

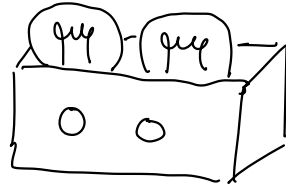
Lecture 3

Oct 1, 2025



Deutsch-Jozsa game

Say there is a box



two lights & two buttons.

Press a button and one of the two lights turns on
(deterministically).

How many (button presses) are required to decide if
actions of the two buttons are same/different?

Classically 2. Also learn which light.

Quantumly 1. Don't know which light.

Need to be able to query box in superposition

Box: $f: \{0,1\} \rightarrow \{0,1\}$.

Want to know if $f(0) = f(1)$ or not.

Consider unitary U_f s.t. $\forall x, y \in \{0,1\}$

$$|x\rangle|y\rangle \xrightarrow{U_f} |x\rangle|y \oplus f(x)\rangle$$

U_f is defined everywhere by linearity.

Algorithm

① Start with $|0\rangle \otimes |1\rangle$.

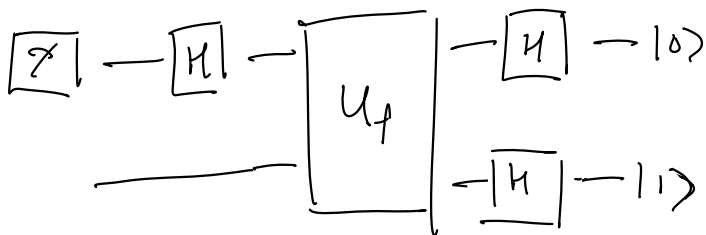
② Apply $H \otimes H$

$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

③ Apply U_f .

④ Apply H to 11

⑤ Measure 1st qubit.



$$|0\rangle|1\rangle \xrightarrow{H \otimes H} \frac{1}{2} \sum_{x \in \{0,1\}} |x\rangle \otimes (|0\rangle - |1\rangle)$$

$$\xrightarrow{U_f} \frac{1}{2} \sum_{x \in \{0,1\}} |x\rangle \otimes \underbrace{(|f(x)\rangle - |1 \oplus f(x)\rangle)}$$

$$\begin{cases} f(x)=0 \Rightarrow |0\rangle - |1\rangle \\ f(x)=1 \Rightarrow |1\rangle - |0\rangle \end{cases}$$

$$= \frac{1}{\sqrt{2}} \sum_{x \in \{0,1\}} (-1)^{f(x)} |x\rangle \otimes \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}} \right)$$

$$\frac{(-1)^{f(0)} |0\rangle + (-1)^{f(1)} |1\rangle}{\sqrt{2}}$$

$$= (-1)^{f(0)} \left(\frac{|0\rangle + (-1)^{f(0)+f(1)} |1\rangle}{\sqrt{2}} \right)$$

$$H \otimes I \mapsto (-1)^{f(0)} |f(0) + f(1)\rangle \otimes \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

Measurement outputs $f(0) + f(1)$ with certainty!

Today: The power of entanglement.

$|\psi\rangle$ a state on $(n_1 + n_2)$ - qubits.

$$\in (\mathbb{C}^2)^{\otimes n_1} \otimes (\mathbb{C}^2)^{\otimes n_2}.$$

Does $|\psi\rangle$ necessarily equal $|\alpha\rangle \otimes |\beta\rangle$?
 $\uparrow \qquad \qquad \qquad \uparrow$
 $(\mathbb{C}^2)^{\otimes n_1} \quad (\mathbb{C}^2)^{\otimes n_2}$

No!

$$\text{Ex. } |\text{EPR}\rangle = \frac{|0\rangle|0\rangle + |1\rangle|1\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \end{pmatrix}.$$

$$\text{Pf. } \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \otimes \begin{pmatrix} \gamma \\ \delta \end{pmatrix} = \begin{pmatrix} \alpha\gamma \\ \alpha\delta \\ \beta\gamma \\ \beta\delta \end{pmatrix}.$$

$$\text{If } |\text{EPR}\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \otimes \begin{pmatrix} \gamma \\ \delta \end{pmatrix}, \text{ then}$$

$$\alpha \gamma = \frac{1}{\sqrt{2}} \Rightarrow \alpha \neq 0, \gamma \neq 0.$$

$$\beta \delta = \frac{1}{\sqrt{2}} \Rightarrow \beta \neq 0, \delta \neq 0.$$

So then $\alpha \delta \neq 0, \beta \gamma \neq 0$, a contradiction.

