\}$$

It is true that $P(\mathcal{E}_1 \cap \mathcal{E}_2) \geq 1 - \delta$, thus

$$\text{w.p.} \geq 1 - \delta \quad \Delta_2 T_2 \leq \Delta_2 \beta.$$

$$\leq 1$$

$$\begin{aligned} &= \mathbb{P}(\cup_{i=1}^n \{E_i^c\}) \\ &\leq \delta \end{aligned}$$

$$\leq \Delta_2 \mathfrak{S} + \Delta_2 T \delta$$

$$\leq 1 + 8\Delta_2^{-1} \sigma^2 \log(2n/\delta) + \Delta_2 T \delta$$

Idea: choose $\delta = 1/T$ ($\Delta_2 \leq 1$)

$$\hookrightarrow \leq 1 + 8\Delta_2^{-1} \sigma^2 \log(2nT) + 1$$

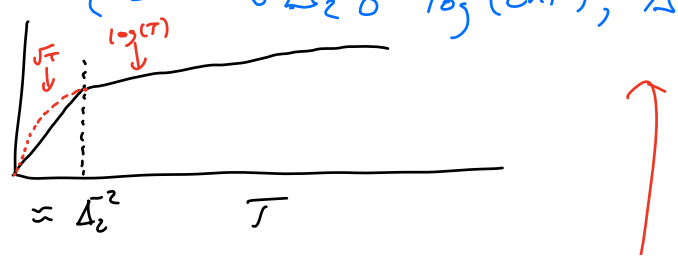
We wanted $R_T = o(T)$ - we have it! $\frac{\log(T)}{T} \rightarrow 0$

Note: If treatments were nearly identical so that $\Delta_2 = 10^{-10}$ then our regret bound says $R_T \leq 10^{10}$.

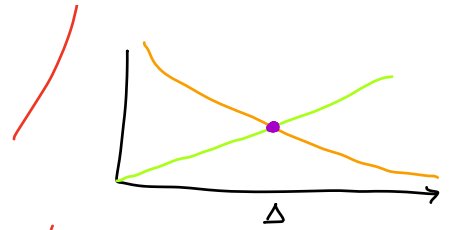
But this is nonsense, b/c it shouldn't matter what is prescribed, they act nearly the same.

We know: $R_T = \Delta_2 \mathbb{E}[T_2] \leq \Delta_2 T$

$$R_T \leq \min \left\{ 2 + 8\Delta_2^{-1} \sigma^2 \log(2nT), \Delta_2 T \right\}$$



maximize over Δ_2



$$\Rightarrow R_T \leq Z + \max_{\Delta \geq 0} \min \{ \underbrace{8\Delta^{-1} \sigma^2 \log(2nT)}_{\text{green}}, \underbrace{\Delta T}_{\text{blue}} \}$$

$$\Delta = \sqrt{\frac{8\sigma^2 \log(2nT)}{T}}$$

$$= Z + \sqrt{8\sigma^2 T \log(2nT)}$$

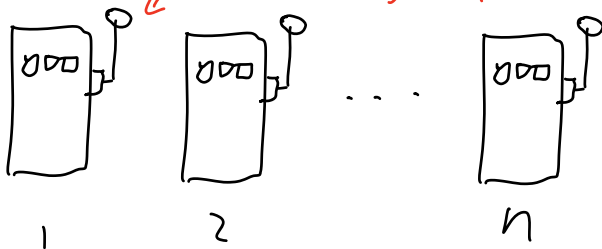
Note $R_T/T \rightarrow 0$

↑
"instance - indep" regret bound

Multi-armed Bandit problem.

Imagine Casino

← arm that you pull to play. "1-armed bandit"



Each machine i pays out random amount w/ mean θ_i

Elimination-style algorithm

Input: n arms $\mathcal{X} = \{1, \dots, n\}$, confidence level $\delta \in (0, 1)$.

Let $\hat{\mathcal{X}}_1 \leftarrow \mathcal{X}, l \leftarrow 1$

while $|\mathcal{X}_l| > 1$ **do**

$$\epsilon_l = 2^{-l}$$

Pull each arm in \mathcal{X}_l exactly $\tau_l = \lceil 2\epsilon_l^{-2} \log(\frac{4l^2|\mathcal{X}_l|}{\delta}) \rceil$ times

Compute the empirical mean of these rewards $\hat{\theta}_{i,l}$ for all $i \in \mathcal{X}_l$

$$\mathcal{X}_{l+1} \leftarrow \mathcal{X}_l \setminus \{i \in \mathcal{X}_l : \max_{j \in \mathcal{X}_l} \hat{\theta}_{j,l} - \hat{\theta}_{i,l} > 2\epsilon_l\}$$

$$l \leftarrow l + 1$$

Output: $\hat{\mathcal{X}}_{l+1}$ (or play the last arm forever in the regret setting)

using pulls only
from round l .

Assume observation $X_{i,t}$ at time t is ^{from arm i} 1-sub-Gaussian

and iid and $\mathbb{E}[X_{i,t}] \in [0, 1]$

Assume wlog $\theta_1 > \theta_2 \geq \dots \geq \theta_n$, $\Delta_i = \theta_1 - \theta_i$, $\Delta = \min_{i: \Delta_i > 0} \Delta_i$
 $= \Delta_2$

Lemma With probability at least $1 - \delta$,

arm $1 \in \mathcal{X}_l$ and $\max_{i \in \mathcal{X}_l} \Delta_i \leq 8\epsilon_l \quad \forall l \in \mathbb{N}$.

Proof For all $l \in \mathbb{N}$ and $i \in [n]$ define

$$E_{i,l} = \{ |\hat{\theta}_{i,l} - \theta_i| \leq \epsilon_l \}$$

Note that $\epsilon_l \geq \sqrt{\frac{2 \log(4l^2|\mathcal{X}_l|/\delta)}{5_l}}$. Thus

$$\mathbb{P}\left(\bigcup_{l=1}^{\infty} \bigcup_{i \in \mathcal{X}_l} E_{i,l}^c\right) \leq \mathbb{P}\left(\bigcup_{l=1}^{\infty} \bigcup_{i=1}^n E_{i,l}^c\right)$$

$$\leq \sum_{l=1}^{\infty} \sum_{i=1}^n \mathbb{P}(E_{i,l}^c)$$

$$\leq \sum_{l=1}^{\infty} \sum_{i=1}^n \frac{\delta}{2l^2|\mathcal{X}_l|} \leq \delta \quad (|\mathcal{X}_l| = n)$$

fact: $\sum_{l=1}^{\infty} \frac{1}{l^2} = \frac{\pi^2}{6} \leq 2$

$P\left(\bigcap_{l=1}^{\infty} \bigcap_{i=1}^n \{E_{i,l}\}\right) \geq 1 - \delta$. Thus, assume these events hold.

If arm l is kicked out of \mathcal{X}_l , then \exists arm $j \in \mathcal{X}_l$:

$$\hat{\theta}_{j,l} - \hat{\theta}_{l,l} > 2\varepsilon_l.$$

Assume $l \in \mathcal{X}_l$ and show $l \in \mathcal{X}_{l+1}$.

For all $j \in \mathcal{X}_l$ we have

$$\hat{\theta}_{j,l} - \hat{\theta}_{l,l} = (\hat{\theta}_{j,l} - \theta_j) + (\theta_j - \hat{\theta}_{l,l}) - \Delta_j$$

$$\leq \varepsilon_l + \varepsilon_l + 0$$

$$= 2\varepsilon_l.$$

$$\Rightarrow l \in \mathcal{X}_{l+1} \Rightarrow \text{arm } l \in \mathcal{X}_l \quad \forall l \in \mathbb{N}.$$

Fix any $i \in \mathcal{X}_l$ s.t. $\Delta_i > 4\varepsilon_l$. \star

Recall $l \in \mathcal{X}_l \quad \forall l \in \mathbb{N}$.