A stronger intuition from counting states:

Recall from part 1 that for any $\epsilon > 0$,

we can find 2^{2^n} n -qubit phase states of which
 $\ni$ subset of size $2^{\Omega(\epsilon^2 2^n)}$ which are ϵ -orthogonal.

within a $(2n)$ qubit space there are $2^{\Omega(\epsilon^2 2^{2n})}$
 such states. How many states are in tensor product?

$$\left[\# \text{ of states of type } |\psi\rangle \otimes |\psi'\rangle \right] = 2^{2^n} \cdot 2^{2^n} = 2^{2^{n+1}}.$$

$$\begin{aligned} \frac{\# \text{ of tensor product states}}{\text{size of } \epsilon\text{-orthogonal set}} &= \frac{2^{2^{n+1}}}{2^{\Omega(\epsilon^2 2^{2n})}} = 2^{2^{n+1} - \Omega(\epsilon^2 2^{2n})} \\ &= 2^{-\Omega(2^n)} \text{ for } \epsilon < 2^{-n/10}. \end{aligned}$$

Def. A state $|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B$ is entangled if $|\psi\rangle$ cannot be expressed as a product state.

What is so special about an entangled state?

Case Study: The EPR state.

What are the only measurements possible?

In standard basis:

$$\frac{|00\rangle + |11\rangle}{\sqrt{2}} \Rightarrow \begin{array}{l} |00\rangle \text{ w pr } \frac{1}{2} \\ |11\rangle \text{ w pr } \frac{1}{2}. \end{array}$$

measure first qubit in $|+\rangle, |-\rangle$ basis.

$$\begin{aligned} \frac{|+\rangle|+\rangle + |-\rangle|-\rangle}{\sqrt{2}} &= \frac{(|0\rangle + |1\rangle)(|0\rangle + |1\rangle) + (|0\rangle - |1\rangle)(|0\rangle - |1\rangle)}{2\sqrt{2}} \\ &= \frac{2|0\rangle|0\rangle + 0|0\rangle|1\rangle + 0|1\rangle|0\rangle + 2|1\rangle|1\rangle}{2\sqrt{2}} \\ &= \frac{|0\rangle|0\rangle + |1\rangle|1\rangle}{\sqrt{2}} = |\text{EPR}\rangle \end{aligned}$$

So $\Pr[\text{measuring } |+\rangle] = \frac{1}{2}$ and resulting state $\propto$

$$(\langle + | \otimes \mathbb{1}) \frac{|+\rangle|+\rangle + |-\rangle|-\rangle}{\sqrt{2}} = \frac{|+\rangle}{\sqrt{2}} + 0 \Rightarrow \text{2nd qubit is then } |+\rangle.$$

Likewise, $\Pr[\text{measuring } |-\rangle] = \frac{1}{2}$ and 2nd qubit is then $|-\rangle$.

Remark EPR $\neq$ mixture over $|0,0\rangle$ and $|1,1\rangle$.

Pf. If we measure first qubit of $|0,0\rangle$ in $|+\rangle, |-\rangle$ basis then $\Pr(X = |+\rangle) = \frac{1}{2}$ but 2nd qubit state is $|0\rangle$.
and $\Pr(X = |-\rangle) = \frac{1}{2}$ but 2nd qubit state is $|0\rangle$.

If we measure first qubit of $|1,1\rangle$ in $|+\rangle, |-\rangle$ basis then $\Pr(X = |+\rangle) = \frac{1}{2}$ but 2nd qubit state is $|1\rangle$
and $\Pr(X = |-\rangle) = \frac{1}{2}$ but 2nd qubit state is $|1\rangle$.