Want to show that arm i does not enter \mathcal{X}_{l+1} .

We need to show $\max_{j \in \mathcal{X}_l} \hat{\theta}_{j,l} - \hat{\theta}_{i,l} > 2\varepsilon_l$.

$$\begin{aligned}
\max_{j \in \mathcal{X}_\ell} \hat{\theta}_{j,\ell} - \hat{\theta}_{i,\ell} &\geq \hat{\theta}_{1,\ell} - \hat{\theta}_{i,\ell} \\
&= (\hat{\theta}_{1,\ell} - \theta_1) + (\theta_i - \hat{\theta}_{i,\ell}) + \Delta_i \\
&\stackrel{(\varepsilon_\ell, \eta \varepsilon_\ell)}{\geq} -\varepsilon_\ell - \varepsilon_\ell + \Delta_i \\
&\stackrel{(*)}{>} -2\varepsilon_\ell + 4\varepsilon_\ell \\
&\geq 2\varepsilon_\ell.
\end{aligned}$$

$$\Rightarrow \max_{j \in \mathcal{X}_\ell} \hat{\theta}_{j,\ell} - \hat{\theta}_{i,\ell} > 2\varepsilon_\ell \Rightarrow i \notin \mathcal{X}_{\ell+1}$$

$$\max_{j \in \mathcal{X}_{\ell+1}} \Delta_j \leq 4\varepsilon_\ell = 8\varepsilon_{\ell+1} \quad (\varepsilon_\ell = 2^{-\ell})$$

$$\Rightarrow \max_{j \in \mathcal{X}_\ell} \Delta_j \leq 8\varepsilon_\ell \quad \forall \ell \in \mathbb{N}. \quad \square$$

Sample complexity. Let T_i be the total pulls from arm i at the end of game.

$$\text{Total \# pulls} = \sum_{i=1}^n T_i.$$

Claim: $|\mathcal{X}_\ell| = 1$ for all ℓ : $8\epsilon_\ell \leq \Delta$

$$8\epsilon_\ell \leq \Delta \Leftrightarrow 8\Delta^{-1} \leq 2^\ell \Leftrightarrow \ell \geq \lceil \log_2(8\Delta^{-1}) \rceil$$

$$\epsilon_\ell = 2^{-\ell} \Rightarrow |\mathcal{X}_\ell| = 1 \text{ for } \ell \geq \lceil \log_2(8\Delta^{-1}) \rceil$$

For any i

$$T_i = \sum_{\ell=1}^{\lceil \log_2(8\Delta^{-1}) \rceil} \mathbb{1}_{\{i \in \mathcal{X}_\ell\}}$$

$$\leq \sum_{\ell=1}^{\lceil \log_2(8\Delta^{-1}) \rceil} \mathbb{1}_{\{\Delta_i \leq 8\epsilon_\ell\}}$$

$$= \sum_{\ell=1}^{\lceil \log_2(8\Delta_i^{-1}) \rceil} 1$$

$$\lceil \log_2(a) \rceil \leq \log_2(a) + 1 = \log_2(2a)$$

$$= \sum_{\ell=1}^{\lceil \log_2(8\Delta_i^{-1}) \rceil} \lceil 2\epsilon_\ell^{-2} \log\left(\frac{4\ell^2 |\mathcal{X}|}{\delta}\right) \rceil$$

$$\leq \lceil 2 \log\left(\frac{4 \lceil \log_2(8\Delta_i^{-1}) \rceil^2 |\mathcal{X}|}{\delta}\right) \rceil \sum_{\ell=1}^{\lceil \log_2(8\Delta_i^{-1}) \rceil} \epsilon_\ell^{-2}$$

$$\leq 2 \log\left(\frac{8 \log_2^2(16\Delta_i^{-1}) |\mathcal{X}|}{\delta}\right) \underbrace{\sum_{\ell=1}^{\lceil \log_2(8\Delta_i^{-1}) \rceil} 4^\ell}_{C \Delta_i^{-2}}$$