Therefore, $\mathcal{X}$ is uniformly random but on outcome $|+\rangle$,
 2^{nd} qubit is a prob. mixture over $|0\rangle$ and $|1\rangle$ and
 not $|+\rangle$. So, we can distinguish $|EPR\rangle$ from $|0,0\rangle$ & $|1,1\rangle$
 mixture.

Alice & Bob and EPR



Then

Alice measures her qubit in basis $|b_0\rangle, |b_1\rangle$. Her measurement
 will be uniformly random $i \in \{0,1\}$, but Bob's state will be $|b_i\rangle$.

No, this does not allow faster than light communication.

Superdense coding.

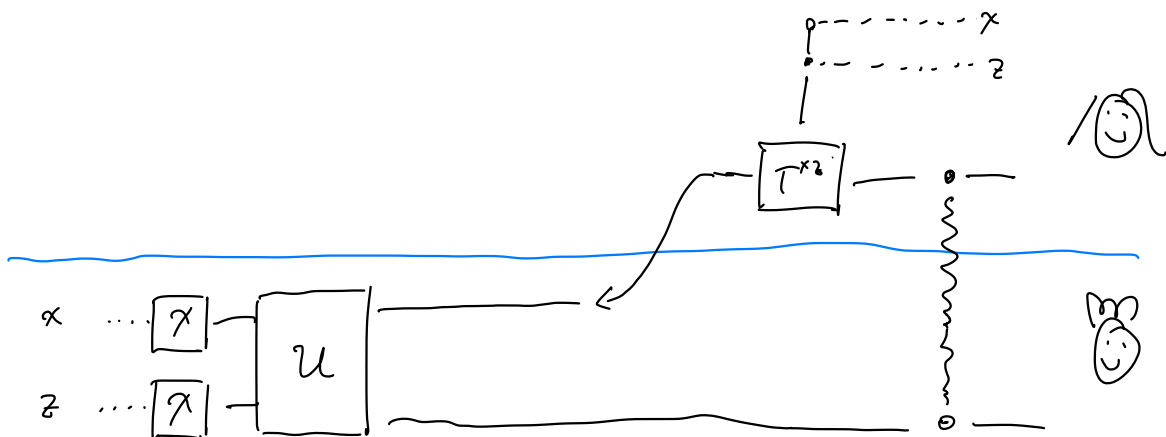
 wants to send messages to , who is leaving

for a trip far away. Can they save on postage using q.m.?

Classically m bits require m communication.

Quantum With some preparation, only $n/2$ bits of communication.

2 bits for the price of 1.

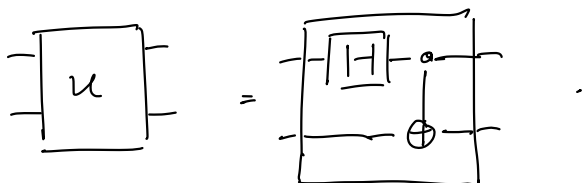


Claim $T^{xz} = X^x Z^z$.

$$|\Phi_{00}\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}, \quad |\Phi_{01}\rangle = \frac{|00\rangle - |11\rangle}{\sqrt{2}}$$

$$|\Phi_{10}\rangle = \frac{|01\rangle + |10\rangle}{\sqrt{2}}, \quad |\Phi_{11}\rangle = \frac{|01\rangle - |10\rangle}{\sqrt{2}}.$$

orthogonal states.



state of Alice & Bob is a 2-qubit state $\in \mathbb{C}^4$.

There are 4 orthogonal states $\in \mathbb{C}^4$.

Alice can switch the global state to any of the 4 of them.

Since they are orthogonal, there exists a distinguishing measurement.

Lesson $|EPR\rangle \neq \underbrace{|+\rangle \otimes |+\rangle}_{\text{unentangled.}}$

$|EPR\rangle \neq \begin{cases} \frac{1}{2} \text{ prob } |00\rangle \\ \frac{1}{2} \text{ prob } |11\rangle \end{cases}$ shared randomness.

Dealing with entanglement and randomness at the same time: density matrix.

For a state $|\psi\rangle$, corresponding density matrix is $|\psi\rangle\langle\psi|$.