$$\leq c' \Delta_i^{-2} \log\left(\frac{\log_2(\Delta_i^{-1}) |\chi|}{\delta}\right)$$

$$\lceil \log_2(8\Delta_i^{-1}) \rceil \leq 1 + \log_2(8\Delta_i^{-1}) = \log_2(16\Delta_i^{-1})$$

$$\sum_{e=0}^N a^e = \frac{1-a^{N+1}}{1-a}$$

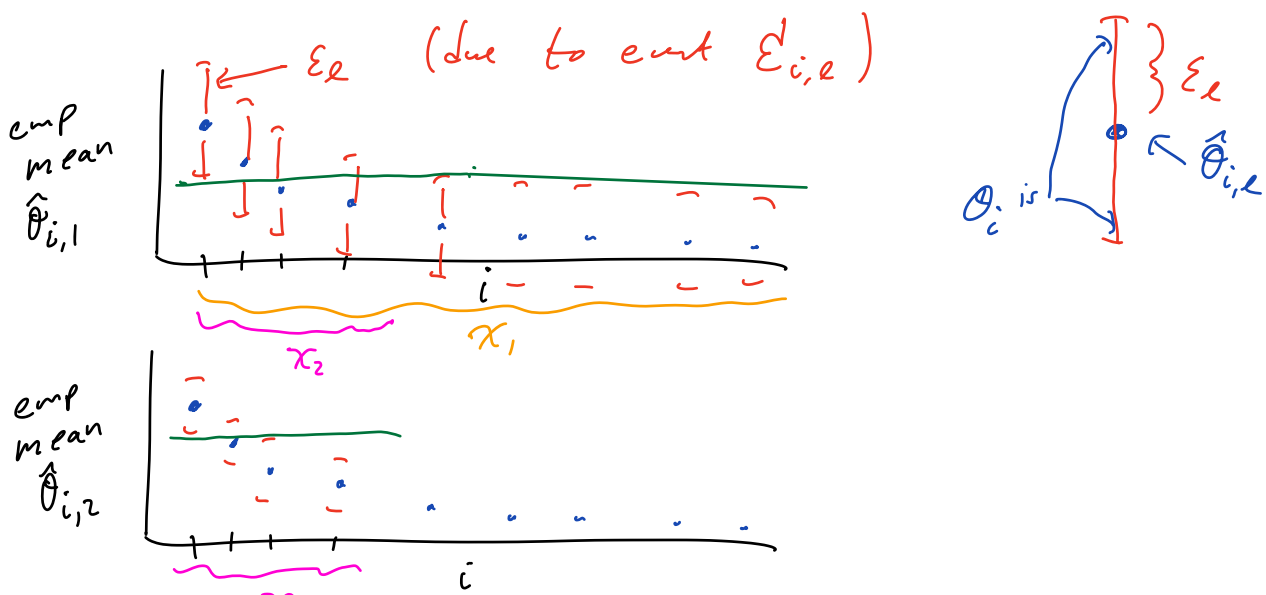
For any $c > 0$, $\log(a) \geq 1 \quad \exists c' > 0$
 s.t. $\log(ca) \leq c' \log(a)$

$$T_1 = \max_{i > 1} T_i \quad \sum_{i=1}^n T_i = T_1 + \sum_{i=2}^n T_i \leq 2 \sum_{i=2}^n T_i$$

Thm) W.p. $\geq 1-\delta$ after at most

$$c'' \sum_{i=2}^n \Delta_i^{-2} \log\left(\frac{\log_2(\Delta_i^{-1}) |\chi|}{\delta}\right) \text{ total pulls we}$$

have that alg exits w/ best arm 1.



x_2
 x_3

$$E_{i,t} = \left\{ |\hat{\theta}_{i,t} - \theta_i| \leq \sqrt{\frac{2 \log(4t^2/\delta)}{t}} \right\}$$