For a prob mixture p_i of $|\psi_i\rangle$ and

p_i of $|\psi_i\rangle$, density matrix

$$\rho = p_0 |\psi_0\rangle\langle\psi_0| + p_1 |\psi_1\rangle\langle\psi_1|.$$

In general
$$\rho = \sum_a p_a |\psi_a\rangle\langle\psi_a|.$$

Properties:

① $\text{tr}(\rho) = 1.$

② $\rho = \rho^\dagger$ and $\rho \geq 0.$

Not.

If $\rho = |\psi\rangle\langle\psi|$ (rank 1)
called "pure state" else
"mixed state".

Strong converse: given every d.m. ρ , $\exists$ a prob. mixture $\{p_a\}$ and states $\{|\psi_a\rangle\}$ s.t.

$$\rho = \sum p_a |\psi_a\rangle\langle\psi_a|.$$

PA: eigendecomposition of ρ gives p_a and $|\psi_a\rangle$.

$\sum$ eigenvalues = 1 and positivity gives $p_a \geq 0$.

Examples

$$|0\rangle \Rightarrow \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|1\rangle \Rightarrow \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$$|+\rangle \Rightarrow \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \quad |-\rangle \Rightarrow \frac{1}{2} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$$

$$\text{conf. p.} \Rightarrow \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}.$$

$$\text{EPR} \Rightarrow \frac{1}{2} \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{pmatrix}$$

$$\text{shared randomness} \Rightarrow \frac{1}{2} \begin{pmatrix} 1 & & & \\ & 0 & & \\ & & 0 & \\ & & & 1 \end{pmatrix}.$$

$$|+\rangle \otimes |+\rangle \Rightarrow \frac{1}{4} (\text{all ones}).$$

The derived consequence of q. computation.

- ① n -qubit state is a density matrix in $\mathcal{L}(\mathbb{C}^2)^{\otimes n}$.
- ② update by unitary transformations

$$\rho \mapsto U \rho U^\dagger$$

③ measurement of a single qubit:

$$\rho \mapsto (|0\rangle\langle 0| \otimes \mathbb{1}) \rho (|0\rangle\langle 0| \otimes \mathbb{1}) + (|1\rangle\langle 1| \otimes \mathbb{1}) \rho (|1\rangle\langle 1| \otimes \mathbb{1})$$

$$\text{Pr}(\text{measuring } 0) = \text{tr}(|0\rangle\langle 0| \otimes \mathbb{1} \rho)$$

$$\text{Pr}(\text{measuring } 1) = \text{tr}(|1\rangle\langle 1| \otimes \mathbb{1} \rho).$$

[State conditioned on outcome 0] =

$$\frac{(|0\rangle\langle 0| \otimes \mathbb{1}) \rho (|0\rangle\langle 0| \otimes \mathbb{1})}{\text{tr}(|0\rangle\langle 0| \otimes \mathbb{1} \rho)}.$$

In class: verify that this state has trace 1.

④ If with probability p_a , first state is ρ_a and second state is σ_a , then overall state is

$$\sum p_a \cdot (\rho_a \otimes \sigma_a).$$

PF that density matrix computations follow from the axioms.

Consider mixture of states $\{|\psi_a\rangle\}$ with prob. p_a
s.t. $\sum_a p_a = 1$.

① The corresponding matrix is $\rho = \sum_a p_a |\psi_a\rangle\langle\psi_a|$.

① $\rho \geq 0$ b.c. sum of positive matrices.

② $\rho = \rho^\dagger$ b.c. sum of Herm. matrices.

$$\begin{aligned}\textcircled{3} \quad \text{tr}(\rho) &= \text{tr}\left(\sum_a p_a |\psi_a\rangle\langle\psi_a|\right) \\ &= \sum_a p_a \text{tr}\left(|\psi_a\rangle\langle\psi_a|\right) \\ &= \sum_a p_a \text{tr}\left(\langle\psi_a|\psi_a\rangle\right) \\ &= \sum_a p_a = 1.\end{aligned}$$

↙ linearity of tr
↙ cyclicity of tr

So ρ is a density matrix.