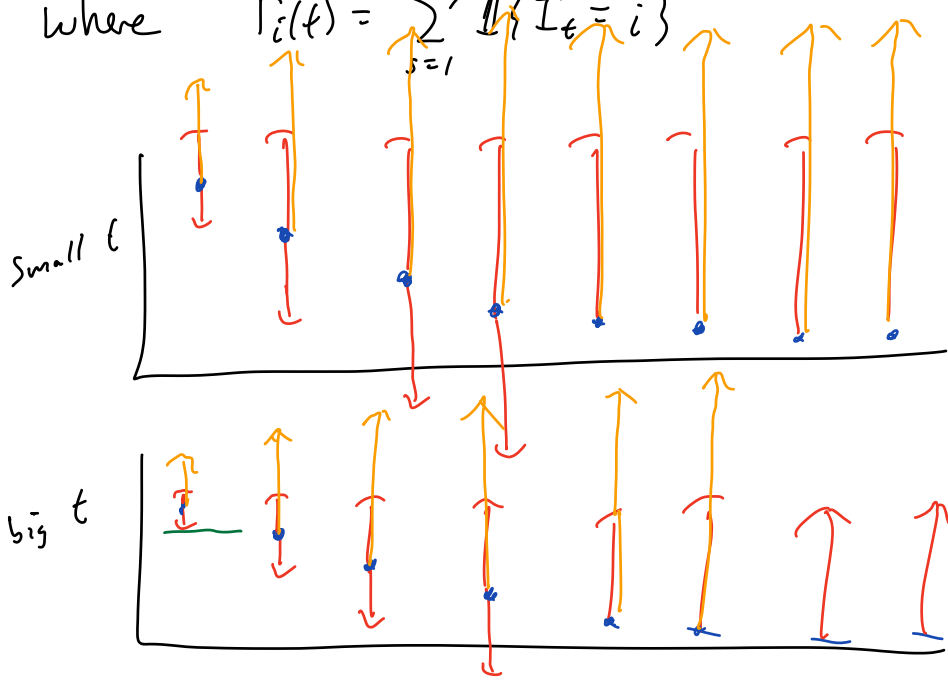
Claim: $P\left(\bigcap_{t=1}^{\infty} \bigcap_{i=1}^n E_{i,t}\right) \geq 1 - \delta.$

Consider UCB (upper confidence bound alg)

Pull every arm once
for $t = n+1, n+2, \dots$

pull $I_t = \operatorname{arg\,max}_{i \in [n]} \hat{\theta}_{i, T_i(t)} + 2 \sqrt{\frac{2 \log(4 T_i(t)^2 / \delta)}{T_i(t)}}$

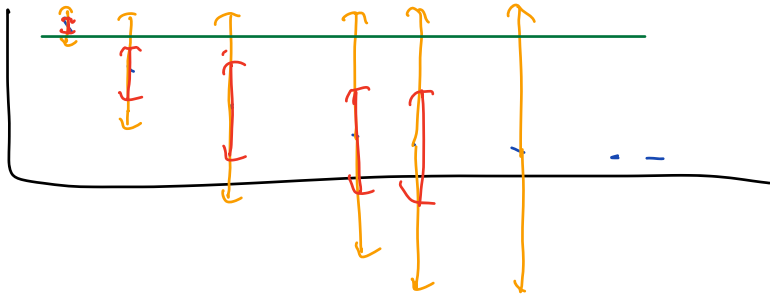
where $T_i(t) = \sum_{s=1}^{t-1} \mathbb{1}\{I_s = i\}$ $\geq \theta_i$



On your homework, you show that $\bar{T}_i(t) \leq c \Delta_i^{-2} \log\left(\frac{T}{\delta}\right)$

$\forall t \in [T]$

$$R_T = \sum_{i=2}^n \Delta_i \mathbb{E}[T_i]$$



We've demonstrated the existence of an algorithm that satisfies $\mathbb{E}[T_i(t)] \leq c \Delta_i^{-2} \log(T) \Rightarrow R_T \leq c \sum_{i=2}^n \Delta_i^{-1} \log(T)$

?

↓

(Lai-Robbins '83)

Thm] Any strategy that satisfies $\mathbb{E}[T_i(t)] = o(t^a)$ for $a > 0$

w/ $\Delta_i > 0$, we have $\lim_{T \rightarrow \infty} \frac{R_T}{\log(T)} \geq \sum_{i=2}^n \frac{2}{\Delta_i}$

\Rightarrow Opt algorithm achieves regret $R_T \leq \sum_{i=2}^n \frac{2}{\Delta_i} \log(T) + c$

$$I_t = \arg \max_{i \in [n]} \hat{\theta}_{i, T_i(t)} + \sqrt{\frac{2\alpha \log(t)}{T_i(t)}}$$

Then A_α $T \rightarrow \infty$, $\alpha \downarrow 1$ UCB- α

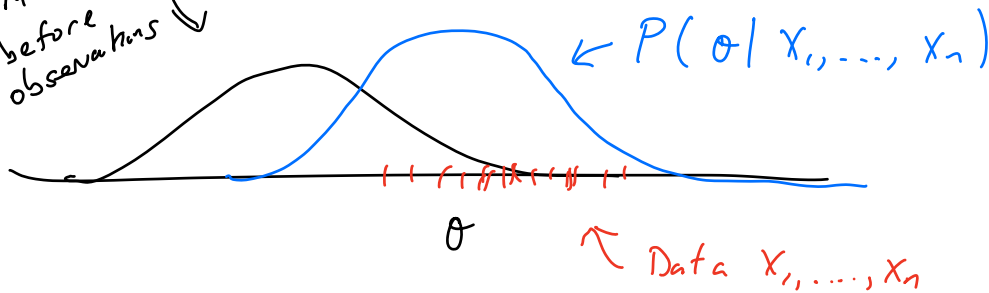
satisfies
$$\lim_{T \rightarrow \infty} \frac{R_T}{\log(T)} \leq \sum_{i=2}^n \frac{2}{\Delta_i}$$

+ Thompson Sampling

Thompson Sampling (1930's Thompson).

$\theta \sim P(\theta)$ Observe $X_1, \dots, X_n \sim \mathcal{N}(\theta, 1)$

Prior belief before observations



$$P(\theta | \{x_i\}_{i=1}^n) \propto P(\{x_i\}_{i=1}^n | \theta) P(\theta)$$

Given posterior distribution for i th arm $P(\theta_i | \{X_{i,1}, \dots, X_{i,T_i(t)}\})$

I can answer the question $P(i = \arg \max_j \theta_j | \{\{X_{i,t}\}_{t=1}^{T_i(t)}\}_{i=1}^n)$

Thompson Sampling

for $t=1, 2, \dots$

Pull arm $I_t = i$ w/ prob $P(i = \arg \max_j \theta_j | \{\{X_{i,t}\}_{t=1}^{T_i(t)}\}_{i=1}^n)$



↙ equivalent

$$\equiv \theta'_t \sim P(\theta \mid \left\{ \left\{ X_{i,t} \right\}_{t=1}^{T_i(t)} \right\}_{i=1}^n)$$

$$I_t = \underset{i \in \{n\}}{\operatorname{argmax}} [\theta'_t]_i$$

$$P(i = \underset{j}{\operatorname{argmax}} \theta_j \mid \left\{ \left\{ X_{i,t} \right\}_{t=1}^{T_i(t)} \right\}_{i=1}^n) = \int \mathbb{1}_{\{i = \underset{j}{\operatorname{argmax}} \theta_j\}} dP(\theta \mid \left\{ \left\{ X_{i,t} \right\}_{t=1}^{T_i(t)} \right\}_{i=1}^n)$$

In MAB, typically no sharing of information:

$$P(\theta) = \prod_{i=1}^n P_i(\theta_i)$$